This work is in a very early stage. The Goal of this Document is

1. combine the collective works from:
 The Hitchhikers to GEOS
 by Berkely Softworks 1988

Geos Programmer's Reference Guide by
by Alexander Donald Boyce 1986
Revised by Bo Zimmerman 1997

The Official GEOS Programmers Reference Guide by Berkeley Softworks 1987

Information now available from the Dissembled GEOS Kernal Information Obtained from my Disassembly of GEOS Applications. If possible. All new KERNAL additions from Wheels 4.4 and from MP3+ MP3 is still being actively developed and it is open source so this should not be a problem. Wheels 4.4 internals has been dissembled so this should be manageable as well but I have not yet started on either of these additions. Any and all other relevant information from others sources including from my own experience developing for GEOS. With appropriate credits given for all source Document s.

- 2. Make all information in the document Hyperlinked.
- 3. Remove all Apple Information from this collection. When I get geoAssembler ported into the Apple GEOS world There will be another document made from this one with all the apple information in it.

This will provide a one stop resource for programming information.

The slowest part of the process is converting the OCR Scanned originals back into editable document form that can be easily edited.

Volunteers are welcome to assist...

My goal is too add 1 page a day to this document until it is completed. Paul B Murdaugh

Writer of Dual Top and the Landmark Series for GEOS paulbmurdaugh@gmail.com

			Status
Chapter	1.	GEOS Kernal	(Growing)
Chantan	2	E	/ <del></del>

(growing)

(Almost ready to merge in)

Chapter 2. Examples
Chapter 3. Memory Map
Chapter 4. 128 BackRAM Memory Map
Chapter 5. REU Bank 0 Memory Map
Appendix I Disk Errors (planned) (planned) (done)

## GEOS Kernal by Name

Name		Description	Category	Page
AllocateBlock	9048	Mark a disk block as in-use.	disk mid-level	31
AppendRecord	C289	Insert a new VLIR record after the current record.	disk VLIR	70
BBMult	C160	Byte by byte (single-precision) unsigned multiply.	math	103
Bell	N/A	Play a bell sound	utility	147
BldGDirEntry	C1F3	Build a GEOS directory entry in memory.	disk mid-level	32
BlkAllOC	C1FC	Allocate sectors for a file.	disk mid-level	33
BlockProcess	C10C	Block process from running. Does not freeze timer.	process	138
Bmult	C163	Byte by word unsigned multiply.	math	104
BootGeos	C000	Reboot GEOS	internal	99
CalcBlksFree	C1DB	Calculate total number of free disk blocks.	disk mid-level	34
CallRoutine	C1D8	pseudo-subroutine call. \$0000 aborts call.	utility	148
ChangeDiskDevice	C2BC	Change disk drive device number.	disk very low-level	17
ChkDkGEOS	C1DE	Check if a disk is GEOS format.	disk mid-level	35
ClearRam	C178	Clear memory to \$00.	memory	115
CloseRecordFile	C277	Close/Save currently open VLIR file.	disk VLIR	71
CmpFString	C26E	Compare two fixed-length strings.	memory	116
CmpString	C26B	Compare two null-terminated strings.	memory	117
CopyFString	C268	Copy a fixed-length string.	memory	118
CopyString	C265	Copy a null-terminated string.	memory	119
CRC	C20E	Cyclic Redundancy Check calculation.	utility	149
Dabs	C16F	Double-precision signed absolute value.	memory	105
DeleteFile	C238	Delete file.	disk high-level	56
DeleteRecord	C283	Delete current VLIR record.	disk VLIR	72
Ddec	C175	Double-precision unsigned decrement.	math	106
Ddiv	C169		math	107
DMult	C166	Double-precision unsigned multiply.	math	109
Dnegate	C172	Double-precision signed negation.	math	110
DoBOp	C2EC	1 ' '	memory	120
DoDlgBox	C256	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	dialog box	14
DoIcons	C15A		icon/menu	92
DoInlineReturn	C2A4	Return from inline subroutine.	utility	150
DoRAMOp	C2D4	RAM-expansion unit access primitive.	memory	127
DoneWithIO	C25F	Restore system after serial I/O.	disk very low-level	18
DrawPoint	C133	Draw, clear, or recover a single screen point.	graphics	89

DSDiv	C16C	Double-precision signed division.	math	111
DShiftLeft	C15D	Double-precision left shift (zeros shifted in).	math	112
DShiftRight	C262	Double-precision right shift (zeros shifted in).	math	113
EnableProcess	C109	Make a process runnable immediately.	process	141
EnterDeskTop	C22C	Leave application and return to GEOS deskTop.	disk high-level	57
EnterTurbo	C214	Activate disk turbo on current drive.	disk very low-level	19
ExitTurbo	C232	Deactivate disk turbo on current drive.	disk very low-level	20
FastDelFile	C244	Quick file delete (requires full track/sector list).	disk mid-level	36
FetchRAM	C2CB	Transfer data from RAM-expansion unit.	memory	128
FillRam	С17В	Fill memory with a particular byte.	memory	121
FindBAMBit	C2AD	Get allocation status of particular disk block.	disk mid-level	37
FindFile	C20B	Search for a particular file.	disk high-level	58
FindFTypes	C23B	Find all files of a particular GEOS type.	disk high-level	59
FirstInit	C271	Initialize GEOS variables.	internal	100
FollowChain	C205	Follow chain of sectors, building track/sector table.	disk mid-level	38
FreeBlock	C2B9	Mark a disk block as not-in-use in BAM.	disk mid-level	39
FreeFile	C226	Free all blocks associated with a file.	disk mid-level	40
FreezeProcess	C112	Pause a process countdown timer.	process	139
GetlstDirEntry	9030	Loads in the first directory block.	disk mid-level	40
GetBlock	C1E4	Read single disk block into memory.	disk low-level	28
GetCharWidth	C1C9	Calculate width of char without style attributes.	text	154
GetDirHead	C247	Read directory header into memory.	disk mid-level	43
GetFile	C208	Load GEOS file.	disk high-level	61
GetFHdrInfo	C229	Read a GEOS file header into fileHeader.	disk mid-level	44
GetFreeDirBlk	C1F6	Find an empty directory slot.	disk mid-level	45
GetNextChar	C2A7	Get next character from keyboard queue.	text	155
GetNxtDirEntry	9033	Get directory entry other than first.	disk mid-level	42
GetOffPageTrSc	9036	Get track and sector of off-page directory.	disk mid-level	47
GetRandom	C187	Calculate new random number.	utility	151
GetRealSize	C1B1	Calculate actual character size with attributes.	text	156
GetSerialNumber	C196	Return GEOS serial number.	internal	101
GetString	C1BA	Get string input from user.	text	157
i_FillRam	C1B4	Inline FillRam.	memory	121
i_MoveData	C1B7	Inline MoveData.	memory	124
InitForIO	C25C	Prepare system for serial I/O.	disk very low-level	21
InitMouse	FE80	Initialize input device.	input driver	96
InitProcesses	C103	Initialize processes.	process	140

InitRam	C181	Initialize memory areas from table.	memory	122
InsertRecord	C286	Insert new VLIR record in front of current record.	disk VLIR	73
LdApplic	C21D	Load GEOS application.	disk mid-level	48
LdDeskAcc	C217	Load GEOS desk accessory.	disk mid-level	50
LdFile	C211	Load GEOS data file.	disk mid-level	52
MoveBData	C2E3	128 BackRAM memory move routine.	memory	123
MoveData	C17E	Intelligent memory block move.	memory	124
NewDisk	C1E1	Initialize a drive.	disk mid-level	53
NextRecord	C27A	Make next VLIR the current record.	disk VLIR	74
OpenDisk	C2A1	Open disk in current drive.	disk high-level	63
OpenRecordFile	C274	Open VLIR file on current disk.	disk VLIR	75
PointRecord	C280	Make specific VLIR record the current record.	disk VLIR	76
PreviousRecord	C27D	Make previous VLIR record the current record.	disk VLIR	77
PutBlock	C1E7	Write single disk block from memory.	disk low-level	29
PutDecimal	C184	Format and display an unsigned double-precision nbr.	text	160
ReadBlock	C21A	Get disk block primitive.	disk mid-level	23
ReadByte	C2B6	Read a File 1 byte at a time.	disk mid-level	54
ReadLink	904B	Read track/sector link.	disk mid-level	24
ReadRecord	C28C	Read current VLIR record into memory.	disk VLIR	78
RenameFile	C259	Rename GEOS disk file.	disk mid-level	64
RestartProcess	C106	Unblock, unfreeze, and restart process.	process	142
SaveFile	C1ED	Save Memory to create a GEOS file.	disk high-level	66
SetDevice	C2B0	Establish communication with a new serial device.	disk high-level	67
SetGEOSDisk	C1EA	Convert normal CBM disk into GEOS format disk.	disk high-level	68
SetMouse	FE89	Reset input device scanning circuitry.	input driver	97
SetMsePic	C2DA	Set and preshift new soft-sprite mouse picture.	mouse/sprite	133
SetNewMode	C2DD	Change GEOS 128 graphics mode (40/80 switch).	graphics	90
Sleep	C199	Put current subroutine to sleep for a specified time.	process	143
SmallPutChar	C202	Fast character print routine.	text	159
StartASCII	7912	Begin ASCII mode printing.	print driver	135
StashRAM	C2C8	Transfer memory to RAM-expansion unit.	memory	129
SwapBData	C2E6	128 memory swap between front/back ram.	memory	125
SwapRAM	C2CE	RAM-expansion unit memory swap.	memory	130
RstrFrmDialog	C2BF	Exits from a dialog box.	dialog box	15
TempHideMouse	C2D7	Hide soft-sprites before direct screen access.	mouse/sprite	134
ToBasic	C241	Pass Control to Commodore BASIC.	utility	152
UnblockProcess	C10F	Unblock a blocked process, allowing it to run again.	process	143

UnfreezeProcess	C115	Unpause a frozen process timer.	process	145
UpdateRecordFile	C295	Update currently open VLIR file without closing	disk VLIR	79
VerifyBData	C2E9	128 BackRAM verify.	memory	126
VerifyRAM	C2D1	RAM-expansion unit verify.	memory	131
VerWriteBlock	C223	Disk block verify primitive.	disk very low-level	25
WriteBlock	C220	Write disk block primitive.	disk very low-level	26
WriteRecord	C28F	Write current VLIR record to disk.	disk VLIR	80

categories

## GEOS Kernal by Category

dialog box			
		Display and begin interaction w/dialog box. Exits from a dialog box.	14 15
disk very Low level	1		
ChangeDiskDevice	C2BC	Change disk drive device number.	1
DoneWithIO	C25F	Restore system after serial I/O.	18
EnterTurbo	C214	Activate disk turbo on current drive.	1
ExitTurbo	C232	Deactivate disk turbo on current drive.	2
InitForIO	C25C	Prepare system for serial I/O.	2
PurgeTurbo	C235	Remove disk turbo from current drive.	2
ReadBlock	C21A	Get disk block primitive.	2
ReadLink	904B	Read track/sector link.	
VerWriteBlock WriteBlock	C223 C220	Disk block verify primitive. Write disk block primitive.	2
disk low level			
		Read single disk block into memory.	28
PutBlock	C1E7	Write single disk block from memory.	29
disk mid-level			
AllocateBlock	9048	Mark a disk block as in-use.	3
BldGDirEntry	C1F3	Build a GEOS directory entry in memory.	3
BlkAllOC	C1FC	Allocate sectors for a file.	3
CalcBlksFree	C1DB	Calculate total number of free disk blocks.	3
ChkDkGEOS	C1DE	Check if a disk is GEOS format.	3
FastDelFile	C244	Quick file delete (requires full track/sector list)	
FindBAMBit	C2AD	Get allocation status of particular disk block.	3
FollowChain	C205	Follow chain of sectors, building track/sector table	
FreeBlock	C2B9	Mark a disk block as not-in-use in <b>BAM</b> .	3
FreeFile	C226	Free all blocks associated with a file.	4
GetlstDirEntry	9030	Loads in the first directory block. Get directory entry other than first.	4
GetNxtDirEntry GetDirHead	9033 C247	Read track 18 sector 0.	4 4
GetFHdrInfo	C247	Read a GEOS file header into fileHeader.	4
GetFreeDirBlk	C1F6	Find an empty directory slot.	4
GetOffPageTrSc		Get track and sector of off-page directory.	4
LdApplic		Load GEOS application.	4
LdDeskAcc		Load GEOS desk accessory.	5
LdFile		Load GEOS data file.	5
NewDisk	C1E1	Initialize a drive.	5
NxtBlkAlloc	\$C24D	Allocate sectors for a file.	-31
PutDirHead			-30
ReadByte		Read a File 1 byte at a time.	5
ReadFile		2	-41
SetGDirEntry			-41
SetNextFree			-31
StartAppl WriteFile			-80 -40
disk high level			
DeleteFile		Delete file.	<b></b> -
			_
EnterDeskTop	C22C	Leave application and return to GEOS deskTop.	5

	Table of Contents	
FindFile FindFTypes GetFile GetPtrCurDkNm OpenDisk RenameFile RstrAppl SaveFile SetDevice SetGEOSDisk	C20B Search for a particular file. C23B Find all files of a particular GEOS type. \$C208 Load GEOS file. \$C298 Compute address of disk's name C2A1 Open disk in current drive. C259 GEOS disk file. \$C23E Load the SWAPFILE \$C1ED Save Memory to create a GEOS file. C2B0 Establish communication with a new serial device C1EA Convert normal CBM disk into GEOS format disk.	categories
disk VLIR		
CloseRecordFile DeleteRecord InsertRecord NextRecord OpenRecordFile PointRecord PreviousRecord ReadRecord	C289 Insert a new VLIR record after the current record C277 Close/Save currently open VLIR file. C283 Delete current VLIR record. C286 Insert new VLIR record in front of current record C27A Make next VLIR the current record. C274 Open VLIR file on current disk. C280 Make specific VLIR record the current record. C27D Make previous VLIR record the current record. C28C Read current VLIR record into memory. C295 Update currently open VLIR file without closing. C28F Write current VLIR record to disk.	71 72 73 74 75 76 77 78
icon/menu		
DoIcons DoMenu DoPreviousMenu GotoFirstMenu RecoverAllMenus RecoverMenu ReDoMenu InitMouse input driver	\$C154 Erase the current menu \$C193 Draw the current menu FE80	92 1-701 1-702 1-703 1-706 1-705 1-704
InitMouse SetMouse SetMousePic SlowMouse UpdateMouse internal	FE80 Initialize input device. FE89 Reset input device scanning circuitry. \$c2da Set and preshift new soft-sprite mouse pictur \$fe83 Reset mouse velocity variables. \$fe86 Update mouse variables from input device.	
BootGeos FirstInit GetSerialNumber InterruptMain MainLoop Panic ResetHandle graphics	C000 Reboot GEOS. C271 Initialize GEOS variables.	99 100 101 1-805 1-804 1-711
BitmapUp i_BitmapUp BitmapClip 8	\$C1AB Draw a click box \$C1AB Draw a click box / inline \$C2AA Draw a coded image	-621 -622 -623

		Table of Contents	
			categories
BitOtherClip	\$C2C5	Draw a coded image with user patches	-624
DisablSprite		Turn off a sprite	-628
DrawLine		Draw/Erase/Copy an arbitrary line	-615
DrawPoint		Draw, clear, or recover a single screen point.	89
DrawSprite		Copy a sprite data block	-625
EnablSprite		Turn on a sprite	-627
FrameRectangle		Draw an outline in a pattern	-604
i FrameRectangle		Draw a solid outline with inline data	-605
GetScanLine		Compute memory address of a row on the screen	-618
GraphicsString		Process a graphic command table	-601
			-602
i_GraphicsString		Process a graphic command table / inline	
HorizontalLine		Draw a horizontal line in a pattern	-616
InvertLine		Reverse video a horizontal line	-614
ImprintRectangle		Copy a box from screen 2 to screen 1	-610
i_ImprintRectangle		Copy a box from screen 2 to screen 1 / inline	-611
InvertRectangle		Reverse video a box	-612
PosSprite		Position a sprite	-626
RecoverLine		Copy a line from screen 2 to screen 1	-613
Rectangle		Fill a box with a pattern	-606
i_Rectangle	\$C19F	Fill a box with a pattern / inline	-607
RecoverRectangle	\$C12D	Copy a box from screen 1 to screen 2	-608
		Copy a box from screen 1 to screen 2 / inline	-609
SetNewMode		Change GEOS 128 graphics mode (40/80 switch).	90
SetPattern		Select a fill pattern	-603
TestPoint		Test the value of a pixel	-619
VerticalLine		Draw a vertical line in a pattern	-617
10101001110	10111	zian a verezear zine in a paecein	01.
math			
BBMult	C160	Byte by byte (single-precision) unsigned multipl	y. 103
Bmult	C160	Byte by word unsigned multiply.	104
Dabs	C165		105
		Double-precision signed absolute value.	
Ddec	C175	Double-precision unsigned decrement.	106
Ddiv	C169	Double-precision unsigned division.	107
DMult	C166	Double-precision unsigned multiply.	109
Dnegate	C172	Double-precision signed negation.	110
DSDiv	C16C	Double-precision signed division.	111
DShiftLeft	C15D	Double-precision left shift (zeros shifted in).	112
DShiftRight	C262	Double-precision right shift (zeros shifted in).	113
memory			
ClearRam		Clear memory to \$00.	115
CmpFString		Compare two fixed-length strings.	116
CmpString		Compare two null-terminated strings.	117
CopyFString	\$C268	Copy a fixed-length string.	118
CopyString	\$C265	Copy a null-terminated string.	119
DoBOp	\$C2EC	(128) Back-RAM memory primitive	120
DoRAMOp		RAM-expansion unit access primitive.	127
FetchRAM		Transfer data from RAM-expansion unit.	128
FillRam		Fill memory with a particular byte.	121
i FillRam		Inline FillRam.	121
InitRam		Initialize memory areas from table.	122
MoveBData		128 BackRAM memory move routine.	123
MoveData		Intelligent memory block move	124
i MoveData		Inline MoveData.	124
I_MOVeData StashRAM		Transfer memory to RAM-expansion unit.	129
SwapBData		128 memory swap between front/back ram.	125
_			
SwapRAM	γC∠CE	Swap memory with an REU memory block.	130
9			

	cate	egories
VerifyBData	\$C2E9 128 BackRAM verify.	126
VerifyRAM	\$C2D1 RAM-expansion unit verify.	131
- 4		
mouse/Sprite		
ClearMouseMode	\$C19C Reset the mouse	-815
HideOnlyMouse	\$C2F2 (128) Temporarily remove soft-sprite mouse pointer.	
IsMseInRegion	\$C2B3 Check if mouse is inside a window	-710
MouseUp	\$C18A Turn on the mouse	-813
MouseOff	\$C18D Turn off the mouse	-814
SetMsePic	C2DA Set and preshift new soft-sprite mouse picture.	133
StartMouseMode		-812
TempHideMouse	\$C2D7 Hide soft-sprites before direct screen access.	134
printer driver		
princer driver		
GetDimensions	\$790C	
InitForPrint	\$7900	
PrintASCI	\$790F	
PrintBuffer	\$7906	
SetNLO	\$7915	
StartASCII	7912 Begin ASCII mode printing.	135
StartPrint	\$7903	100
StopPrint	\$7909	
Scopiline	4,303	
process		
BlockProcess	C10C Block process from running. Does not freeze timer.	138
EnableProcess	C109 Make a process runnable immediately.	141
FreezeProcess	C112 Pause a process countdown timer.	139
InitProcesses	C103 Initialize processes.	140
RestartProcess	C106 Unblock, unfreeze, and restart process.	142
Sleep	C199 Put current routine to sleep for a specified time.	143
UnblockProcess	C10F Unblock a blocked process, allowing it to run again	. 143
UnfreezeProcess	C115 Unpause a frozen process timer	145
text		
GetCharWidth	\$C1C9 Calculate width of char without style attributes.	154
GetNextChar	\$C2A7 Get next character from keyboard queue.	155
GetRealSize	\$C1B1 Calculate actual character size with attributes.	156
GetString	C1BA Get string input from user.	157
InitTextPrompt	\$C1C0 Create the text cursor sprite	-208
LoadCharSet	\$C1CC Load and activate a new font	-212
PromptOff	\$C29E Turn off the text cursor	-210
PromptOn	\$C29B Turn on the text cursor	-209
PutChar	\$C145 Display a character	-201
PutDecimal	\$C184 Format and display an unsigned double-precision nbr	
PutString	\$C148 Display a text string	-204
i PutString	\$C1AE Display a text string with inline data	-205
UseSystemFont	\$C14B Select the BSW font	-211
SmallPutChar	\$C202 Fast character print routine.	159
-	<u>.</u>	

			categories
utility			_
CRC DoInlineReturn GetRandom	C1D8 C20E C2A4 C187	Play a bell sound pseudo-subroutine call. \$0000 aborts call. Cyclic Redundancy Check calculation. Return from inline subroutine. Calculate new random number. Pass Control to Commodore BASIC.	147 148 149 150 151
Wheels Kernal			
		Load New Kernal Group B Unload Kernal Group	
KG_REU	0		
GetRAMBam PutRAMBam AllocA11RAM AllocRAMBlock FreeRAMBlock GetRAMInfo RamBlkAlloc RemoveDrive SvRamDevice DelRamDevice RamDevinfo	\$500 \$500 \$500 \$501	3 6 9 C F 2 5 8 B	
KG_DEVICE 1 DevNumChange			
SwapDrives KG_DISK 2	\$500	3	
NSetGEOSDisk DBFormat FormatDisk DBEraseDisk EraseDisk	\$500 \$500 \$500 \$500 \$500	3 6 9	
KG_READFILE 3			
OReadFile	\$500	0	
KG_WRITEFILE	4		
OWriteFile	\$500	0	
KG_DIRECTORY	5		
ChgParType ChPartition ChSubdir ChDiskDirectory GetFEntries TopDirectory UpDirectory	\$500 \$500 \$500 \$500 \$500 \$500 \$501	3 6 9 C F	

DownDirector GoPartition ChPartOnly FindRamLink	сy	\$5015 \$5018 \$501E \$5027
KG_MKDIR	6	
MakeDirector MakeSysDir	J.A	\$5000 \$5003
KG_VALDISK	7	
ValDisk		\$5000
KG_CPYDISK	8	
CopyDisk TestCompatik	oility	\$5000 \$5003
KG_COPY	9	
CopyFile	-	\$5000
KG_DESKTOP	10	
NewDesktop OEnterDeskto InstallDrive FindDesktop FindAFile	-	\$5000 \$5003 \$5006 \$5009 \$500c
KG_ToBasic	11	
K <b>ToBasic</b>		\$5000

## Structures

disk

\_\_\_\_\_

## Directory Entry

## Constants

errors

\_\_\_\_\_

examples

disk

-----

 ${\tt CheckDiskSpace}$ 

dialog box

## Chapter 1 GEOS Kernal

dialog box			
DoDlgBox	C256	Display and begin interaction w/dialog box.	14
RstrFrmDialog	C2BF	Exits from a dialog box.	15

dialog box

**DoDlgBox:** (C64, C128) C256

Function: Initializes, displays, and begins interaction with a dialog box.

Parameters: r0 DIALOG - pointer to dialog box definition (word).

r5-r10 can be used to send parameters to a dialog box.

When using **DBGetFileS** 

r5 BUFFER Ptr to buffer to store returned filename.

r7L FILETPE GEOS file type to search for (byte). (NULL for all)
r10 PERMNAME GEOS file type to search for (byte). (NULL for all)

Wheels: When using DBGetFileS and bit 7 of r7L is set.

r5 FILTER Ptr to Filter Procedure. Called once for every file before adding to the list of files.

r7L FILETPE GEOS file type to search for (byte). (NULL for all) r10 PERMNAME GEOS file type to search for (byte). (NULL for all)

Returns: rOL return code: typically the number of the system icon clicked

on to exit.

Note: returns when dialog box exits through RstrFrmDialog.

Destroys: n/a

Description: DoDlgBox saves off the current state of the system, places GEOS in a near warm start state, displays the dialog box according to the definition table (whose address is passed in r0), and begins tracking the user's interaction with the dialog box. When the dialog box finishes, the original system state is restored, and control is returned to the

application.

Simple dialog boxes will typically contain a few lines of text and one or two system icons (such as **OK** and **CANCEL**). When the user clicks on one of these icons, the GEOS system icon routine exits the dialog box with an internal call to **RstrFrmDialog**, passing the number of the system icon selected in **sysDBData**. **RstrFrmDialog** restores the system state and copies **sysDBData** to **rOL**.

More complex dialog boxes will have application-defined icons and routines that get called. These routines, themselves, can choose to load a value into **sysDBData** and call **RstrFrmDialog**.

a varie into Diezzada ana carr incultanzazag.

Note: Part of the system context save within DoDlgBox saves the current stack

pointer. Dialog boxes cannot be nested.  ${\tt DoDlgBox}$  is not reentrant. That

is, a dialog box should never call DoDlgBox.

Structure: DIALOG

Example:

See also: RstrFrmDialog

RstrFrmDialog: (C64,C128) C2BF

Function: Exits from a dialog box, restoring the system to the state prior to the

call to DoDlgBox.

Parameters: none.

Returns: Returns to point where DoDlgBox was called. System context is restored.

rOL contains sysDBData return value.

Destroys: assume a, x, y, r0H-r15

Description: RstrFrmDialog allows a custom dialog box routine to exit from the a

dialog box.  ${\tt RstrFrmDialog}$  is typically called internally by the GEOS system icon dialog box routines. However, it may be called by any dialog

box routine to force an immediate exit.

RstrFrmDialog first restores the GEOS system state (context restore) and then calls indirectly through recoverVector to remove the dialog box rectangle from the screen. The routine in recoverVector is called with the r2-r4 loaded for a call to RecoverRectangle. By default recoverVector points to RecoverRectangle, which will automatically recover the foreground screen from the background buffer. However, if the application is using background buffer for data, it will need to intercept the recover by placing the address of its own recover routine in recoverVector. If there is no shadow on the dialog box, then recoverVector is only called through once with r2-r4 holding the coordinates of the dialog box rectangle. However, if the dialog box has a shadow, then recoverVector will be called through two times: first for the patterned shadow rectangle and second for the dialog box rectangle. The application may want to special-case these two recovers when recovering.

Note:

RstrFrmDialog restores the sp register to value it contained at the call to DoDlgBox just before returning. This allows RstrFrmDialog to be called with an arbitrary amount of data on top of the stack (as would be the case if called from within a subroutine). GEOS will restore the stack pointer properly.

Structure: DIALOG

Example

# disk very low level

## disk very low level

ChangeDiskDevice	C2BC	Change disk drive device number.	17
DoneWithIO	C25F	Restore system after serial I/O.	18
EnterTurbo	C214	Activate disk turbo on current drive.	19
ExitTurbo	C232	Deactivate disk turbo on current drive.	20
InitForIO	C25C	Prepare system for serial I/O.	21
PurgeTurbo	C235	Remove disk turbo from current drive.	21
ReadBlock	C21A	Get disk block primitive.	23
ReadLink	904B	Read track/sector link.	24
VerWriteBlock	C223	Disk block verify primitive.	25
WriteBlock	C220	Write disk block primitive.	2.6

disk very low level

ChangeDiskDevice: (C64, C128) C2BC

Function: Instruct a drive to change its serial device number.

Parameters: a NEWDEVNUM - new device number to give current drive (byte).

curDrive drive whose device number will change.

Uses: curDrive drive whose device number will change.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Alters: curDrive NEWDEVNUM

Destroys: a,y

 $\textbf{Description: Change Disk Device} \ \ \text{requests the turbo software to change the serial device}$ 

number of the current drive. Most applications have no need to call this routine, as it is in the realm of low-level disk utilities. **ChangeDiskDevice** is used primarily by the deskTop and Configure programs

to add, rearrange, and remove drives.

Be aware that changing the device number merely instructs the turbo software in the drive to monitor a different serial bus address. Many internal GEOS variables and disk drivers expect the original device

number to remain unchanged.

Note: If ChangeDiskDevice is used on a RAMdisk, curDrive and curDevice both

change. However, because of the nature of the RAMdisk driver, the RAMdisk

does not respond as this new device.

Example:

See also: SetDevice

disk very low level

DoneWithIO: (C64, C128) C25F

Function: Restore system after I/O across the serial bus.

Parameters: none.

Returns: nothing.

Destroys: a, y

Description: DoneWithIO restores the state of the system after a call to InitForIO.

It restores the interrupt status, turns sprite DMA back on, returns the 128 to its original clock speed, and switches out the ROM and I/O banks

if appropriate (only on C64).

Disk and printer routines access the serial bus between calls to

InitForIO and DoneWithIO.

Example: MyPutBlock

See also: InitForIO

disk very low level

**EnterTurbo:** (C64, C128) C214

Function: Activate disk drive turbo mode

Parameters: none.

Uses: curDrive currently active disk drive.

curType v1.3+: checks disk type because not all use turbo

software.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y

 $\textbf{Description: EnterTurbo} \ \ \text{activates the turbo software in the current drive.} \ \ \text{If the}$ 

turbo software has not yet been downloaded to the drive,  ${\bf EnterTurbo}$  will download it. The turbo software allows GEOS to perform high-speed serial

disk access.

**EnterTurbo** treats different drive types appropriately. A RAMdisk, for example, does not use turbo code so **EnterTurbo** will not attempt to

download the turbo software.

The very-low level Commodore GEOS read/write routines, such as ReadBlock, WriteBlock, VerWriteBlock, and ReadLink, expect the turbo software to be

active. Call **EnterTurbo** before calling one of these routines.

Example: MyPutBlock

See also: WriteBlock, ExitTurbo, PurgeTurbo.

disk very low level

**ExitTurbo:** (C64, C128) C232

Function: Activate disk drive turbo mode.

Parameters: none.

Uses: curDrive currently active disk drive.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y

Description: ExitTurbo deactivates the turbo software in the current drive so that

the serial bus may access another device. SetDevice automatically calls

this before changing devices.

Note: If the turbo software has not been downloaded or is already inactive,

ExitTurbo will do nothing.

Example:

See also: EnterTurbo, PurgeTurbo.

disk very low level

InitForIO: (C64, C128) c25c

Function: Prepare for I/O across the serial bus

Parameters: none.

Returns: nothing.

Destroys: a, y

Description: InitForIO prepares the system to perform I/O across the Commodore serial

bus. It disables interrupts, turns sprite DMA off, slows the 128 down to lMhz, switches in the ROM and I/O banks if necessary, and performs  $\,$ 

anything other initialization needed for fast serial transfer.

Call InitForIO before directly accessing the serial port (e.g., in a printer driver) or before using ReadBlock, WriteBlock, VerWriteBlock, or ReadLink. To restore the system to its previous state, call DoneWithIO.

Example: MyPutBlock

See also: DoneWithIO, SetDevice

disk very low level

**PurgeTurbo:** (C64, C128) **c235** 

Function: Completely deactivate and remove disk drive turbo code from current

drive, returning to standard Commodore DOS mode.

Parameters: none

Uses: curDrive currently active disk drive.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y

Description: PurgeTurbo deactivates and removes the turbo software from the current

drive, returning control of the device to the disk drive's internal ROM software. This allows access to normal Commodore DOS routines. An application may want to access the Commodore DOS to perform disk

functions not offered by the GEOS Kernal such as formatting.

Example:

See also: EnterTurbo, ExitTurbo.

disk very low level

ReadBlock: (C64, C128) C12A

Function: Very low-level read block from disk.

Parameters: rlL TRACK-valid track number (byte),

rlH SECTOR-valid sector on track (byte).

r4 BUFFER - address of buffer of BLOCKSIZE bytes to read block into

(word).

Uses: curDrive currently active disk drive.

curType GEOS 64 vl.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y

Description: ReadBlock reads the block at the specified TRACK and SECTOR into BUFFER.

If the disk is shadowed, **ReadBlock** will read from the shadow memory. **ReadBlock** is a pared down version of GetBlock. It expects the application to have already called **EnterTurbo** and **InitForIO**. By removing this overhead from GetBlock, multiple sector reads can be accomplished without the redundant initialization. This is exactly what happens in many of the higher-level disk routines that read multiple blocks at once, such

as **ReadFile**.

**ReadBlock** is useful for multiple-sector disk operations where speed is an issue and the standard GEOS routines don't offer a decent solution. **ReadBlock** can function as the foundation of specialized, high-speed disk

routines.

Example: MyGetBlock

See also: GetBlock, WriteBlock, VerWriteBlock.

disk very low level

**ReadLink:** (C64, C128) 904B

Function: Read link (first two bytes) from a disk block

Parameters: rlL TRACK - track number (byte),

rlH SECTOR - sector on track (byte).

r4 BUFFER — address of buffer of at least BLOCKSIZE bytes, usually

points to diskBlkBuf (word).

Uses: curDrive currently active disk drive.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y

Description: ReadLink returns the track/sector link from a disk block as the first

two bytes in BUFFER. The remainder of BUFFER (BLOCKSIZE-2 bytes) may or

may not be altered.

**ReadLink** is useful for following a multiple-sector chain in order to build a track/sector table. It mainly of use on 1581 disk drives, which walk through a chain significantly faster when only the links are read. Routines such as **DeleteFile** and **FollowChain** will automatically take

advantage of this capability of 1581 drives.

Note: Disk drives that do not offer any speed increase through ReadLink will

simply perform a ReadBlock.

Note: Does not work in 1541 Drivers. Use ReadBlock instead.

Example:

See also: ReadBlock, FollowChain

disk very low level

VerWriteBlock: (C64, C128) c223

Function: Very low-level verify block on disk.

Parameters: rlL TRACK - track number (byte).

rlH SECTOR - valid sector on track (byte).

r4 BUFFER - address of buffer of BLOCKSIZE bytes that contains data

that should be on this sector (word).

Uses: curDrive currently active disk drive.

**curType** GEOS 64 vl.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y

Description: VerWriteBlock verifies the validity of a recently written block. If the block does not verify, the block is rewritten by calling WriteBlock. VerWriteBlock is a low-level disk routine and expects the application to

have already called EnterTurbo and InitForIO.

VerWriteBlock can be used to accelerate the verifies that accompany multiple sector writes by first writing all the sectors and then verifying them. This is often faster than verifying a sector immediately after writing it because when writing sequential sectors, the GEOS turbo code will catch the sector interleave. If a sector is written and then immediately verified, the turbo code will need to wait for the disk to make one complete revolution before the newly-written sector will again pass under the read/write head. By writing all the sectors first and catching the interleave, then verifying all the sectors (again, catching the interleave), the dead time when the turbo code is Waiting for the disk to spin around is minimized. Many of the higher-level disk routines that write multiple blocks do just this.

**VerWriteBlock** is useful for multiple-sector disk operations where speed is an issue and the standard GEOS routines don't offer a decent solution. **VerWriteBlock** can function as the foundation of specialized, high-speed disk routines.

**VerWriteBlock** does not always do a byte-by-byte compare with the data in BUFFER. Some devices, such as the Commodore 1541, can do a cyclic redundancy check on the data in the block, and this internal checksum is sufficient evidence of a good write. Other devices, such as RAM-expansion units, have built-in byte-by-byte verifies.

Example: MyPutBlock

See also: WriteBlock, PutBlock

disk very low level

WriteBlock: (C64, C128) c220

Function: Very low-level write block to disk.

Parameters: rlL TRACK - valid track number (byte).

rlH SECTOR-valid sector on track (byte).

r4 BUFFER — address of buffer of BLOCKSIZE bytes that contains data

to write out (word).

