Lab 1: The Instruction Fetch Stage

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October 17, 2018

1 Introduction

The objective of this lab is to implement and test the Instruction Fetch (IF) pipeline stage of the MIPS five stage pipeline. There are five modules that needs to be implemented which are: the program counter, the incrementer, the MUX, the memory, and the instruction fetch ID latch. All of these modules will make up the Instruction Fetch stage.

2 Interface

The program counter module takes in a 32-bit number that represents the New Program Counter (NPC), updates the current PC to the NPC value, and outputs a 32-bit number to the rest of the dependent modules.

Table 1: pc mod Inputs

Name	Function
npc	The current 32-bit PC value

Table 2: pc_mod Outputs

Name	Function
pc_out	The updated 32-bit PC value

The incrementer module takes in a 32-bit constant 1 and a 32-bit PC value. Then the two inputs are added together, and outputs the result to another 32-bit wire.

Table 3: incrementer Inputs

Name	Function
pc_in	The 32-bit
	number that is
	the current PC

value	7	value
-------	---	-------

Table 4: incrementer Outputs

Name	Function
	The updated 32-bit number that is the result
	of incrementing pc_in by 1.

The MUX module takes in the next sequential PC value and the PC value that is calculated in the Execute stage depending on the value of the select line set in the Memory stage.

Table 5: mux Inputs

Name	Function
a	The 32-bit number that holds the next sequential PC value
b	The 32-bit number that holds the PC value that is calculated in the Execute stage
sel	The line that determines if inputs a or b will be selected to be the output

Table 6: mux Onputs

Name	Function
у	The 32-bit wire that holds the value that is specified via the select line

The instruction memory module takes in the current 32-bit PC value as its input and outputs the 32-bit instruction which that PC value is currently pointing at.

Table 7: mem Inputs

Name	Function
addr	The current 32-bit PC value

Table 8: mem Outputs

		- 110 - 0 0 1 - 1 - 1 - 1 - 1 - 1 - 1 -
Nan	ne	Function
data	ı	The 32-bit register which the PC value is currently pointing at

The Instruction Fetch-Decode latch (IF_ID_Latch) module takes in the next sequential PC value and address and outputs them to the next stage.

Table 9: IF ID Inputs

Name	Function
instruction_in	The 32-bit wire from the Fetch stage
npc_in	The 32-bit wire from the Fetch stage

Table 10: IF_ID outputs

Name	Function
instruction_out	The updated 32-bit PC value from instruction_in
npc_out	The updated 32-bit PC value from npc_in

The Instruction Fetch module for the Instruction Fetch stage is where all the individual modules are connected to each other. It takes in a PC value and the sequential PC value, and outputs the IF_ID latch instructions and IF_ID latch PC value.

Table 11: ifetch Inputs

Name	Function
EX_MEM_PCSrc	The PC value from the instruction fetch-decode latch stage
EX_MEM_NPC	The 32-bit sequential PC value from the instruction fetch-decode latch stage

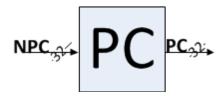
Table 12: ifetch Outputs

Name	Function
IF_ID_INSTR	The instructions from the instruction fetch-decode latch stage
IF_ID_NPC	The PC value from the instructions fetch-decode latch stage

3 Design

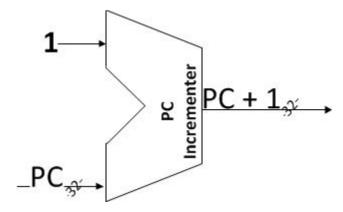
The design of the program counter takes in a 32-bit New Program Counter (NPC), updates the current PC value to the NPC value, and outputs the results to all the dependent modules.

Figure 1: Program Counter design



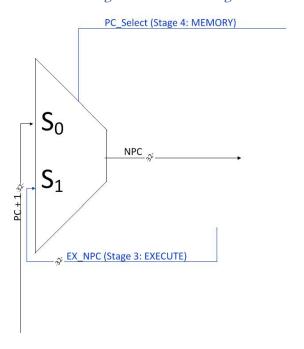
The design of the incrementer takes in the current PC value and outputs the PC value incremented by 1.

Figure 2: Incrementer design



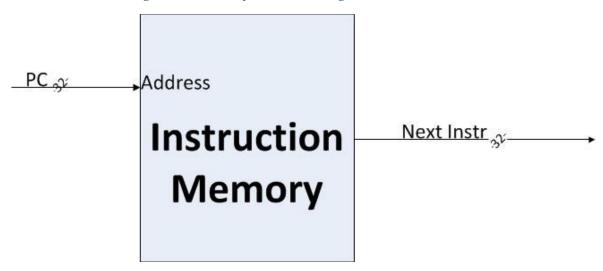
The design of the MUX module takes in a s0 and s1 wire in which a select line determines the next PC value depending if the line goes high (1) or low (1).

Figure 3: MUX design



The design for the instruction memory module takes in a 32-bit PC value, and outputs the 32-bit instruction depending on where the PC value is currently pointing at.

