This is a C code project for learning about Ethereum communication.

For example, consider the following documentation from the Ethereum project website.  $^{1}$ 

Alice wants to send an encrypted message that can be decrypted by Bob's static private key  $k_B$ . Alice knows about Bob's static public key  $K_B$ .

To encrypt the message m, Alice generates a random number r and corresponding elliptic curve public key R = r \* G and computes the shared secret  $S = P_x$  where  $(P_x, P_y) = r * K_B$ . She derives key material for encryption and authentication as  $k_E \parallel k_M = \text{KDF}(S, 32)$  as well as a random initialization vector iv. Alice sends the encrypted message  $R \parallel iv \parallel c \parallel d$  where  $c = \text{AES}(k_E, iv, m)$  and  $d = \text{MAC}(\text{sha256}(k_M), iv \parallel c)$  to Bob.

Let

```
r = {\sf ephemeral\_private\_key} 32 bytes R = {\sf ephemeral\_public\_key} 64 bytes S = {\sf shared\_secret} 32 bytes K_B = {\sf peer\_public\_key} 64 bytes k_E = {\sf aes\_key} 16 bytes k_M = {\sf hmac\_key} 32 bytes
```

Then this is the code for r and R = r \* G.

```
ec_genkey(ephemeral_private_key, ephemeral_public_key);
```

This is the code for  $S = P_x$  where  $(P_x, P_y) = r * K_B$ .

ec\_ecdh(shared\_secret, ephemeral\_private\_key, peer\_public\_key);

And this is the code for  $k_E \parallel k_M = \text{KDF}(S, 32)$ .

kdf(aes\_key, hmac\_key, shared\_secret);

<sup>1</sup>https://github.com/ethereum/devp2p/blob/master/rlpx.md

## Review of shared secrets

Let  $k_a$  and  $k_b$  be private keys and let  $K_a$  and  $K_b$  be public keys such that

$$K_a = k_a G, \quad K_b = k_b G$$

where G is the generator point.

It follows that

$$\frac{K_a}{k_a} = G, \quad \frac{K_b}{k_b} = G$$

Hence

$$\frac{K_a}{k_a} = \frac{K_b}{k_b}$$

and

$$k_b K_a = k_a K_b$$

After A and B exchange public keys, both can compute shared secret S.

$$k_a K_b = S, \quad k_b K_a = S$$

Since S is a point, by convention  $S_x$  is used for the actual shared secret.