Uses: curDrive currently active disk drive.

**curType** GEOS 64 vl.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y

Description: WriteBlock writes the block at BUFFER to the specified TRACK and SECTOR.

If the disk is shadowed, **WriteBlock** will also write the data to the shadow memory. **WriteBlock** is pared down version of **PutBlock**. It expects the application to have already called **EnterTurbo** and **InitForIO**, and it

does not verify the data after writing it.

WriteBlock can be used to accelerate multiple-sector writes and their accompanying verifies by writing all the sectors first and then verifying them. This is often faster than verifying a sector immediately after writing it because when writing sequential sectors, the GEOS turbo code will catch the sector interleave. If a sector is written and then immediately verified, the turbo code will need to wait for the disk to make one complete revolution before the newly written sector will again pass under the read/write head. By writing all the sectors first and catching the interleave, then verifying all the sectors (again, catching the interleave), the dead time when the turbo code is waiting for the disk to spin around is minimized. Many of the higher-level disk routines that write multiple blocks do just this.

WriteBlock is useful for multiple-sector disk operations where speed is an issue and the standard GEOS routines don't offer a decent solution. WriteBlock can function as the foundation of specialized, high-speed disk routines.

Example: MyPutBlock

See also: PutBlock, ReadBlock, VerWriteBlock.

## disk low-level

disk low level							
GetBlock	C1E4	Read single disk block into memory.	28				
PutBlock	C1E7	Write single disk block from memory.	29				

disk low-level

GetBlock: (C64, C128) C1E4

Function: General purpose routine to get a block from current disk.

Parameters: r4 BUFFER - address of buffer to place block; must be at least

BLOCKSIZE bytes (word).

rlL TRACK - track number (byte).

r1H SECTOR - sector number on track (byte).

Uses: curDrive currently active disk drive.

curType GEOS 64 vl.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

rl, r4 unchanged

Destroys: a,y

(1581 drive, r1, r4)

AAnote: Need to confirm is this is still true

Description: GetBlock reads a block from the disk into BUFFER. GetBlock is useful for

implementing disk utility programs and new file structures.

GetBlock is a higher-level version of ReadBlock. It calls InitForIO, EnterTurbo, ReadBlock, and DoneWithIO. If an application needs to read many blocks at once, ReadBlock may offer a faster solution. If the disk is shadowed, GetBlock will read from the shadow memory, resulting in a

faster transfer.

The Commodore 1581 driver has a bug that causes its GetBlock to trash rl

and r4.

Example:

See also: PutBlock, WriteBlock, BlkAllOC.

disk low-level

**PutBlock:** (C64, C128) **C1E7** 

Function: General purpose routine to write a block to disk with verify.

Parameters: r4 BUFFER - address of buffer to get block from;

rlL TRACK - valid track number (byte).
rlH SECTOR - valid sector on track (byte).

Uses: curDrive currently active disk drive.

curType GEOS 64 vl.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

rl, r4 unchanged

Destroys: a,y

Description: PutBlock writes a block from BUFFER to the disk. PutBlock is useful for

implementing disk utility programs and new file structures.

PutBlock is a higher-level version of WriteBlock. It calls InitForIO, EnterTurbo, ReadBlock, and DoneWithIO. If an application needs to write many blocks at once, WriteBlock may offer a faster solution. If the disk is shadowed, PutBlock will also write the data to the shadow memory..

Note<sup>3</sup>: PutBlock does no boundary check on the buffer. If the buffer is less the

then **BLOCKSIZE** (\$100) bytes, **PutBlock** will write the buffer and the memory contents that are after the buffer. This normally will not cause any problems as the size of data in the data block is stored in offset

1 of the block when the block is not full.

Example:

See also: GetBlock, WriteBlock, BlkAllOC.

AllocateBlock	9048	Mark a disk block as in-use.	31
BldGDirEntry	C1F3	Create a directory entry in memory.	32
BlkAllOC	C1FC	Allocate sectors for a file.	33
CalcBlksFree	C1DB	Compute number of free blocks on disk	34
ChkDkGEOS	C1DE	Check if a disk is GEOS format.	35
FastDelFile	\$C244	Delete a temporary file.	-407
FindBAMBit	\$C2AD	Check if a disk sector is in use.	-316
FollowChain	\$C205	Create a list of sectors used by a file.	-420
FreeBlock	\$C2B9	Free a BAM bit in memory.	
FreeFile	\$C226	Free a file's sectors.	-408
GetlstDirEntry	\$9030	Loads in the first directory block.	
GetNxtDirEntry	9033	Get directory entry other than first.	
GetDirHead	\$C247	Read track 18 sector 0.	-308
GetFHdrInfo	\$C229	Get a file's load address.	-419
GetFreeDirBlk	\$C1F6	Find a hole in the directory.	-416
GetOffPageTrSc	9036	Get track and sector of off-page directory.	
LdApplic	\$C21D	Load and run a file, given a directory entry.	-412
LdDeskAcc	\$C217	Load a file with memory swapping.	-413
LdFile	\$C211	Load a file, given a directory entry.	-411
NewDisk	C1E1	Initialize a drive.	
NxtBlkAlloc	\$C24D	Allocate sectors for a file.	-319
PutDirHead	\$C24A	Write to track 18 sector 0.	-309
ReadByte	\$C2B6	Read a File 1 byte at a time.	
ReadFile	\$C1FF	Load a chain of blocks into memory.	-415
SetGDirEntry	\$C1F0	Create a directory entry on disk.	-417
SetNextFree	\$C292	Find and allocate a disk block.	-317
StartAppl	\$C22F	Run a program that is in memory.	-807
WriteFile	\$C1F9	Save memory to preallocated sectors.	-404

#### AllocateBlock: (C64, C128) 9048

Allocate a disk block, marking it as in use.. Function:

Parameters: r6L track number of block (byte).

r6H sector number of block (byte).

curDrive drive that disk is in. Uses:

curDirHead this buffer must contain the current directory header.

dir2Head2† (BAM for 1571 and 1581 drives only)

dir3Head3<sup>†</sup> (BAM for 1581 drive only)

 $;^{\dagger}$ used internally by GEOS disk routines; applications generally don't use.

x error (\$00 = no error). Returns:

> BAD BAM **r6** unchanged

Alters: curDirHead BAM updated to reflect newly allocated blocks.

dir2Head† (BAM for 1571 and 1581 drives only)
dir3Head† (BAM for 1581 drive only)

Destroys: a,y, **r7**, **r8**H

Description: AllocateBlock allocates a single block on this disk by setting the

appropriate flag in the block allocation map (BAM).

If the sector is already allocated then a BAD BAM error is returned. AllocateBlock does not automatically write out the BAM. See PutDirHead for more information on writing out the BAM. The Commodore 1541 device drivers do not have a jump table entry for AllocateBlock. All other device drivers, however, do. The following subroutine will properly

allocate a block on any device, including the 1541.

NewAllocateBlock

Example: CallNewAlloc

BldGDirEntry: (C64, C128) C1F3

Function: Builds a directory entry in memory for a GEOS file using the

information in a file header.

Parameters: r2 NUMBLOCKS - number of blocks in file (word).

r6 TSTABLE - pointer to a track/sector list of unused blocks (unused but allocated in the BAM), usually a pointer to fileTrScTab; BlkAllOC

can be used to build such a list (word).

r9 FILEHDR - pointer to GEOS file header (word).

Uses: curDrive drive that disk is in.

Returns: r6 pointer to first non-reserved block in track/sector table

(BldGDirEntry reserves one block for the file header and a second

block for the index table if the file is a VLIR file).

Alters: dirEntryBuf contains newly-built directory entry.

Destroys: a,y, r5

Description: Given a GEOS file header, BldGDirEntry will build a system specific

directory entry suitable for writing to an empty directory slot.

Most applications create new files by calling **SaveFile**. **SaveFile** calls **SetGDirEntry**, which calls **BldGDirEntry** as part of its normal processing.

Example: MySetGDirEntry

**BlkAllOC:** (C64, C128) **C1FC** 

Function: Allocate enough disk blocks to hold a specified number of bytes.

Parameters: r2 BYTES — number of bytes to allocate space for. Commodore version can allocate up to 32,258 bytes (127 Commodore blocks).

r6 TSTABLE - pointer to buffer for building out track and sector table
 of allocated blocks, usually points to fileTrScTab (word).

Uses: curDrive drive that disk is in.

curDirHead this buffer must contain the current directory header.

 ${\tt dir2Head2^{\dagger}}$  (BAM for 1571 and 1581 drives only)

dir3Head3† (BAM for 1581 drive only)

interleave desired physical sector interleave (usually 8); used by
setNextFree. Applications need not set this explicitly -

will be set automatically by internal GEOS routines.

;  $^{\dagger}$ used internally by GEOS disk routines; applications generally don't use.

Returns:

r2 number of blocks allocated to hold BYTES amount of data.

r3L track of last allocated block.
r3H sector of last allocated block.

Alters: curDirHead BAM updated to reflect newly allocated blocks.

dir2Head\* (BAM for 1571 and 1581 drives only)

dir3Head<sup>†</sup> (BAM for 1581 drive only)

Destroys: a, y, r4-r8

Description: BlkAllOC calculates the number of blocks needed to store BYTES amount of data, taking any standard overhead into account (such as the two-byte track/sector link required in each Commodore block), then calls CalcBlksFree to ensure that enough free blocks exist on the disk. If there are not enough free blocks to accommodate the data, BlkAllOC returns an INSUFFICIENT\_SPACE error without allocating any blocks. Otherwise, BlkAllOC calls SetNextFree to allocate the proper number of

unused blocks.

BlkAllOC builds out a track and sector table in the buffer pointed to by TSTABLE. The 256 bytes at fileTrScTab are usually used for this purpose. When BlkAllOC returns, the table contains a two-byte entry for each block that was allocated: the first byte is the track and the second byte is the sector. The last entry in the table has its first byte set to \$00, indicating the end of the table. The second byte of the last entry is an index to the last byte in the last block. This track/sector list can be passed directly to WriteFile for use in writing data to the blocks.

Note: For more information on the scheme used to allocate successive blocks,

refer to **SetNextFree**.

Example: GrabSomeBlocks

CalcBlksFree: (C64, C128) C1DB

Function: Calculate total number of free blocks on disk.

Parameters: r5 DIRHEAD - address of directory header, should always point to

curDirHead (word).

Uses: curDrive drive that disk is in.

dir2Head2<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head3† (BAM for 1581 drive only)

;tused internally by GEOS disk routines; applications generally don't use.

Returns: r4 number of free blocks.

r5 unchanged.

r3 in GEOS v1.3 and later: total number of available blocks on empty disk. This is useful because v1.3 and later support disk devices other than the 1541. GEOS versions earlier than v1.3 leave r3 unchanged.

Destroys: a, y

Description: CalcBlksFree calculates the number of free blocks available on the disk.

An application can call CalcBlksFree, for example, to tell the user the amount of free space available on a particular disk. GEOS disk routines

that allocate multiple blocks at once, such as **BlkAlloC**, call **CalcBlksFree** to ensure enough free space exists on the disk to prevent a surprise **ENSUFFICENT\_SPACE** error, midway through the allocation. (This is why it is usually not necessary to check for sufficient space before saving a file or a VLIR record—the higher level GEOS disk routines handle

this checking automatically.)

CalcBlksFree looks at the BAM in memory and counts the number of unallocated blocks. The BAM is stored in the directory header and the directory header is stored in the buffer at curDirHead. Calling CalcBlksFree requires first loading r5 with the address of curDirHead.

LoadW r5, #curDirHead jsr CalcBlksFree

When checking the total number of blocks (both allocated and free) on a particular disk device, call **CalcBlksFree** with r3 loaded with the number of blocks on a 1541 disk device. On GEOS v1.3 and above, this number is changed to reflect the actual number of blocks in the device. On previous versions of GEOS, r3 comes back unchanged.

```
N1541 BLOCKS = 664 ; total number of blocks on 1541 devices
```

```
LoadW r3, #N1541_BLOCKS; assume 1541 block count for v1.2 Kernal's
```

LoadW r5, #curDirHead ; point to the directory header

jsr CalcBlksFree ; r3 comes back with total number of blocks

; on this device

Example: CheckDiskSpace

See also: NxtBlkAlloc, SetNextFree, GetFreeDirBlk, FreeBlock.

ChkDkGEOS: (C64, C128) C1DE

Function: Check Commodore disk for GEOS format.

 $\textbf{Parameters:} \quad \textbf{r5} \quad \textbf{DIRHEAD} \ - \ \textbf{address} \quad \textbf{of directory header, should always point to}$ 

curDirHead (word).

Returns: a TRUE/FALSE matching isGEOS.

Z flag=0 GEOS Disk
Z flag=1 Non GEOS Disk

Alters: isGEOS set to TRUE if disk is a GEOS disk, otherwise set to FALSE.

Destroys: a, y

Description: ChkDkGEOS checks the directory header for the version string that flags it as a GEOS disk (at OFF\_GEOS\_BD). The primary difference between a GEOS disk and a standard Commodore disk is the addition of the off-page directory and the possibility of GEOS files on the disk. GEOS files have an additional file header block that holds the icon image and other information, such as the author name and permanent name string. To convert a non-GEOS disk into a GEOS disk, use SetGEOSDisk.

**OpenDisk** automatically calls **ChkDkGEOS**. As long as **OpenDisk** is used before reading a new disk, applications should have no need to call **ChkDkGEOS** 

### Example:

```
isr
         GetDirHead
                            ; read in the directory header
                             ; check status
   txa
   bne
         99$
                             ; exit on error
   LoadW r5, #curDirHead
                            ; point to directory header
   jsr
         ChkDkGEOS
                             ; Check for GEOS disk
   beq 50$
                             ; if not a GEOS disk, branch
   ; code here to handle GEOS disk
   bra 90$
                             ; jump to exit
50$
   ; code here to handle non-GEOS disk
90$
   clc
                             ; Success Exit
   rts
99$
   sec
                             ; error exit
   rts
```

FastDelFile: (C64, C128) C244

Function: Special Commodore version of DeleteFile that quickly deletes a

sequential file when the track/sector table is available.

Parameters: r0 FILENAME - pointer to null-terminated file name (word).

r3 TSTABLE - pointer to track and sector table of file, usually points

to fileTrScTab (word).

Uses: curDrive

curType GEOS 64 v 1.3 and later for detecting REU shadowing.

curDirHead BAM updated to reflect newly freed blocks.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

**Returns:** x = (\$00 = no = rror).

Destroys: a,y, r0, r9

 $\textbf{Description: FastDelFile} \ \ \text{quickly deletes a sequential file by taking advantage of an}$ 

already existing track/sector table. It first removes the directory entry determined by FILENAME and calls FreeBlock for each block in a track/sector table at TSTABLE. The track/sector table is in the standard format, such as that returned from ReadFile, where every two-byte entry constitutes a track and sector. A track number of \$00 terminates the

table.

FastDelFile is fast because it does not need to follow the chain of sectors to delete the individual blocks. It can do most of the deletion by manipulating the BAM in memory then writing it out with a call to

PutDirHead when done.

FastDelFile will not properly delete VLIR files without considerable work on the application's part. Because there is no easy way to build a track/sector table that contains all the blocks in all the records of a VLIR file, it is best to use DeleteFile or FreeFile for deleting VLIR

files or **DeleteRecord** for deleting a single record.

FastDelFile calls GetDirHead before freeing any blocks. This will

overwrite any BAM and directory header in memory.

Note: FastDelFile can be used to remove a directory entry without actually

freeing any blocks in the file by passing a dummy track/sector table, where the first byte (track number) is \$00 signifying the end of the

table: See Example DeleteDirEntry

Examples: DeleteDirEntry, ReadAndDelete

# FindBAMBit: (C64, C128) C2AD

Function: Get disk block allocation status.

Parameters: r6L TRACK -track number of block (byte).

 ${f r6H}$  SECTOR — sector number of block (byte).

Uses: curDrive

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

Returns: st z flag reflects allocation status (1 = free; 0 = allocated).

**r6** unchanged

1541 drives only:

x offset from **curDirHead** for BAM byte.

r8H mask for isolating BAM bit. a BAM byte masked with r8H.

r7H offset from curDirHead of byte that holds free blocks on track total.

Destroys: non-1541 drives:

a, y, r7H, r8H.

1541 drives:

y (a, r7H, and r8H all contain useful values).

Description: FindBAMBit accesses the BAM of the current disk "in curDirHead) and returns the allocation status of a particular block. If the BAM bit is zero, then the block is in-use; if the BAM bit is one, then the block is free. FindBAMBit returns with the z flag set to reflect the status of the BAM so that a subsequent bne or beq branch instructions can test the

status of a block after calling FindBAMBit.

bne BlockIsFree ;branch if block is free

- or -

beq BlockInUse ;branch if block is in-use

Note: FindBAMBit will return the allocation status of a block on any disk device, even those with large or multiple BAMs (such as the 1571 and 1581 disk drives). Only the 1541 driver, however, will return useful information in a, y, r7H, and r8H. For an example of using these extra

1541 return values, refer to AllocateBlock.

Examples:

LoadB r6L, #TRACK ; get track and sector number

LoadB r6H, #SECTOR

jsr FindBAMBit ; get allocation status
beq BlocklnUse ; branch if already in use

FollowChain: (C64, C128) C205

Function: Follow a chain of Commodore disk blocks, building out a track/sector

table.

Parameters: rlL START\_TRACK - track number of starting block (byte).

rlH START\_SEC - sector number of starting block (byte).

 ${f r3}$  TSTABLE — pointer to buffer for building out track and sector

table of chain, usually points to fileTrScTab (word).

Uses: curDrive

curType GEOS 64 v 1.3 and later for detecting REU shadowing.

**Returns:** x = (\$00 = no = no).

r3 unchanged

track/sector built-out in buffer pointed to by TSTABLE.

Alters: diskBlkBuf used for temporary block storage.

Destroys: a, y, r1, r4

Description: FollowChain constructs a track/sector table for a list of chained blocks

on the disk. It starts with the block passed in START\_TR and START\_SC and follows the links until it encounters the last block in the chain. Each block (including the first block at START\_TR, START\_SC) becomes a

part of the track/sector table.

Commodore disk blocks are linked together with track/sector pointers. The first two bytes of each block represent a track/sector pointer to the next block in the chain. Each sequential file and VLIR record on the disk is actually a chained list of blocks. FollowChain follows these track/sector links, adding each to the list at TSTABLE until it encounters a track pointer of \$00, which terminates the chain. FollowChain adds this last track pointer (\$00) and its corresponding sector pointer (which is actually an index to the last valid byte in the block) to the track/sector table and returns to the caller.

FollowChain builds a standard track/sector table compatible with routines such as WriteFile and FastDelFile.

disk mid-level

FreeBlock: (C64, C128) C2B9

Function: Free an allocated disk block.

Parameters: r6 track number of block to free (byte).

r6H sector number of block to free (byte).

Uses: curDrive

curDirHead must contain the current directory header.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

BAD BAM if block already free.

r6L, r6H unchanged.

Alters: curDirHead BAM updated to reflect newly allocated block.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

Destroys: a,y,r7,r8H

Description: FreeBlock tries to free (deallocate) the block number passed in r6. If

the block is already free, then FreeBlock returns a BAD BAM error.

Note: FreeBlock was not added to the Commodore GEOS jump table until v1.3, but

it can be accessed directly under GEOS v1.2. The following routine will check the GEOS version number and act correctly under GEOS v1.2 and

later. See Example MyFreeBlock

Example: MyFreeBlock

**FreeFile:** (C64, C128) **c226** 

Function: Free all the blocks in a GEOS file (sequential or VLIR) without

deleting the directory entry. The GEOS file header and any index

blocks are also deleted.

Parameters: r9 DIRENTRY - pointer to directory entry of file being freed, usually

points to dirEntryBuf (Apple GEOS: must be in main memory.)

(word).

Uses: curDrive

curType GEOS 64 v 1.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

r6L, r6H unchanged.

Alters: diskBlkBuf used for temporary block storage.

curDirHead BAM updated to reflect newly allocated block.

dir2Head<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head† (BAM for 1581 drive only)

fileHeader temporary storage of the index table when deleting a

VLIR file.

;tused internally by GEOS disk routines; applications generally don't use.

Destroys: a,y,r0-r9

Description: Given a valid directory entry, FreeFile will delete (free) all blocks

associated with the file. The GEOS file header and any index blocks associated with the file are also be freed. The directory entry on the

disk, however, is left intact

The directory entry is a standard GEOS data structure returned by routines such as FindFile, GetlstDirEntry and GetNxtDirEntry. FreeFile

is called automatically by **DeleteFile**.

FreeFile tries to free (deallocate) the block number passed in r6. If the block is already free, then FreeBlock returns a BAD BAM error.

FreeFile calls GetDirHead to get the current directory header and BAM into memory. It then checks at OFF GHDR PTR in the directory entry for

a GEOS file header block, which it then frees.

If the file is a sequential file, **FreeFile** walks the chain pointed at by the **OFF\_DE\_TR\_SC** track/sector pointer in the directory header and frees all the blocks in the chain. **FreeFile** then calls **PutDirHead** to write out

the new BAM.

When using **GetlstDirEntry** and **GetNxtDirEntry**, do not pass **FreeFile** a pointer into **diskBlkBuf**. Copy the full directory entry (**DIRENTRY\_SIZE** bytes) from **diskBlkBuf** to another buffer (such as **dirEntryBuf**) and pass **FreeFile** the pointer to that buffer. Otherwise when **FreeFile** uses

diskBlkBuf it will corrupt the directory entry.

Because **FreeFile** deletes a block at a time as it follows the chains, it is capable of deleting files with chains larger than 127 blocks, which is the standard GEOS limit imposed by the size of TrScTable.

See also: DeleteFile, FreeDir, FreeBlock.

GetlstDirEntry: (C64, C128) 9030

Function: Loads in the first directory block of the current directory and

returns a pointer to the first directory entry within this block.

Parameters: none.

Uses: curDrive

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

r5 pointer to first directory entry within diskBlkBuf.

Alters: diskBlkBuf directory block.

Destroys: a,y,r1,r4

Description: GetlstDirEntry reads in the first directory block of the current

directory and returns with r5 pointing to the first directory entry. **GetlstDirEntry** is called by routines like **FindFTypes** and **FindFile**.

To get a pointer to subsequent directory entries, call GetNxtDirEntry.

Since Commodore GEOS does not support a hierarchical file system, the "current directory" is actually the entire disk.

**GetlstDirEntry** did not appear in the jump table until version 1.3. An application running under version 1.2 can access **GetlstDirEntry** by calling directly into the Kernal. The following subroutine will work on Commodore GEOS v1.2 and later:

; MyGetlstDirEntry - Call instead of GetlstDirEntry

; to work on GEOS vl.2 and later

; EQUATE: v1.2 entry point directly into Kernal. Must

; do a version check before calling.

MyGetlstDirEntry:

lda **version** ; check version number

cmp #\$13

bcc 10\$ ; branch < v1.3
jmp Get1stDirEntry ; direct call</pre>

10\$

jmp o GetlstDirEntry ; go through jump table

#### Example:

See also: GetNxtDirEntry, FindFTypes.

GetNxtDirEntry: (C64, C128) 9033

Function: Given a pointer to a directory entry returned by GetlstDirEntry or

 ${\tt GetNxtDirEntry,\ returns\ a\ pointer\ to\ the\ next\ directory\ entry.}$ 

Parameters: r5 CURDIRENTRY - pointer to current directory entry as returned from

GetlstDirEntry or GetNxtDirEntry; will always be a pointer into

diskBlkBuf (word).

Uses: curDrive

diskBlkBuf must be unaltered from previous call to GetlstDirEntry

or GetNxtDirEntry.

**curType** GEOS 64 v 1.3 and later for detecting REU shadowing.

**Returns:** x = error (\$00 = no error).

r5 pointer to next directory entry within diskBlkBuf.

y non-zero if end of directory reached

Alters: diskBlkBuf directory block.

Destroys: a,y,r1,r4

Description: GetNxtDirEntry increments r5 to point to the next directory entry in diskBlkBuf. If diskBlkBuf is exceeded, the next directory block is read

in and **r5** is returned with an index into this new block. Before calling

GetNxtDirEntry for the first time, call GetlstDirEntry.

**GetNxtDirEntry** did not appear in the jump table until version 1.3. An application running under version 1.2 can access **GetNxtDirEntry** by calling directly into the Kernal. The following subroutine will work on Commodore GEOS v1.2 and later:

```
; MyGetNxtDirEntry - Call instead of GetNxtDirEntry
; to work on GEOS vl.2 and later
; EQUATE: vl.2 entry point directly into Kernal. Must
; do a version check before calling.
o GetNxtDirEntry = $ca10
                         ; exact entry point
MyGetNxtDirEntry:
  lda version
                           ; check version number
  cmp #$13
  bcc 10$
                          ; branch < v1.3
  jmp GetNxtDirEntry
                          ; direct call
10$
  jmp o GetNxtDirEntry
                          ; go through jump table
```

# Example:

See also: Get1stDirEntry, FindFTypes.

disk mid-level

GetDirHead: (C64, C128) C247

Function: Read directory header from disk. GEOS also reads in the BAM

Parameters: none.

Uses: curDrive

**Returns:** x = error (\$00 = no error).

r4 pointer to curDirHead.

Alters: curDirHead contains directory header

dir2Head<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head† (BAM for 1581 drive only)

;tused internally by GEOS disk routines; applications generally don't use.

Destroys: a,y,r1

Description: GetDirHead reads the full directory header (256 bytes) into the buffer

at  ${\tt curDirHead}$ . This block also includes the BAM (block allocation map)

for the entire disk.

GEOS disks, like the standard Commodore disks upon which they are based, have one directory header. The directory header occupies one full block on the disk. The Commodore directory header contains information about the disk, such as the location of the directory blocks, the disk name, and the GEOS version string (if a GEOS disk). The Commodore directory header also contains the disk BAM, which flags particular sectors as used or unused

**GetDirHead** calls **GetBlock** to read in the directory header block into the buffer at **curDirHead**. The directory header block contains the directory header and the disk BAM (block allocation map). Typically, applications don't call **GetDirHead** because the most up-to-date directory header is almost always in memory (at **curDirHead**), **OpenDisk** calls **GetDirHead** to get it there initially. Other GEOS routines update it in memory, some calling **PutDirHead** to bring the disk version up to date.

Because Commodore disks store the BAM information in the directory header it is important that the BAM in memory not get overwritten by an outdated BAM on the disk. An application that manipulates the BAM in memory (or calls GEOS routines that do so), must be careful to write the BAM back out (with <code>PutDirHead</code>) before calling any other routine that might overwrite the copy in memory. <code>GetDirHead</code> is called by routines such as <code>OpenDisk</code>, <code>SetGEOSDisk</code>, and <code>OpenRecordFile</code>, etc.

### Example:

See also: PutDirHead

GetFHdrInfo: C229 (C64, C128)

Loads the GEOS file header for a particular directory entry. Function:

Parameters: r9 DIRENTRY - pointer to directory entry of file, usually points to

dirEntryBuf (Apple GEOS: must be in main memory) (word).

Uses: curDrive

> GEOS 64 v 1.3 and later for detecting REU shadowing. curType

error (\$00 = no error). Returns:

r7 load address copied from the O GHST ADDR word of the GEOS file

header.

track/sector copied from bytes +1 and +2 of the directory entry (DIRENTRY). This is the track/sector of the first data block of a sequential file (OFF DE TR SC) or the index table block of a VLIR file

(OFF INDEX PTR).

contains 256-byte GEOS file header. Alters: fileHeader

> track/sector of header added to first two bytes of this fileTrScTab

table; a subsequent call to **ReadFile** or similar routine will augment this table beginning with the third byte

(fileTrScTab+2) so as not to disrupt this value.

Destroys: a,y,r4

Description: Given a valid directory entry, GetFHdrInfo will load the GEOS file header

into the buffer at fileHeader.

The directory entry is a standard GEOS data structure returned by routines such as FindFile, GetlstDirEntry and GetNxtDirEntry. GetFHdrInfo is called by routines such as LdFile just prior to calling ReadFile (to load in a sequential file or record zero of a VLIR).

GetFHdrInfo gets the block number (Commodore track/sector) of the GEOS file header by looking at the OFF GHDR PTR word in the directory entry.

disk mid-level

GetFreeDirBlk: (C64, C128) C1F6

Function: Search the current directory for an empty slot for a new directory

entry. Allocates another directory block if necessary.

Parameters: rlOL DIRPAGE - directory page to begin searching for free slot; each

directory page holds eight files and corresponds to one notepad

page on the GEOS deskTop. The first page is page one.

Uses: curDrive

curType GEOS 64 v 1.3 and later for detecting REU shadowing.
curDirHead this buffer must contain the current directory header.

 ${\tt dir2Head2^{\dagger}}$  (BAM for 1571 and 1581 drives only)

dir3Head3† (BAM for 1581 drive only)

interleave<sup>†</sup> desired physical sector interleave (usually 8); Applications

need not set this explicitly — will be set automatically by internal GEOS routines. Only used when new directory block  $\,$ 

is allocated.

; tused internally by GEOS disk routines; applications generally don't use.

**Returns:** x error (\$00 = no error).

FULL\_DIRECTORY

rl0L page number of empty directory slot.

rl block (track/sector) number of directory block in diskBlkBuf.

y index to empty directory slot in diskBlkBuf.

Alters: curDirHead contains directory header

dir2Head<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head† (BAM for 1581 drive only)

Destroys: a,r0,r3,r5,r7-r8.

**Description: GetFreeDirBlk** searches the current directory looking for an empty slot for a new directory entry. A single directory page has eight directory slots, and these eight slots correspond to the eight possible files that

can be displayed on a single GEOS deskTop notepad page.

**GetFreeDirBlk** starts searching for an empty slot beginning with page number **DIRPAGE**. If **GetFreeDirBlk** reaches the last directory entry without finding an empty slot, it will try to allocate a new directory block. If DIRPAGE doesn't yet exist, empty pages are added to the directory structure until the requested page is reached.

\$01 will most often be passed as the DIRPAGE starting page number, so that all possible directory slots will be searched, starting with the first page. If higher numbers are used, **GetFreeDirBlk** won't find empty directory slots on lower pages and extra directory blocks may be allocated needlessly.

 ${\tt GetFreeDirBlk}$  is called by  ${\tt SetGDirEntry}$  before writing out the directory entry for a new GEOS file.

Since GEOS 2.0 does not support a hierarchical file system, the "current directory" is actually the entire disk. A directory page corresponds exactly to a single sector on the directory track. There is a maximum of 18 directory sectors (pages) on a Commodore disk. If this 18th page is exceeded, GetFreeDirBlk will return a FULL DIRECTORY error.

**GetFreeDirBlk** allocates blocks by calling **SetNextFree** to allocate sectors on the directory track. **SetNextFree** will special-case the directory track allocations. Refer to **SetNextFree** for more information.

GetFreeDirBlk does not automatically write out the BAM. See PutDirHead for more information on writing out the BAM.

Example: MySetGDirEntry

GetOffPageTrSc: (C64, C128) 9036

Function: Get track and sector of off-page directory.

Parameters: none.

Uses: curDrive drive that disk is in.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

y \$ff if the disk is not a GEOS disk and therefore has no off-page

directory block, otherwise \$00.

rlL track of off-page directory.
rlH sector of off-page directory.

r4 pointer to curDirHead.

Destroys: a,y, r5

Description: Commodore GEOS disks have an extra directory block somewhere on the disk

called the off-page directory. The GEOS deskTop uses the off-page directory block to keep track of file icons that have been dragged off of the notepad and onto the border area of the deskTop. The off-page

directory holds up to eight directory entries.

GetOffPageTrSc reads the directory header into the buffer at curDirHead and calls ChkDkGEOS to ensure that the disk is a GEOS disk. If the disk is not a GEOS disk, it returns with \$ff in the y register. Otherwise, GetOffPageTrSc copies the off-page track/sector from the OFF\_OP\_TR\_SC

word in the directory header to rl and returns \$00 in y.

```
; Put off-page block into diskBlkBuf
```

```
GetOffPageTrSc
                                ; get off-page directory block
      jsr
      txa
                                ; check for error
      bne
              99$
                                ; check for GEOS disk
      tya
      tax
                                ; put in x in case error
              99$
      bne
      LoadW r4, #diskBlkBuf ; get off-page block
                                ; return with error in x
      jsr
99$
      rts
```

**LdApplic:** (C64, C128) **C21D** 

**Function:** Load and (optionally) run a GEOS application, passing it the standard application startup flags as if was launched from the deskTop.

**r0L** LOAD OPT:

bit 0: 0 load at address specified in file header; application
 will be started automatically

1 load at address in r7; application will not be started automatically.

bit 7: 0 not passing a data file.

1 r2 and r3 contain pointers to disk and data file names,

bit 6: 0 not printing data file.

1 printing data file; application should print file and exit

r7 LOAD\_ADDR - optional load address, only used if bit 0 of LOAD\_OPT
 is set (word).

PATA\_DISK - only valid if bit 7 or bit 6 of LOAD\_OPT is set: pointer
to name of the disk that contains the data file, usually a pointer
to one of the DrxCurDkNm buffers (word).

r3 DATA\_FILE - only valid if bit 7 of LOAD\_OPT is set: pointer to name
 of the data file (word).

Uses: curDrive drive that disk is in.

curType GEOS 64 v 1.3 and later for detecting REU shadowing.

Returns: only returns if alternate load address or disk error.

 $\mathbf{x}$  error (\$00 = no error).

Passes: usually doesn't return, but warmstarts GEOS and passes the following:

r0 as originally passed to LdApplic.

r2 as originally passed to LdApplic (use dataDiskName).
r3 as originally passed to LdApplic.(use dataFileName).

Alters: GEOS brought to a warmstart state.

 ${\tt dataDiskName}$  contains name of data disk if bit 7 of  ${\tt r0}$  is set.  ${\tt dataFileName}$  contains name of data file if bit 6 of  ${\tt r0}$  is set

Destroys: a,x,y,r0-r15

Description: LdApplic is a mid-level application loading routine called by the higher

level **GetFile**. Given a directory entry of a GEOS application file, **LdApplic** will attempt load it into memory and optionally run it. **LdApplic** calls **LdFile** to load the application into memory: a sequential file is loaded entirely into memory but only record zero of a VLIR file is loaded. Based on the status of bit 0 of **LOAD\_OPT**, optionally runs the application by calling it through **StartAppl**.

Most applications will not call **LdApplic** directly but will go indirectly through **GetFile**.

Note:

Only in extremely odd cases will an alternate load address be specified for an application. Loading an application at another location is not particularly useful because it will most likely not run at an address other than its specifiec load address. When **LdApplic** returns to the caller, it does so before calling **StartAppl** to warmstart GEOS.

Example:

See also: GetFile, LdDeskAcc, StartAppl

disk mid-level

**LdDeskAcc:** (C64, C128) **c217** 

Function: Load and run a .GEOS desk accessory.

Parameters: r9 DIRENTRY - pointer to directory entry of file, usually points to

dirEntryBuf (word).

rOL RECVR\_OPTS - should be set to \$00 (see below for explanation) (byte).

Uses: curDrive drive that disk is in.

Returns: returns when desk accessory exits with a call to RstrAppl.

 $\mathbf{x}$  error (\$00 = no error).

Passes: warmstarts GEOS and passes the following to the desk accessory:

r10L as originally passed to LdDeskAcc (should be \$00; see below).

Alters: nothing directly; desk accessory may alter some buffers that are not

saved.

