

Integer Value Optimization

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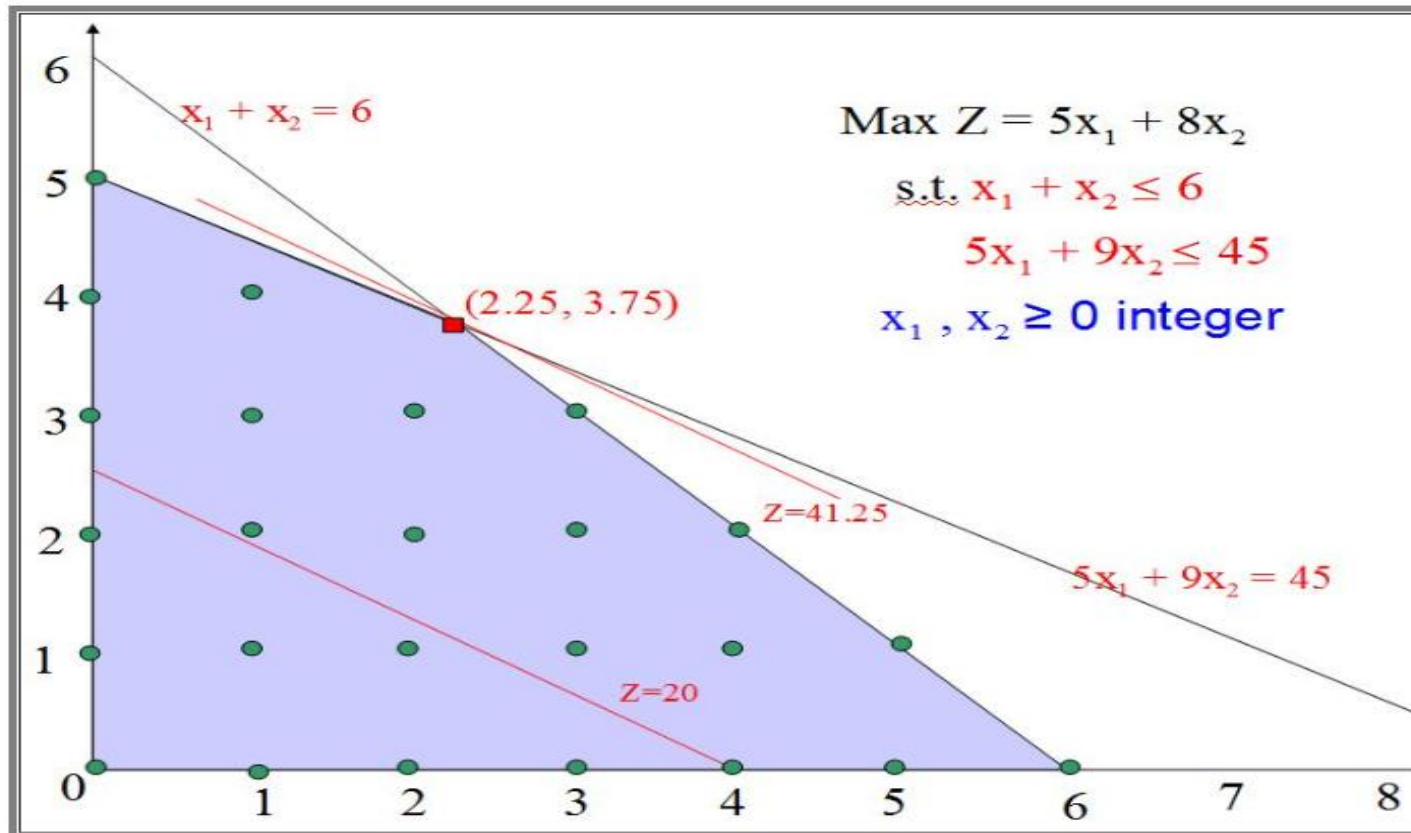
MND 404

Systems

◦ Example:

Linear Program	Corresponding Integer Program
$\text{Max } Z = 5x_1 + 8x_2$ $\text{s.t. } x_1 + x_2 \leq 6$ $5x_1 + 9x_2 \leq 45$	$\text{Max } Z = 5x_1 + 8x_2$ $\text{s.t. } x_1 + x_2 \leq 6$ $5x_1 + 9x_2 \leq 45$ $x_1, x_2 \geq 0 \text{ integer}$

The feasible regions of the **Linear Program (LP)** and **Integer Program (IP)** are shown in the following figure:



Explanation:

- The **feasible region** defined by the **constraints** of the **Linear Program** is the *entire shaded region*
- The **integer program** imposes *more constraints*

Therefore:

- A *subset* of the **feasible region** of the LP will satisfy the **constraints** of the **integer program**

Specifically:

- The *"integral" points* inside the **feasible (shaded) region** satisfy the **integer value constraints**

Feasible "region" of the **Integer Program**:

- *All* the **dots** inside the *shaded region* form the **feasible region** for the **Integer Program**.