Figure 4: Memory module design



4 Implementation

Listing 1: Implementation for pc_mod

```
`timescale 1ns / 1ps

module pc_mod( output reg [31:0] pc_out, input wire [31:0] npc
);

initial
  begin
     pc_out <= 0;
  end

always @ ( npc )
  begin
     #1 pc_out <= npc;
  end

endmodule</pre>
```

Listing 2: Implementation for incrementer

```
`timescale 1ns / 1ps

module incrementer ( input wire [31:0] pc_in, output wire
[31:0] pc_out );

   assign pc_out = pc_in + 1; // Increment PC by 1
endmodule
```

Listing 3: Implementation for mem

```
module mem (
    output reg [31:0] data,
    input wire [31:0] addr
    // Regsiter Declarations
    reg [31:0] MEM[0:127];
    // Initialize Registers
    initial begin
       MEM[0] <= 'hA00000AA;</pre>
       MEM[1] <= 'h10000011;
       MEM[2] <= 'h20000022;
       MEM[3] <= 'h30000033;
       MEM[4] \le 'h40000044;
       MEM[5] <= 'h50000055;
       MEM[6] <= 'h60000066;
       MEM[7] <= 'h70000077;
       MEM[8] <= 'h80000088;
       MEM[9] <= 'h90000099;
     end
     always @ ( addr )
     begin
        data <= MEM[addr];</pre>
     end
endmodule
```

Listing 4: Implementation for if_id module

Listing 5: Implementation for the top-level Instruction Fetch module

```
module I FETCH (
    input EX MEM PCSrc,
    input wire [31:0] EX MEM NPC,
    output wire [31:0] IF ID INSTR, IF ID NPC
);
    // signals
    wire [31:0] PC;
    wire [31:0] data out;
    wire [31:0] npc, npc mux;
    // instantiations
    mux mux1 ( .y(npc mux),
               .a(EX MEM NPC),
               .b(npc),
               .sel(EX MEM PCSrc));
   pc_mod pc_mod1 ( .pc_out(PC),
                     .npc(npc mux));
    mem mem1 ( .data(data out),
               .addr(PC));
    if id if id1 ( .instructions out(IF ID INSTR),
                   .npc out(IF ID NPC),
                   .instructions in (data out),
                   .npc in(npc));
    incrementer incrementer1 ( .pc out(npc),
                                .pc in(PC));
endmodule
```

5 TestBench Design

This test bench is designed to simulate instruction fetch routines over a 20ns simulated run. The PC value starts at 0 initially and jumps to 1 after 9 cycles, and then jumps the PC value to 5 in which it outputs all the instruction fetch routines that were in the instruction memory module (Listing 3).

Listing 6: Instruction Fetch Test Bench

```
// Engineer:
//
// Create Date: 10/17/2018 12:15:19 PM
// Design Name:
// Module Name: pipeline
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
module pipeline();
   // Inputs
   reg EX_MEM_PCSrc;
   reg [31:0] EX MEM NPC;
   // Outputs
   wire [31:0] IF ID INSTR;
   wire [31:0] IF ID NPC;
   // Instantiate the Unit Under Test (UUT)
   I FETCH uut (
       .EX MEM PCSrc(EX MEM PCSrc),
       .EX_MEM_NPC(EX_MEM_NPC),
       .IF ID INSTR(IF ID INSTR),
       .IF_ID_NPC(IF_ID_NPC)
    );
    initial
    begin
       EX MEM NPC = 0;
       EX MEM PCSrc = 0;
       EX MEM PCSrc = 1;
       EX MEM NPC = 5;
       EX MEM PCSrc = 0;
    #10;
    $stop;
    end
endmodule
```

6 Simulation

Figure 5: Timing diagram from 0ns – 8ns

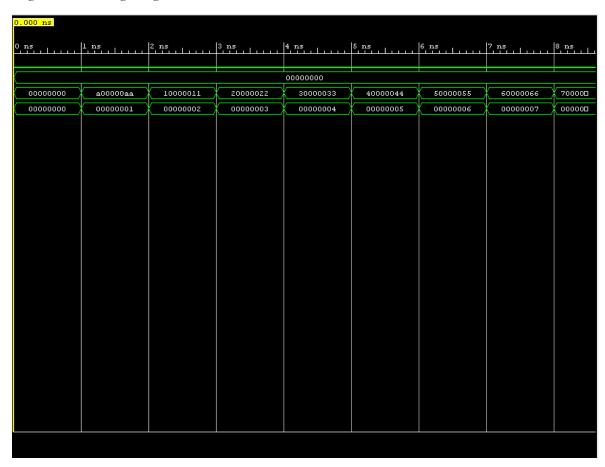


Figure 6: Timing diagram from 9ns – 17ns

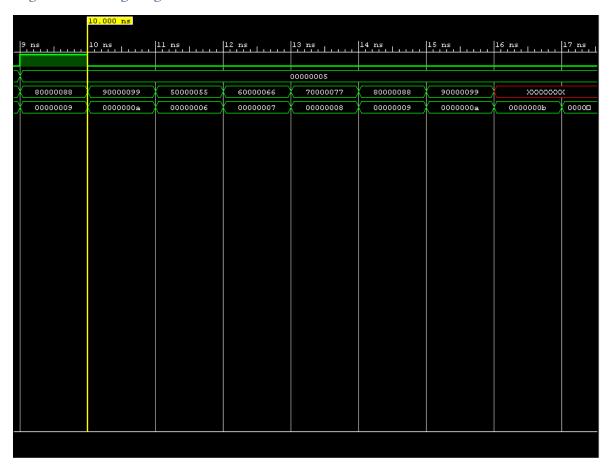


Figure 7: Timing diagram from 13ns - 20ns

7 Conclusions

The instruction fetch stage for this lab was successfully implemented. The main thing I learned in doing is how the instruction fetch-decode latch is designed and implemented in Verilog code. I never thought about using a top-level module that connects all the modules together into one since I just instantiate all the modules in the test bench. What I would do differently in this lab was to get more familiar with the VIvado software beforehand to save myself time to figure out how to add source files and debug the Verilog code.