Destroys: a,x,y,r0-r15

Description: LdDeskAcc is a mid-level desk accessory loading routine called by the

higher level **GetFile**. Given a directory entry of a GEOS desk accessory file, **LdDeskAcc** will attempt load it into memory and run it. When the user closes the desk accessory, control returns to the calling

application.

LdDeskAcc first loads in the desk accessory's file header to get the start and ending load address. Under GEOS 64 and Apple GEOS, it will then save out the area of memory between these two addresses to a file on the current disk named SWAP FILE. The GEOS 128 version saves this area to the 24K desk accessory swap area in back RAM. Desk accessories larger than 24K cannot be used under GEOS 128 (to date, there are none); a BFR OVERFLOW error is returned.

After saving the overlay area, the dialog box and desk accessory save-variables are copied to a special area of memory, the current stack pointer is remembered, and the desk accessory is loaded and executed. When the desk accessory calls RstrAppl to return to the application, this whole process is reversed to return the system to a state similar to the one it was in before the desk accessory was called. The SWAP FILE file is deleted.

Most applications will not call  ${\bf LdDeskAcc}$  directly, but will go indirectly through  ${\bf GetFile.}$ 

C64: GEOS versions 1.3 and above have a GEOS file type called TEMPORARY. When the deskTop first opens a disk, it deletes all files of this type. The

SWAP FILE is a TEMPORARY file.

#### Note:

The RECVR OPTS flag originally carried the following significance:

- bit 7: 1 force desk accessory to save foreground screen area and restore it on return to application.
  - 0 not necessary for desk accessory to save foreground.
- bit 6: 1 force desk accessory to save color memory and restore it on return to application.
  - 0 not necessary for desk accessory to save foreground.

#### Note:

It was found that the extra code necessary to make desk accessories save the foreground screen and color memory provided no real benefit because this context save can just as easily be accomplished from within the application itself. The RECVR\_OPTS flag is set to \$00 by all Berkeley Softworks applications, and desk accessories can safely assume that this will always be the case. (In fact, future versions of GEOS may force rlOH to \$00 before calling desk accessories just to enforce this standard!)

The application should always set  ${\bf r10H}$  to \$00 and bear the burden of saving and restoring the foreground screen and the color memory. (Color memory only applicable to GEOS 64 and GEOS 128 in 40-column mode.)

disk mid-level

**LdFile:** (C64, C128) **c211** 

Function: Given a directory entry, loads a sequential file or record zero of a

VLIR record.

Parameters: r9 DIRENTRY - pointer to directory entry of file, usually points to

dirEntryBuf (word).

Uses: curDrive drive that disk is in.

curType GEOS 64 v 1.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

r7 pointer to last byte read into BUFFER plus one.

Alters: fileHeader contains 256-byte GEOS file header. (This is a 512-byte

buffer in Apple GEOS, although only 256 bytes are used

in the GEOS file header for compatibility).

fileTrScTab track/sector of header in first two bytes of this table

(fileTrScTab+0 and fileTrScTab+1); As the file is loaded, the track/sector pointer to each block is added to the file track/sector table starting at fileTrScTab+2

and fileTrScTab+3.

Destroys: Not Listed in source material. LdFile is in an Unusable state already

so this is to be expected.

Description: LdFile is a mid-level file handling routine called by the higher level

**GetFile.** Given a directory entry of a sequential file, **LdFile** will load it into memory. Given the directory entry of a VLIR file, **LdFile** will

load its record zero into memory.

Most applications will not call  $\mathbf{LdFile}$  directly, but will go

indirectly through GetFile.

All versions of  $\mathbf{LdFile}$  to date under Commodore GEOS are unusable because the load variables that are global under Apple GEOS (loadOpt and loadAddr) are local to the Kernal and inaccessible to applications.

Fortunately this is not a problem because applications can always go through **GetFile** to achieve the same effect.

disk mid-level

**NewDisk:** (C64, C128) **C1E1** 

Function: Tell the turbo software that a new disk has been inserted into the

drive.

Parameters: r1L1 Track to position the disk drive head at.

rlH¹ Sector to position the disk drive head at.

Uses: curDrive drive that disk is in.

curType GEOS 64 v 1.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Destroys: a,y, r0-r3

Description: NewDisk informs the disk drive turbo software that a new disk has been

inserted into the drive. It first calls **EnterTurbo** then sends an initialize command to the turbo code. If the disk is shadowed, the shadow

memory is also cleared.

NewDisk gets called automatically when OpenDisk opens a new disk. An application that does not deal with anything but the low-level disk

routines might want to call NewDisk instead of OpenDisk to avoid the unnecessary overhead associated with reading the directory header and

initializing internal file-level variables.

Note: NewDisk has no effect on a RAMdisk. Also, some early versions of the

1541 turbo code leave the disk in the drive spinning after it is first loaded. A call to **NewDisk** during the application's initialization will

stop the disk.

Note: 1 It also positions the head over a particular sector.

Calls: 2 EnterTurbo, InitForIO, DoneWithIO

**ReadByte:** (C64, C128) **C2B6** 

Function: Special version of ReadFile that allows reading a chained list of

blocks a byte at a time.

Parameters: on initial call only:

r1 START\_TRSC - track/sector of first data block (word).

 ${f r4}$  BLOCKBUF — pointer to temporary buffer of  ${\bf BLOCKSIZE}$  bytes for use by

ReadByte, usually a pointer to diskBlkBuf (word).

**r5** \$0000 (word).

Uses: curDrive drive that disk is in.

curType GEOS 64 v 1.3 and later for detecting REU shadowing.

Returns: a byte returned

x = (\$00 = no = ror).

 ${f r1},\ {f r4},\ {f r5}$  contain internal values that must be preserved between calls

to ReadByte.

Destroys: y

Description: ReadByte allows a chain of blocks on the disk to be read a byte at a

time. The first time **ReadByte** is called, **r1**, **r4**, and **r5** must contain the proper parameters. When **ReadByte** returns without an error, the a register will contain a single byte of data from the chain. To read another byte, call **ReadByte** again. Between calls to **ReadByte**, the application must

preserve  ${\tt r1}$ ,  ${\tt r4}$ ,  ${\tt r5}$ , and the data area pointed to by BLOCKBUF.

**ReadByte** loads a block into BLOCKBUF and returns a single byte from the buffer at each call. After returning the last byte in the buffer, **ReadByte** loads in the next block in the chain and starts again from the beginning of BLOCKBUF. This process continues until there are no more bytes in the

file. A BFR OVERFLOW error is then returned.

ReadByte is especially useful for displaying very large bitmaps with

BitOtherClip

Note: Reading a chain a byte at a time under Commodore GEOS involves finding

the first data block and passing its track/sector to **ReadFile**. The track/sector of the first data block in a sequential file is returned in **rl** by **GetFHdrlnfo**. The first data block of a VLIR record is contained

in the VLIR's index table.

Example:

See also: OpenDisk, SetDevice

# disk high-level

#### disk high level \_\_\_\_\_ Delete a file. DeleteFile C238 56 C22C EnterDeskTop Load and run DESKTOP. 57 \$C20B Lookup a file in the directory FindFile FindFTypes \$C23B Create a table of file names -401 GetFile \$C208 Load a file, given a file name -410 GetPtrCurDkNm \$C298 Compute address of disk's name -315 OpenDisk \$C2A1 Open a disk \$C259 Rename a file -405 RenameFile \$C23E Load the SWAPFILE \$C1ED Save memory to a file -409 RstrAppl SaveFile SetDevice \$C2BO Select a drive SetGEOSDisk C1EA Convert a disk to GEOS format

disk high-level

DeleteFile: (C64, C128) c238

Function: Delete a GEOS file by deleting the its directory entry and freeing all

its blocks. Works on both sequential and VLIR files.

Parameters: r0 FILENAME - pointer to null-terminated name of file to delete

Uses: curDrive

**curType** GEOS 64 v 1.3 and later for detecting REU shadowing.

**Returns:** x = (\$00 = no = rror).

Alters: diskBlkBuf used for temporary block storage

fileHeader temporary storage of index table when deleting a VLIR

file.

Written to Disk:

curDirHead BAM updated to reflect newly freed blocks.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

Destroys:

a, y, r0-r9.

Description: Given a null-terminated filename, DeleteFile will remove it from the

current directory by deleting its directory entry and calling FreeFile

to free all the blocks in the file.

DeleteFile first calls FindFile to get the directory entry and ensure the file does in fact exist. If the file specified with FILENAME is not

found, a FILE NOT FOUND error is returned.

The directory entry is deleted by setting its OFF CFILE TYPE byte to

\$00.

disk high-level

EnterDeskTop: (C64, C128) C22C

Function: Standard application exit to GEOS deskTop.

Parameters: none.

Returns: never returns to application.

Description: EnterDeskTop takes no parameters and looks for a copy of the file DESK

TOP on each drive. Later versions of GEOS are only compatible with the correspondingly later revision of the deskTop and will check the version number in the permanent name string of the DESK TOP file to ensure that it is in fact a newer version. If after all drives are searched no valid copy of the deskTop is found, **EnterDeskTop** will prompt the user to insert

a disk with a copy of the deskTop on it.

### Example:

See also: RstrAppl, GetFile.

**FindFile:** (C64, C128) C20B

Function: Search for a particular file in the current directory.

Parameters: r6 FILENAME - pointer to null-terminated name of file of a maximum

of 16 bytes (not counting null terminator). (word).

Uses: curDrive

curType GEOS 64 vl.3 and later for detecting REU shadowing

\$886E<sup>1</sup> Flag byte

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

r1 track/sector of directory block containing entry;
r5 Pointer to directory entry within diskBIkBuf.

Alters: dirEntryBuf directory entry of file if found.

diskBlkBuf contains directory block where FILENAME found.

Destroys: a,y,r4,r6

Description: Given a null-terminated filename, FindFile searches through the current

directory and returns the directory entry in **directryBuf**. If the file specified with **FILENAME** is not found, a **FILE NOT FOUND** error is returned.

Since Commodore GEOS 2.0 does not support a hierarchical file system, the current directory is actually the entire disk. The directory entry

is deleted by setting its OFF CFILE TYPE byte to \$00.

Note: 1 If the flag byte at \$886E is \$FF, then both drives 8 and 9 will be

scanned if necessary. If the flag is \$00, then the lookup is only to the current drive. If there is only one drive, then this flag has no effect.

Note: Wheels, Gateway and MP3 All support a hierarchical file system. As of

this writing, the Author or this document (PBM) does not yet know the details of this support. This section, (and many others I am sure) will

be updated when I have researched these systems.

ANote: (Note to Author: Confirm behavior of and Give/Get a Name if confirmed

the 886E Flag address)

Example: LoadBASIC

**FindFTypes:** (C64, C128) C23B

Function: Builds a list of files of a particular GEOS type from the current

directory.

Parameters: r6 BUFFER - pointer to buffer for building-out file list; allow

ENTRY SIZE+1 bytes for each entry in the list (word).

 ${f r10}$  PERMNAME — pointer to permanent name string to match or \$0000 to

ignore permanent name string (word).

r7H MAXFUJES - maximum number of filenames to return, usually used to

prevent overwriting buffer.

**r7L** FILETYPE - GEOS file type to search for (byte).

Uses: curDrive

curType GEOS 64 vl.3 and later for detecting REU shadowing

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

r7H decremented once for each file name (Apple GEOS: high-bit is always

cleared).

Alters: diskBlkBuf used as temporary buffer for directory blocks.

Destroys: a, y, rO-r2L, r4, r6.

Description: FindFTypes build a list of files that match a particular GEOS file type

and, optionally, a specific permanent name string.

The data area at BUFFER, where the list is built-out, must be large enough to accommodate MAXFILES filenames of ENTRY SIZE+1 bytes each.

FindFTypes first clears enough of the area at BUFFER to hold MAXFILES filenames then calls **GetlstDirEntry** and **GetNxtDirEntry** to go through each directory entry in the current directory. When the GEOS file type of a directory entry matches the FILETYPE parameter, **FindFTypes** goes on to check for a matching permanent name string.

If the PERMNAME parameter is \$0000, then this check is bypassed and the filename is added to the list. If the PERMNAME parameter is non-zero, the null terminated string it points to is checked, character-by-character, against the permanent name string in the file's header block. Although the permanent name string in the GEOS file header is 16 characters long, the comparison only extends to the character before the null-terminator in the string at PERMNAME.

Since permanent name strings typically end with Vx.x, where x.x is a version number (e.g., 2.1), a shorter string can be passed so that the specific version number is ignored. For example, a program called geoQuiz version 1.3 might use "geoQuiz V1.3" as the permanent name string it gives its data files. When geoQuiz version 3.0 goes searching for its data files, it can pass a PERMNAME string of "geoQuiz V" so data files for all versions of the program will be added to the list

When a match is found, the filename is copied into the list at BUFFER. The filenames are placed in the buffer as they are found (the same order they appear on the pages of the deskTop notepad). With a small buffer, matching files on higher-numbered pages may never get added to the list.

Note:

Since Commodore GEOS does not support a hierarchical file system, the "current directory" is actually the entire disk. The filenames appear in the list null terminated even though they are padded with a0 in the directory.

Example:

See also: FindFile, GetlstDirEntry, GetNxtDirEntry.

**GetFile:** (C64, C128) C208

**Function:** General-purpose file routine that can load an application, desk accessory, or data file.

Parameters: r6 FILENAME- pointer to null-terminated filename (word).

#### When loading an application:

#### **r0L** LOAD OPT:

- bit 0: 0 load at address specified in file header; application
   will be started automatically
  - 1 load at address in r7; application will not be started automatically.
- bit 7: 0 not passing a data file.
  - 1 r2 and r3 contain pointers to disk and data file names.
- bit 6: 0 not printing data file.
  - 1 printing data file; application should print file and exit
- r7 LOAD\_ADDR optional load address, only used if bit 0 of LOADJOPT
   is set (word).
- PATA\_DISK only valid if bit 7 or bit 6 of LOAD\_OPT is set:
   pointer to name of the disk that contains the data file, usually
   a pointer to one of the DrxCurDkNm buffers (word).
- PATA\_FILE only valid if bit 7 or bit 6 of LOAD\_OPT is set:
  pointer to name of the data file (word).

# When loading a desk accessory:

#### Uses: curDrive

curType GEOS 64 vl.3 and later for detecting REU shadowing

# Returns: When loading an application:

only returns if alternate load address or disk error.

x = rror (\$00 = no error).r0, r2, r3, and r7 unchanged.

# When loading a desk accessory:

returns when desk accessory exits with a call to RstrAppL x error (\$00 = no error).

# When loading a data file:

x = (\$00 = no = ror).

# Passes: When loading an application:

warmstarts GEOS and passes the following to the application

r0 as originally passed to GetFile.

 ${\bf r2}$  as originally passed to  ${\bf GetFile}$  (use dataDiskName).

r3 as originally passed to GetFile. (use dataFileName).

dataDiskName contains name of data disk if bit 7 of r0 is set
dataFileName contains name of data file if bit 6 of r0 is set

# When loading a desk accessory:

warmstarts  ${\tt GEOS}$  and passes the following:

rl0L as originally passed to GetFile.

When loading a data file: not applicable.

Alters: When loading an application:

GEOS brought to a warmstart state.

Destroys: a,x,y,r0-r10 (only applies to loading a data file).

Description: GetFile is the preferred method of loading most GEOS files, whether a data file, application, or desk accessory. (The only exception to this is a VLIR file, which is better handled with the VLIR routines such as

OpenRecordFile and ReadRecord). Most applications will use GetFile to load and execute desk accessories when the user clicks on an item in the geos menu. Some applications will use GetFile to load other applications. The GEOS deskTop, in fact, is just another application like any other. Depending on the user's choice of actions — open an application, open an application's data file, print an applications' data file — the deskTop sets LOAD OPT, DATA DISK, DATA FILE appropriately and calls GetFile.

GetFile first calls FindFile to locate the file at FILENAME, then checks the GEOS file type in the directory entry. If the file is type DESK ACC, then GetFile calls LdDeskAcc. If the file is type APPLICATION or type AUTO\_EXEC, GetFile calls LdApplic. All other file types are loaded with the generic LdFile.

The following GEOS constants can be used to set the  $LOAD\_OPT$  parameter when loading an application:

ST_LD_AT_ADDR	Load at address: load application at the address passed
	in <b>r7</b> as opposed to the address in the file header.
ST_LD_DATA	Load data file: application is being passed the name of a data file to load.
ST_PR_DATA	Print data file: application is being passed the name of a data file to print.

Note:

The RECVR\_OPTS flag used when loading desk accessories originally carried the following significance:

- bit 7: 1 force desk accessory to save foreground screen area and restore it on return to application.
  - 0 not necessary for desk accessory to save foreground.
- bit 6: 1 force desk accessory to save color memory and restore it on return to application.
  - 0 not necessary for desk accessory to save color memory.

The application should always set r10H to \$00 and bear the burden of saving and restoring the foreground screen and the color memory. (Color memory only applicable to GEOS 64 and GEOS 128 in 40-column mode.) See LdDeskAcc Note for more information.

### Example:

See also: LdFile, LdDeskAcc, LdApplic

disk high-level

**OpenDisk:** (C64, C128) C2A1

Function: Open the disk in the current drive

Parameters: None:

Uses: curDrive drive that disk is in. Set by call SetDevice

Calls: NewDisk, GetDirHead , ChkDkGEOS, GetPtrCurDkNm

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

r5 pointer to disk name buffer as returned from GetPtrCurDkNm. This

is a pointer to one of the DrxCurDkNm arrays.

Alters: DnxCurDkNm current disk name array contains disk name

isGEOS set to TRUE if disk is a GEOS disk, otherwise set to

FALSE.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

Destroys: a, y, r0-r4.

Description: OpenDisk initiates access to the disk in the current drive. OpenDisk is

meant to be called after a new disk has been inserted into the disk drive. It prepares the drive and disk variables for dealing with a new disk. An application will usually call **OpenDisk** immediately after calling

SetDevice

Note: Because GEOS uses the same allocation and file buffers for each drive,

it is important to close all files and update the BAM if necessary (use

PutDirHead) before accessing another disk.

OpenDisk first calls NewDisk to tell the disk drive a new disk has been inserted (if the disk is shadowed, the shadow memory is also cleared). GetDirHead is then called to load the disk's header block and BAM into curDirHead. With a valid header block in memory, ChkDkGEOS is called to check for the GEOS I.D. string and set the isGEOS flag to TRUE if the

disk is a GEOS disk. Finally, OpenDisk copies the disk name string from

curDirHead to the disk name buffer returned by GetPtrCurDkNm.

Note: This Routine calls GetDirHead which loads in the BAM from disk.

PutDirHead should be called prior to this routine if the BAM has been

modified by Freeing or allocating blocks.

disk high-level

**RenameFile:** (C64, C128) C259

Function: Renames a file that is in the current directory.

Parameters: r6 OLDNAME - pointer to null-terminated name of file as it appears on

the disk (word).

r0 NEWNAME - pointer to new null-terminated name (word).

Uses: curDrive drive that disk is in. Set by call SetDevice

Calls: NewDisk, GetDirHead , ChkDkGEOS, GetPtrCurDkNm

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Alters: diskBlkBuf used for temporary block storage.

curDirHead BAM updated to reflect newly freed blocks.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

Destroys: a, y, r4-r6.

Description: RenameFile searches the current directory for OLDFILE and changes the

name string in the directory entry to NEWFILE.

RenameFile first calls FindFile to get the directory entry and ensure the OLDFILE does in fact exist. (If it doesn't exist, a FILE NOT FOUND

error is returned.)

The directory entry is read in, the new file name is copied over the old

file name, and the directory entry is rewritten. The date stamp of the

file is not changed.

When using **GetlstDirEntry** and **GetNxtDirEntry** to establish the old file name, do not pass **RenameFile** a pointer into **diskBlkBuf**. Copy the file

name from **diskBlkBuf** to another buffer (such as **dirEntryBuf**) and pass **FreeFile** the pointer to that buffer. Otherwise when **FreeFile** uses

diskBlkBuf it will corrupt the file name.

Note3: This Routine calls FindFile which loads in the BAM in from disk. it is

important to close all VLIR files and update the BAM if necessary (use

PutDirHead) before using RenameFile.

Example:

See also: FreeFile, FreeBlock.

disk high-level

**SaveFile:** (C64, C128) \$C1ED

Function: create a GEOS sequential OR VLIR file and save a region of memory.

Parameters: r9 HEADER pointer to GEOS file header for file.

Required Offsets into Header to Set

.word \$00 = Pointer to Filename

.byte \$69 = DOS File Type

.byte \$70 = GEOS Filetype

.byte \$71 = GEOS Structure (SEQ or VLIR)
.word \$72 = Memory to Save Start Address

(Set to \$0000 for an empty file)

.word \$74 = Memory to Save End Address

(Set to \$FFFF for an empty file)

r10L DIRPAGE Directory page to begin searching for an empty directory

slot.

Uses: curDrive

year, month, day, hours, minutes for date-stamping file.
curType GEOS 64 vl.3 and later for detecting REU shadowing
interleave desired physical sector interleave (usually 8).

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

r1 Track and Sector of last block written

r9 Unchanged

r6 pointer to fileTrScTab

Alters: dirEntryBuf contains newly-built directory entry.

diskBlkBuf contains contents of last block written

fileTrScTab \$00-\$01 contain T/S of File Header.

End of Table is marked with Track=0

curDirHead BAM updated to reflect newly allocated block.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

Destroys: a,y, r0-r8

**Description:** SaveFile is the most general purpose write data routine in GEOS. It creates a new file, either sequential or VLIR with a Header Block. VLIR

files will have all of the memory written to Record 0 of the VLIR.

SaveFile calls SetGDirEntry and BlkAllOC to construct the file, then calls WriteFile to put the data into it. After the file is saved, the

BAM is written to disk

Note<sup>1,3</sup>: If the Start Address = \$0000 and the End Address = \$FFFF (Or if Start Address = End Address) no data blocks are written. A VLIR's VLIR block will have all empty records. An empty SEO file's directory entry will

will have all empty records. An empty SEQ file's directory entry will have a start T/S of 00/FF. (This is not a normal valid state for a SEQ  $\,$ 

file and should have at least one block added to it).

The HEADER holds all the information needed to create the file. All of the information listed as Required must be populated in the HEADER.

Example:

See also: GetFile, OpenRecordFile.

SetDevice: (C64, C128) C2B0

Function: Establish communication with a new peripheral

Parameters: a DEVNUM - 8,9,10,11 (DRIVE A through DRIVE D) for disk drives,

PRINTER for serial printer, or any other valid serial device bus

address (byte).

Uses: curDevice currently active device.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Alters: curDevice new current device number.

driveType table).

Destroys: a,y

**Description:** SetDevice changes the active device and is used primarily to switch from one disk drive to another. SetDevice also allows a printer driver to

gain access to the serial bus by using a DEVNUM value of PRINTER.

Each I/O device has an associated device number that distinguishes its I/O from devices. At any given time only one device is active. The active device is called the current device and to change the current device an application calls **SetDevice**.

**SetDevice** is designed to switch between serial bus devices, DEVNUM reflects the architecture of serial bus: disk drives are numbered 8 through 11 and the printer is numbered 4. However, not all I/O devices are actual serial bus peripherals. A RAMdisk, for example, uses a special device driver to make a cartridge port RAM-expansion unit emulate a Commodore disk drive. **SetDevice** switches between these devices just as if they were daisy chained off of the serial bus.

GEOS up through v1.2 supports two disk devices, DRIVE A and DRIVE B. Commodore GEOS v1.3 and later supports up to four disk devices, DRIVE $\sim$ A through DRIVE-D. Desktop Only Supports 3 Devices.

Note:

SetDevice calls ExitTurbo so that the old device is no longer actively sensing the serial bus, then installs the new device driver as necessary to make the new device (DEVNUM) the current device. With more than one type of device attached (e.g., a 1541 and a 1571), GEOS must switch the device drivers, making the driver for the selected device active. GEOS stores inactive device drivers in the Commodore 128 back RAM and in special system areas in an REU. GEOS applications must use SetDevice to change the active device. An application should never directly modify curDrive or curDevice.

SetGEOSDisk: (C64, C128) C1EA

Function: Convert Commodore disk to GEOS format.

Parameters: none.

Uses: curDrive

curType GEOS 64 vl.3 and later for detecting REU shadowing.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Alters: curDirHead directory header is read from disk.

dir2Head (BAM for 1571 and 1581 drives only)

dir3Head (BAM for 1581 drive only)

Destroys: a,y

Description: SetGEOSDisk converts a standard Commodore disk into GEOS format by

writing the GEOS ID string to the directory header (at OFF\_GEOS ID) and creating an off-page directory block. An application can call SetGEOSDisk after OpenDisk returns the isGEOS flag set to FALSE. Typically the user

is prompted before the conversion.

SetGEOSDisk expects the disk to have been previously opened with OpenDisk. It first calls GetDirHead to read the directory header into memory then calls CalcBlksFree to see if there is block available for the off-page directory (if there isn't, an INSUFFICIENT\_SPACE error is returned). SetNextFree is then called to allocate the off-page directory block. The off-page directory block is written with empty directory entries and a pointer to it is placed in the directory header (at OFF OP TR SC). Finally PutDirHead is called to write out the new BAM and

directory header.

AppendRecord	C289	Add a VLIR chain
CloseRecordFile	C277	Close a VLIR file
DeleteRecord	C283	Remove a VLIR chain
InsertRecord	C286	Insert a VLIR chain
NextRecord	C27A	Move to next VLIR chain
OpenRecordFile	C274	Open a VLIR file
PointRecord	C280	Go to a specific VLIR chain
PreviousRecord	C27D	Move to previous VLIR chain
ReadRecord	C28C	Load a VLIR chain
UpdateRecordFile	C295	Update a VLIR file
WriteRecord	C28F	Save memory to a VLIR chain

AppendRecord: (C64, C128) c289

Function: Adds an empty record after the current record in the index table, moving

all subsequent records down one slot to make room.

Parameters: none.

Uses: curDrive drive that disk is in. Set by call SetDevice

fileWritten if FALSE, assumes record just opened (or updated) and

reads BAM/VBM into memory. ANOTE: Confirm

curType GEOS 64 vl.3 and later for detecting REU shadowing

curDirHead BAM updated to reflect newly allocated block.

dir2Head<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head† (BAM for 1581 drive only)

'used internally by GEOS disk routines; applications generally don't use.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

OUT OF RECORDS

Alters: curRecord new record number

last updated.

fileHeader buffer contains VLIR index table.

note: When making manual changes to the VLIR setting fileWritten to TRUE

will cause CloseRecordFile to write the changes to disk.

Destroys: a,y, r0L, r1L, r4

Preparatory routines1: OpenRecordFile

Description: AppendRecord inserts an empty VLIR record following the current record in the index table of an open VLIR file, moving all subsequent records down in the record list. The new record becomes the current record. A VLIR file can have a up to MAX\_VLIR\_RECS records (127 on the Commodore and 254 on the Apple). If adding a Record exceeds this value, then an

**OUT OF RECORDS** error is returned.

A record added with **AppendRecord** occupies no disk space until data is written to it. The new record is marked as empty in the VLIR index table (\$00 \$FF). When a VLIR file is first created by SaveFile, all records are marked as unused (\$00 \$00). Some applications call **AppendRecord** repeatedly after creating a new file until an **OUT\_OF\_RECORDS** error is returned This marks all the records as used and prepares them to accept

data with calls to WriteRecord.

Note: AppendRecord does not write the VLIR index table out to the disk. Call

CloseRecordFile or UpdateRecordFile to save the index table when all

modifications are complete.

Note: Use PointRecord to check the status of a particular record (unused,

empty, or filled).

Example: SaveRecord

See also: InsertRecord, DeleteRecord, PointRecord

CloseRecordFile: (C64, C128) c277

Function: Close the current VLIR file (updating it in the process) so that another

may be opened

Parameters: none.

Uses: curDrive drive that disk is in. Set by call SetDevice

fileWritten if FALSE, assumes record just opened (or updated) and

reads BAM/VBM into memory. ANOTE: Confirm

fileHeader VLIR index table.

fileSize total number of disk blocks used in file (includes index

block, GEOS file header, and all records).

curType GEOS 64 vl.3 and later for detecting REU shadowing

curDirHead BAM updated to reflect newly allocated block.

dir2Head† (BAM for 1571 and 1581 drives only)

dir3Head† (BAM for 1581 drive only)

'used internally by GEOS disk routines; applications generally don't use.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Alters: fileWritten<sup>†</sup> set to TRUE to indicate the file has been altered since

last updated.

fileHeader buffer contains VLIR index table.

note: When making manual changes to the VLIR setting fileWritten to TRUE

will cause CloseRecordFile to write the changes to disk.

Destroys: a,y, r1, r4, r5.

Preparatory routines1: OpenRecordFile

Description: CloseRecordFile first calls UpdateRecordFile then closes the VLIR file

so that another may be opened.

Because Commodore GEOS stores the BAM in global memory, the application must be careful not to corrupt it before the VLIR file is updated or

closed. For more information, refer to UpdateRecordFile.

Example: SaveRecord

DeleteRecord: (C64, C128) C283

Function: Removes the current VLIR record from the record list, moving all

subsequent records upward to fill the slot and freeing all the data

blocks associated with the record.

Parameters: none.

Uses: curDrive drive that disk is in. Set by call SetDevice

fileWritten<sup>†</sup> if FALSE, assumes record just opened (or updated) and

reads BAM into memory.

curRecord Current record number

fileHeader VLIR index table.

curDirHead BAM updated to reflect newly allocated block.

dir2Head† (BAM for 1571 and 1581 drives only)

dir3Head<sup>†</sup> (BAM for 1581 drive only)

'used internally by GEOS disk routines; applications generally don't use.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

Alters: curRecord only changed if deleting the last record in the table, in

which case it becomes the new last record.

fileWritten set to TRUE to indicate the file has been altered since

last updated.

fileHeader Record Marked as empty (\$00 \$FF)

dir2Head† (BAM for 1571 and 1581 drives only)

dir3Head† (BAM for 1581 drive only)

Destroys: a,y, r0-r9

Description: DeleteRecord removes the current record from the record list by moving

all subsequent records upward to fill the current record's slot. Any

data blocks associated with the record are freed.

DeleteRecord does not update the BAM and VLIR file information on the

disk. Call CloseRecordFile or UpdateRecordFile to update the file when

done modifying.

InsertRecord: (C64, C128) C286

Function: Adds an empty record before the current record in the index table, moving

all subsequent records (including the current record) downward.

Parameters: none.

Uses: curDrive drive that disk is in. Set by call SetDevice

fileWritten if FALSE, assumes record just opened (or updated) and

reads BAM/VBM into memory.

curRecordCurrent record numberfileHeaderVLIR index table.

dir2Head<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head<sup>†</sup> (BAM for 1581 drive only)

'used internally by GEOS disk routines; applications generally don't use.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

MAX\_VLIR\_RECS OUT OF RECORDS

Alters: curRecord new record number

fileWritten' set to TRUE to indicate the file has been altered since

last updated.

fileHeader buffer contains VLIR index table.

usedRecords number of records in file that are currently in use.

Destroys: a,y, r0L

Preparatory routines1: OpenRecordFile

Description: InsertRecord attempts to insert an empty VLIR record in front of the current record in the index table 7of an open VLIR file, moving all subsequent records downward in the record list. The new record becomes the current record. A VLIR file can have a maximum of MAX\_VLIR\_RECS records. If adding a record will exceed this value, an OUT\_OF\_RECORDS

error is returned. In the index table, the new record is marked as used but empty  $(\$00,\$FF)^{1}$ .

Note: An application can create an empty VLIR file with SaveFile.

Note: GEOS up to 2.0 does not support a hierarchical file system, the

"current directory" is actually the entire disk.

Note: 3 This Routine calls GetDirHead which loads in the BAM from disk.

PutDirHead should be called prior to this routine if the BAM has been

modified by Freeing or allocating blocks.

NextRecord: (C64, C128) C27A

Function: Makes the next record the current record.

Parameters: none

Uses: fileHeader index table checked to establish whether record exists.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

INV RECORD

Y Track of VLIR Chain. A value of \$00 here means record is allocated

but not in use (has no data blocks).

a new current record number.

r1L Track of VLIR Chain
r1H Sector of VLIR Chain

Alters: curRecord new record number

Destroys: nothing

Preparatory routines1: OpenRecordFile

Description: NextRecord makes the current record plus one the new current record. A

subsequent call to ReadRecord or WriteRecord will operate with this

record.

If the record does not exist, then NextRecord returns an then NextRecord

returns an INV RECORD (invalid record) error.

OpenRecordFile: (C64, C128) C274

Function: Open an existing VLIR file for access.

Parameters: r0 FILENAME-pointer to null-terminated name of file (word).

Uses: curDrive drive that disk is in. Set by call SetDevice

curType GEOS 64 vl.3 and later for detecting REU shadowing

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

STRUCT MISMATCH

rlL track/Sector of VLIR Block

r1H

r5 pointer into diskBlkBuf to start of directory entry.

Alters: fileHeader buffer contains VLIR index table.

usedRecords number of records in file that are currently in use.

fileWritten set to FALSE to indicate VLIR file has not been written

to.

fileSize total number of disk blocks used in file (includes

index block, GEOS file header, and all records).

Destroys: a,y,r1,r4-r6

Preparatory routines: SetDevice, OpenDisk

Description: Before accessing the data in a VLIR file, an application must call

OpenRecordFile. OpenRecordFile searches the current directory for FILENAME and, if it finds it, loads the index table into fileHeader. OpenRecordFile initializes the GEOS VLIR variables (both local and global) to allow other VLIR routines such as WriteRecord and ReadRecord

to access

the file. Only one VLIR file may be open at a time. A previously opened

VLIR file should be closed before opening another.

If an application passes a FILENAME of a non-VLIR file, OpenRecordFile

will return a **STRUCT MISMATCH** error.

Note: An application can create an empty VLIR file with SaveFile.

Note: GEOS up to 2.0 does not support a hierarchical file system, the

"current directory" is actually the entire disk.

Note: 3 This Routine calls **GetDirHead** which loads in the BAM from disk.

PutDirHead should be called prior to this routine if the BAM has been

modified by Freeing or allocating blocks.

PointRecord: (C64, C128) C280

Function: Make a particular record the current record.

Parameters: a RECORD - record number to make current (byte).

Uses: fileHeader index table checked to establish whether record exists.

usedRecords Number of Currently Used Records in the VLIR file

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

y Track of VLIR Chain. A value of \$00 here means record is allocated

but not in use (has no data blocks).

a new current record number.

r1L Track of VLIR Chain
r1H Sector of VLIR Chain

Alters: curRecord new record number

Destroys: nothing

Preparatory routines1: OpenRecordFile

Description: PointRecord makes RECORD the current record so that a subsequent call to

ReadRecord or WriteRecord will operate with RECORD. VLIR records are

numbered zero through MAX\_VLIR\_RECS-1.

If the record does not exist (you pass a record number that is larger

than the number of currently used records), then PointRecord returns an

INV RECORD (invalid record) error.

PreviousRecord: (C64, C128) C27D

Function: Makes the previous record the current record.

Parameters: none

Uses: fileHeader index table checked to establish whether record exists.

**Returns:**  $\mathbf{x}$  error (\$00 = no error).

INV RECORD

Y Track of VLIR Chain. A value of \$00 here means record is allocated

but not in use (has no data blocks).

a new current record number.

r1L Track of VLIR Chain
r1H Sector of VLIR Chain

Alters: curRecord new record number

Destroys: nothing

Preparatory routines1: OpenRecordFile

Description: PreviousRecord makes the current record minus one the new current record.

