

Problem 1 =

1) $10000000000n^2 \leq 10000000001n^2$ | $O(g(n)) = O(n^2)$

$n^3 \leq 2n^3$ | $O(g(n)) = O(n^3)$

$O(n^3)$ is asymptotically greater since when we are solving for big O seeing that n^3 is greater than $10000000000n^2$

Another way is because the exponent is bigger

2) $n^2 \log(n)$ vs $n^{((\log(n))^{10})}$

$n^2 \log(n)$ is greater because of log law, since when multiply by log instead of n^2 we get $\log(n^2) + \log(\log(n))$ vs $\log(n) + \log(\log(n)^{10})$. $2 \log(n)$ vs $\log(n)$ means that the left is greater than the right.

3) $n^{\log(n)}$ vs $2^{\sqrt{2}}$

$\log(n)(\log(n))$ vs $\sqrt{n} \log(2)$

$\log((\log(n))^2)$ vs $\log(\sqrt{n} \log(2))$

$2 \log(\log(n))$ vs $\log(\sqrt{n}) + \log(\log(2))$ hence we state that $2^{\sqrt{2}}$ is greater.

4) These are equal since when we simplify both into big O We see that both functions are under the same $O(2^n)$ hence when compared their constants which shows 2^{2n} is greater than 2^n

Problem 2 =

Best case = n is $O(1)$ since if the case is even or multiple of 2 it only runs once and then it will exit the method

Worst case = $O(\sqrt{n})$ because

- $i * i = n$

- $i^2 = n$

- $i = \sqrt{n}$

Average Case = $O(\sqrt{n})$ because for defining theta we need for $g(n)$ to be the same hence why theta takes the worst case which is the same.