Optimal solution of IP \leq Optimal solution of LP

- Relationship between a LP and its derived IP

- Fact 1:

- Optimal solution of IP \leq Optimal solution of LP

Reason:

- We know that:

- A solution x_1, x_2, \dots, x_n that satisfies the constraints of the **IP problem** will *also* satisfy the constraints of the **LP problem**

- Therefore:

- The **optimal solution** $x_1^{opt}, x_2^{opt}, \dots, x_n^{opt}$ of the **IP problem** is a **solution** of the **LP problem**

- In other words:

$$\text{OptimalSolution(IP)} \leq \text{OptimalSolution (LP)}$$

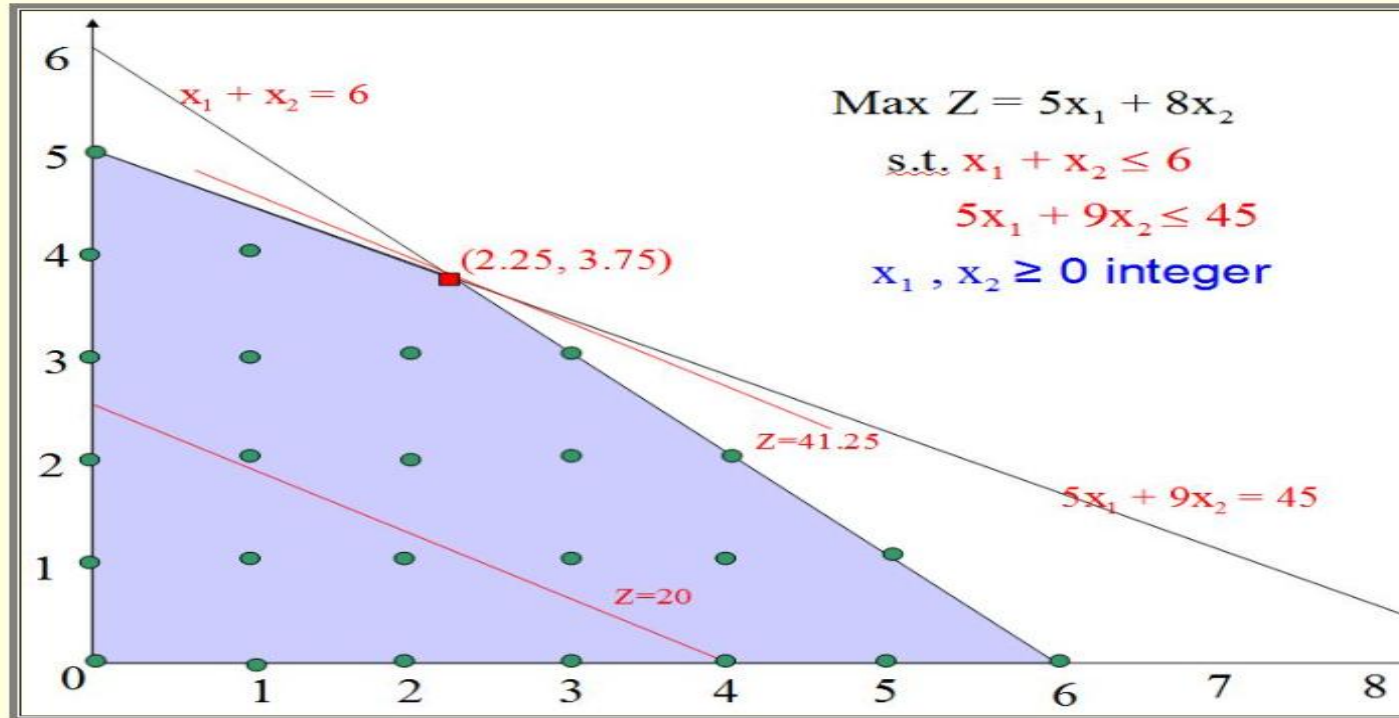
Optimal Solution of IP may not be equal to the Optimal Solution of LP

Fact 2:

- Optimal solution of LP may *not* be equal to the optimal solution of the derived IP

Example:

Consider:



The optimal solution of the Linear Program (LP) is:

$x_1 = 2.25$ (not integer)
 $x_2 = 3.75$ (not integer)

This solution is *not feasible* for the IP problem

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• **IF** the optimal solution of LP is *integral* (= all values are integers), **then** this solution is *also* the optimal solution of the derived IP

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• **IF** the Linear Program LP is *infeasible* (i.e., the *feasible region* is empty), then the Integer Program IP is *also infeasible*

- **Complexity of Integer Programming**

- **Integer Program** is **hard** to solve:

- The **Integer Program** problem is a **known NP-hard problem**

Wikipedia link: [click here](#)

- "**NP-hard**" problems (and the related "**NP-complete**" problems) are problems where:

- There are **no known $O(n^k)$** algorithm for solving the problem

(Algorithm to solve these problems have **exponential order** of **running time** (i.e., $O(2^n)$)

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- **Practice:**

- **Integer Programs** are solved using **heuristics**

Class Problem

Consider the following optimization model where the decision variables are x_1 and x_2 and where α is a parameter of the problem that is fixed, *i.e.*, it is not a variable:

$$\begin{array}{ll} \max & x_1 + \alpha x_2 \\ \text{s.t.} & x_1 + 2x_2 \leq 5 \\ (\mathcal{P}) & x_2 \geq x_1 - 5 \\ & x_2 \geq 0. \end{array}$$

1. Is this optimization model linear or nonlinear, continuous or mixed integer?
2. Represent the feasible region of this problem.
3. For what value of α does (\mathcal{P}) have multiple optimal solutions? Explain.
4. For what value α does (\mathcal{P}) have a single optimal solution of value 5? Explain.
5. For each one of the following constraints, describe if the optimization model obtained by adding this constraint to (\mathcal{P}) is (i) linear or nonlinear, (ii) feasible or infeasible:
 - (a) $x_1^2 + x_2^2 \leq 1$.
 - (b) $x_1 \geq 3$.
 - (c) $x_1 + x_2^2 \geq 4$.
6. If the objective is changed from maximization to minimization, is the optimization model (\mathcal{P}) unbounded? Why?