A subsequent call to **ReadRecord** or **WriteRecord** will operate with this

record.

If the record does not exist, then PreviousRecord returns an INV RECORD

(invalid record) error.

ReadRecord: C28C (C64, C128)

Function: Read in the current VLIR record.

Parameters: r7 BUFFER - pointer to start buffer where data will be read into (word).

r2 BUFSIZE - size of buffer: Commodore version can read up to 32,258 bytes (127 Commodore blocks); Apple version can read up to the maximum two-byte number that can be passed in r2: 65,535 (\$FFFF) bytes (word).

drive that disk is in. Set by call SetDevice Uses: curDrive

> curRecord Current record number

VLIR index table. Table holds Track / Sector of first fileHeader

block of each record.

curType GEOS 64 vl.3 and later for detecting REU shadowing

error (\$00 = no error). Returns: Х

> \$00 = empty record, no data read. \$ff = record contained data.

pointer to last byte read into BUFFER plus one if not an empty record, otherwise unchanged.

rl if BFR OVERFLOW error returned, contains the track/sector of the block that, had it been copied from diskBlkBuf to the application's buffer space, would have exceeded the size of BUFFER. The process of copying any extra data from diskBlkBuf to the end of BUFFER is left to the application. The data starts at diskBlkBuf+2. If no

error, then **rl** is destroyed

fileTrScTab As the chain blocks in the record is followed, the Alters: track/sector pointer of each block is added to the file track/sector table. The track and sector of the first block in the record is added at

fiieTrScTab+2 and fileTrScTab+3. Refer to ReadFile for more information.

y, (r1), r2-r4 (see above for r1) Destroys:

Preparatory routines1: OpenRecordFile

Description: ReadRecord reads the current record into memory at BUFFER. If the record

contains more than BUFSIZE bytes of data, then a BFR OVERFLOW error is

returned.

ReadRecord calls ReadFile to load the chain of blocks into memory.

UpdateRecordFile: (C64, C128)

Function: Update the disk copy of the VLIR index table, BAM and other VLIR

information such as the file's time/date-stamp. This update only takes

place if the file has changed since opened or last updated.

Parameters: none.

Uses: curDrive drive that disk is in. Set by call SetDevice

fileWritten
 if FALSE, assumes record just opened (or updated) and
fileHeader
 VLIR index table. Table holds Track / Sector of first

block of each record.

fileSize total number of disk blocks used in file (includes

index block, GEOS file header, and all records).

year, month, day, hours, minutes for date-stamping file.

curType GEOS 64 vl.3 and later for detecting REU shadowing

curDirHead BAM updated to reflect newly allocated block.

dir2Head† (BAM for 1571 and 1581 drives only)

dir3Head<sup>†</sup> (BAM for 1581 drive only)

'used internally by GEOS disk routines; applications generally don't use.

**Returns:** x = error (\$00 = no error).

Alters: fileWritten set to FALSE to indicate that file hasn't been altered

since last updated

Destroys: a, y, r1, r4, r5

Description: UpdateRecordFile checks the fileWritten flag. If the flag is TRUE, which

indicates the file has been altered since it was last updated, **UpdateRecordFile** writes the various tables kept in memory out to disk (e.g., index table, BAM) and time/date-stamps the directory entry. If

the fileWritten flag is FALSE, it does nothing.

UpdateRecordFile writes out the index block, adds the time/date-stamp and fileSize information to the directory entry, and writes out the new

BAM with a call to PutDirHead.

Because Commodore GEOS stores the BAM in global memory, the application must be careful not to corrupt it before the VLIR file is updated. If the **fileWritten** flag is **TRUE** and the BAM is reread from disk, the old copy (on disk) will overwrite the current copy in memory. In the normal use of VLIR disk routines, where a file is opened, altered, then closed before any other disk routines are executed, no conflicts will arise.

WriteRecord: (C64, C128) C28F

Function: Write data to the current VLER record.

Parameters: none.

Uses: curDrive drive that disk is in. Set by call SetDevice

fileWritten if FALSE, assumes record just opened (or updated) and

curType GEOS 64 vl.3 and later for detecting REU shadowing

curDirHead BAM updated to reflect newly allocated block.

dir2Head<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head<sup>†</sup> (BAM for 1581 drive only)

\*used internally by GEOS disk routines; applications generally don't use.

**Returns:** x = error (\$00 = no error).

Alters: fileWritten set to FALSE to indicate that file hasn't been altered

since last updated

fileHeader index table adjusted to point to new chain of blocks

for current record.

fileSize adjusted to reflect new size of file.

fileTrScTab Contains track/sector table for record as returned from

**BlkAlloC.** The track and sector of the first block in the record is at fileTrScTab+0 and fileTrScTab+1 The end of the table is marked with a track value of \$00.

curDirHead BAM updated to reflect newly freed and allocated

blocks.

dir2Head<sup>†</sup> (BAM for 1571 and 1581 drives only)

dir3Head<sup>†</sup> (BAM for 1581 drive only)

Destroys: a, y, r1, r4, r5

Description: WriteRecord writes data to the current record All blocks previously

associated with the record are freed. **BlkAllOC** is then used to allocate enough new blocks to hold BYTES amount of data The data is then written to the chain of sectors by calling **WriteFile**. The fileSize variable is

updated to reflect the new size of the file.

WriteRecord does not write the BAM and internal VLIR file information to disk. Call CloseRecordFile or UpdateRecordFile when done to update the

disk with this information.

Note: WriteRecord correctly handles the case where the number of bytes to write

(BYTES, R2) is zero. The record is freed and marked as allocated but not

in use.

AddDirBlock	9039	Needs Documenting
CallDrvRoutine:	9042	Needs Documenting
CheckDrvStatus:	9045	Needs Documenting
Get <b>diskBlkBuf</b>	903C	Needs Documenting
Put <b>diskBlkBuf</b>	903F	Needs Documenting
<u>UpdateRecordFile</u> JmpI	ndX:	9D80 Jump Table Needs Documenting

**AddDirBlock:** (C64, C128) 9042

Function: Call Directly into Disk Driver

Parameters:

Uses:

Returns:

Alters:

Destroys:

Description:

Note: Needs Documenting...

 CallDrvRoutine:
 (C64, C128)
 9042

Function: Call Directly into Disk Drive

Parameters:

Uses:

Returns:

Alters:

Destroys:

Description:

Note: Needs Documenting...

GetdiskBlkBuf:	(C64, C128)	9045

Function: Call Directly into Disk Driver

Parameters:

Uses:

Returns:

Alters:

Destroys:

Description:

Note: Needs Documenting...

PutdiskBlkBuf: (C64, C128)

9045

Function: Call Directly into Disk Driver

Parameters:

Uses:

Returns:

Alters:

Destroys:

Description:

Note: Needs Documenting...

CheckDrvStatus: (C64, C128) 9045

Function: Call Directly into Disk Driver

Parameters:

Uses:

Returns:

Alters:

Destroys:

Description:

Note: Needs Documenting...

UpdateRecordFileJmpIndX: (C64, C128)

9D80

Function: Jump Table

Parameters:

Uses:

Returns:

Alters:

Destroys:

Description:

Note: Needs Documenting...

# graphics

BitmapUp	\$C1AB	Draw a click box	-621
i BitmapUp	\$C1AB	Draw a click box / inline	-622
 BitmapClip	\$C2AA	Draw a coded image	-623
BitOtherClip	\$C2C5	Draw a coded image with user patches	-624
DisablSprite	\$C1D5	Turn off a sprite	-628
DrawLine	\$C130	Draw/Erase/Copy an arbitrary line	-615
DrawPoint	\$C133	Draw/Erase/Copy a point on the screen	89
DrawSprite	\$C1C6	Copy a sprite data block	-625
EnablSprite	\$C1D2	Turn on a sprite	-627
FrameRectangle	\$C127	Draw an outline in a pattern	-604
i_FrameRectangle	\$C1A2	Draw a solid outline with inline data	-605
GetScanLine	\$C13C	Compute memory address of a row on the screen	-618
GraphicsString	\$C136	Process a graphic command table	-601
i_GraphicsString	\$C1A8	Process a graphic command table / inline	-602
HorizontalLine	\$C118	Draw a horizontal line in a pattern	-616
InvertLine	\$C11B	Reverse video a horizontal line	-614
ImprintRectangle	\$C250	Copy a box from screen 2 to screen 1	-610
i_ImprintRectangle	\$C253	Copy a box from screen 2 to screen 1 / inline	-611
InvertRectangle	\$C12A	Reverse video a box	-612
PosSprite	\$C1CF	Position a sprite	-626
RecoverLine		Copy a line from screen 2 to screen 1	-613
Rectangle		Fill a box with a pattern	-606
i_Rectangle	\$C19F	Fill a box with a pattern / inline	-607
RecoverRectangle	\$C12D	Copy a box from screen 1 to screen 2	-608
i_RecoverRectangle	\$C1A5	Copy a box from screen 1 to screen 2 / inline	-609
SetNewMode		Set GEOS 128 to 40 or 80 column mode	90
SetPattern		Select a fill pattern	-603
TestPoint	\$C13F	Test the value of a pixel	-619
VerticalLine	\$C121	Draw a vertical line in a pattern	-617

graphics

**DrawPoint:** (C64, C128) \$**c133** 

Function: Set, clear, or recover a single screen point (pixel).

Parameters: r3 XI - x-coordinate of pixel (word).

**r11L** Y1 - y-coordinate of pixel (byte).

st MODE:

N C Operation

1 X recover pixel from background screen to foreground

0 1 Draw set pixel using dispBufferOn.
0 0 Erase clear pixel using dispBufferOn.

Uses: when setting or clearing pixels:

### dispBuffer0n:

bit 7 — write to foreground screen if set. bit 6 — write to background screen if set.

### Destroys:

a, x, y, r5-r6

#### Description:

**DrawPoint** is a very versatile routine. It can copy a point from one screen to another, as well as draw or erase it. This routine is called by DrawLine to draw lines.

The carry (c) flag and sign (n) flag in the processor status register (s) are used to pass information to **DrawPoint**. The following tricks can be used to set or clear these flags appropriately:

- Use sec and clc to set or clear the carry (c) flag.
- Use lda # to set the sign (n) flag.
- Use 1da #0 to clear the sign (n) flag.
- 128: DOUBLE W, ADD1 W can be used on r3

graphics

SetNewMode: (C128) \$C2DD

Function: Set 128 mode to 40 or 80 column mode.

**Uses:** graphMode GRMODE — new graphics mode to change to:

40-Column: GR\_40 80-Column: GR 80

Returns: nothing.

Destroys:

a, x, y, r0-r15

Description: SetNewMode the Operation mode of the Commodore 128.

### 40-column mode (graphMode == GR 40)

- 1: 8510 clock speed is slowed down to lMhz because VIC chip cannot operate at 2Mhz.
- 2: rightMargin is set to 319.
- 3: UseSystemFont is called to begin using the 40-column font.
- 4: 40-column VIC screen is enabled.
- 5: 80-column VDC is set to black on black, effectively disabling it.

### 80-column mode (graphMode == GR 80)

- 1: 8510 clock speed is raised to 2Mhz.
- 2: rightMargin is set to 639.
- 3: UseSystemFont is called to begin using the 80-column font.
- 4: 40-column VIC screen is disabled.
- 5: 80-column VDC screen is enabled.

Example: Change Mode

# icon/menu

Dolcons	C15A	Display and activate an icon table.	92
DoMenu	\$C151	Menu processor	1-701
DoPreviousMenu	\$C190	Close current menu	1-702
GotoFirstMenu	\$C1BD	Close all menu levels	1-703
RecoverAllMenus	\$C157	Erase all menus	1-706
RecoverMenu	\$C154	Erase the current menu	1-705
Re <b>DoMenu</b>	\$C193	Draw the current menu	1-704

icon/menu\_

**DoIcons:** (C64, C128) C15A

Function: Display and activate an icon table.

Parameters: r0 ICONTABLE - pointer to the icon table to use.

Uses: dispBufferOn:

bit 7 - draw icons to foreground screen if set. bit 6 - draw icons to background screen if set.

Destroys: rO-rl5, a, x, y

Description: Dolcons takes an icon, draws the enabled icons (those whose OFF\_I\_PIC word is non-zero) and instructs MainLoop to begin tracking the user's interaction with the icons. This routine is the only way to install icons. Every application should install at least one icon, even if only a dummy icon.

If **DoIcons** is called while another icon table is active, the new icons will take precedence. The old icons are not erased from the screen before the new ones are displayed.

**DoIcons** is a complex routine which affects a lot of system variables and tables. The following is an outline of its major actions:

- 1: All enabled icons in the table are drawn to the foreground screen and/or the background buffer based on the value in **dispBufferOn**.
- 2: StartMouseMode is called. If the OFF\_IC XMOUSE word of the icon table header is non-zero, then StartMouseMode loads mouseXPosition and mouseYposition with the values in the OFF\_IC\_XMOUSE and the OFF\_IC\_YMOUSE parameters of the icon table header (see StartMouseMode for more information).
- 4: faultData is cleared to \$00, indicating no faults.
- 5: If the MOUSEON\_BIT of mouseOn is clear, then the MENUON BIT is forced to one. This is because GEOS assumes that it is in a power-up state and that mouse tracking should be fully enabled. If the MOUSEON\_BIT bit is set, GEOS leaves the menu-scan alone, assuming that the current state of the MENUON BIT is valid.
- 6: The ICONSON\_BIT and MOUSEON\_BIT bits of mouseOn are set thereby enabling icon-scanning.

When an icon event handler is given control, **rOL** contains the number of the icon clicked on (beginning with zero) and **rOH** contains **TRUE** if the event is a double-click or FALSE if the event is a single click.

•

Example: IconsUp

icon/menu

**DoMenu:** (C64, C128) C15A

Function: Display and activate a menu structure.

Parameters: r0 MENU - pointer to the menu structure to display.

a POINTER\_OVER — which menu item (numbered starting with zero) to center the pointer over.

Destroys: a, x, y, ro-r13

**Description:** DoMenu draws the main menu (the first menu in the menu structure) and instructs MainLoop to begin tracking the user's interaction with the menu. This routine is the only way to install a menu.

If **DoMenu** is called while another menu structure is active, the new menu will take precedence. The old menu is not erased from the screen before the new menu is displayed. If the new menu is smaller (or at a different position) than the old menu, parts of the old menu may be left on the screen. A typical way to avoid this is to erase the old menu with a call to Rectangle, passing the positions of the main menu rectangle and drawing in a white pattern. However, a more elegant solution involves calling **RecoverAllMenus**, which will erase any extant menus by recovering from the background buffer.

**DoMenu** is a complex routine which affects a lot of system variables and tables. The following is an outline of its major actions:

- 1: Menu level 0 (main menu) is drawn to the foreground screen.
- 2: StartMouseMode is called. mouseXPosition and mouseYposition are set so that the pointer is centered over the selection number passed in a. Under Apple GEOS, if the CallRoutine POINTER\_OVER number in the accumulator has its high-bit set, then the mouse will not be repositioned Under GEOS 64 and GEOS 128, DoMenu always forces the mouse to a new position. If you do not want the mouse moved, surround the call to DoMenu with code to save and restore the mouse positions. The following code fragment will install menus without moving the mouse.

```
DoMenu2:
   php
                              ; Save Processor Status Register
    sei
                              ; disable interrupts around call
    PushW
             mouseXPos
                             ; save mouse x
    PushB
             mouseYPos
                             ; save mouse y
    lda
             #0
                             ; dummy menu value
                            ; install menus (mouse will move)
    jsr
             DoMenu
    PopB
             \begin{tabular}{ll} \textbf{mouseYPos} & \textbf{;} & \textbf{restore original mouse y} \\ \end{tabular}
    PopW
             mouseXPos
                            ; restore original mouse x
    plp
                              ; Restore Interrupts to their saved state
    rts
```

**3: SlowMouse** is called. With a joystick this will kill all accumulated speed in the pointer, requiring the user to reaccelerate. With a proportional mouse, this will have no effect.

- 4: faultData is cleared to \$00, indicating no faults.
- 5: If the MOUSEON\_BIT of **mouseOn** is clear, then the ICONSON\_BIT is forced to one. This is because GEOS assumes that it is in a power-up state and that mouse tracking should be fully enabled. If the MOUSEON\_BIT bit is set, GEOS leaves the icon-scan alone, assuming that the ICONSON\_BIT is valid.
- 6: The MENUON\_BIT and MOUSEON\_BIT bits of mouseOn are set, thereby enabling menu-scanning.
- 7: The mouse fault variables (mouseTop, mouseBottom, mouseLeft, and mouseRight) are set to the full screen dimensions.

# input driver

InitMouse	\$fe80	Initialize input device.
SetMouse	\$fe89	Reset input device scanning circuitry.
<b>SetMouse</b> Pic	\$c2da	Set and preshift new soft-sprite mouse picture.
SlowMouse	\$fe83	Reset mouse velocity variables.
UpdateMouse	\$fe86	Update mouse variables from input device.

#### GEOS Kernal

input driver

InitMouse: (c64,C128)

Function: Initialize the input device.

Parameters: none.

Returns: nothing.

Alters: mouseXPos initialized (typically 8).

mouseYPos initialized (typically 8).

mouseData initialized (typically reflects a released button).

pressFlag initialized (typically set to \$00).

Destroys: assume a,x,y,r0-r15

Description: GEOS calls InitMouse after first loading an input driver. The input

driver is expected to initialize itself and begin tracking the input

device. An application should never need to call InitMouse.

Example:

See also: SlowMouse, UpdateMouse, SetMouse, StartMouseMode, MouseUp.

#### GEOS Kernal

SetMouse: input driver

(C128)

FE89

Function: Input device scan reset.

Parameters: none.

Returns: nothing.

Destroys: assume a,x,y,r0-r15

Description: GEOS 128 calls SetMouse during Interrupt Level, immediately after the

keyboard is scanned for a new key, to reset the pot (potentiometer) scanning lines so that they will recharge with the new value of. It is primarily of use with the Commodore 1351 mouse, which requires having the pot lines reset regularly. Other input drivers will have a **SetMouse** routine that merely performs an rts. An application should never need to

call SetMouse.

Example:

See also: SlowMouse, UpdateMouse, Initmouse.

### GEOS Kernal

internal

### internal

C000	Reboot GEOS.	99
\$C271	Initialize GEOS variables.	1-803
C196	Return GEOS serial number.	101
\$C100	Main interrupt level processing.	1-805
\$C1C3	GEOS's main loop.	1-804
\$C2C2	Report system errors.	1-711
\$C003	internal Bootstrap entry point	
	\$C271 C196 \$C100 \$C1C3 \$C2C2	C000 Reboot GEOS. \$C271 Initialize GEOS variables. C196 Return GEOS serial number. \$C100 Main interrupt level processing. \$C1C3 GEOS's main loop. \$C2C2 Report system errors. \$C003 internal Bootstrap entry point

internal

**BootGeos:** (c64,C128) c000

Function: Restart GEOS from a non-GEOS application.

Parameters: none.

Returns: Does not return.

**Destroys:** n/a

Description: BootGeos provides a method for an non-GEOS to run in the GEOS

environment—starting up from the deskTop and returning to GEOS when done. The non-GEOS application need only preserve the area of memory between  ${\bf BootGeos}$  (\$C000) and  ${\bf BootGeos}+\$7f$  (\$C07f). The rest of the GEOS Kernal may be overwritten. To reboot GEOS, simply jmp  ${\bf BootGeos}$ , which completely reloads the operating system (either from disk in a "boot11 procedure or from a RAM-expansion unit in an "rboot11 procedure) and returns to the

GEOS deskTop.

A program can check to see if it was loaded by GEOS by checking the memory starting at \$c006 (bootName) for the ASCII (not CBMASCII) string "GEOSBOOT". If loaded by GEOS, the program can check bit 5 of \$c012 (sysFlgCopy): if this bit is clear, ask the user to insert their GEOS boot disk before continuing, otherwise a boot disk is not needed because GEOS will rboot from the RAM expansion unit. To actually return to GEOS, set CPU\_DATA to \$37 (KRNL BAS 10 IN) on a Commodore 64 and set config to \$40 (CKRNL BAS IO IN) on a Commodore 128, then jump to BootGeos

Example: RoadTrip

internal

**FirstInit:** (c64, C128) **c271** 

Function: Simulates portions of the GEOS coldstart procedure without actually

rebooting GEOS or destroying the application in memory.

Parameters: none.

Returns: GEOS variables and system hardware in a coldstart state; stack and

application space unaffected.

Destroys: a, x, y, rO-r2

Description: FirstInit is part of the GEOS coldstart procedure. It initializes nearly

all GEOS variables and data structures (both global and local), including those which are usually only done once, when GEOS is first booted, such as setting the configuration variables to a default, power-up state.

GEOS calls this routine internally. Applications will not find it

especially useful.

Note: The GEOS font variables are not reset by FirstInit; a call to

UseSystemFont may be necessary.

internal

GetSerialNumber: (c64,C128) c196

**Function:** Return the 16-bit serial number or pointer to the serial string for

the current GEOS kernal

Parameters: none.

Returns: r0 16-bit serial number

**Destroys:** a

 $\textbf{Description: GetSerialNumber} \ \ \text{gives an application access to an unencrypted copy of} \\$ 

the GEOS serial number or serial string for comparison purposes. You cannot change the actual serial string or number by altering this copy.

math

III.a LII
-----------

BBMult	C160	Byte by byte (single-precision) unsigned multiply.	103
Bmult	C163	Unsigned 16 bit by 8 bit multiply	104
Dabs	C16F	16 bit absolute value	105
Ddec	C175	Decrement a 16 bit integer	106
Ddiv	C169	Unsigned 16 bit division	107
DMult	C166	Unsigned 16 bit by 16 bit multiply	109
Dnegate	C172	Negate a 16 bit integer	110
DSDiv	C16C	Signed 16 bit division	111
DShiftLeft	C15D	Multiple 16 bit arithmetic shift left	112
DShiftRight	C262	Multiple 16 bit logical shift right	113

<u>math</u>

**BBMult:** (c64,C128) C160

Function: Unsigned byte-by-byte multiply: multiplies two unsigned byte operands

to produce an unsigned word result.

low-byte of a word variable (byte pointer to a word variable).

y OPERAND2 - zero-page address of the byte multiplier (byte pointer

to a byte variable).

Note: result = OPERAND1 (word) \* OPERAND2 (word).

Returns: x, y, and byte pointed to by OPERAND2 unchanged. word pointed to by

OPERAND1 contains the word result.

Destroys: a,r7L,r8

Description: BBMult is an unsigned byte-by-byte multiplication routine that

multiplies two bytes to produce a 16-bit word result (low/high order). The byte in OPERAND1 is multiplied by the byte in OPERAND2 and the result is stored as a word back in OPERAND1. Note OPERAND1 starts out as a byte parameter but becomes a word result with the high-byte at OPERAND 7+1.

Note: Because  ${\tt r7}$  and  ${\tt r8}$  are destroyed in the multiplication process, they

cannot be used to hold either operand.

No overflow can occur when multiplying two bytes because the result

always fits in a word(ff\*f = fe01).

Example: 8BitMultiply

math\_

**BMult:** (c64, C128) **C163** 

Function: Unsigned word-by-byte multiply: multiplies an unsigned word and an

unsigned byte to produce an unsigned word result

to word variable).

y OPERAND2 - zero-page address of multiplier (byte pointer to a

word variable — use a word variable; only the low-byte is used in the multiplication process, but the high-byte of the word is

destroyed).

Note: result = OPERAND1 (word) \* OPERAND2 (byte).

Returns: x, y unchanged.

word pointed to by OPERAND2 has its high-byte set to \$00, and its low-

byte unchanged word pointed to by OPERAND1 contains the word result.

Destroys: a, r6-r8.

Description: BMult is an unsigned word-by-byte multiplication routine that multiplies

the word at one zero-page address by the byte at another to produce a 16-bit word result. **Bmult** operates by clearing the high-byte of OPERAND2 and calling **Bmult**. The result is stored as a word back in OPERAND 1.

Note: r6, r7 and r8 are destroyed in the multiplication process, they cannot

be used to hold the operands.

Overflow in the result (beyond 16-bits) is ignored.

Example: 16x8Multiply, ConvToUnits

math

Compute absolute value of a two's-complement signed word

(c64,C128)

 ${\tt OPERAND}$  - zero-page address of word to operate on (byte Parameters: X

pointer to a word variable).

Returns: unchanged. x,y

word pointed to by OPERAND contains the absolute value result.

Destroys:

Dabs:

Function:

Description: Dabs takes a signed word at a zero-page address and returns its absolute

value. The address of the word (OPERAND) is passed in x. The absolute

C16F

value of OPERAND is returned in OPERAND.

The equation involved is: if (value < 0) then value = -value.

Example: DSmult

See also: **DNegate**  <u>math</u>

**Ddec:** (c64,C128) C175

Function: Decrement a word

Parameters: X OPERAND - zero-page address of word to decrement (byte pointer

to a word variable).

Returns: x,y unchanged.

st z flag is set if resulting word is \$0000.

zero page word pointed to by OPERAND contains the decremented word.

Destroys: a

Description: Ddec is a double-precision routine that decrements a 16-bit zero-page

word. The absolute address of the word is passed in x. If the result of the decrement is zero, then the z flag in the status register is set and can be tested with a subsequent beq or bne. **Ddec** is useful for loops

which require a two-byte counter.

Note<sup>3</sup>: the Macro **DecW** should be used in cases where speed is more important

then code size. Inner loops should always use **DecW** if space allows. **Ddec** should be used when space is at a premium as it is only cost only 5 bytes to use. The kernal uses **Ddec** in **CRC** because space in the kernal is more valuable then the speed of the **CRC** procedure that is not normally ever

used in an inner loop. See Example DdecvsDecW

Example: Kernal\_CRC, DdecvsDecW, DecCounter

math\_

**Ddiv:** (c64,C128) **C169** 

Function: Unsigned word-by-word (double-precision) division: divides one unsigned

word by another to produce an unsigned word result.

Parameters: x OPERAND1 — zero-page address of word dividend (byte

pointer to a word variable).

y OPERAND2 - zero-page address of word divisor (byte

pointer to a word variable).

Note: result = OPERAND1 (word) / OPERAND2 (word).

Returns: x,y and word pointed to by OPERAND2 unchanged,

word pointed to by OPERAND1 contains the result **r8** contains the fractional remainder (word).

Destroys: a,r9

Description: Ddiv is an unsigned word-by-word division routine that divides the word

at one zero-page address (the dividend) by the word at another (the divisor) to produce a 16-bit word result and a 16-bit word fractional remainder The word in OPERAND 1 is divided by the word in OPERAND2 and the result is stored as a word back in OPERAND1. The remainder is returned

in **r8.** 

Note: Because r8 and r9 are used in the division process, they cannot be used

to hold operands.

If the divisor (OPERAND2) is greater than the dividend (OPERAND1), then the fractional result will be returned as \$0000 and OPERAND1 will be

returned in r8.

Although dividing by zero is an undefined mathematical operation, **Ddiv** makes no attempt to flag this as an error condition and will simply return incorrect results. If the divisor might be zero, the application should check for this situation before dividing as in:

```
lda zpage,y ; get low byte of divisor
ora zpage+l,y ; get high byte of divisor
```

bne 10\$; if either non-zero, go divide

jmp DivideByZero ; else, flag error

10\$

jmp Ddiv

There is no possibility of overflow (a result which cannot fit in 16 bits).

Example: ConvToUnits, CheckDiskSpace

See also: DSDiv, DMult, BBMult, BMult

math\_

# DivideBySeven: (Apple)

**Function:** Divide a byte value by 7

Parameters: r0L OPERAND1 - byte to divide/7

Returns: a result

Destroys: a

Description: Bonus Code Page CBM GEOS has no DivideBySeven in the Kernal like Apple.

So here is a block to do a similar operation on an 8 bit value

DvBy7:

lda **r0L** 

lsr

lsr

lsr

adc rOL

ror

lsr

lsr

adc rOL

ror

lsr

lsr

rts

<u>math</u>

Function: Unsigned word-by-word (double-precision) multiply: multiplies two

unsigned words to produce an unsigned word result.

Parameters:  $\mathbf{x}$  OPERAND 1 — zero-page address of word multiplicand (byte

pointer to a word variable).

y OPERAND2 - zero-page address of word multiplier (byte pointer

C166

to a word variable).

(c64, C128)

Note: results OPERAND1 (word) \* OPERAND2 (word).

Returns: x,y, word pointed to by OPERAND2 unchanged

word pointed to by OPERAND contains the word result.

Destroys: a, r6-r8

DMult:

Description: DMult is an unsigned word-by-word multiplication routine that multiplies

the word at one zero-page address by the word at another to produce a 16-bit word result (all stored in low/high order). The word in OPERAND1 is multiplied by the word in OPERAND2 and the result is stored as a word

back in OPERAND1.

Note: Because r6, r7 and r8 are destroyed in the multiplication process, they

cannot be used to hold the operands.

Overflow in the result (beyond 16-bits) is ignored.

Example: DSmult

math\_

**Dnegate:** (c64, C128) **C172** 

Function: Negate a signed word (two's complement sign-switch).

Parameters: X OPERAND - zero-page address of word to operate on (byte

pointer to a word variable).

Returns: x,y unchanged.

Destroys: a

Description: Dnegate negates a zero-page word. The absolute address of the word

OPERAND) is passed in x. The absolute value of OPERAND is returned in

OPERAND.

The operation of this routine is: value = (value A \$FFFF) + 1.

Example: DSmult

See also: Dabs

 $\mathtt{math}_{}$ 

**DSDiv:** (c64,C128) C16C

Function: Signed word-by-word (double-precision) division: divides one two's

complement

Parameters: X OPERAND1 - zero-page address of signed word dividend (byte

pointer to a word variable).

y OPERAND2 - zero-page address of signed word divisor (byte pointer

to a word variable).

Returns: x,y unchanged.

r8 the fractional remainder (word).

word pointed to by OPERAND2 equals its absolute value. word pointed to by OPERAND 1 contains the word result.

Destroys: a,r9

Description: DSDiv is a signed, two's complement word-by-word division routine that

divides the word in one zero-page pseudo register (the dividend) by the word in another (the divisor) to produce a 16-bit word signed result and a 16-bit word fractional remainder The word in OPERAND 1 is divided by the word in OPERAND2 and the result is stored as a word back in OPERAND1

with the remainder in r8.

The remainder is always positive regardless of the sign of the dividend. This will cause problems with some mathematical operations that expect a signed remainder. The following code fragment will fix this problem:

Note: Because r8 and r9 are used in the division process, they cannot be used

as the operands.

Although dividing by zero is an undefined mathematical operation, **DSDiv** makes no attempt to flag this as an error condition and simply returns incorrect results. If the divisor might be zero, the application should check for this situation before dividing:

check for this situation before dividing

zpage = \$00

DSDivPre:

lda zpage,y
ora zpage+1,y
bne 10\$; get low byte of divisor
;get high byte of divisor
;if either non-zero, go divide

jmp DivideByZero ;else, flag error

10\$

jmp DSDiv ;divide

math\_

DShiftLeft: (c64,C128) C15D

Function: Arithmetically left-shift a zero-page word.

Parameters: X OPERAND - address of the zero-page word to shift (byte pointer

to a word variable).

Y COUNT - number of times to shift the word left (byte).

Returns: a,y unchanged.

**y** #\$ff

st c (carry flag) is set with last bit shifted out of word.

zero page address pointed to by OPERAND contains the shifted word.

Destroys: nothing

Description: DShiftLeft is a double-precision math routine that arithmetically left-

shifts a 16-bit zero-page word (low/high order). The address of the word is passed in  ${\bf x}$  and the number of times to shift the word is passed in

y. Zeros are shifted into the low-order bit.

An arithmetic left-shift is useful for quickly multiplying a value by a power of two. One left-shift will multiply by two, two left-shifts will multiply by four, three left-shifts will multiply by eight, and so on:

value = value \* 2count.

Note: If a COUNT of \$00 is specified, the word will not be shifted.

Carry Flag <- High Byte <- Low Byte

C 7-6-5-4-3-2-1-0 7-6-5-4-3-2-1-0 <- 0

math\_

DShiftRight: (c64,C128) c262

Function: Arithmetically right-shift a zero-page word.

Parameters: X OPERAND - address of the zero-page word to shift (byte pointer

to a word variable).

Y COUNT - number of times to shift the word right (byte).

Returns: a,x unchanged.

**y** #\$ff

st c (carry flag) is set with last bit shifted out of word.

zero page address pointed to by OPERAND contains the shifted word.

Destroys: nothing

Description: DShiftRight is a double-precision math routine that arithmetically

right-shifts a 16-bit zero-page word (low/high order). The address of the word is passed in x and the number of times to shift the word is

passed in y. Zeros are shifted into the high-order bit.

An arithmetic left-shift is useful for quickly multiplying a value by a power of two. One left-shift will multiply by two, two left-shifts will multiply by four, three left-shifts will multiply by eight, and so on:

value = value \* 2count.

Note: If a COUNT of \$00 is specified, the word will not be shifted.

High Byte -> Low Byte -> Carry Flag

0 -> 7-6-5-4-3-2-1-0 7-6-5-4-3-2-1-0 C

Example: MseToCardPos, ConvToUnits

# GEOS Kernal

memory

m	no	-	. 7

ClearRam	\$C178	Fill a memory region with zeroes	115
CmpFString	\$C26E	Memory block comparison	118
CmpString	\$C26B	String compare	117
CopyFString	\$C268	String copy	118
CopyString	\$C265	Memory block move	119
DoBOp	\$C2EC	(128) Back-RAM memory action primitive	120
DoRAMOp	\$C2D4	Perform any of the Below REU commands	127
FetchRAM	C2CB	Retrieve memory from an REU	128
FillRam	\$C17B	Memory block fill	121
$\mathtt{i}_{\texttt{FillRam}}$	\$C1B4	Memory block fill with inline data	121
InitRam	\$C181	Multiple memory location initialization	122
MoveBData	\$C2E3	(128) Move data between front and back RAM	123
MoveData	\$C17E	Intelligent memory block move	124
i_MoveData	\$C1B7	Inline Version inline data	124
StashRAM	\$C2C8	Stash memory into an REU	129
SwapBData	\$C2E6	(128) Swaps memory between Front and Back RAM	125
SwapRAM	\$C2CE	Swap memory with an REU memory block	130
VerifyBData	\$C2E9	(128) Compares two regions of memory	126
VerifyRAM	\$C2D1	Verify (compare) memory with REU	131

\$C187

Function: Clear a region of memory to \$00.

Parameters: ClearRam COUNT - number of bytes to clear (0 - 64K) (word).

(C64,C128)

r1 ADDR - address of area to clear (word).

Returns: nothing.

ClearRam:

Destroys: a,y, r0,r1,r2L

Description: ClearRam clears COUNT bytes starting at ADDR to ADDRESS+COUNT. It useful

for initializing .ramsect variable and data sections.

Note: Do not use ClearRam to initialize r0-r2L.

Note: Also, for when space is at a premium, it actually takes fewer bytes to

call  $i_Fill_{Ram}$  with a fill value of \$00.

Note: 1 ClearRam sets r2L to \$00 and calls FillRam.

Example: InitBuffers

See also: FillRam, InitRam.

 CmpFString:
 (C64,C128)
 \$c26E

Function: Compare two fixed-length strings.

**Parameters:**  $\mathbf{x}$  SOURCE — zero-page address of pointer to source string (byte pointer to a word pointer).

p DEST - zero-page address of pointer to destination string (byte pointer to a word pointer).

a LEN - length of strings (1-255). A LEN value of \$00 will cause CmpFString to function exactly like CmpString, expecting a null terminated source string.

Returns: st status register flags reflect the result of the comparison.

Destroys: a,x,y

**Description: CmpFString** compares the fixed-length string pointed to by SOURCE to the string of the same length pointed to by DEST.

CmpFString with a LEN value of \$00 causes the routine to act exactly like CmpString.

**CmpFString** compares each character in the strings until there is a non-matching pair. The result of the comparison between the non-matching pair is passed back in the processor status register (st). If the strings match, the z flag is set. This allows the application to test the result of a string comparison with standard test and branch operations:

bne branch if strings don't match

beq branch if strings match

bcc branch if source string is less than DEST string

Note: The strings may contain internal NULL'S. These will not terminate the

comparison.

Example: Find

**CmpString**: (C64,C128) \$**c26**B

Function: Compare two null-terminated strings.

string.

y DEST - zero-page address of pointer to destination null

terminated string.

Returns: st status register flags reflect the result of the comparison.

Destroys: a,x,y

Description: CmpString compares the null-terminated source string pointed to by SOURCE

to the destination string pointed to by DEST. The strings are compared a byte at a time until either a mismatch is found or a null is encountered  $\frac{1}{2}$ 

in both strings.

**CmpString** compares each character in the strings until there is a non-matching pair. The result of the comparison between the non-matching pair is passed back in the processor status register (st). If the strings match, the z flag is set. This allows the application to test the result of a string comparison with standard test and branch operations:

bne branch if strings don't match

beg branch if strings match

bcs branch if source string is greater than or equal to DEST string

bcc branch if source string is less than DEST string

Note: CmpString cannot compare strings longer than 256 bytes (including the

null). The compare process is aborted after 256 bytes.

Example: Find2

**CopyFString**: (C64, C128) \$C268

Function: Copy a fixed-length string.

pointer to a word pointer).

y DEST - zero-page address of pointer to destination buffer (byte

pointer to a word pointer).

**a** LEN - length of strings (1-255)

Returns: Buffer pointed to by DEST contains copy of source string.

Destroys: a,x,y

Description: CopyFString copies a fixed-length string pointed to by SOURCE to the

buffer pointed to by DEST If the source and destination areas overlap,

the source must be lower in memory for the copy to work properly.

Note: Because the LEN parameter is a one-byte value, CopyFString cannot copy

a string longer than 255 bytes. A LEN value of \$00 causes CopyFString

to act exactly like CopyString.

Note: The source string may contain internal NULL'S. These will not terminate

the copy operation.

Example: CopyBuffer

See also: CmpFString, CopyString.

See also: CopyString, CmpFString, MoveData.

**CopyString:** (C64,C128) \$**c268** 

Function: Copy a null-terminated string.

Parameters: x SOURCE - zero-page address of pointer to a NULL terminated

source string (byte pointer to a word pointer).

y DEST - zero-page address of pointer to destination buffer (byte

pointer to a word pointer).

Returns: Buffer pointed to by DEST contains copy of source string, including

the terminating NULL

Destroys: a,x,y

Description: CopyString copies a null terminated string pointed to by SOURCE to the

buffer pointed to by DEST. All Characters in the string are copied, including the null-terminator. If the source and destination areas overlap, the source must be lower in memory for the copy to work properly.

CopyString cannot copy more than 256 bytes. The Copy process is aborted

after 256 bytes.

Note: NULL terminated Strings can be an arbitrary size including > 255

Example: CopyBuffer

**DoBOp:** (C128) \$**c2EC** 

Function: Primitive for communicating with REU (RAM-Expansion Unit) devices.

Parameters: r0 ADDR1 - address of first block in application memory (word).

r1 ADDR2 - address of second block in application memory (word).

r2 COUNT-number of bytes to operate on (word).

r3L A1BANK-ADDR1 bank: 0 = front RAM; 1 = back RAM (byte).

r3H A2BANK-ADDR2 bank: 0 = front RAM; 1 \* back RAM (byte).

**y** MODE - operation mode:

#### b1 b0 Description

0	0	move from memory at ADDR1 to memory at ADDR2.
0	1	move from memory SAADDR2 to memory at ADDR1.
1	0	swap memory at ADDR1 with memory at ADDR2.
1	1	verify (compare) memory at ADDR1 against memory at ADDR2.

Note: the **DoBOp** MODE parameter closely matches the low nibble of the Do RAM Op CMD parameter.

Returns: r0-r3 unchanged.

When verifying:

x \$00 if data matches; \$ff if mismatch.

DEV\_NOT\_FOUND if bank or REU not available.

Destroys: a,x,y

Description: DoBOp is a generalized memory primitive for dealing with both memory

banks on the Commodore 128. It is used by MoveBData, SwapBData, and

VerifyBData.

Note: DoBOp should only be used on designated application areas of memory. When

moving memory within the same bank the destination address must be less than source address. When swapping memory within the same bank, ADDR1

must be less than ADDR2.

FillRam: ,I FillRam: (C64,C128) \$C17B,\$C1B4

Function: Fills a region of memory with a repeating byte value.

Parameters: Normal:

r0 COUNT - number of bytes to clear (0 - 64K) (word).

rl ADDR - address of area to clear (word).
r2L FILL - byte value to fill with (byte).

Inline:

.word COUNT - number of bytes to clear (0 - 64K) (word).

.word ADDR - address of area to clear (word).
.byte FILL - byte value to fill with (byte).

Returns: r2L unchanged.

Destroys: a,y, r0,r1

Description: FillRam fills COUNT bytes starting at ADDR with the FILL byte. This

routine is useful for initializing a block of memory to any desired

value.

Note: Do not use FillRam to initialize r0-r2L.

Note:

Example: InitBuffers

See also: ClearRam, InitRam.

InitRam: (C64, C128) \$C181

Function: Table driven initialization for variable space and other memory areas.

Parameters: r0 TABLE -address of initialization table (word).

Returns: nothing.

Destroys: a,x,y, r0-r2L

Description: InitRam uses a table of data to initialize blocks of memory to preset

values. It is useful for setting groups of variables to specific values. It is especially good at initializing a group of noncontiguous variables

in a "two bytes here, three bytes there" fashion.

The initialization table that is pointed to by the TABLE parameter is a data structure made up from the following repeating pattern:

```
.word address ;start address of this block
.byte count ;number of bytes to initialize
```

.byte byte1,byte2,... ; count bytes of data

.word address ;start address of next block

. . .

The table is made of blocks that follow the above pattern, count bytes starting at address are initialized with the next count bytes in the table. (A count value of \$00 is treated as 256.) To end the table, use

.word NULL

where InitRam expects the next address parameter.

Note: Do not use InitRam to initialize r0-r2L.

MoveBData: (C128) \$C2E3

Function: Special version of MoveData that will move data within either front RAM

or back RAM (or from one bank to the other).

Parameters: r0 SOURCE - address of source block in memory (word)

r2 COUNT - number of bytes to move (word)

r3L SRCBANK - source bank: 0 = back RAM; 1 = front RAM (byte)

r3H DSTBANK - destination bank: 0 = back RAM; 1 = front RAM (byte)

Returns: r0-r3 unchanged.

Destroys: a,x,y

Description: MoveBData is a block move routine that allows data to be moved in either

front RAM, back RAM, or between front and back (bank 1, the front bank, is the normal GEOS application area). If the SOURCE and DEST areas are

in the same bank and overlap, DEST. must be less than SOURCE.

MoveBData is especially useful for copying data from front RAM to back

RAM or from back RAM to front RAM.

MoveBData uses the DoBOp primitive by calling it with a MODE parameter

of \$00.

Note: MoveBData should only be used to move data within the designated

application areas of memory. MoveBData is significantly slower than MoveData and should be avoided if the move will occur entirely within

front RAM.

Example:

See also: MoveBData, SwapBData, VerifyBData, DoBOp.

MoveData: ,I MoveData: (C64,C128)

\$C2E3

Function: Moves a block data from one area to another.

Parameters: Normal:

r0 SOURCE - address of source block in memory (word)
r1 DEST - address of destination block in memory (word)

r2 COUNT - number of bytes to move (word)

Inline:

.word SOURCE
.word DEST
.word COUNT

Returns: r0,R1,R2 unchanged.

Destroys: a,y

Description: MoveData will move data from one area of memory to another. The source and destination blocks can overlap in either direction, which makes this routine ideal for scrolling, insertion sorts, and other applications that need to move arbitrarily large areas of memory around. The move is actually a copy in the sense that the source data remains unaltered

unless the destination area overlaps it.

64 & 128: If the DMA MoveData option in the Configure program is enabled (GEOS vl.3 and later), MoveData will use part of bank 0 of the installed RAM-expansion unit for an ultrafast move operation. An application that calls

MoveData in the normal manner will automatically take advantage of this selection. An application that relies upon a slower MoveData (for timing or other reasons) can disable the DMA-move by temporarily clearing bit 7 of sysRAMFlg. This bit can also be used to read the status of the DMA-

move configuration.

Due to insufficient error checking in GEOS, do not attempt to move more than 30,976 (\$7900) bytes at one time when the DMA-move option is

enabled. Break the move up into multiple calls to MoveData.

Due to insufficient error checking in GEOS, do not attempt to move more than 14,592 (\$3900) bytes at one time when the DMA-move option is enabled. Break the move up into multiple calls to MoveData. MoveData

should only be used to move data within the standard front RAM application space. Use MoveData to move memory within back RAM or between

front RAM and back RAM.

Because the RAM-expansion unit DMA follows the VIC chip bank select, an application that is displaying a 40-column screen from back RAM must either disable DMA-moves or temporarily switch the VIC chip to front RAM before the **MoveData** call. (Note to **Author:** Confirm information here. I see no reason REU would not function equally well with back ram. And also be able to do ultra fast transfers between front to back. Using

stash ram, bank switch, then fetch ram. Needs testing.)

Note: Do not use MoveData on r0-r6.

Example:

See also: MoveBData, CopyString.

SwapBData: (C128) \$c2E6

Function: Swaps two regions of memory within either front RAM or back RAM (or

between one bank and the other).

Parameters: r0 ADDR1 - address of first block in application memory (word).

r1 ADDR2 - address of second block in application memory (word).

r2 COUNT - number of bytes to swap (word).

r3L AlBANK - ADDR1 bank: 0 = front RAM; 1 = back RAM (byte)

r3H A2BANK - ADDR2 bank: 0 = front RAM; 1 = back RAM (byte).

Returns: r0-r3 unchanged.

Destroys: a,x, y

Description:

**SwapBData** is a block swap routine that allows data to be swapped in either front RAM, back RAM, or between front and back. If the ADDR1 and ADDR2 areas are in the same bank and overlap, ADDR2. must be less than

ADDR1.

SwapBData is especially useful for swapping data from front RAM to back

RAM or from back RAM to front RAM.

SwapBData uses the DoBOp primitive by calling it with a MODE parameter

of \$02.

Note: SwapBData should only be used to swap data within the designated

application areas of memory.

VerifyBData: (C128) \$c2E9

Function: Compares (verifies) two regions of memory against each other. The

regions may either be in front RAM or back RAM (or one in front and the

other in back).

Parameters: r0 ADDR1 - address of first block in application memory (word).

r1 ADDR2 - address of second block in application memory (word).

r2 COUNT - number of bytes to swap (word).

r3L A1BANK - ADDR1 bank: 0 = front RAM; 1 = back RAM (byte)r3H A2BANK - ADDR2 bank: 0 = front RAM; 1 = back RAM (byte).

Returns: r0-r3 unchanged.

x \$00 if data matches; \$FF if mismatch.

Destroys: a, y

Description: VerifyBData is a block verify routine that allows the data in one region

of memory to be compared to the data in another region in memory. The regions may be in either front RAM, back RAM, or in front and back. The ADDR1 and ADDR2 areas may overlap even if they are in the same bank.

VerifyBData uses the DoBOp primitive by calling it with a MODE parameter

of \$03.

Note: VerifyBData should only be used to compare data within the designated

application areas of memory.

**DoRAMOp:** (c64 v1.3+,C128) \$C2D4

Function: Primitive for communicating with REU (RAM-Expansion Unit) devices.

r1 REUDEST - address in REU bank (word).

r2 COUNT - number of bytes to operate with (word).

r3L REUBANK - REU bank number to use (byte).
y CMD - command to send to REU (byte).

Returns: r0-r3 unchanged.

x error code: \$00 (no error) or

DEV NOT FOUND if bank or REU not available.

a REU status byte and ed with \$60

Destroys: y

Description: DoRAMOp is a very low-level routine for communicating with a RAM

expansion unit on a C64 or C128. This routine is a "use at your own risk"

GEOS primitive

**DoRAMOp** operates with the with the RAM-expansion unit directly and handles all the necessary communication protocols and clock-speed save/restore (if necessary).

The CMD parameter is stuffed into the REC Command Register (EXP\_BASE+\$01). Although **DoRAMOp** does no error checking on this parameter, it expects the high-nibble to be \$1001 (transfer with current configuration and disable FFOO decoding). The lower nibble can be one of the following:

%00 Transfer from Commodore to REU.

%10 Swap. %11 Verify.

Note: the low nibble of the  ${\it DoRAMOp}$  CMD parameter closely matches the

DoBOp MODE parameter.

Note: Note: On a Commodore 128, if the VIC chip is mapped to front RAM (with

the MMU VIC bank pointer), the REU will read/write using front RAM. Similarly, if the VIC chip is mapped to back RAM, the REU will read/write using back RAM. The REU ignores the standard bank selection controls on

the 8510. GEOS 128 defaults with the VIC mapped to front RAM.

For more information on the Commodore REU devices, refer to the Commodore 1764 RAM Expansion Module User's Guide or the 170011750 RAM Expansion

Module User's Guide.

**FetchRAM:** (c64 v1.3+,C128) **C2CB** 

Function: Primitive for transferring data from an REU

Parameters: r0 CBMDEST - address in Main Memory to start writing (word).

r1 REUSRC - address in REU bank to start reading (word).

r2 COUNT - number of bytes to fetch (word).

r3L REUBANK - REU bank number to fetch from (byte)

Returns: r0-r3 unchanged.

x error code: \$00 (no error) or

 ${\tt DEV\_NOT\_FOUND}$  if bank or REU not available.  ${\tt a}$  REU status byte and ed with \$60 (\$40 = success).

Destroys: y

Description: FetchRAM moves a block of data from a REU BANK into Commodore memory.

FetchRAM uses the DoRAMOp primitive by calling it with a CMD parameter

of %10010001. \$91

Note: Refer to DoRAMOp for notes and warnings.

Example:

See also: StashRAM, SwapRAM, VerifyRAM, DoRAMOp, MoveBData

**StashRAM:** (c64 v1.3+,C128) \$c2c8

Function: Primitive for transferring data to an REU

Parameters: r0 CBMSRC - address in Main Memory to start reading (word).

r1 REUDEST - address in REU bank to stash data (word).

Returns: r0-r3 unchanged.

x error code: \$00 (no error) or

 ${\tt DEV\_NOT\_FOUND}$  if bank or REU not available.  ${\tt a}$  REU status byte and ed with \$60 (\$40 = success).

Destroys: y

Description: StashRAM moves a block of data from Commodore memory into an REU bank.

This routine is a "use at your own risk" low-level GEOS primitive

StashRAM uses the DoRAMOp primitive by calling it with a CMD parameter

of %10010000. \$90

Note: Refer to DoRAMOp for notes and warnings.

Example:

See also: SwapRAM, FetchRAM, VerifyRAM, DoRAMOp, MoveBData

**SwapRAM:** (c64 v1.3+,C128) \$c2CE

Function: Primitive for swapping data between Commodore memory and an REU.

Parameters: r0 CBMADDR - address in Commodore to swap (word).

r1 REUADDR - address in REU to swap (word).
r2 COUNT - number of bytes to swap (word).

r3L REUBANK - REU bank number to fetch from (byte).

Returns: r0-r3 unchanged.

x error code: \$00 (no error) or

DEV NOT FOUND if bank or REU not available.

a REU status byte and'ed with \$60 (\$40 = successful swap).

Destroys: y

Description: SwapRAM swaps a block of data in an REU bank with a block of data in

Commodore memory.

SwapRAM uses the DorAMOp primitive by calling it with a CMD parameter

of % 10010010. \$92

Note: Refer to DoRAMOp for notes and warnings.

Example:

See also: StashRAM, FetchRAM, VerifyRAM, DoRAMOp, SwapBData

**VerifyRAM:** (c64 v1.3+,C128) \$**c2D1** 

Function: Verify (compare) data in main memory with data in an REU.

Parameters: r0 CBMADDR - address in Commodore to start (word)

r1 REUADDR - address in REU bank to start (word).
r2 COUNT - number of bytes to verify (word)
r3L REUBANK - REU bank number to verify with (byte).

Returns: r0-r3 unchanged.

x error code: \$00 (no error) or

DEV NOT FOUND if bank or REU not available.

a REU status byte and'ed with \$60

\$40 = data match \$20 = data mismatch

Destroys: y

Description: VerifyRAM Compares a block of data in Commodore memory with a block of

data in an REU bank to Verify the contents match. If bit 5 of the a register is set, there was an failed comparison during validation.

VerifyRAM uses the DoRAMOp primitive by calling it with a CMD parameter

of % 10010011. \$93

Note: Refer to DoRAMOp for notes and warnings.

# GEOS Kernal

mouse/sprite

# mouse/sprite

ClearMouseMode	\$C19C	Reset the mouse	1-815
IsMseInRegion	\$C2B3	Check if mouse is inside a window	1-710
MouseUp	\$C18A	Turn on the mouse	1-813
MouseOff	\$C18D	Turn off the mouse	1-814
SetMsePic	C2DA	Set and preshift new soft-sprite mouse picture.	133
StartMouseMode	\$C14E	Initialize the mouse	1-812
TempHideMouse	\$C2D7	Hide soft-sprites before direct screen access.	134

mouse/sprite

SetMsePic: \$C184 (C64, C128)

Uploads and pre-shifts a new mouse picture for the software sprite Function:

handler.

Parameters: r0 MSEPIC - pointer to 32 bytes of mouse sprite image data or one of

the following special codes:

ARROW

Returns: nothing

Destroys: a, x, y, r0-r15

Description: The software sprite routines used by GEOS 128 in 80-column mode treat the mouse sprite (sprite #0) differently than the other sprites. Sprite#0 is optimized and hardcoded to provide reasonable mouse-response while minimizing the flicker typically associated with erasing and redrawing a fastmoving object. The mouse sprite is limited to a 16x8 pixel image. The image includes a mask of the same size and both are stored in a pre-shifted form within internal GEOS buffers. For these reasons, a new mouse picture must be installed with SetMsePic (as opposed to a normal DrawSprite). SetMsePic pre-shifts the image data and lets the soft-sprite mouse routine know of the new image.

> SetMsePic accepts one parameter: a pointer to the mask and image data or a constant value for one of the predefined shapes. If a user-defined shape is used, the data that MSEPIC points to is in the following format:

16 bytes	16x8 "cookie cutter" mask. Before drawing the software			
	mouse sprite, GEOS and's this mask onto the foreground			
	screen. Any zero bits in the mask, clear the			
	corresponding pixels. One bits do not affect the screen.			
16 bytes	16x8 sprite image. After clearing pixels with the mask			
	data, the sprite image is or'ed into the area. Any one			
	bits in the sprite image set the corresponding pixels.			
	Zero bits do not affect the screen.			

Note: SetMsePic calls HideOnlyMouse. ARROW Equate - ARROW = \$00 Note<sup>3</sup>:

mouse/sprite

TempHideMouse: (C64, C128) \$c184

Function: Temporarily removes soft-sprites and the mouse pointer from the graphics

screen.

Parameters: nothing.

Uses: graphMode

Destroys: a, x

Description: TempHideMouse temporarily removes all soft-sprites (mouse pointer and

sprites 2-7) unless they are already removed. This routine is called by all GEOS graphics routines prior to drawing to the graphics screen so that software sprites don't interfere with the graphic operations. An application that needs to do direct screen access should call this

routine prior to modifying screen memory.

The sprites will remain hidden until the next pass through MainLoop.

Note: In 40-column mode (bit 7 of graphMode is zero), TempHideMouse exits

immediately without affecting the hardware sprites.

Example:

See also: HideOnlyMouse

print driver

print driver

**StartASCII:** (C64, C128) **7912** 

Function: Enable ASCII text mode printing. 7912

Parameters: rl WORKBUF - pointer to a 640-byte work buffer for use by the printer

driver. (word). PrintASCII uses this work area as an intermediate buffer,

the buffer must stay intact throughout the entire page.

Returns: x STATUS - printer error code; \$00 = no error.

Destroys: a, y, r0-r15

Description: StartASCII enables ASCII text mode printing. An application calls

StartASCII at the beginning of each page. It assumes that InitForPrint

has already been called to initialize the printer.

**StartASCII** takes control of the serial bus by opening a fake Commodore file structure and requests the printer (device 4) to enter listen mode. It then sends the proper control sequences to place the printer into

text mode.

Example:

See also: PrintASCII, StopPrint, StartPrint

# GEOS Kernal

process

pro	~~	~~

BlockProcess	C10C	Prevent a recurring timed event from running	138
EnableProcess	C109	Force a recurring timed event to run	141
FreezeProcess	C112	Stop a recurring timed event's timer	139
InitProcesses	C103	Initialize a table of recurring timed events	140
RestartProcess	C106	Enable a recurring timed event	142
Sleep	C199	Set up a time delay	143
UnblockProcess	C10F	Allow a recurring timed event to execute	143
UnfreezeProcess	C115	Start a recurring timed event's timer	145

BlockProcess: (C64, C128) C103

Function: Block a processes events.

Parameters: x PROCESS - process to block (0 to n-1, where n is the number of

processes in the table) (byte).

Returns: x unchanged.

**Destroys:** a

Description: BlockProcess causes MainLoop to ignore the runnable flag of a particular

process so that if a process timer reaches zero (causing the process to become runnable) no process event is generated until the process is subsequently unblocked with a call to <code>UnblockProcess</code>. <code>BlockProcess</code> stops the process the <code>MainLoop</code> level. Refer to <code>FreezeProcess</code> to stop the

process at the Interrupt Level.

BlockProcess does not stop the countdown timer, which continues to decrement at Interrupt Level (assuming the process is not frozen). When the timer reaches zero, the runnable flag is set and the timer is restarted. As long as the process is blocked, though, MainLoop ignores this runnable flag and, therefore, never generates an event. When a blocked process is later unblocked, MainLoop checks the runnable flag. If the runnable flag was set during the time the process was blocked, this pending event generates a call to the appropriate service routine. Only one event is generated when a process is unblocked, even if the

timer reached zero more than once.

Note: If a process is already blocked, a redundant call to BlockProcess has

no effect.

Example:

SuspendClock:

ldx #CLOCK\_PROCESS ; x <- process number of the clock
jmp BlockProcess ; block that particular process</pre>

FreezeProcess: (C64, C128) C112

Function: Freeze a process's countdown timer at its current value.

Parameters: x PROCESS - process to freeze (0 to n-1, where n is the number of

processes in the table) (byte).

Returns: x unchanged.

**Destroys:** a

Description: FreezeProcess halts a process's countdown timer so that it is no longer

decremented every vblank. Because a frozen timer will never reach zero, the process will not become runnable except through a call to **EnableProcess**. When a process is unfrozen with **UnfreezeProcess**, its

timer again begins counting from the point where it was frozen.

Note: If a process is already frozen, a redundant call to FreezeProcess has

no effect.

Example:

See also: UnfreezeProcess, BlockProcess

InitProcesses: (C64, C128) c103

Function: Initialize and install a process data structure.

Parameters: a NUM PROC - number of processes in table (byte).

**r0 PTABLE** - pointer to process data structure to use (word).

Returns: r0 unchanged.

Destroys: a, x, y, rl

Description: InitProcesses installs and initializes a process data structure. All

processes begin as frozen, so their timers are not decremented during vblank. Processes can be started individually with **RestartProcess** after

the call to InitProcesses.

InitProcesses copies the process data structure into an internal area of memory hidden from the application. GEOS maintains the processes within this internal area, keeping track of the event routine addresses, the timer initialization values (used to reload the timers after they time-out), the current value of the timer, and the state of each process (i.e., frozen, blocked, runnable). The application's copy of the process data structure is no longer needed because GEOS remembers this information until a subsequent call to InitProcesses.

Note:

Although processes are numbered starting with zero, NUM\_PROC should be the actual number of processes in the table. To initialize a process table with four processes, pass a  $NUM\_PROC$  value of \$04. When referring to those processes (i.e., when calling routines such as UnblockProcess), use the values \$00-\$03. Do not call InitProcesses with a  $NUM\_PROC$  value of \$00 or a  $NUM\_PROC$  value greater than  $MAX\_PROCESSES$  (the maximum number of processes allowable).

EnableProcess: (C64, C128) c109

Function: Makes a process runnable immediately.

Parameters: x PROCESS - process to enable (0 - n-l, where n is the number of

processes in the table) (byte).

Returns: x unchanged.

**Destroys:** a

Description: EnableProcess forces a process to become runnable on the next pass

through MainLoop, independent of its timer value.

**EnableProcess** merely sets the runnable flag in the process table. When **MainLoop** encounters an unblocked process with this flag set, it will attempt to generate an event just as if the timer had decremented to zero.

**EnableProcess** has no privileged status and cannot override a blocked process. However, because it doesn't depend on or affect the current timer value, the process can become runnable even with a frozen timer.

EnableProcess is useful for making sure a process runs at least once, regardless of the initialized value of the countdown timer. It is also useful for creating application-defined events which run off of MainLoop: a special process can be reserved in the data structure but never started with RestartProcess. Any time the desired event-state is detected, a call to EnableProcess will generate an event on the next pass through MainLoop. EnableProcess can be called from Interrupt Level, which allows a condition to be detected at Interrupt Level but processed during MainLoop.

RestartProcess: (C64, C128) C106

Function: Reset a process's timer to its starting value then unblock and unfreeze

the process.

Parameters: x PROCESS - process to restart (0 - n-1 where n is the number of

processes in the table) (byte).

Returns: r0 unchanged.

Destroys: a, x, y, rl

Description: RestartProcess sets a process's countdown timer to its initialization

value then unblocks and unfreezes it Use **RestartProcess** to initially start a process after a call to **InitProcesses** or to rewind a process to

the beginning of its cycle.

Note: RestartProcess clears the runnable flag associated with the process,

thereby losing any pending call to the process.

RestartProcess should always be used to start a process for the first time because InitProcesses leaves the value of the countdown timer in

an unknown state.

Example:

See also: InitProcesses, EnableProcess, UnfreezeProcess, UnblockProcess.

**Sleep:** (C64, C128) C199

Function: Pause execution of a subroutine ("go to sleep") for a given time interval.

Parameters: r0 DELAY - number of vblanks to sleep (word),

Returns: does not return directly to caller (see description below).

Destroys: a, x, y

Description: Sleep stops executing the current subroutine, forcing an early rts to the routine one level lower, putting the current routine "to sleep." At Interrupt Level, the DELAY value associated with each sleeping routine is decremented. When the associated DELAY value reaches zero, MainLoop removes the sleeping routine from the sleep table and performs a jsr to the instruction following the original jsr Sleep, expecting a subsequent rts to return control back to MainLoop. For example, in the normal course of events, MainLoop might call an icon event service routine (after an icon is clicked on). This service routine can perform a jsr Sleep. Sleep will force an early rts, which, in this case, happens to return control to MainLoop. When the routine awakes (after DELAY vblanks have occurred), MainLoop performs a jsr to the instruction that follows the original jsr Sleep. When this wake-up jsr occurs, it occurs at some later time the contents of the processor registers and GEOS pseudoregisters are uninitialized. A subsequent rts will return to MainLoop.

Sleeping in Detail:

- 1: The application calls **Sleep** with a jsr **Sleep**. The jsr places a return address on the stack and transfers the processor to the **Sleep** routine.
- 2: **Sleep** pulls the return address (top two bytes) from the stack and places those values along with the DELAY parameter in an internal sleep table.
- 4: **Sleep** executes an rts. Since the original caller's return address has been pulled from the stack and saved in the sleep table, this rts uses the next two bytes on the stack, which it assumes comprise a valid return address. (Note: it is imperative that this is in fact a return address; do not save any values on die stack before calling **Sleep**.)
- 5: At Interrupt Level GEOS decrements the sleep timer until it reaches zero.
- 6: On every pass, MainLoop checks the sleep timers. If one is zero, then it removes that sleeping routine from the table, adds one to the return address it pulled from the stack (so it points to the instruction following the jsr Sleep), and jsr's to this address. Because no context information is saved along with the Sleep address, the awaking routine cannot depend on any values on the stack, in the GEOS pseudoregisters, or in the processor's registers.

Note: A DELAY value of \$0000 will cause the routine to sleep only until the next pass through MainLoop.

When debugging an application, be aware that **Sleep** alters the normal flow of control.

Example: BeepThrice

See also: InitProcesses.

UnblockProcess: (C64, C128) C10F

Function: Allow a process's events to go through.

Parameters: x PROCESS - number of process (0 - n-1, where n is the number of

processes in the table) (byte).

Returns: x unchanged.

**Destroys:** a

Description: UnblockProcess causes MainLoop to again recognize a process's runnable

flag so that if a process timer reaches zero (causing the process to

become runnable) an event will be generated.

Because the GEOS Interrupt Level continues to decrement the countdown timer as long as the process is not frozen, a process may become runnable while it is blocked. As long as the process is blocked, however, MainLoop will ignore the runnable flag. When the process is subsequently unblocked, MainLoop will recognize a set runnable flag as a pending event and call the appropriate service routine. Multiple pending events are ignored: if a blocked process's timer reaches zero more than once, only one event will be generated when it is unblocked. To prevent a pending event from happening, use RestartProcess to unblock the process.

Note: If a process is not blocked, an unnecessary call to UnblockProcess will

have no effect.

Example:

See also: BlockProcess, UnfreezeProcess, EnableProcess, RestartProcess.

process

UnfreezeProcess: (C64, C128) c115

Function: Resume (unfreeze) a process's countdown timer.

**Parameters:** x PROCESS - number of process (0 - n-1, where n is the number of

processes in the table) (byte).

Returns: x unchanged.

**Destroys:** a

 $\textbf{Description: UnfreezeProcess} \ \ \texttt{causes} \ \ \texttt{a} \ \ \texttt{frozen} \ \ \texttt{process's} \ \ \texttt{countdown} \ \ \texttt{timer} \ \ \texttt{to} \ \ \texttt{resume}$ 

decrementing. The value of the timer is unchanged; it begins decrementing again from the point where it was frozen. If a process is

not frozen, a call to UnfreezeProcess will have no effect.

Note: If a process is not frozen, a call to UnfreezeProcess will have no

effect.

Example:

See also: FreezeProcess, BlockProcess

# GEOS Kernal

# utility

# Utility

Bell	N/A	Play a bell sound	147
CallRoutine	C1D8	pseudo-subroutine call. \$0000 aborts call.	148
CRC	C20E	Cyclic Redundancy Check calculation.	149
DoInlineReturn	C2A4	Return from inline subroutine.	150
GetRandom	C187	Calculate new random number.	151
ToBasic	C241	Pass Control to Commodore BASIC.	152

N/A

Function: Makes a brief beeping sound

(Apple)

Parameters: none.

Returns: nothing.

Destroys: a

Bell:

Description: Bell sounds a 1000 Hz signal. The sound lasts approximately 1/10th of a

second.

Note: Bell does not exists in Commodore Geos. This code provides the behavior

of the Apple Bell.

```
; Author: Dan Kaufman (w Chris Hawley)
```

FCLo = voice1Regs + 7 + \$07 FCHi = voice1Regs + 7 + \$08 res\_filt = voice1Regs + 7 + \$09 mode\_vol = voice1Regs + 7 + \$0A pulse = %01000001

 $\begin{array}{ccc} & & & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ &$ 

Bell:

PushB CPU\_DATA ;switch to I/O space LoadB CPU\_DATA, #IO\_IN controlReg1, #0 sta att\_dec1 LoadB mode\_vol, #\$18

LoadB sus rell, #SOUND ON LoadW PWLol, #\$800 LoadB FCLo,#0 sta FCHi sta res filt att\_decl,#6 sus\_rell,#0 LoadB LoadB LoadB freqLol, #\$DF freqHil/#\$25 LoadB

LoadB controlRegl, #pulse
PopB CPU\_DATA ; return to memory space

rts

Example: BeepThrice

See also: Ddec

# CallRoutine: (C64, C128) C1D8

Function: Perform a pseudo-subroutine call, checking first for a null address

(which will be ignored).

Parameters: a [ADDRESS - low byte of subroutine to call.

x ]ADDRESS — high byte of subroutine to call.

where ADDRESS is the address of a subroutine to call.

Returns: depends on subroutine at ADDRESS.

**Destroys:** depends on subroutine at ADDRESS.

 $\textbf{Description: CallRoutine} \ \ \text{offers a clean and simple way to perform an indirect jsr}$ 

through a vector or call a subroutine with an address from a jump table. Before simulating a jsr to the address in the x and a registers, it also checks for a null address (\$0000). If the address is \$0000 (x=\$00 and a=\$00), **CallRoutine** performs rts without calling any subroutine address. This makes it easy to nullify a vector or an entry in a jump table by

using a \$0000 value.

GEOS frequently uses **CallRoutine** when calling through vectors. This is why placing a \$0000 into **keyVector**, for example, causes GEOS ignore the vector. Other examples of this usage are **intTopVector**, **intBotVector**, and

mouseVector.

Note: CallRoutine modifies the st register prior to performing the jsr. It,

therefore, cannot be used to call routines that expect processor status flags as parameters (flags may be returned in the st register, however). CallRoutine may be called from Interrupt Level (off of routines in IntTopVector and IntBotVector). Do not use CallRoutine to call inline

(i) routines, as it will not return properly.

Example: HandleCommand:

CRC: C20E (c64, C128)

16-bit cyclic redundancy check (CRC). Function:

- pointer to start of data (word). Parameters: r0 DATA

> r1 LENGTH - of bytes to check (word).

CRC value for the specified range (word). Returns: r2

Destroys: a, y, r0-r3L

Description: CRC calculates a 16-bit cyclic-redundancy error-checking value on a range

of data. This value can be used to check the integrity of the data at a later time. For example, before saving off a data file, and application might perform a CRC on the data and save the value along with the rest of the data. Later, when the application reloads the data, it can perform another CRC on it and compare the new value with the old value. If the

two are different, the data has unquestionably been corrupted.

Given the same data, CRC will produce the same value under all versions Note:

of GEOS.

 $Note^1$ : This routine is called by the bootup routines to compute the checksum

> of GEOS BOOT. This checksum is used to create the interrupt vector address. The reason for this was to prevent piracy. This can be used to

check the integrity of a memory region.

#### Example: Kernal CRC

```
MAGIC VALUE = $0317
                                ; CRC value that we're looking for
DATA SIZE=$2434
                                 ; Size of data
```

.ramsect

buffer .block DATA SIZE

.psect

Checksum:

```
LoadW r0, #buffer ; r0 <- data area to checksum
LoadW r1, #DATA_SIZE ; r1 <- bytes in buffer to check
jsr
      CRC
                                ; r2 <- CRC value for data area
CmpWI r2, MAGIC VALUE
                               ; return status to caller
```

rts

DoInlineReturn: (c64,C128) C2A4

Function: Return from an inline subroutine.

Parameters: a DATABYTES - number of inline data bytes following the jsr plus

one (byte) .

stack top byte on stack is the status register to return (execute a

php just before calling).

Returns: (to the inline jsr) x, y unchanged from the jmp DolnlineReturn. st

register is pulled from top of stack with a plp.

Destroys: a

Description: DoInlineReturn simulates an rts from an inline subroutine call, properly skipping over the inline data. Inline subroutines (such as the GEOS routines which begin with i ) expect parameter data to follow the subroutine call in memory. For example, the GEOS routine i Rectangle is

called in the following fashion:

jsr i\_Rectangle ;subroutine call
.byte y1, y2 ;inline data

.word x1, x2

jsr FrameRectangle ;returns to here

Now if i\_Rectangle were to execute a normal rts, the program counter would be loaded with the address of the inline data following the subroutine call. Obviously, inline subroutines need some means to resume processing at the address following the data. DoInlineReturn Provides this facility. The normal return address is placed in the global variable returnAddress. This is the return address as it is popped off the stack, which means it points to the third byte of the inline jsr (an rts increments the address before resuming control). The status registers is pushed onto the stack with a php, DoInlineReturn is called with the number of inline data bytes plus one in the accumulator, and control is returned at the instruction following the inline data.

Inline subroutines operate in a consistent fashion. The first thing one does is pop the return address off of the stack and store it in returnAddress. It can then index off of returnAddress as in Ida (returnAddress), y to access the inline parameters, where the y-register contains \$01 to access the first parameter byte, \$02 to access the second, and so on (not \$00, \$01, \$02, as might be expected because the address actually points to the third byte of the inline jsr). When finished, the inline subroutine loads the accumulator with the number of inline data bytes and executes a jmp **DoInlineReturn**.

Note:

**DoInlineReturn** must be called with a jmp (not a jsr) or an unwanted return address will remain on the stack. The x and y registers are not modified by **DoInlineReturn** and can be used to pass parameters back to the caller. Inline calls cannot be nested without saving the contents of returnAddress. An inline routine will not work correctly if not called directly through a jsr (e.g., **CallRoutine** cannot be used to call an inline subroutine).

Example: i VerticalLine

See also: Ddec

GetRandom: (C64, C128) c187

Function: Creates a 16-bit random number.

Parameters: none.

Uses: random random seed for next random number.

Alters: random random contains a new 16-bit random number.

Returns: depends on subroutine at ADDRESS.

Destroys: a

Description: GetRandom produces a new pseudorandom (not truly random) number using

the following linear congruential formula:

random = (2\*(random+1) // 65521)

(remember: // is the modulus operator)

The new random number is always less than 65221 and has a fairly even

distribution between 0 and 65521.

Note: GEOS calls GetRandom during Interrupt Level processing to automatically

keep the random variable updated. If the application needs a random number more often than random can be updated by the Kernal, then

GetRandom must be called manually.

Example:

**ToBasic:** (C64, C128) **c187** 

Function: Removes GEOS and passes control to Commodore BASIC with the option of

loading a non-GEOS program file (BASIC or assembly-language) and/or

executing a BASIC command.

> r5 DIR\_ENTRY - pointer to the directory entry of a standard Commodore file (PRG file type), which itself can be either a BASIC or ASSEMBLY GEOS-type file. If this parameter is \$0000, then no file will be loaded.

r7 LOADADDR — if r5 is non-zero, then this is the file load address. For a BASIC program, this is typically \$801. If r5 is zero and a tokenized BASIC program is already in memory, then this value should point just past the last byte in the program. If r5 is zero and no program is in memory, this value should be \$803, and the three bytes at \$800-\$802 should be \$00.

Returns: N/ADestroys: N/A

Description: ToBasic gives a GEOS application the ability to run a standard Commodore

assembly-language or  ${\tt BASIC}$  program. It removes GEOS, switches in the  ${\tt BASIC}$  ROM and I/O bank, loads an optional file, and sends an optional

command to the BASIC interpreter.

Once **ToBasic** has executed, there is no way to return directly to the GEOS environment unless the RAM areas from \$C000\$ through <math>\$C07F\$ are preserved (those bytes may be saved and restored later). To return to GEOS, the called program can execute a jump to <math>\$C000\$ (BootGEOS).

A program in the C64 environment can check to see if it was loaded by GEOS by checking the memory starting at \$C006 for the ASCII (not CBMASII) string "GEOS BOOT" If loaded by GEOS, the program can check bit 5 of \$C012: if this bit is set, ask the user to insert their GEOS boot disk; if this bit is clear, GEOS will reboot from the RAM expansion unit To actually return to GEOS, set CPU\_DATA to \$37 (KRNL\_BAS\_IO\_IN) and jump to \$C000

Example: LoadBASIC

# GEOS Kernal

+0\*\*+

	_		
т.	e	x	τ

GetCharWidth	\$C1C9	Get a character's width	154
GetNextChar	\$C2A7	Read a character from the keyboard	155
GetRealSize	\$C1B1	Get a character's stats	156
GetString	C1BA	Read a line of text from the user	157
InitTextPrompt	\$C1C0	Create the text cursor sprite	-208
LoadCharSet	\$C1CC	Load and activate a new font	-212
PromptOff	\$C29E	Turn off the text cursor	-210
PromptOn	\$C29B	Turn on the text cursor	-209
PutChar	\$C145	Display a character	-201
PutDecimal	\$C184	Display a 16 bit integer	160
PutString	\$C148	Display a text string	-204
i PutString	\$C1AE	Display a text string with inline data	-205
 UseSystemFont	\$C14B	Select the BSW font	-211
Small PutChar	\$0202	Draw a character on the screen	159

GetCharWidth: (C64, C128) \$c1c9

Function: Calculate the pixel width of a character as it exists in the font (in

its plaintext form). Ignores any style attributes.

Parameters: a CHAR - character code of character (byte).

Uses: curlndexTable

**Returns:** a character width in pixels.

Destroys: y

Description: GetCharWidth calculates the width of the character before any style

attributes are applied. If the character code is less than 32, \$00 is returned. Any other character code returns the pixel width as calculated from the font data structure. The sprites will remain hidden until the

next pass through MainLoop.

Because **GetCharWidth** does not account for style attributes, it is useful for establishing the number of bits a character occupies in the font

data structure.

Note: In 40-column mode (bit 7 of graphMode is zero), TempHideMouse exits

immediately without affecting the hardware sprites.

Example:

See also: GetRealSize

GetNextChar: (C64, C128) \$c2A7

Function: Retrieve the next character from the keyboard queue.

Parameters: none.

Returns: a keyboard character code of character or NULL if no characters

available.

Alters: pressFlag if the call to GetNextChar removes the last character

from the queue, then the KEYPRESS BIT is cleared.

Destroys: x

Description: GetNextChar checks the keyboard queue for a pending keypress and returns

a non-zero value if one is available. This allows more than one character

to be processed without returning to MainLoop

Example:

See also: GetString

GetRealSize: (C64, C128) \$C1B1

Function: Calculate the printed size of a character based on any style attributes.

Parameters: a CHAR - character code of character (byte).

Uses: curHeight

baselineOffset

Returns: y character width in pixels (with attributes).

- \* character height in pixels (with attributes).
- a character baseline offset (with attributes).

Destroys: nothing.

Description: GetRealSize calculates the width of the character based any style

attributes The character code must be 32 or greater. If the character code is  ${\tt USELAST}$ , the value in  ${\tt lastWidth}$  is returned. Any other character code returns the pixel width as calculated from the font data structure

and the MODE parameter.

lastWidth is local to the GEOS Kemal and therefore inaccessible to

applications.

### Example:

; Calculate size of largest character in current font

lda #'W' ; capital W is a good choice

ldx #(SET\_BOLD|SET\_OUTLINE) ; widest style combo

jsr GetRealSize ; dimensions come back in x,y

GetString: (C64, C128) C1BA

Function: Get a string from the keyboard using a cursor prompt and echoing

characters to the screen as they are typed. Runs concurrently with

MainLoop.

Parameters: r0 BUFR - pointer to string buffer. When called this buffer can

contain a null-terminated default string (if no default string is used, the first byte of the buffer must be NULL).

This buffer must be at least MAX\_CH+1 bytes long.

rlL FLAG - \$00 = use system fault routine;

\$80 = use fault routine pointed to by r4 (byte).

**r2L** MAX\_CH - maximum number of characters to accept

(not including the null-terminator).

r11 XPOS - x-coordinate to begin input (word).

rlH YPOS - y-coordinate of prompt and upper-left of characters.

To calculate this value based on baseline printing position, subtract the value in  ${\bf baselineOffset}$  from the baseline

printing position (byte).

r4 FAULT - optional (see FLAG) pointer to fault routine.

keyVector STRINGDONE - routine to call when the string is terminated

by the user typing a carriage return.

\$0000 = no routine provided.

Uses: at call to GetString:

curHeight for size of text prompt.

baselineOffset for positioning default string relative to prompt

any variables used by PutString.

while accepting characters:

**keyVector** vectors off of MainLoop through here with characters.

string pointer to start of string buffer.

any variables used by PutChar.

Returns: from call to GetString:

stringX starting prompt x-position.
stringY starting prompt y-position.

**string** BUFR (pointer to start of string buffer).

when done accepting characters:

x length of string / index to null

**string** BUFR (pointer to start of string buffer).

keyVector \$0000
stringFaultVec \$0000

Destroys: at call to GetString:

r0-r13, a, x, y.

Description: GetString installs a character handling routine into GetString and returns immediately to the caller. During MainLoop, the string is built

up a character at a time in a buffer. When the user presses [Return], GEOS calls the STRINGDONE routine with the starting address of the

string in string and the length of the string in the x-register.

The following is a breakdown of what **GetString** does:

- 1: Variables local to the **GetString** character input routine are initialized. Global string input variables such as **string**, **stringX**, and **stringY** are also initialized.
- 2: **PutString** is called to output the default input string stored in the character buffer. If no default input string is desired, the first byte of the buffer should be a **NULL**.
- 3: The **STRINGDONE** parameter in **keyVector** is saved away and the address of the **GetString** character routine (SystemStringService) is put into keyVector.
- 4: If the application supplied a fault routine, install it into **StringFaultVec**, otherwise install a default fault routine.
- 5: The prompt is initialized by calling **InitTextPrompt** with the value in **curHeight**. **PromptOn** is also called.
- 6: Control is returned to the application.

lastWidth is local to the GEOS Kemal and therefore inaccessible to
applications.

Note:

String is not null-terminated until the user presses [Return]. To simulate a [Return], use the following code:

;Simulate a CR to end GetString

This will also terminate the GetString input.

Note:

This note courtesy of Bill Coleman...Because **GetString** runs off of **MainLoop**, it is a good idea to call **GetString** from the top level of the application code and return to **MainLoop** while characters are being input. That is, while at the top level of your code you can call **GetString** like this:

Since the routine specified by the **STRINGDONE** value stored in **keyVector** is called when the user has finished entering the string, that is where your application should again take control and process the input.

Note<sup>2</sup>:

If the user manages to type off the end of the screen, specifically past rightMargin, GetString will stop echoing characters although it will still enter the characters into the buffer.

See also: PutChar, PutString, GetNextChar.

**SmallPutChar:** (C64, C128) \$c202

Function: Print a single character without the PutChar overhead

Parameters: a CHAR - character code (byte).

r11 XPOS — x-coordinate of left of character (word). r1H YPOS — y-coordinate of character baseline (word).

Uses: Same as PutChar

Returns: r11 x-position for next character.

rlH unchanged

Destroys: a, x, y, rlL, r2-rl0, rl2, rl3

**Description:** SmallPutChar is a bare bones version of PutChar. SmallPutChar will not handle escape codes, does no margin faulting, and does not normalize the x coordinates on GEOS 128.

SmallPutChar will assume the character code is a valid and printable character Any portion of the character that lies above windowTop or below windowBottom will not be drawn. If a character lies partially outside of leftMargin or rightMargin, SmallPutChar will only print the portion of the character lies within the margins. SmallPutChar will also accept small negative values for the character x-position, allowing characters to be clipped at the left screen edge.

Note:

Partial character clipping at the **leftmargin**, including negative x-position clipping, is not supported by early versions of GEOS 64 (earlier than vl.4) — the entire character is clipped instead Full **leftmargin** clipping is supported on all other versions of GEOS: GEOS 64 vl.4 and above, GEOS 128 (both in 64 and 128 mode.

Like **PutChar**, 159 is the maximum CHAR value that **SmallPutChar** will handle correctly. Most fonts will not have characters for codes beyond 129.

Example:

128: DOUBLE W, ADD1 W cannot be used on r11

\$C184

**Function:** Format and print a 16-bit positive integer.

Parameters: a FORMAT - formatting codes (byte) - see below.

 ${f r0}$  NUM - 16-bit integer to convert and print (word).

r11 XPOS - x-coordinate of leftmost digit (word).

rlH YPOS - y-coordinate of baseline (word).

(C64, C128)

Uses: Same as PutChar

PutDecimal:

Returns: r11 x-position for next character.

rlH unchanged

Destroys: a, x, y, r0, r1L, r2-r10, r12, r13

**Description: PutDecimal** converts a 16-bit positive binary integer to ASCII and sends the result to PutChar. The number is formatted based on the FORMAT

parameter bytes in the a-registers as follows:

FORMAT:

7 6 5 4 3 2 1 0

, 0	0 1 0	2 1 0	
b7	b6	b0-b5	

b7 justification:

1 = left

0 = right.

b6 leading zeros:

1 = suppress

0 = print.

b5-b0 field width in pixels (only used if right justifying).

The following constants may be used:

SET\_LEFTJUST SET\_RIGHTJUST

SET\_SUPPRESS SET NOSUPPRESS

Note: The maximum 16-bit decimal number is 65535 (\$FFFF), so the printed

number will never exceed five characters.

Example:

### GEOS Kernal

wheels

Chapter 2 Wheels 4.4

### Wheels Kernal

GetNewKernal \$9d80 Load New Kernal Group
RstrKernal \$9d83 Unload Kernal Group

wheels

GetNewKernal: (Wheels 4.4 64,128) \$9D80

Function: Load Modular Kernal Group

Parameters: a GROUPNBR to load | RUNFLAG

RUNFLAG Bit 6 of a.

1 Selected Kernal Group Swapped into memory at 5000-5FFF.

O First Routine in group executed. (Kernal Group swapped back).

Destroys: (unknown)

Return: varies depending on RUNFLAG and GROUPNBR.

Description: GetNewKernal allows access to the Extended Kernal available in 4.4.

If RUNFLAG is 0 **GetNewKernal** behaves as a far jsr to the first routine in the Kernal Group. Performing the following...

Swap the extended kernel group into memory. Execute the first routine in the group. Swap the kernal back out of memory. Control is returned to the caller.

If RUNFLAG is set:

Extended Kernal is swapped into memory at 5000-5FFF. Control is returned to the caller.

(Kernal will remain in memory until a call to **RstrKernal** to swap it back.)

Note: Kernal Groups are loaded from the Last REU bank which is reserved

exclusively for the 4.4 Kernal.

Note: Caller cannot be in the Range 5000-5FFF as that address range is swapped

out with the Kernal Group

Example:

KG\_REU=\$00 NO\_RUN=%01000000 RUN FIRST=%00000000

LoadREUGrp:

lda #KG\_REU|NO\_RUN
jmp GetNewKernal

wheels

**RstrKernal:** (Wheels 4.4 64,128) \$9D83

Function: Unload Extended Kernal group.

Parameters: none

Destroys: a

Return: nothing

Alters: Memory area from 5000-5FFF is restored to its previous contents.

Description: RstrKernal is used to restore the memory area 5000-5FFF after using

GetNewKernal to load in an extended Kernal Group.

### Example:

GRP\_REU=\$00 NO\_RUN=%01000000 RUN FIRST=%00000000

.ramsect

freeBanks: .block 1

.psect

#### GetBanksFree:

lda #GRP REU|NO RUN ; Select REU Group . And don't execute 1st

jsr GetNewKernal ; Load in Kernal Group.

jsr GetRAMInfo ; Call Kernal Group function to get

; number of free REU banks.

MoveB **r4H**, freeBanks ; save the result

jmp RstrKernal ; Remove Kernal Group, restoring 5000-5FFF

; to its previous contents.

Examples

hardware

# Examples

hardware

# GetFPS:

```
;Author PBM
; PASS: Nothing
;Return: a = fps
        minus flag set if known model was not found
        minus return should never happen without a bug in C64Model
models: .byte %00,%01,%10,%11
NBR MODELS=*-models
frates: .byte 50,60,60,50
GetFPS:
    jsr C64Model
10$
    ldx #NBR_MODELS-1
    cmp models,x
    beq
           90$
    dex
    bpl 10$
lda [TRUE
    rts
90$
    lda frates,x
    rts
```

# C64Model:

```
; Detect PAL/NTSC
; Original Name: DetectC64Model
; Author: TWW
; 312 rasterlines -> 63 cycles per line PAL
                => 312 * 63 = 19656 Cycles / VSYNC => #>76 %00
; 262 rasterlines -> 64 cycles per line NTSC V1
                 => 262 * 64 = 16768 Cycles / VSYNC => #>65 %01
; 263 rasterlines -> 65 cycles per line NTSC V2
                \Rightarrow 263 * 65 = 17095 Cycles / VSYNC \Rightarrow #>66 %10
; 312 rasterlines -> 65 cycles per line PAL DREAN
                => 312 * 65 = 20280 Cycles / VSYNC => #>79 %11
C64Model:
   ;-- Use CIA #1 Timer B to count cycled in a frame
   lda #$FF
   sta $DC06
   sta $DC07
                    ; Latch #$FFFF to Timer B
10$
   bit $D011
   bpl 10$
                    ; Wait until Raster > 256
20$
   bit $D011
   bmi 20$
                     ; Wait until Raster = 0
   ldx #%00011001
   stx $DCOF
                     ; Start Timer B (One shot mode
                     ; (Timer stops automatically when underflow))
30$
   bit $D011
   bpl 30$
                     ;Wait until Raster > 256
40$
   bit $D011
   bmi 40$
                     ;Wait until Raster = 0
   sec
   sbc $DC07
                    ; Hibyte number of cycles used
   and #%0000011
   rts
```

DetectC64Model Source from CodeBase64
https://codebase64.org/doku.php?id=base:detect pal ntsc

Note<sup>3</sup>: I believe this will also work on a 128 in 40 column mode. Need to test.

Examples

	-1-	
na	TП	

math

-----

# 8BitMultiply:

```
8BitMultiply- 8 Bit unsigned multiply.
         x - zpage address of multiplicand
          y - zpage address of multiplier
; returns: unsigned result in address pointed to by {\bf x}
          x, y unchanged
; Multiply {\bf rlL} by {\bf rlH} and store the word result in r2
8BitMultiply:
   MoveB rlL, r2L
                         ; r2L <- rlL copy of OPERAND1
   ldx
          #r2L
                         ; x <- source register address
                         ; y <- destination register address
   ldy
         \#r1H
   jsr
          BBMult
                         ; r2 <- r2L * r2H do multiplication
   rts
```

# 16x8Multiply:

```
16x8Multiply - 16x8 Bit unsigned multiply.
        x - zpage address of multiplicand
         y - zpage address of multiplier
; returns: unsigned result in address pointed to by \boldsymbol{x}
         x, y unchanged
; Multiply the value in r9 by 87 and store the result back in r9
; (rl is destroyed)
*********************
16x8Multiply:
                      ; point to OPERAND1 in r9
  ldx
         #r9
         rlL,#87
                      ; rl <- 87 (OPERAND2)
  LoadB
                      ; point to OPERAND2 in rl
  ldy
         #rl
         Bmult ; r9 <- r9 * rlL
  jsr
  rts
```

### ConvToUnits:

```
This routine converts a pixel measurement to inches or, optionally,
  centimeters, at the rate of 80 pixels per inch or 31.5 pixels per
  centimeter.
  pass: r0 - number to convert (in pixels)
  return: r0 - inches / centimeters
          rlL - tenths of an inch / millimeters.
  destroys: a, x, y, r0-r1, r8-r9
; Assembler time decision on whether inches or centimeters is to be used.
.if AMERICAN
   INCHES = TRUE
.else
   :Metric
   INCHES = FALSE
.endif
ConvToUnits:
                         ; First, convert r0 to length in 1/20 of
                         ; standard units
.if INCHES
                         ; For Inches, need to multiply by
                              20 1
                         ; ----- = ---
                         ; 80 dots/inch
                         ; which amounts to a divide by four
   ldx
           #r0
   ldy
           #2
          DShiftRight
   jsr
.else
                         ; For Centimeters, need to multiply by
                          ; 20 1
                          ; ----- = ---
                         ; 31.5 dots/cm
                         ; First multiply by 40
   LoadB
          rl,#40
                         ; (word value)
   ldx
           #r0
                        ; (byte value)
           #rl
   ldv
                        ; r0 * r0*40 (byte by word multiply)
   jsr
          Bmult
   LoadW
          rl,#63
                        ; then divide by 63
   ldx
          #r0
   ldy
          #rl
   jsr
          Ddiv
                        ; r0- r0/63
.endif
   ;-- Start of Common Code ; r0 * result in 1/20ths
          r0
                                   ; add in one more 1/20th, for rounding
   IncW
                        ; now divide by 20 (to move decimal over one)
   LoadW
           rl,#20
   ldx
          #r0
                         ; dividend
                         ; divisor
           #rl
   ldy
          Ddiv
                         ; r0 = r0 /20 (r0 = result in proper unit)
   jsr
   MoveB
          r8L,rlL
                         ; rlL - 1/20ths
   lsr
          rlL
                         ; and convert to 1/10ths (rounded)
   rts
                         ; exit
```

# Kernal CRC:

```
This is the actual Kernal Code for CRC.
        r0 - pointer to start of data
   pass:
          r1 - # of bytes to check
  return: r2 - CRC Checksum
  destroys: a, x, y, r0, r1, r3L
                     **********
Kernal_CRC:
         ldy
              #$FF
         sty
              r2L
         STY
              r2H
         iny
10$
         lda
              #$80
         sta
              r3L
20$
              r2L
         asl
         rol
              r2H
              (r0),y
         lda
         and
              r3L
              30$
         bcc
         eor
              r3L
30$
              40$
         beq
         lda
              r2L
              #%00100001
         eor
              r2L
         sta
         lda
              r2H
         eor
              #%00010000
         sta
              r2H
40$
         lsr
              r3L
         bcc
              20$
         iny
              50$
         bne
         inc
              r0H
50$
         ldx
              #r1
         jsr
              Ddec
         lda
              r1L
         ora
              r1H
         bne
              10$
         rts
```

# DecCounter:

zCounter = \$70 COUNT = \$FFF0

# DecCounter:

LoadW zCounter,COUNT

10\$

Jsr DoSomething
ldx #zCounter
jsr Ddec
bne 10\$

rts

case:4

# DdecvsDecW:

Size in Bytes vs Speed in Cycles of Ddec and DecW

**Ddec** represents a maximum of 7 byte savings over **DecW** every time it is used in your code. If Not needing a zero result after **DecW** then only a 3 byte savings.

DecW takes roughly ½ the time to execute. In and Inner loop executed 1 Million times. DecW will save roughly 20 seconds off the time vs Ddec

```
zCounter=$70
```

.macro DecW dest
 lda dest
 bne dolow
 dec dest+1
dolow:
 dec dest
.endm

### Ddec code block.

Op Code	Instruction	Bytes	Cycles
A2 70	ldx #zCounter	2	2
20 OE C2	jsr <b>Ddec</b>	3	6
	(Kernal Routine)	0	27 - 32
	Total	5	35 - 40

### DecW macro code block.

Op Code	Instruction	Bytes	Cycles	
A9 70	lda zCounter	2	3	
D0 02	bne 10\$	2	2 or 3 or 4	
C6 71	dec zCounter+1	2	5	
	10\$			
C6 70	dec zCounter	2	5	
	Total	8	11 Worst Case 15	
	if branch crosses p	age	12	

;-- When using  $\bf DecW$  on a counter, Add check for word=0 after the DecW macro A9 70 lda zCounter 2 2

05 70		ora zCounter+1	2	3
		Total	12	16 - 20

Kernal <b>Ddec</b>	;Actual Kernal Co	de for <b>Ddec</b>	
Op Code	Instruction	Cycles	
B5 00	lda zpage,X	4	
D0 02	bne 10\$	(1/256ish chance 2) or 3 or Wors	st
D6 01	dec zpage+1,X	6	
	10\$		
D6 01	dec zpage,X	6	
B5 00	lda zpage,X	4	
D6 01	ora zpage+1,X	4	
60	rts	6	
		=======================================	
	- · ·	- · · · · · · · · · · · · · · · · · · ·	

Total Best Case: 27 Worst Case: 32 if branch crosses Page 28 (1/256 chance)

#### DSmult:

```
DSMult - double-precision signed multiply.
          x - zpage address of multiplicand
           y - zpage address of multiplier
   returns: signed result in address pointed to by x
;
           word pointed to by y is absolute-value of the
                  multiplier passed
           x, y unchanged
   Strategy:
;
           Establish the sign of the result: if the signs of the
;
           multiplicand and the multiplier are different, then the result
;
           is negative; otherwise, the result is positive. Make both the
           multiplicand and the multiplier positive, do unsigned
           multiplication on those, then adjust the sign of the result
           to reflect the signs of the original numbers.
   destroys:
                  a, r6 - r8
```

#### DSmult:

```
lda
                            ;get sign of multiplicand (hi-byte)
            zpage+l,x
                            ; and compare with sign of multiplier
   eor
            zpage+l,y
   php
                            ; save the result for when we come back
            Dabs
                             ;multiplicand = abs(multiplicand)
   jsr
   stx
            r6L
                             ; save multiplicand index
                             ; put multiplier index into x
   tya
                             ; for call to Dabs
   tax
            Dabs
                            ;multiplier = abs(multiplier)
   jsr
            r6L
   ldx
                            ;restore multiplier index
   jsr
            DMult
                            ; do multiplication as if unsigned
   plp
                            ;get back sign of result
            90$
                             ; ignore sign-change if result positive
   bpl
                             ;otherwise, make the result negative
   jsr
            Dnegate
90$
   rts
```

Examples

memory

memory

 ${\tt CopyBuffer}$ 

# CopyBuffer:

```
SrcBuff: .byte "Any Values can be in the buffer", NULL, CR
        .byte $0C,"NULLS are just zeros here",CR
LENBUFF = (*-SrcBuff)
.ramsect
     DestBuff .block LENSTRING
•psect
CopyBuffer:
     LoadW r5, #SrcBuff ; point to start of source buffer LoadW r11, #DestBuff ; point to start of destination buffer
     LoadW rll, #DestBuff ldx #r5
                            ; x <- source register address
                            ; y <- destination register address
     ldy #rll
          #LENBUFF
                           ; a <- length of buffer
     lda
     jsr CopyFString ; DestBuff <- SrcBuff (copy)</pre>
     rts
SrcStr: .byte "Any values but null can be in the string", NULL
LENSTRING = (*-SrcStr)
.ramsect
     DestBuff .block LENSTRING
.psect
CopyStr:
     ldx
          #r0
                            ; x <- source register address
     ldy
           #rl
                            ; y <- destination register address
     jsr CopyString ; DestBuff <- SrcStr (copy)</pre>
     rts
```

### Find:

```
REC SIZE = 5 ; size of each record
.ramsect
        Data: .block 1024 ; Table of Zip Code Locations.
.psect
        Key: .byte "65803" ;Zip Code to Find
Find:
                                     ; r2 <- total number of records
        LoadW r2, #NUM RECS
        LoadW r0,#Key
LoadW r1,#Data
                                         ; r0 <- pointer to keyword
                                         ; rl <- pointer to start of search list
10$
                                         ; DO
                                                 \mathbf{x} <- source string - key ,
        ldx
               #r0
                                         ;
        ldy #rl ; y <- destination string - list
lda #REC_SIZE ; a <- length of each record
jsr CmpFString ; compare key with current record
beq 20$ ; if they match, branch to handler
AddVW #REC_SIZE,rl ; otherwise point to the next record
DecW r2 ; r2- (decrement counter)
        bne 10$
                                         ; WHILE (r2 > 0)
        ;---
      jmp NotMatched ; jmp to no match handler
jmp Matched ; jmp to match handler
20$
```

# Find2:

```
Find2:

LoadW r0, #original ; r0 <- pointer to original string
LoadW r1, #copy ; r1 <- pointer to copy
ldx #r0 ; x <- source string =* key
ldy #r1 ; y <- destination string - list
jsr CmpString ;
beq 20$
jmp NotMatched ; jmp to no match handler

20$ jmp Matched ; jmp to match handler

original:
.byte "Mark Charles Heartless", NULL

Copy:
.byte "Mark Charlie Heartless", NULL
```

# InitBuffers:

```
; initialize buffers and variables to zero
InitBuffers:
   LoadW
        r0, #varStart
                                ; clear variable space
   LoadW rl, # (varEnd-varStart)
   isr
          ClearRam
  LoadW
          r0, #heapStart
                                ; clear heap
   LoadW
          rl, # (heapEnd-heapStart)
          ClearRam
   jmp
Alternate version. Using more space efficient i_Fill_{Ram}
InitBuffers:
          i FillRam
                                 ; clear variable space
   jsr
   .word
         varStart
   .word varEnd-varStart
         $AA
                                 ; With any value you choose
   .byte
          i FillRam
   jsr
                                 ; clear heap
   .word heapStart
   .word heapEnd-heapStart
         $00
   .byte
                                ; Heap set to zero's
   rts
```

Examples

disk routines

disk

CheckDiskSpace

# CheckDiskSpace:

```
; DESCRIPTION: Ensures that the current disk has a enough space for a
            minimum number of bytes. Does not take into account any
             index blocks or other blocks needed to maintain the file
            structure. Works with GEOS 64, GEOS 128
          r2
;Pass:
                  number of bytes we need
            x = If not enough space, returns an
;Returns:
                  INSUFFICIENT SPACE error.
             x = 0 Is there is enough space.
             Z Flag follows value of X.
;Destroyed: a, y, r2, r3, r8, r9
; Number of bytes that can be stored in each block on the disk. Accounts for
; two-byte track/sector link on Commodore versions of GEOS.
NO ERROR
          = 0
BLOCK SIZE = $100
BLOCK BYTES = BLOCK SIZE - 2
.macro bgt raddr
  beq label
   bcs
          raddr
label:
.endm
CheckDiskSpace:
                           ; r2 - # of BYTES to check for
   lda
       r2L
          r2H
   ora
                            ; check if zero bytes requested
         80$
   beq
                            ; if so, exit with no error
   LoadW r3, BLOCK BYTES
                           ; r3 <- number of bytes per block.
                            ; divide r2 by r3 to get number of
   ldx
         #r2
         #r3
                            ; blocks to hold BYTES
   ldy
                            ; r2 <- r3/r2
   jsr
         Ddiv
        r8L
                            ; r8L <- remainder
   lda
        r8H
   ora
                            ; Any remainder bytes?
   beq
         10$
                            ; if not, OK
   IncW
                            ; otherwise 1 more block needed
                            ; r2 = BLOCKS needed to hold BYTES
10$
                            ; get number of free blocks on disk
   LoadW r5, #curDirHead
                           ; point to directory header
   jsr
          CalcBlksFree
                           ; r4 <- free blocks on disk
          r2,r4
                            ; are there enough free blocks?
   CmpW
          99$
                            ; if not, assume. correct, branch.
   bgt
   ldx
          #NO ERROR
                      ; otherwise, no error
   rts
99$
          #INSUFFICIENT SPACE ; not enough space
   ldx
   rts
                            ; exit
```

# DeleteDirEntry:

; Pass: r0 pointer to filename

.ramsect

rFileName: .block 17

#### DeleteDirEntry:

LoadW r0, rFileName

LoadW r3, #NullTrScTable ; pass dummy table

jmp FastDelFile

This will also work correctly with a VLIR file. For freeing (deleting) all the blocks in a file without removing the directory entry refer to **FreeFile**.

#### ReadAndDelete:

.if COMMENT

Read sequential file into memory and then delete it from disk

Pass: r6 pointer to filename r7 where to put data

r2 size of buffer (max size of file)

Returns: x error code

Destroys: a, y, r0-r9

#### Implementation:

Call **FindFile** to get the directory entry of the file to load/delete. We pass the directory entry to **GetFHdrlnfo** to get the GEOS header block. We check the header to ensure we're not trying to read in a VLIR file. After **GetFHdrlnfo**, the parameters are already set up correctly to call ReadFile (fileTrScTab+0, **fileTrScTab**+1 contains header block and rl contains first data block). **ReadFile** reads in the file's blocks, building out the remainder of the **fileTrScTab**, which we pass to **FastDelFile** to free all blocks in the file (including the file header block, which is the first entry in the table).

.endif

#### ReadAndDelete:

```
; save pointer for FastDelFile
   MoveW r6.r0
          FindFile
   Jsr
                                        ; find file on disk
                                        ; set status flags
   txa
                                        ; branch on error
          99$
   bne
   LoadW r9, dirEntryBuf
                                        ; get directory entry
          GetFHdrlnfo
                                       ; get GEOS file header
   jsr
   txa
                                        ; set status flags
   bne
                                        ; branch on error
          fileHeader+OFF GSTRUCT TYPE
   lda
   cmp
          #VLIR
                                        ; check filetype
         10$
   bne
                                        ; branch if not VLIR
   ldx
         #STRUCT MISMAT
                                        ; can't load VLIR
   bne
          99$
                                        ; branch always for error
10$
   jsr
         ReadFile
                                        ; read in file
   txa
                                        ; else set status flags
   bne
          99$
                                        ; branch on other error
   LoadW r3, #fileTrScTab
                                        ; track/sector table
                                         ; file read OK, delete it!
   jsr
          FastDelFile
99$
   rts
                                         ; error in x
```

#### GrabSomeBlocks:

```
GrabSomeBlocks - allocate enough disk blocks to hold
;
                data in buffer.
 pass: Nothing
  returns: Carry flag. 1 = Error, 0 = success.
          X = Error Nbr if Carry is set or 0.
K = 1024
                       ;one kilobyte
.ramsect
  buffer: .block 5*K -1 ; 5K buffer .'
  bufferE: .block 1
                       ; End of 5k Buffer
   BUF SIZE = (bufferE - buffer)+1; size of buffer
.psect
GrabSomeBlocks:
  LoadW r2,BUF\_SIZE; number of bytes to allocate
         r6, fileTrSecTab ; buffer to build out table
   LoadW
                   ; allocate the blocks
          BlkAllOC
  jsr
                       ; check status
   txa
        99$
  bne
                       ; and exit on error
   ; more code here
90$
   ldx
         #0
                       ; Success exit
   clc
   rts
99$
                       ; Error Exit
   sec
   rts
```

## MyFreeBlock:

```
*********
   {f MyFreeBlock} — allocate specific block in BAM
;
  with any CBM device driver. And any GEOS Version
  pass:
           r6L = track #
           r6h = sector #
  Note:
           FreeBlock was not added to the
          GEOS jump table until vl.3
MyFreeBlock:
           version ; check GEOS version number
   lda
                         ; version Less then 1.3?
           #$13
   cmp
           10$
   bcc
                         ; if not, go through jump table
   Jmp
           FreeBlock
10$
   jsr
           FindBAMBit
                          ; Returns r8H = mask for BAM byte
                                    r7H = offset to track
                                   \mathbf{x} = offset into bam
                          ;
                                   a = masked value
           99$
                         ; if 1, then not allocated, give error
   bne
   txa
   bne
           99$
   lda
           r8H
                         ; get mask
   eor
           curDirHead, X
                         ; flip BAM bit to make available
   sta
           curDirHead, x
   ldx
           r7H
                          ; one more free block
   inc
           curDirHead, X
   ldx
           #0
                          ; NO ERROR
   rts
99$
          #BAD BAM
   ldx
                         ;
   rts
```

#### MySetGDirEntry:

.if COMMENT

This routine duplicates the function of the Kernal's SetGDirEntry for demonstration purposes. It shows examples of the following routines:

```
BldGDirEntry
GetFreeDirBlk
PutBlock
```

```
Same as SetGDirEntry
Destroys: Same as SetGDirEntry
                                                                         .endif
DIRCOPYSIZE=30
                             ; Size of directory entry for copy
TDSIZE=5
                              ; number of bytes in time/date entry
MySetGDirEntry:
                             ; build directory entry for GEOS file
   jsr
            BldGDirEntry
                              ; get block with free directory entry
   jsr
             GetFreeDirBlk
                              ; r3 = 1st byte of free entry
                              ; block number of block in rl
   txa
                              ; test for error code
   bne
             99$
                              ; if error, exit...
                              ; get offset into diskBlkBuf for dir entry
   tva
   clc
             #[diskBlkBuf ; and get absolute address in buffer
   adc
   sta
            r5L
            # | diskBlkBuf
   lda
            # ()
                              ; (propagate carry)
   adc
            r5H
   sta
   ldy
             #DIRCOPYSIZE
                             ; copy over some bytes
10$
                             ; get byte from directory entry built
   lda
             dirEntryBuf, y
             (r5),y
                             ; store new entry into block buffer
   sta
   dey
   bpl
             10$
                             ; loop till copied
            TimeStampEntry ; stamp the dir entry with time & date r4,#diskBlkBuf ; write out the new directory entry
   jsr
   LoadW
            PutBlock
   jsr
   txa
                              ; get error status
   bne
            99$
                              ; if error, exit
   clc
                              ; Success exit
   rts
99$
   sec
   rts
                             ; Error exit
TimeStampEntry:
   ldy
             #(OFF YEAR+TDSIZE)-1 ; offset to time/date stamp
10$
             dirEntryBuf,y
                             ; get the year/month/day/hour/minute
   lda
   sta
            (r5),y
                              ; store in dir entry
   dey
             10$
   bpl
                              ; Loop until done
   rts
```

## MyPutBlock:

```
***********
   MyPutBlock - Write diskBlkBuf to disk
;
   pass:
            r1L = track #
            r1H = sector #
            r4 = Address of block to write.
            verify = FALSE (0) Do Not Verify
                    <> 0 Verify after Write
   Note:
            If you have multiple blocks to write you should
            write the entire chain and then verify the chain.
            See WriteBlock description for more information
    .ramsect
                                         1
            nextTrack:
                            .block
                                        1
            nextSector:
                             .block
                            .block
                                       $FE
            outbuffer:
            track
                            .block
                                        1
                            .block
            sector
                             .block
                                        1
            verify:
   .psect
CallMyPutB:
            r4, outBuffer-2
   LoadW
   MoveB
            track, r1L,
            sector, r1H
   MoveB
            verify, [#TRUE
   LoadB
            MyPutBlock
   jsr
   bcs
            99$
   rts
                             ;return good status in carry
99$
                             ; Error Handler or let caller handle error
   . . .
   rts
MyPutBlock:
   jsr
            EnterTurbo
                             ; go into turbo mode
                             ; check for error in X
   txa
   bne
            99$
                            ; branch if error found
   jsr
            InitForIO
                            ; prepare for serial I/O
                             ; primitive write block
   jsr
            WriteBlock
                             ; set status flags
   txa
            99$
   bne
                             ; branch if error found
            verify
                            ; check verify flag
   lda
   beq
            80$
                            ; branch if not verifying
                            ; verify block we wrote
   jsr
            VerWriteBlock
   txa
                            ; set status flags
   bne
            99$
                             ; branch if error found
80$
            DoneWithIO
                            ; restore after I/O done
   jsr
   clc
   rts
                             ; No Errors
99$
            DoneWithIO
                            ; restore after I/O done
   jsr
   sec
                             ; Error Status exit
   rts
```

#### MyReadBlock:

```
***********
  MyReadBlock - Read sector from disk into diskBlkBuf
  Demonstrates use of very-low level disk primitives
           r1L = track #
           r1H = sector #
           r4 = Address of block to read into.
.ramsect
           nextTrack:
                         .block
                                    1
                                    1
           nextSector:
                          .block
           outbuffer:
                          .block
                                   ŚFE
                                   $100
           inbuffer
                          .block
                         .block
           track
                         .block
           sector
            .block 1
   verify:
   .psect
CallMyPutB:
   LoadW
          r4, inBuffer
   MoveB
          track, r1L,
   MoveB sector, r1H
          MyReadBlock
   jsr
   bcs
           99$
   rts
                          ;return good status in carry
99$
                          ; Error Handler or let caller handle error
   rts
MyReadBlock:
                         ; go into turbo mode
   jsr
           EnterTurbo
                          ; check for error in X
   txa
                          ; branch if error found
           99$
   bne
                         ; prepare for serial I/O
   jsr
           InitForIO
                         ; primitive read block
   jsr
           ReadBlock
   jsr
           DoneWithIO
                         ; restore after I/O done (x is preserved in DoneWithIO)
   txa
                         ; get error result of ReadBlock
   bne
          99$
                         ; branch if error found
80$
   clc
   rts
99$
   sec
   rts
```

#### NewAllocateBlock:

```
; NewAllocateBlock - allocate specific block in BAM
; with any CBM GEOS device driver.
; Pass: r6L,r6H track, sector to allocate
; Uses: BAM in curDirHead
; Returns: x error status ($00 = success, BAD BAM = block already in use, etc.)
; Destroys: a,y,r7, r8H.
· *********
                       *******
BAD BAM=$0B
DRV 1571=2
NO ERROR=0
NewAllocateBlock:
   ldy
                           ; get current drive
           curDrive
   lda
            driveType-8,y
                           ; get drive type
           #%00001111
                           ; keep only drive format
   and
   cmp
            #DRV 1571
                           ; see if 1571 or above
           1541$
                            ; branch if 1541
   bcc
                           ; else, use driver routine
            AllocateBlock
   qmr
1541$
           FindBAMBit
   jsr
                           ; get BAM bit info
   beq
           110$
                            ; if zero, then it's not free
           r8H
   lda
                            ; get bit mask for BAM
            #$FF
   eor
                            ; convert to clearing mask
   and
            curDirHead, X
                           ; and with BAM byte to clear
                           ; bit and show as allocated
                           ; and store back.
   sta
            curDirHead, x
   ldx
                           ; get base of track9s entry
            r7H
   dec
            curDirHead, x
                           ; dec # free blocks this track
                           ; show no error
   ldx
            #NO ERROR
   rts
                            ; exit
99$
   ldx
            #BAD BAM
                            ; show error - already in use
   rts
                            ; exit
Example Caller Routine;
.ramsect
   diskBlock .block 2
.psect
CallNewAlloc:
   MoveW
           diskBlock, r6
                          ; block to allocate
            NewAllocateBlock ; (see above)
   jsr
            #BAD BAM
                            ; BAD BAM means block in use
   срх
                            ; branch if block already in use
   beq
            95$
   txa
                            ; check for other error
           99$
                            ; branch if error
   bne
            ; code to handle newly allocated block goes here
95$
                            ; block was not free...
            ; code to handle block already allocated goes here
99$
           MyDiskError ; call error handler with error in x
   jmp
```

#### SaveRecord:

```
SaveRecord - Append new record into am existing VLIR
;
           appendPoint = Already set to the last VLIR Record
                     = Buffer populated with VLIR's filename
NAME LENGTH=17
.ramsect
     appendPoint: .block 1
                                          ; record to append to
                    .block NAME LENGTH
                                          ; hold null-terminated filename
     filename:
     bufStart:
                    .block 1023 ; data buffer
                     .block 1
                                           ; length of buffer
     bufEnd:
                    = (BufEnd - BufStart)+l
     BUFLENGTH
.psect
SaveRecord:
     LoadW r0, #filename ; pointer to filename jsr OpenRecordFile ; open VLIR file
                          ; open VLIR file
     txa
                          ; check open status
     bne
                          ; exit on error
                          ; get record to append to
     lda appendPoint
     jsr PointRecord
                          ; go to that record
                          ; check point status
     txa
                          ; exit on error
     bne 99$
     jsr AppendRecord ; append a record at this point
LoadW r7, #bufStart ; point at data buffer
     LoadW r2, #BUFLENGTH ; bytes in buffer (bufEnd-bufStart)
     jsr WriteRecord
                          ; write buffer to record
     txa
                          ; get write status
                          ; exit on error
     bne 99$
     jsr CloseRecordFile ; close VLIR file
     txa
                           ; check point status
     bne 99$
                          ; exit on error
90$
                          ; Clean Exit
     clc
                          ; clear carry for all ok
     rts
                          ; Error handler
99$:
                           ; Set carry to show returning with an error
     sec
     rts
```

NOTE: geoProgrammer1.1 does not support the \* counter in .ramsect. The method above must be used when the assembler needs to calculate the size of a ramsect field.

Examples	: - <b>-</b> 1
	internal

## internal

## RoadTrip:

```
; Show Leaving GEOS to use all of the resources of the machine and returning again via
; rebooting by either REU disk.
BYTESTOSAVE=$80
                               ; no. of bytes to save at BootGeos
RBOOT BIT=5
                               ; bit in sysFlgCopy to check
CKRNL BAS IO IN=$40
config=$FF00
.ramsect
   GEOS save .block BYTESTOSAVE ; save area for GEOS restart block
.psect
RoadTrip:
                             ; Save Kernal Boot strap
; Do anything ... Use all of kernal ram
   jsr OnEntry
   jsr UnEntry
jsr HaveAFunTrip
                               ; just no kernal calls while you are gone
                               ; Reboot the Kernal
   jmp OnExit
OnEntry:
   ldx #BYTESTOSAVE
                           ; save bytes GEOS needs so we can use area
10$
                                ; STARTLOOP
   lda BootGeos-1,x
                                ; copy a byte
   sta GEOS save-1,x
   dex
                                    count--
   bne 10$
                                ; if (count != 0), then loop
                                ; ENDLOOP
   rts
OnExit:
                               ; Get. version of GEOS
   lda version
   cmp #$13
   bcc 64$
                               ; If version < 1.3, then branch
   lda cl28Flag
                               ; else, test for GEOS 128
                               ; If GEOS64, then branch
   bpl 64$
128$
                             ; load 128 memory mapping
   lda CKRNL BAS IO IN
   sta config
   bra 200$
64$
   lda #KRNL BAS IO IN
                              ; load 64 memory mapping
   sta CPU DATA
$200
   ldx
       #BYTESTOSAVE
                               ; restore bytes GEOS needs to restart
                               ; STARTLOOP
10$
   lda
       GEOS save-1,x
                               ; copy a byte
   sta BootGeos-1,x
   dex
                                    count--
   bne 10$
                                ; if (count != 0), then loop
                                ; ENDLOOP
                              ; check for Rboot flag
   lda #(%1<<RBOOT BIT)
   and sysFlgCopy
   bne 90$
                               ; if flag is clear, branch to rboot
   jsr AskForBootDisk
                              ; else, get user to insert boot disk
90$
   jmp BootGeos
```

	Examples	
İ	-	internal

Examples

graphics

graphics

 ${\tt Change Mode}$ 

# ChangeMode:

#### GREYPAT=2

```
ChangeMode:
        jsr GreyScreen ; grey out old screen
lda graphMode ; switch mode by flipping
eor #%1000000 ; 40/80 bit
jsr SetNewMode ; and calling SetNewMode
jsr GreyScreen ; grey out new screen
                                                   ; and calling SetNewMode
; grey out new screen
        rts
                                                     ; exit
GreyScreen:
      jsr i_GraphicsString
      .byte NEWPATTERN, GREYPAT ; set to grey pattern .byte MOVEPENTO ; Put pen in upper left .word 0 ; x
                                                     ; Put pen in upper left
       .byte 0
                                                     ; у
      .byte RECTANGLETO
                                                    ; grey out entire screen
      .word (SC_PIX_WIDTH-1) | DOUBLE_W | ADD1_W
      .byte SC_PIX_HEIGHT-1
      .byte NULL
      rts
```

#### MseToCardPos:

.if COMMENT

\*\*\*\*\*\*\*\*\*\*\*\*

#### MseToCardPos

converts current mouse positions to card position

pass: Nothing

uses: mouseXPos, mouseYPos

Returns: rOL mouse card x-position (byte)
rOH mouse card y-position (byte)

modes sara i posteron

Destroys: a,x,y

\*\*\*\*\*\*\*\*\*\*\*\*\*

.endif

#### MseToCardPos:

; save current interrupt disable status
; disable interrupts so mouseXPos doesn't change*
; copy mouse x-position to zp work reg (r0)
; get mouse y-position
; reset interrupt status asap.
; divide x-position $(\mathbf{r0})$ by 8
; (shift right 3 times)
; this gives us the card x-position in ${ t r0L}$
; shift y-position in ${f a}$ right 3 times
; which is a divide by 8
; and gives us the card y-position in ${f a}$
; set card y-position
; exit

Note:

If you do not disable interrupts prior to getting the value of mouseXPos you could get r0H with Lydia/site and before getting really an interrupt occurs and the mouse position is updated during the interrupt. Now when you do/star for r0L it is for a different mouseXPos reading giving unpredictable results.

Note<sup>3</sup>:

By also getting the Y value while interrupts are disabled, you are guaranteed to also get a consistent reading for all three parts of the mouse position.

Examples

TCOILS	/menus

icons/menu

ChangeMode

## IconsUp:

```
IconsUp:
  LoadB dispBufferOn, # (ST_WR_FORE | ST_WR_BACK) ; draw to both buffers
   LoadW r0,#IconTable
          Dolcons ; exit
   isr
   rts
Important: Due to a limitation in the icon-scanning code, the application must
         always install an icon table with at least one icon. If the application
         is not using icons, create a dummy icon table with one icon (see
         below).
; NoIcons: Install a dummy icon table. For use in applications that
; aren't using icons. Call early in the initialization of the
; application, before returning to MainLoop.
DummylconTable:
                   ; one icon
; dummy mouse x (don't reposition)
; dummy mouse y
     .byte 1
     .word NULL
.byte NULL
                    ; bitmap pointer to $0000 (disabled)
     .word NULL
                    ; dummy x-pos
     .byte NULL
     .byte NULL
                    ; dummy y-pos
     Nolcons:
     LoadW r0, #DummyIconTable ; point to dummy icon table
     jmp DoIcons ; install. Let DoIcons rts
```

Examples

m01160/	sprite

mouse/sprite

-----

 ${\tt Change Mode}$ 

## ArrowUp:

```
; Put up a new mouse picture
ArrowUp:
  LoadW r0, #DnArrow
                       ; point at new image
  jsr
         SetMsePic
                       ;install it
  rts
;macro to store a word value in high/low order
.macro HILO word
  .byte ]word, [word
.endm
; Mouse picture definition for down-pointing arrow
DnArrow:
  HILO %1111111110000000
                       ;mask
  HILO %11111110011111110
  HILO %0001100111111001
  HILO %0110011111100111
  HILO %0111111110011111
  HILO %0111111110011111
  HILO %00000000001111
  ;image
  HILO %000000001111110
  HILO %000000111111000
  HILO %0110011111100000
  HILO %0111111110000000
  HILO %0111111110000000
  HILO %0111111111100000
  HILO %0000000000000
```

#### utility

------

BeepThrice Beep three times. Runs off the MainLoop by using Sleep

HandleCommand Given a command number this routine handles dispatching

control to the appropriate routine.

LoadBASIC Loads a Commodore BASIC program and starts it running.

# BeepThrice:

```
; Beep three times
; Runs off the MainLoop by using Sleep
.if TARGET NTSC
  FRAME RATE=60
.else
  FRAME RATE=50
.endif
BELL INTERVAL = (FRAME RATE/10) ;approximately. 1/10 second.
BeepThrice:
                        ; sound the bell
    jsr
          Bell
    LoadW r0, BELL_INTERVAL ;
                       ; pause a bit
    jsr Sleep
        Bell
                       ; sound the bell again
    jsr
    LoadW r0, BELL INTERVAL
                       ; pause a bit
    jsr Sleep
                        ; sound the bell again and let bell rts
    jmp
        Bell
```

Note<sup>3</sup>: see **GetFPS** for detecting Frame Rate for portability between hardware.

#### HandleCommand:

```
; HandleCommand
; DESCRIPTION: Given a command number this routine handles dispatching
; control to the appropriate routine.
; Pass: y command number
; Returns: depends on command
; Destroyed: depends on command
UNIMPLEMENTED = $0000
HandleCommand:
                          ; check command # against last cmd num
           #TOT CMDS
   сру
           99$
                           ; exit if command is invalid
   bcs
                          ; get low byte of routine address
   lda
           CMDtabL,y
                          ; get high byte routine address
   ldx
           CMDtabH, y
   jsr
                          ; call the routine
           CallRoutine
99$
   rts
                           ; exit
; The table below is a collection of the the high/low bytes of the routine
; associated with each command number. If a command is not yet implemented
; use the UNIMPLEMENTED constant
CMDtabL:
   .byte
          [UNIMPLEMENTED ; Low Byte of command 0
   .byte
          [Cmdl
                          ; Low Byte of command 1
           [Cmd2
   .byte
                          ; etc...
   .byte
           [Cmd3
           [Cmd4
   .byte
   .byte
           [Cmd5
          ;low bytes
CMDtabH:
   .byte ]UNIMPLEMENTED ; High Byte of command 0
   .byte
          ]Cmdl
                          ; High Byte of command 1
   .byte
          1Cmd2
                          ; etc...
           ]Cmd3
   .byte
           ]Cmd4
   .byte
   .byte
           ]Cmd5
TOT CMDS = (CMDtabH-CMDtabL) ; Total Number of commands
   ; Perform some action here.
Cmd2:
   ; Perform some action here.
Cmd3:
   ; Perform some action here.
   rts
Cmd4:
   ; Perform some action here.
   rts
Cmd5:
   ; Perform some action here.
   rts
```

## i VerticalLine:

```
; Inline version of VerticalLine.
; Pass:
; .word xl
; .word x2
; .byte y1
IVERT_BYTES = 5 number of inline bytes in call
i_VerticalLine:
     ;--- Save away the inline return address
     PopW returnAddress
     ;--- Load up VerticalLine's parameters
     ldy #VJBYTES
     lda
        sta rllL
10$
                            ; load other params in a loop
    dey
        (returnAddress),y
                            ; They occupy consecutive GEOS
    lda
    sta r3L-1, y
                            ; pseudoregisters, so this will,
     cpy #1
                             ; work correctly
        10$
    bne
     ;--- Now call VerticalLine with registers loaded
     jsr VerticalLine
     ;--- and do an inline return
    php ; save st reg to return
lda #IVERT_BYTES +1 ; # of bytes + 1
jmp DoInlineReturn ; jump to inline return. Do not jsr!
```

#### LoadBASIC:

```
Loads a Commodore BASIC program and starts it
; running. Assumes that the program is a standard BASIC
; file that loads at $801. This example does little
; error checking.
; Pass:
           Nothing
*******
UNIMPLEMENTED = $0000
basicProg:
   .byte "GodZilla", NULL
runCommand:
   .byte "RUN", NULL
LoadBASIC:
   LoadW
           r6,basicProg ; Find Basic Program to run
   jsr
           FindFile
                           ; r5 will now point to programs DIR Entry
   txa
           99$
                           ; If FILE NOT FOUND or other Disk Errors exit.
   bne
           r0, runCommand ; point at command string
   LoadW
           r7, #$801
                            ; assume standard address
   LoadW
   jmp
           ToBasic
99$
   sec
   rts
```

# Icons, Menus, and Other Mouse Presses

When the user clicks the mouse button, GEOS determines whether the mouse pointer was positioned over an icon, a menu item, or some other region of the screen. GEOS has a unique method of handling a mouse press for each of these cases. If the user pressed on an icon, GEOS calls the appropriate icon event routine. If the user pressed on a menu, GEOS opens up a submenu or calls the appropriate menu event routine, whichever is applicable. And if the user pressed somewhere else, GEOS calls through **otherPressVector**, letting the application handle (or ignore) these "other" mouse presses.

#### **Icons**

When you open a disk by clicking on its picture, delete a file by dragging it to the trash can, or click on the CANCEL button in a dialog box, you are dealing with *icons*, small pictorial representations of program functions. A GEOS icon is a bitmapped image, whether die picture of a disk or a button-shaped rectangle, that allows the user to interact with the application. When the application enables icons, GEOS draws them to the screen and then keeps track of their positions. When the user clicks on an icon, an icon event is generated, and the application is given control with information concerning which icon was selected.

#### **Icon Table Structure**

The information for all active screen icons is stored in a data structure called the *icon table*. GEOS only deals with one icon table at a time. The icon table consists of an *icon table header* and a number of *icon entries*. The whole table is stored sequentially in memory with the header first, followed by the individual icon entries.

## **Icon Table Header**

The icon table header is a four byte structure which tells GEOS how many icons to expect in the structure and where to position the mouse when the icons are enabled. It is in the following format:

#### **Icon Table Header:**

Index	Constant	Size	Description
+0	OFF_NM_ICNS	byte	Total number of icons in this table.
+1	OFF_IC_XMOUSE	word	Initial mouse x-position. If \$0000, mouse position will not be altered.
+3	OFF_IC_YMOUSE	byte	Initial mouse y-position.

This first byte reflects the number of icon entries in the icon table (and, hence, the number of icons that can be displayed). The table can specify up to MAX\_ICONS icons.

The next word (bytes 1 and 2) is an absolute screen x-coordinate and the following byte (byte 3) is an absolute screen y-coordinate. The mouse will be positioned to this coordinate when the icons are first displayed. If you do not want the mouse positioned, set the x-coordinate word to \$0000, which will signal Dolcons to leave the mouse positions alone.

#### **Icon Entries**

Following the icon table header are the icon entries, one for each specified in the OFF\_I\_NUM byte in the icon table header. Each icon entry is a seven-byte structure in the following format:

#### **Icon Entries:**

Index	Constant	Size	Descr	ription
+0	OFF_I_PIC		word	Pointer to compacted bitmap picture data for this Icon. If set
				to \$0000, icon is disabled.
+2	OFF_I_X		byte	Card x-position for icon bitmap.
+3	OFF_I_Y		byte	Y-position of icon bitmap.
+4	OFF_I_WIDTH		byte	Card width of icon bitmap.
+5	OFF_I_HEIGHT		byte	Pixel height of icon bitmap.
+6	OFF_I_EVENT		word	Pointer to icon event routine to call if this icon is selected.

Note: OFF\_I\_NEXT=8 Offset to Next Icon in structure if it exists.

The first word (**OFF\_I\_PIC**) is a pointer to the compacted bitmap data for the icon. The icon can be of any size (up to the full size of the screen). If this word is set to **NULL** (\$0000), the icon is disabled.

The third byte (**OFF\_I\_X**) is the x byte-position of the icon. The x byte-position is the x-position in bytes. Icons are placed on the screen by **BitmapUp** and so must appear on an eight-pixel boundary. The byte-position can be calculated by dividing the pixel-position by eight ( $x_byte_position = x_pixel_position/8$ ).

The fourth byte (**OFF\_I\_WIDTH**) is the pixel position of the top of the icon. The icon will be placed at (x\_byte\_position\*8, y\_pixel\_position).

The next two bytes (**OFF\_I\_WIDTH** and **OFF\_I\_HEIGHT**) are the width in bytes and height in pixels, respectively. These values correspond to the geoProgrammer internal variables **PicW** and **PicH** when they are assigned immediately after a pasted icon image.

The final word (**OFF\_I\_EVENT**) is the address of the icon event handler associated with this icon.

# **Sample Icon Table**

The following data block defines three icons which are placed near the middle of the screen. The mouse is positioned over the first icon:



PAINTW = PicW PAINTH = PicH PAINTX = 16/8 PAINTY = 80

PaintIcon:

#### Writelcon:



WRITEW = PicW WRITEH = PicH

WRITEX = PAINTX + PAINTW + I SPACE

WRITEY = PAINTY

#### Publishlcon:



PUBLISHW = PicW PUBLISHH = PicH

PUBLISHX = WRITEX + WRITEW + I SPACE

PUBLISHY = WRITEY

; The actual icon data structure to pass to Dolcons follows IconTable:

#### I header:

.byte NUMOFICONS

.byte PAINTY + PAINTH/2

#### I entries:

#### PaintIStruct:

.word PaintEvent ; event handler

#### WriteIStruct:

. WOIG WILLELCON ; pointer to bitmap
.byte WRITEX, WRITEY ; icon position
.byte WRITEW, WRITEH ; icon width, height
.word WriteEvent ; icon position ; icon width, height .word WriteEvent ; event handler

#### PublishIStruct:

; event handler .word PublishEvent

NUMOFICONS = (\*-I\_entries)/IESIZE ; number of icons in table

; Dummy icon event routines which do nothing but return PaintEvent:

WriteEvent:

PublishEvent:

rts

# **Installing Icons**

When an application is first loaded, GEOS will not have an active icon structure. GEOS must be given the address of the applications icon table before MainLoop can display and track the user's interaction with them. GEOS provides one routine for installing icons

• **DoIcons** Display and activate an icon table.

**DoIcons** draws the enabled icons and instructs MainLoop to begin watching for a single- or double-click on one. The icon table stays activated and enabled until the ICONS\_ON\_BIT of mouseOn is cleared or another icon table is installed by calling Dolcons with the address of a different icon structure. In either case, the old icons are not erased from the screen by GEOS.

Dolcons will draw to the foreground screen and background buffer depending on the value of dispBufferOn. Icons are usually permanent structures in a display and so often warrant being drawn to both screens. If icons are only drawn to the foreground screen, they will not be recovered after a menu or dialog box.

Example: **IconsUp** 

**Important:** Due to a limitation in the icon-scanning code, the application must always install an icon table with at least one icon. If the application is not using icons, create a dummy icon table with one icon (see below).

```
; NoIcons Install a dummy icon table. For use in applications that
         aren't using icons. Call early in the initialization of
         the application, before returning to MainLoop.
Nolcons:
    LoadW r0, #DummyIconTable ; point to dummy icon table
           DoIcons
                            ; install. Let Dolcons rts
    qmp
DummyIconTable:
    .byte 1
                            ; one icon
    .word $0000
                            ; dummy mouse x (don't reposition)
    .byte $00
                            ; dummy mouse y
    .word $0000
                            ; bitmap pointer to $0000 (disabled)
    .byte $00
                            ; dummy x-pos
    .byte $00
                            ; dummy y-pos
    .byte 1,1
                            ; dummy width and height
     .word $0000
                            ; dummy event handler
```

# **MainLoop and Icon Event Handlers**

When the user clicks the mouse button on an active icon, GEOS **MainLoop** will recognize this as an icon event and call the icon event handler associated with the particular icon. The icon event handler is given control with the number of the icon in r0L (the icon number is based on the icon's position in the table: the first icon is icon 0). Before the event handler is called, though, **MainLoop** might flash or invert the icon depending on which of the following values is in **iconSelFlag**:

## Constants for iconSelFlag:

ST_NOTHING	\$00	The icon event handler is immediately called; the icon image is untouched
ST_FLASH	\$80	The icon is inverted for selectionFlash vblanks and then reverted to its normal state before the
		event handler is called.
ST_NVERT	\$40	The icon is inverted (foreground screen image only) before the event handler is called. The event
		handler will usually want to revert the image before returning to <b>MainLoop</b> by calculating the
		bounding rectangle of the icon, loading <b>dispBufferOn</b> with <b>ST_WR_FORE</b> , and calling
		InvertRectangle.

# **Detecting Single- and Double-clicks on Icons**

When the user first clicks on an icon, GEOS loads the global variable **dblClickCount** with the GEOS constant **CLICK\_COUNT**. GEOS then calls the icon event handler with r0H set to **FALSE**, indicating a single-click. **dblClickCount** is decremented at interrupt level every vblank. If the icon event handler returns to MainLoop and the icons user again clicks on the icon before **dblClickCount** reaches zero, GEOS calls the icon event handler a second time with **r0H** set to **TRUE** to indicate a double-click.

Checking for a double-click or a single-click (but not both) on a particular icon is trivial: merely check **r0H**. If **r0H** is **TRUE** when you're looking for a single-click or its **FALSE** when you're looking for a double-click, then return to MainLoop immediately. Otherwise, process the click appropriately. This way, if the user single-clicks on an icon which requires double-clicking or double-clicks on an icon which requires single-clicking, the event will be ignored.

However, checking for both a double- or a single-click on the same icon (and performing different actions) is a bit more complicated because of the way double-clicks are processed: during the brief interval between the first and second clicks of a double-click, the icon event handler will be called with r0H set to **FALSE**, which will appear as a single-click; when the second press happens before **dblClickCount** hits zero, the icon event handler is called a second time with **r0H** set to TRUE, which will appear as a double-click. There is no simple way (using the GEOS double click facility) to distinguish a single-click which is part of a double-click from a single-click which stands alone.

There are two reliable ways to handle single- and double-click actions on icons: the additive function method and the polled mouse method. The additive function method relies on a simple single-click event which toggles some state in the application and a double-click event (usually more complicated) which happens in addition to the single-click event. The GEOS deskTop uses the additive function method for selecting (inverting) file icons on a single-click and selecting and opening them on a double-click. The icon event handler first checks the state of **r0H**. If it is **FALSE** (single-click) then the icon (and an associated selection flag) is inverted. If it is **TRUE** (double-click) then the file is opened. If the user single-clicks, the icon is merely inverted. If the user double-clicks, the icon is inverted (on the first click) and then processed as if opened (on the second click).

#### Example:

```
; Icon double-click handler
; additive function method
IconEventl:
    1 da
        r0H
                     ; check double-click flag .
        10$
                     ; branch if second click of a double-click
    bne
                     ; else, this is a single-click or the
                     ; first push of a double-click,
    jsr
        InvertIcon
                    ; so just invert the selection
    bra
        90$
10$
        OpenIcon
                     ; double-click detected, go process it
    jsr
90$
                     ; exit
    rts
```

The polled-mouse method can be used when the single-click and double-click functions are mutually exclusive. When a single-click is detected the icon event handler, rather than returning to **MainLoop** and letting GEOS manage the double-click, handles it manually by loading **dblClickCount** with a delay and watching **mouseData** for a release followed by a second click.

#### Example:

```
* ***********************
; Icon double-click handler
; polled mouse method Open Icon
 ******************
IconEvent2:
 ;--- User pressed mouse once, start double-click counter going
     LoadB dblClickCount, #CLICK COUNT ; start delay
;--- Loop until double-click counter times-out or button is released
10$
     lda
          dblClickCount
                                      ; check double-click timer
                                      ; If timed-out, no double-click
          40$
     beq
     lda
          mouseData
                                      ; Else, check for release
          10$
                                      ; loop until released
     bpl
 ; --- mouse was released, loop until double-click counter times-out or
;--- button is pressed a second time.
20$
     lda
          dblClickCount
                                      ; check double-click timer
     beq
          30s
                                      ; If timed-out, no double-click
     lda mouseData
                                      ; Else, check for second press
     bmi
          20$
                                      ; loop until pressed
;--- Double-click detected (no single-click)
30$
    jmp DoDoubleClick
                                      ; do double-click stuff
;--- Single-click detected (no double-click)
40$ jmp DoSingleClick
                                     ; do single-click stuff
```

Note:

These techniques for handling single- and double-clicks are described here as they pertain to icons; they are not directly applicable to applications that detect mouse clicks through otherPressVector. When control vectors through otherPressVector, the value in r0H is meaningless. For more information on otherPressVector, refer to "Other Mouse Presses" in this chapter.

# Other Things to Know About Icons

#### Icon Releases and otherPressVector

When the user clicks on an active icon, MainLoop will call the proper icon event routine rather than vectoring through otherPressVector. However, the routine pointed to by otherPressVector will get called when the mouse is released. Applications that aren't using otherPressVector can disable this vectoring by storing a. \$0000 into OtherPressVector (\$0000 is actually its default value). Applications that depend on otherPressVector, however, can check mouseData and ignore all releases.

#### **Example:**

```
;OtherPressVector routine that ignores releases (high bit of mouseData is set on
releases)
MyOtherPress:
                                    ; control comes here from otherPressVector
            lda
                  mouseData
                                    ; check state of the mouse button
                  90$
                                    ; ignore it if it's a release
            bmi
            jsr
                  PressDown
                                    ; otherwise process the press
90$:
            rts
                                    ; exit
```

For more information on **otherPressVector**, refer to "Other Mouse Presses" in this chapter.

#### **Icon Precedence**

GEOS draws icons sequentially. Therefore, if icons overlap, the ones which are drawn later will be drawn on top. When the user clicks somewhere on the screen, GEOS scans the icon table in this same order, looking for an icon whose rectangular boundaries enclose the coordinates of the mouse pointer. If more than one icon occupies the coordinate position, the icon that is defined first in the icon table (and therefore drawn on bottom) will be given the icon event. If an active menu and an icon overlap, the menu will always be given precedence.

#### **Disabling Icons**

An application can disable an icon in the current icon structure by clearing the **OFF\_I\_PIC** word of the icon (setting it to \$0000). If an icon is disabled prior to a call to **Dolcons**, the icon will not be drawn. If an icon is disabled after the call to Dolcons, the icon will remain on the screen but will be ignored during the icon scan. The application can reenable the icon by restoring the **OFF\_I\_PIC** word to its original value. (Actually, any non-zero value will do because reenabling an icon does not redraw it, it only restores the coordinates to MainLoop's active search list.)

## **GEOS 128 Icon Doubling**

As with bitmaps, special flags in the icon data structure can be set to automatically double the xposition and/or icon width when GEOS 128 is running in 80-column mode. To have an position icon's x-position automatically doubled in 80-column mode, bitwise-or the **OFF\_I\_X** parameter with **DOUBLE\_B**. To double an icon's width in 80-column mode, bitwise-or the **OFF\_I\_WIDTH** parameter with **DOUBLE\_B**. These bits will be ignored when GEOS 128 is running in 40-column mode. Do not, however, use these doubling bits when running under GEOS 64. GEOS 64 will try to treat the doubling bit as part of the coordinate or width value rather than a special-case flag. For more information, refer to "GEOS 128 X-position and Bitmap Doubling" in Chapter "Graphics Routines".

#### Example:

```
; ***********************************
; SAMPLE GEOS 128 ICON TABLE THAT USES AUTOMATIC DOUBLING FEATURE
; using compiler flags for conditional assembly between C128 and C64
; Note: You can build programs that work on both the 128 in 80cols
; and the 64.
 ******************
C128=TRUE
C64=FALSE
.if !C128
     .echo Error: cannot assemble GEOS 128 specific code without C128 flag set
.else
PaintIcon:
PAINTW = PicW
PAINTH = PicH
PAINTX = 16/8
PAINTY = 80
; The actual icon data structure to pass to Dolcons follows
IconTable:
I header:
  .byte NUMOFICONS
   .word ((PAINTX*8) + (PAINTW*8/2)) | DOUBLE_W ; position mouse over paint icon
  .byte PAINTY + PAINTH/2
I entries:
PaintIStruct:
                                   ; pointer to bitmap
   .word PaintIcon
                                      ; x card position (dbl in 80-column mode)
  .byte PAINTX | DOUBLE B
                                     ; y-position
  .byte PAINTY
  .byte PAINTW | DOUBLE B
                                      ; icon width (dbl in 80-column mode)
  .byte PAINTH
                                      ; icon height
                                      ; event handler
  .word PaintEvent
NUMOFICONS - (*-I entries)/OFF_I_NEXT ; number of icons in table
; Dummy icon event routines which do nothing but return
PaintEvent:
  rts
.endif
213
```

	Overviews
	Icons, Menus, and Other Mouse Presses
214	

# **Menus**

Menus, one of the most common and powerful user-interface facilities provided by GEOS, allow the application to offer lists of items and options to the user. The familiar menus of the GEOS desktop, for example, provide options for selecting desk accessories, manipulating files, copying disks, and opening applications. Virtually every GEOS-based program will take advantage of these capabilities, providing a consistent interface across applications.

GEOS menus come in two flavors: horizontal and vertical. The main menu, the menu which is always displayed, is usually of the horizontal type and is typically placed at the top of the screen. Each selection in the main menu usually has a corresponding vertical sub-menu that opens up when an item in the main menu is chosen. These sub-menus can contain items that trigger the application to perform some action. They can also lead to further levels of sub-menus. For example, a horizontal main menu item can open up to a vertical menu, which can have items which then open up other horizontal sub-menus, which can then lead to other vertical menus, and so on.

#### **Division of Labor with Menus**

GEOS divides the labor of handling menus between itself and the application. The GEOS Kernal handles all of the user's interaction with the menus. This includes drawing the menu items, opening up necessary sub-menus, and restoring the Screen area from the background buffer when the menus are retracted. MainLoop manages the menus, keeping track of which items the user selects. If the user moves off of the menu area without making a selection, GEOS automatically retracts the menus without alerting the application.

If the user selects a menu item which generates a menu event, the application's menu event handler is called with the menus left open. Leaving the menus open allows the application to choose when and how to retract them: all the way back to the main menu, up one or more levels (for multiple sub-menus), or up no levels (keeping the current menu open). This lets the application choose the menu level which is given control upon return, thereby allowing multiple selections from a sub menu without forcing the user to repeatedly traverse the full menu tree for each option.

#### Menu Data Structure

The main menu, all its sub-menus, their individual selectable items, and various attributes associated with each menu and each item are all stored in a hierarchical data structure called the menu tree. Conceptually, a menu tree with multiple sub-menus might have the following layout:

## Examples

structures

structures

dialog/Icons/Menus/Graphics

#### DIALOG:

#### Note<sup>2</sup>:

The first entry in a DB table is a command byte defining its position. This can either be a byte indicating a default position for the DB, DEF\_DB\_POS (%10000000), or a byte indicating a user defined position, SET\_DB\_POS (%0000000) which must be followed by the position information.

The position command byte is or'ed with a system pattern number to be used to fill in a shadow box. The shadow box is a rectangle of the same dimensions as the DB and is filled with one of the system patterns. The shadow box appears underneath the Dialog Box, Offset 1 card right and 1 card down.

Start of Default Dialog

.byte DEF\_DB\_POS | pattern

.byte top ; (0-199)

.byte bottom ; (0-319 or 0-639)

.word right ; (0-319 or 0-639)

Note<sup>1</sup>: standard window size: columns 72-263 rows 40-135

Note1: If the shadow pattern is zero, then no shadow is drawn.

Note: Icon descriptors are stored in a table at \$880C

Note<sup>3</sup>: Maximum # of Dialog Icons is 8. This can be worked around by drawing your own images and detecting mouse clicks over the images.

**Note**: The following is a list of global variables stored by the window processor:

baselineOffset curPattern string curSetWidth curHeight curingerial
dispBufferOn mouseOn msePicPtr window10
windowBottom leftMargin rightMargin appMain
intTopVector ioBotVector mouseVector keyVector
inputVector mouseFaultVec otherPressVector alarmTmt
selectionFlash alphaFla curIndexTable cardDataPntr currentMode windowTop keyVector otherPressVector alarmTmtVector alphaFlag iconSelFlag faultData menuNumber mouseTop mouseBottom mouseLeftmouseRight stringX,stringY

I/O address's \$D000-\$D010 \$D01B-\$D01D \$D025-\$D026 \$D015 \$D028-\$D02E

#### Position Commands:

After the position byte (or bytes) may appear a number of icon or command bytes. Most require position coordinates. The  ${\bf x}$  and  ${\bf y}$  positions are an offset from the upper left corner of the DB.

Icons x position uses bytes (cards) 0-40 x\_boffset Text x position uses pixels 0-319 x\_poffset y position is always in pixels 0-199.

## Note<sup>3</sup>:

GEOS 128 always doubles the x positions in a dialog box when the system is in 80 column mode. Do not try to use **DOUBLE\_W** as this will be a VERY large x coordinate. **DBUSERICON** Structures DO need **DOUBLE\_B** for width if the user icon is not a native 80 col icon.

Dialog Box Icons

Icon	Value	Example	Description
OK	1	<pre>.byte OK .byte x_boffset .byte y_offset</pre>	Draw OK Icon
CANCEL	2		Draw CANCEL Icon
YES	3		etc
NO	4		
OPEN	5		
DISK	6		
NOT-USED	7-10		Marked for future use. When is the future?

Dialog Commands

Dialog Commands				
Command	Value	Example	Description	
DBTXTSTR	11	.byte <b>DBTXTSTR</b>	Put tTextStr	
		.byte x_poffset		
		.byte y_offset		
		.word ptrtTextStr		
DBVARSTR	12	.byte DBVARSTR	Put text @@zPgPtr	
		.byte x_poffset	zPgPtr is an address of a zero	
		.byte y_offset	page ptr to string	
		.byte zPgPtr		
DBGETSTRING	13	.byte <b>DBGETSTRING</b>	Get typed user input. ZpgPtr	
		.byte x_poffset	points to address of a buffer to	
		.byte y_offset	use for the input that is	
		.word ZPgPtr	BUFFERSIZE bytes.	
		.byte BUFFERSIZE		
SBSYSOPV	14	.byte <b>SBSYSOPV</b>	Closes DB when the mouse is	
			pressed anywhere other then over	
			an Icon	
DBGRPHSTR	15	.byte DBGRPHSTR	Draws a <b>GraphicsString</b>	
		_		
DBGetFileS	16	.word gGraphicsString	Presents a File Selection box	
DBGetFileS		<pre>.word gGraphicsString .byte DBGetFileS</pre>	Presents a File Selection box	
DBGetFileS		<pre>.word gGraphicsString .byte DBGetFileS .byte x_boffset</pre>		
DBGetFileS DBOPVEC		<pre>.word gGraphicsString .byte DBGetFileS</pre>	Presents a File Selection box	
	16	<pre>.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset</pre>	Presents a File Selection box for the user to pick from.	
	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y offset .byte DBOPVEC	Presents a File Selection box for the user to pick from.  Vector to call when mouse button	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y offset .byte DBOPVEC .word MsePressVector	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset .byte y_offset	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData .word NULL	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset .byte y_offset	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData .word NULL .byte width in bytes	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset .byte y_offset	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData .word NULL .byte width in bytes .byte Height in Pixels	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset .byte y_offset	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData .word NULL .byte width in bytes .byte Height in Pixels	
DBOPVEC	16	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y_offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset .byte y_offset	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData .word NULL .byte width in bytes .byte Height in Pixels .word ptrIconAction	
DBUSERICON	16 17 18	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset .byte y_offset .word UserIcon	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData .word NULL .byte width in bytes .byte Height in Pixels .word ptrIconAction  Note: (width   DOUBLE_B for 128)	
DBUSERICON	16 17 18	.word gGraphicsString .byte DBGetFileS .byte x_boffset .byte y offset .byte DBOPVEC .word MsePressVector .byte DBUSERICON .byte x_boffset .byte y_offset .word UserIcon  .byte DB_USR_ROUT	Presents a File Selection box for the user to pick from.  Vector to call when mouse button is pressed.  UserIcon Table .word ptrIconData .word NULL .byte width in bytes .byte Height in Pixels .word ptrIconAction  Note: (width   DOUBLE_B for 128)  Call User_Vector after the DB	

## Menu

## Menu/Sub Menu Header

## Size

+0	OFF_M_Y_TOP	byte	Top edge of menu rectangle (yl pixel position).
+1	OFF_M_Y_BOT	byte	Bottom edge of menu rectangle (y2 pixel position).
+2	OFF_M_X_LEFT	word	Left edge of menu rectangle (xl pixel position).
+4	OFF_M_X_RIGHT	word	Right edge of menu rectangle (x2 pixel position).
+6	OFF_NUM_M_ITEMS	byte	Menu type bitwise-or'ed with number of items in this
			menu/sub-menu.

## Menu/Sub-menu Types (use in attribute byte):

HORIZONTAL	Arrange menu items in this menu/sub-menu horizontally.
VERTICAL	Arrange menu items in this menu/sub-menu vertically.
CONSTRAINED	Constrain the mouse to the menu/sub-menu. If the menu is a sub-menu, the mouse can still be moved off to the parent menu (off the top of a vertical sub menu or off the left of a horizontal menu).
UNCONSTRAINED	Do not constrain the mouse to the menu/sub-menu. If the user moves off of the menu, GEOS will retract it

## Bitwise Breakdown of the Attribute Byte:

7	6	5	4	3	2	1	0
b7	b6			b5-k	o 0		

b7 orientation: 1 = vertical; 0 = horizontal,

b6 constrained: 1 = yes; 0 = no.

b5-b0 number of items in menu/sub-menu (up to  ${\tt MAX\_M\_ITEMS}$ ).

Examples

_	 ru	_	┺-	_	_	_	_
S	 411	ıc:	<b>L</b> .	11	1	е.	8

disk

-----

## structures

# Directory Entry:

Offset	Hex Dec	Constant	# Bytes	Description	
00		OFF CFILE TYPE	1	DOS file type	
			Bit 7	1=File Closed	
			Bit 6	Write Protect	bit
				<b>ST WR PR</b> %0100	0000
			Bit 0-2		
				DEL=0	- 1 F -
				SEQ=1	
				<b>PRG</b> =2	
				USR=3 (GEOS u	ses USR)
				,	rmitted under GEOS)
				<b>CBM</b> =5	,
01		OFF INDEX PTR	2	Index table po	inter (VLIR file T/S)
		OFF DE TR SC			or file's 1st data block
03		OFF FNAME	16		ed with hard spaces \$A0
\$13	19	OFF GHDR PTR	2	_	f GEOS Header block
\$15	21	OFF GSTRUC TYPE	1	GEOS file stru	cture type
	·			SEQUENTIAL=0	
				VLIR=1	
\$16	22	OFF GFILE TYPE	1	GEOS file type	indicator
				NOT GEOS=0	
				BASIC=1	;C-64 Basic w/Header
				ASSEMBLY=2	;C-64 Assembly w/Header
				<b>DATA</b> =3	;C-64 DATA File w/Header
				SYSTEM=4	;GEOS System File
				DESK ACC=5	;GEOS desk accessory
				APPLICATION=6	;GEOS application
				APPL_DATA=7	;GEOS data file
				FONT=8	;GEOS font
				PRINTER=9	;GEOS Print Driver
					0 ;GEOS mouse etc.
				DISK_DEVICE=11	;GEOS DISK driver
				SYSTEM_BOOT=12	;GEOS boot file
					;GEOS Swap File
					ill automatically delete
				all temporary	files when opening a
				disk.)	
				AUTO_EXEC=14	;Application to
					be ran just after
					efore deskTop runs.
				_	;128 Input driver
\$17	-	OFF_YEAR	5	Y/M/D/H/M	
\$1C	28	OFF_SIZE	2	File Size in b	locks

Examples	
	structures

Appendex

hardware

Appendex hardware

dialog

**6510 data register:** (64,128) 01

## CPU\_DATA

KRNL_BAS_IO_IN	RAM_64K	IO_IN
8k KERNAL ROM	8K RAM	8K RAM
I/O	4K RAM	I/O
4K RAM	4K RAM	4K RAM
8K BASIC	8K RAM	8K RAM
24K RAM	24K RAM	24K RAM
Zero Page	Zero Page	Zero Page

See also: TestPoint, DrawLine, GetScanLine.

## 17XX RAM Expansion:

```
EXP BASE:
$DF00: STATUS REGISTER
       Bit 7: INTERRUPT PENDING (1 = interrupt waiting to be served)
              Not Used by GEOS
       Bit 6: END OF BLOCK (1 = transfer complete)
               unnecessary
       Bit 5: FAULT (1 = block verify error)
               Set if a difference between C64- and REU-memory areas was found
               during a compare-command.
       Bit 4: SIZE (1 = 256 \text{ KB}) set on 1764 and 1750 and clear on 1700.
       Bits 3..0: VERSION
$DF01: COMMAND REGISTER. Write to this register to start operation
       Bit 7: EXECUTE (1 = transfer per current configuration)
               Set this bit to execute a command.
       Bit 6: reserved (normally 0)
       Bit 5: LOAD (1 = enable autoload option)
               With autoload enabled the address and length registers (see
               below) will be unchanged after a command execution.
               Otherwise the address registers will be counted up to the
               address off the last accessed byte of a DMA + 1,
               and the length register will be changed (normally to 1).
       Bit 4: FF00
               If this bit is set command execution starts immediately
               after setting the command register.
               Otherwise command execution is delayed until write access to
              memory position $FF00
       Bits 3..2: reserved (normally 0)
       Bits 1..0: TRANSFER TYPE
          00 = transfer C64 -> REU
          01 = transfer REU -> C64
          10 = swap C64 < -> REU
          11 = compare C64 - REU
$DF02: .word C64 BASE ADDRESS
$DF04: .word REU BASE ADDRESS
$DF05: .byte
             BANK
$DF07: .word Transfer Size
$DF09: Interrupt Mask Register.
                                 Not used by GEOS
$DF0A: Address Control Register.
       Bit 7: C64 ADDRESS CONTROL (1 = fix C64 address)
       Bit 6:
                 REU ADDRESS CONTROL (1 = fix REU address)
       Bits 5..0: unused
```

Note: By using a fixed address in the REU as a source you can very quickly initialize large blocks of ram

## Full Reference

http://www.zimmers.net/anonftp/pub/cbm/documents/chipdata/programming.reu
Richard Hable

#### C128 MMU:

```
Configuration Register MMUReg=$FF00 ;Mirror of D500. FF00 is Always Visible
```

```
;MMUReg Bits
                 ;Zone 5
                          $D000-DFFF
;- Bit 0
MIO =%0
                 ; I/O
MCROM =%1
                ;Character ROM
                ;Zone 2 ;$4000-7FFF
;- Bit 1
                ;Basic ROM
MBASIC =%00
                ;External Function ROM
MEXTROM=%10
;- Bits 2,3
                ;Zone 3 $8000-BFFF
MUBASIC=%0000
                ;Basic ROM
MUIROM =%0100
                 ; Internal Function ROM
                 ;External Function ROM
MUEROM =%1000
                ;RAM
MURAM =%1100
                            $C000-CFFF,
;- Bits 4,5
                ;Zone 4
                            $E000-FEFF
MHKERNAL = %000000 ; KERNAL ROM
MHIROM=%010000 ;Internal Function ROM
                 ;External Function ROM
MHEROM=%100000
                 ;RAM
MHERAM=%110000
                ;Bank Select
;- Bits 6,7
                ;Bank 0
MBANK0=%0000000
MBANK1=%01000000 ;Bank 1
MBANK2=%10000000 ;Bank 2
MBANK3=%11000000 ;Bank 2
```

	Configuration Register
Bits	Description
0	D000-DFFF
	0 I/O
	1 1 RAM or Character ROM
1	4000-7FFF
	0 BASIC ROM low
	1 RAM
2-3	8000-BFFF
	00 Basic ROM
	01 Internal Function ROM
	10 External Function ROM
	11 RAM
4-5	C000-CFFF,E000-EFFF
	00 Kernal ROM
	01 Internal Function ROM
	10 External Function ROM
	11 RAM
6-7	Bank Select
	00 Bank 0
	01 Bank 1
	10 Bank 2
	11 Bank 3

```
BANK 0 = MBANK0 | MHERAM | MURAM | MEXTROM | MCROM ; NO ROMS, RAM 0
BANK 0 = %001111111; No ROMs, RAM 0
BANK 1 = %01111111 ; No ROMs, RAM 1
BANK 2 = $10111111 ; No ROMs, RAM 2 ; Requires 512k expanded 128.
                                      ;Otherwise same as bank 0
BANK 3 = %11111111 ; No ROMs, RAM 3 ; Requires 512k expanded 128.
                                      ;Otherwise same as bank 0
BANK 4 = MBANKO | MHIROM | MUIROM | MEXTROM | MIO
BANK 5 = MBANK1 | MHIROM | MUIROM | MEXTROM | MIO
BANK 6 = MBANK2 | MHIROM | MUIROM | MEXTROM | MIO
BANK 7 = MBANK3 | MHIROM | MUIROM | MEXTROM | MIO
BANK 8 = MBANKO | MHEROM | MUEROM | MEXTROM | MIO
BANK 9 = MBANK1 | MHEROM | MUEROM | MEXTROM | MIO
BANK 10 = MBANK2 | MHEROM | MUEROM | MEXTROM | MIO
BANK 11 = MBANK3 | MHEROM | MUEROM | MEXTROM | MIO
BANK 12=%00000110 ;int function ROM, Kernal and IO, RAM 0
BANK 13=%00001010;
BANK 14 = %00000001; all ROMs, char ROM ram 0
BANK 15=%00000000; all ROMs, RAMO power on default
BANK 99=$00001110 ; IO, KERNAL, RAM 0 48K
```

```
Ram Configuration Register
MMURCR=$FF06; Mirror of D506. FF06 is Always Visible
```

```
;MMUReg Bits
;- Bit 0
                  ;Zone 5 $D000-DFFF
MIO
    =%0
                  ;I/O
MCROM =%1
                  ;Character ROM
;- Bit 1
                  ;Zone 2 ;$4000-7FFF
MBASIC =%00
                 ;Basic ROM
MEXTROM=%10
                  ;External Function ROM
;- Bits 2,3
                  ;Zone 3
                            $8000-BFFF
MUBASIC=%0000
                  ;Basic ROM
MUIROM =%0100
                  ; Internal Function ROM
MUEROM =%1000
                  ;External Function ROM
MURAM =%1100
                  ;RAM
                  ;Zone 4
                             $C000-CFFF,
;- Bits 4,5
                             $E000-FEFF
MHKERNAL = %000000 ; KERNAL ROM
MHIROM=%010000
                  ; Internal Function ROM
MHEROM=%100000
                  ;External Function ROM
MHERAM=%110000
                  ;RAM
;- Bits 6,7
                 ;Bank Select
MBANK0=%0000000 ;Bank 0
MBANK1=%01000000 ;Bank 1
                  ;Bank 2
MBANK2=%1000000
                 ;Bank 2
MBANK3=%11000000
```

F	RAM Configuration Register				
Bits	Description				
0-1	Size of Common Ram				
	00 1k				
	01 4k				
	10 8k				
	11 16K				
2-3	Common Ram Location				
	00 Disabled				
	01 Bottom				
	10 Top				
	11 Both				
4-5	Not Used				
6-7	Bank Select for VIC				
	00 Bank 0				
	01 Bank 1				
	10 Bank 2				
	11 Bank 3				

```
CMRAM 1K =%00
CMRAM 4K =%01
CMRAM 8K =%10
CMRAM 16K=%11
;-- Set Shared RAM size to 16K
lda MMURCR
    #%11111100
and
ora CMRAM 16K
sta MMURCR
.macro SetBankConfiguration(id) {
 .if(id==0) {
 lda #%00111111 // no ROMs, RAM0
.if(id==1) {
 lda #%01111111 // no ROMs, RAM1
 .if(id==12) {
 lda #%00000110 // int.function ROM, Kernal and IO, RAMO
 .if(id==14) {
 lda #%00000001 // all ROMs, char ROM, RAM0
227
```

## hardware

```
if(id==15) {
  lda #%00000000  // all ROMs, RAM0. default setting.
}
.if(id==99) {
  lda #%00001110  // IO, kernal, RAM0. 48K RAM.
}
sta MMUCR
}
.endm

.macro SetVICBank (bank) {
  lda $DD00
  and #%11111100
  ora #3 - bank
  sta $DD00
.endm
```

Appendex	
memory	maps

memory maps

-----

GEOS Memory Map:

## All address Values in Hex

## C64 Memory Regions

00	100	200	400	6000	8000	A000	BF40	D000	E000
Zero	Stack	deskTopVars	AppRAM	Back	Disk	Fore	Kernal	I/0 or	Kernal
Page	Page			Screen	Buffers	Screen	Low	Kernal	High

## Zero Page

======		
00	CPU DDR	6510 data direction register
01	CPU DATA	Built-in 6510 I/O port, bit oriented
02-21	r0-r15	GEOS Kernal zero Page pseudoregisters
22-23	curPattern	Pointer to fill pattern data
24-25	string	Pointer to input buffer
	_	;- Current Font Settings
26	baselineOffset	Font pixels below line of print
27-28	curSetWidth	Size in bytes of bit Font stream
29	curHeight	Point size of the font
2A-2B	curIndexTable	Address of bit stream indices table
2C-2D	curDataPntr	Address of the first bit stream ; Font End
2E	currentMode	Defines the current print style
2F	dispBufferOn	Controls the screen to draw too. Fore/Back or Both.
30	mouseOn	Mouse control flag.
31-32	msePicPtr	Pointer Mouse sprite, default= 84C1
		;- Text Clipping
33	windowTop	Top margin, usually 0 (Top of screen)
34	windowBottom	Bottom margin, usually 199
35-36	leftMargin	Left margin
37-38	<del>-</del>	Right margin
39	pressFlag	Input control flags
3A-3B	mouseXPos	Mouse's X position
3C	mouseYPos	Mouse's Y position
3D-3E	returnAddress	Address for inline return
3F	graphMode	40 / 80 column mode flag on 128
40	returnAddress	Pointer to click box data table
42-43	callRouVector	Jump vector used by CallRoutine
44-45	dialogBoxVec	DoDlgBox pointer to window descriptor block.
45-6F	•	Used by Kernal
70-7F	A2-A9	Generically Named. Application zPage area
80-FA	zKerIO	Kernal I/O*
BA	curDevice	current serial device number
FB-FE	A0-A1	Generically Named. Application zPage area
FF		Used by Kernal
		-

\*Note: 80-FA is only used by the kernal during IO. See  $\mathbf{SwZp}$  for how to make safe use of this area in your applications.

## Stack Page

========	:=======:	
0100-01FF		6510 Hardware Stack Area. (GEODEBUGGER uses bottom of stack as a data area)
0200-02FF 0300-03FF	deskTopVars	;Application may freely use this block

## 128 BackRAM:

GEOS Primary Bank is Bank 1.

BackRAM is bank 0. This allows common RAM to be turned on and have parts of bank 0 then Appear into the memory space of bank 1 as shared ram is always Bank 0 ram and is always visible to the CPU when active.

Bank 0:

0000-03FF: ?

0400-1FFF: Soft Sprites

2000-9FFF: Swap area for Desk Accessories

If your application does not use Desk Accessories this may be used

as Application data area.

A000-FFFF: ??

\$0000	\$400	\$FF00	\$FF05
	BANK 0	MMU	ROM

### Bank 1 GEOS Address Space

\$0000	\$400	\$FF00	\$FF05
	BANK 1 GEOS APPLICATION SPACE	MMU	ROM

### Bank 2

\$0000	\$400	\$FF00	\$FF05
Common RAM	BANK 2 (bank 0 if 128 is not expanded)	MMU	ROM

### Bank 3

\$0000	\$400	\$FF00	\$FF05
Common RAM	BANK 3 (bank 1 if 128 is not expanded)	MMU	ROM

### Bank 14

\$0000	\$400	\$4000	\$D000	\$E000	\$FF00	\$FF05
Common RAM	RAM 0	Basic ROM	Char Rom	Kernal ROM	MMU	ROM

## Bank 15

\$0000	\$400	\$4000	\$D000	\$E000	\$FF00	\$FF05
Common RAM	RAM 0	Basic ROM	I/O	Kernal ROM	MMU	ROM

## **REU-BANKO**

## Constants

Zero Page

Constants

Zero Page

-----

## pseudoregisters:

pseudoregisters are used when calling into the GEOS kernal. Each call will have a list of pseudoregisters to setup. Registers have common uses across the GEOS API but none are exclusively used for only one thing. r12-r15 are very rarely used and make for very safe temporary zpage use. Never use other data areas for temporary storage unless you have already used all of the available options in r0-r15 that do not conflict with your current kernal interaction.

```
.zsect
         $02
   r0
          .block
                  2
                        ; Pointer
                  $02
   r0L
                        ; holds result after DoDlgBox
   r0H
                  $03
         ==
                        ; Used in RAM operations
   r1
         .block
                2
                       ; Track Number in Disk I/O
   r1T.
                  $04
                 $05
                        ; Sector Number in Disk I/O, Y Position for PutChar
   r1H
                        ; Ptr to diskname , Buffer Size during Disk I/O
         .block 2
   r2
                       ; Top Margin, Pixel Width, Str Length
    r2L
                  $06
    r2H
                 $07
                       ; Bottom Margin, Pixel Height
                        ; Left Margin, Ptr dataFileName
   r3
         .block 2
                 $08
   r3L
                       ; Top Margin
                                     Track for Allocate Block.
                 $09
                       ; Bottom Margin Sector for Allocate Block
   r3H
         .block 2
                        ; Ptr to Disk Buffers, margins on boxes
   r4
                        ; Sprite Number
                 $0A
   r3L
         =
                        ; Dest Bank on Move operations.
                  $0B
   r3H
   r5
         .block
                  2
                        ; Ptr to DirEntry
                  $0C
   r5L
         =
   r5H
         ==
                  $0D
                        ; Ptr to T/S List for block allocates
         .block
    r6T.
                  SOE
                  $0F
    r6H
         ==
                        ; Start address of Read/Write buffer
         .block 2
   r7
                 $10
    r7L
                        ; FileType to find with FindFTypes
                 $11
                        ; Number of files to get from FindFTypes
    r7H
         ==
   r8
         .block 2
                        ; Internal Kernal use during some kernal calls
                  $12
    r8L
    r8H
                  $13
         ==
                 2
   r9
          .block
                        ; Pointer to disk structures. DirEntrys/ Info Sector etc.
   r9L
                  $14
   r9H
         ==
                  $15
                2
   r10
                        ; Class Pointer.
         .block
   r10L =
                 $16
                       ; Desk Top Page number
   r10H ==
                 $17
         .block 2
   r11
                        ; x Position for PutChar
    r11L =
                  $18
                        ; row Number in DrawPoint
    r11H ==
                  $19
   r12
         .block
                        ; Not Used by Kernel as a parameter
   r12L
                 $1A
   r12H ==
                 $1B
   r13
                 2
         .block
    r13L =
                  $10
    r13H ==
                  $1D
         .block
   r14
                        ; Not Used by Kernel as a parameter
    r14L =
                  $1E
    r14H ==
                  $1F
                        ; Not Used by Kernel. Commonly used in GEOS Applications
         .block 2
   r15L =
                 $20
                        ; This is the first Goto for temp zpage use.
                 $03
   r15H ==
```

Disk

\_\_\_\_\_

## Disk Errors:

GEOS I/O Routines return errors in the X register

Standard Constant			Description
NO ERROR			No Error Occurred
NO BLOCKS	1	\$01	Not Enough Blocks On Disk
INV TRACKS	2		Invalid Track or Sector
INSUFF SPACE	3		Disk Full, Insufficient Space
	4	\$03	Directory is Full
FILE NOT FOUND	5	\$05	File Not Found
BAD BAM	6	\$06	File Not Found Bad Bam: Attempt to deallocate
_			an unallocated block. (Or the reverse)
UNOPENED_VLIR	7	\$07	VLIR file not open
_			Illegal VLIR chain number.
INV_RECORD	8	\$08	Invalid VLIR Record. Bad Track/Sector
OUT_OF_RECORDS	9	\$09	Out of Records: Too many VLIR chains
STRUCT_MISMATCH	10	\$0A	Geos Structure Mismatch
			File is not a VLIR file.
BFR_OVERFLOW	11	\$0B	Buffer Overflow: ReadRecord max read size
			exceeded.
CANCEL_ERR	12		Deliberate Cancel Error
DEV_NOT_FOUND	13	\$0D	Device Not Found
INCOMPATIBLE	14		Incompatible 40/80
HDR_NOT_THERE	32	\$20	Disk Block Read error:
			No Header Block sync character.
NO_SYNC	33		Unformatted or Missing Disk
DBLK_NOT_THERE			
	35	•	Data Block Checksum Error
WR_VER_ERR	37		Write Verify Error
WR_PR_ON	38		Write Protect On
	39	\$27	Disk Block Write: Header Checksum Error
DSK_ID_MISMAT			Disk ID Mismatch
			Drive Speed Read error
DOS_MISMATCH	115	\$73	Wrong DOS Indicator

Data

variables Address (hex)

64 128 Size Default Saved Description Chapter 3 Data †128 BackRAM Name

variables

By Name:

## Address (hex)

Name	64	128	Size	Default	Saved	Description †128 BackRAM
alarmSetFlag:	851C	851C	1	FALSE	No	TRUE if the alarm is set for geos to monitor, else FALSE
alarmTmtVector:	84AD	84AD	2	0	Yes	address of a service routine for the alarm clock time-out
						(ringing, graphic etc.) that the application can use if
						necessary.
alphaFlag:	84B4	84B4	1	0	Yes	Flag for alphanumeric string input
						0 if not getting text input
						llxx xxxx if getting text input.
						bit 0-5 - Counter before prompt flashes
						bit 6 - Flag indicating prompt is visible
						bit 7 - Flag indicating alphanumeric input is on
appMain:	849B	849B	2	0	No	Vector that allows applications to include their own main
						loop code. The code pointed to by appMain will run at the
						end of every GEOS MainLoop.
backBufPtr:	_	131B <sup>†</sup>	16	None	No	Screen pointer where the back buffer came from. Resides in
						back ram of C128.
bakclr0: [0-3]	D021	D021	1	?	No	Background colors 0-3. 1 Byte each, 4 Total Bytes.
	:	:				Hardware Registers
	D024	D024				
backXBufNum:	-	132B <sup>†</sup>	8	None	No	For each sprite, there is one byte here for how many bytes
						wide the corresponding sprite is. Used by C128 soft sprite
						routines and resides in back ram.
backYBufNum:	_	1333†	8	None	No	For each sprite, there is one byte here for how many
						scanlines high the corresponding sprite. Used by soft
						sprite routines and resides in back ram.
bootName:	C006	C006	9	GEOS	No	This is the start of the "GEOS BOOT" string.
				BOOT		
BRKVector:	84AF	84AF	2	CF85	Yes	Vector to the routine that is called when a BRK
						instruction is encountered. The default is to the
						operating system
						System Error dialog box routine.
bkvec:	0316	0316	2	?	No	BRK instruction vector when ROMs are switched in.
baselineOffset:	26	26	1	\$06	Yes	Offset from top line to baseline in character set. i.e. it
						changes as fonts change. Default \$06 - for BSW 9 Font
callRouVector	42	42	2	None	No	
CPU_DATA:	01	01	1	RAM_64K	No	Address of 6510 data register that controls the hardware
				_		memory map of the C64.
CPU_DDR:	00	0.0	1	%101111	No	address of 6510 data direction register
_						Note: Writing \$00 to this address will disable output to
						CPU DATA register. This may cause unexpected results.

## Address (hex)

Name	64	128	Size	Default		Description †128 BackRA
curDataPntr:	2C	2C	2	D2DC (BSW 9)	Yes	This is a pointer to the actual card graphic data for the current font in use.
curDirHead:	8200	8200	256	\$00	No	buffer containing header information for the disk in currently selected drive.
curDevice:	ВА	BA	1	\$08	No	current serial device number. See <b>curDrive</b> for more information
curDrive:	8489	8489	1	\$08	No	device number of the currently active disk drive. For Commodore, allowed values are 8 - 11.
curEnable:	_	1300 <sup>†</sup>	1	None	No	This is an image of the C64 mobenble register.
currentHeight:	29	29	1	\$09	Yes	card height in pixels of the current font in use.
curIndexTable:	2A	2A	2	D218	Yes	pointer to the table of sizes, in bytes, of each card in of the current font.
curmobx2:	-	1302 <sup>†</sup>	1	None	No	Image of the C64 mobx2 register. Used for C128 soft sprites. Resides in back ram
curmoby2:	-	1301 <sup>†</sup>	1	None	No	Image of C64 moby2 register. Used for C128 soft sprites. Resides in back ram.
curPattern:	22	22	2	D010	Yes	Pointer to the first byte of the graphics data for the current pattern in use.  Note: Each pattern is 1 byte wide and 8 bytes high, to give an 8 by 8 bit pattern.
curRecord:	8496	8496	1	0	No	Current record number for an open VLIR file.  Note: When a VLIR file is opened, using OpenRecordFile.  curRecord is set to 0 if there is at least 1 record in the file, or -1 if their are no records.
currentMode:	2E	2E	1	\$00	Yes	current text drawing node. Each bit is a flag for a drawing style. If set, that style is active, if clear it is inactive. The bit usage and constants for manipulating these bits are as follows.  Bit Style Constant

## Address (hex)

Name	64	128	Size	Default	Saved	Description †128 BackRAM
curSetWidth: curType:	3c 88C6	3c 88C6	2 1	\$00 Drive 8 Type	Yes	b0: Unused Clears all flags (plain text) SET_PLAINTEXT = %00000000 Any combination of flags can be set or clear. If current node is plaintext, all flags are clear.  Constants that can be used within text strings themselves that affect currentMode are: UNDERLINEON, UNDERLINEOFF, REVERSEON, REVERSEOFF, BOLDON, ITALICON, OUTLINEON, PLAINTEXT  Card width in pixels for the current font Holds the current disk type. This value is copied from driveType for quicker access to the current drive b7: Set if the disk is a RAM disk b6: Set if using disk shadowing  Only one of bit 6 or 7 may be set. Other constants used with curType are DRV_NULL = 0 No drive present at this device address DRU_1541 = 1 Drive type Commodore 1541 DRV_1571 = 2 Drive type Commodore 1571 DRV_1581 = 3 Drive type Commodore 1581
curXpos0:	_	1303 <sup>†</sup>	16	None	No	The current X positions of the C128 soft sprites. BackRAM
curYpos0:	_	1313†	8	None	No	The current Y positions of the C128 soft sprites. BackRAM
diskBlkBuf:	8000	8000	256	\$00	No	General disk block buffer. Initialized to all zeros
doRestFlag:	-	1B54 <sup>†</sup>	1	\$00	No	Flag needed because of overlapping soft sprite problems on C128. Set to TRUE if we see a sprite that needs to be redrawn and therefore all higher numbered sprites need to be redrawn as well. Resides in BackRAM.
dir2Head:	8900	8900	256	None	No	1571,1581 Second BAM block
dir3Head:	9C80	9C80	256	None	No	1581 Third BAM block
dlgBoxVector:	44	44	2	None	No	

## Appendex

## wheels

- ;Dumping Ground for Wheels info until it gets organized
- ; Wheels
- ; these are addresses to routines that are in the extended  $% \left( 1\right) =\left( 1\right) \left( 1\right)$
- ; kernal that get loaded in at \$5000 in groups.