Matrix approximation

COMS 4771 Fall 2019

Overview

- ightharpoonup Compression using k-means
- ► Compression using SVD
- Power iteration
- ► Latent semantic analysis
- ► Matrix completion

0 / 15

Encoder/decoder interpretation of k-means clustering

- ▶ Given data $x_1, \dots, x_n \in \mathbb{R}^d$, we learn a lossy encoding/decoding (compression) scheme
- lacksquare Scheme is represented by $oldsymbol{c}_1,\ldots,oldsymbol{c}_k\in\mathbb{R}^d$
- ► Encoder: $\boldsymbol{x} \mapsto \arg\min_{i} \|\boldsymbol{x} \boldsymbol{c}_{i}\|_{2}^{2}$
 - ▶ Need $\log_2 k$ bits to represent encoding of x!
- ▶ Decoder: $j \mapsto c_j$

Matrix view

▶ Let
$$m{A} = \begin{bmatrix} \leftarrow & m{x}_1^{\mathsf{T}} & \rightarrow \\ & \vdots & \\ \leftarrow & m{x}_n^{\mathsf{T}} & \rightarrow \end{bmatrix} \in \mathbb{R}^{n \times d}$$
 (forget the $1/\sqrt{n}$ scaling)

▶ Let $m{C} = \begin{bmatrix} \leftarrow & m{c}_1^{\mathsf{T}} & \rightarrow \\ & \vdots & \\ \leftarrow & m{c}_k^{\mathsf{T}} & \rightarrow \end{bmatrix} \in \mathbb{R}^{k \times d}$

▶ We try to approximate $m{A}$ with $m{BC}$ where $m{B} \in \{0,1\}^{n \times k}$ is

- ▶ We try to approximate A with BC, where $B \in \{0,1\}^{n \times k}$ is matrix where $B_{i,j} = 1$ iff x_i is assigned to group j.
- ▶ The k-means cost is $\| \boldsymbol{A} \boldsymbol{B} \boldsymbol{C} \|_F^2$
 - ▶ Here, $\|\cdot\|_F$ is a matrix norm called <u>Frobenius norm</u>, which treats the $n \times d$ matrix as a vector in nd-dimensional Euclidean space
- ▶ BC is a particular rank $\leq k$ approximation of the data matrix A

1 / 1:

2 / 15

3 / 1

Matrix factorization

- lacktriangle Try to approximate A with BC, where $B \in \mathbb{R}^{n \times k}$ and $C \in \mathbb{R}^{k \times d}$, to minimize $\|\boldsymbol{A} - \boldsymbol{B}\boldsymbol{C}\|_F^2$.
 - lacktriangle Think of B as the encodings of the data in A
 - \blacktriangleright At least get "dimension reduction" when k < d
 - ► Maybe not quite as space-efficient as k-means
- ▶ Theorem (Schmidt, 1907; Eckart-Young, 1936): Optimal solution is given by truncating the singular value decomposition (SVD) of A

Singular value decomposition

lacktriangle Every matrix $A \in \mathbb{R}^{n \times d}$ —say, with rank r—can be written as

$$oldsymbol{A} = \sum_{i=1}^r \sigma_i oldsymbol{u}_i oldsymbol{v}_i^{\scriptscriptstyle\mathsf{T}}$$

where

- $ightharpoonup \sigma_1 \ge \cdots \ge \sigma_r > 0$ (singular values)
- $lackbox{} u_1,\ldots,u_n\in\mathbb{R}^n$ (orthonormal left singular vectors)
- $lackbox{v}_1,\ldots,v_n\in\mathbb{R}^d$ (orthonormal right singular vectors)
- Can also write as

$$oldsymbol{A} = oldsymbol{U} oldsymbol{S} oldsymbol{V}^{\mathsf{T}}$$

where

- $lackbox{lackbox{\lorenthing}} lackbox{lackbox{U}} = [oldsymbol{u}_1|\cdots|oldsymbol{u}_r] \in \mathbb{R}^{n imes r}$, satisfies $oldsymbol{U}^{\mathsf{T}}oldsymbol{U} = oldsymbol{I}$
- \triangleright $S = \operatorname{diag}(\sigma_1, \ldots, \sigma_r) \in \mathbb{R}^{r \times r}$
- $V = [v_1] \cdots | v_r] \in \mathbb{R}^{d \times r}$, satisfies $V^{\mathsf{T}}V = I$
- \blacktriangleright This is sometimes called the *thin SVD*, since the matrices U and Vare as thin as possible.

Truncated SVD

- ▶ Let A have SVD $A = \sum_{i=1}^r \sigma_i u_i v_i^{\mathsf{T}}$ (rank of A is r)
- ▶ Truncate at rank k (for any $k \le r$): rank-k SVD

$$oldsymbol{A}_k \coloneqq \sum_{i=1}^k \sigma_i oldsymbol{u}_i oldsymbol{v}_i^{\scriptscriptstyle \mathsf{T}}$$

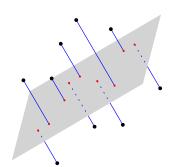
- lacktriangle Can write as $A_k := U_k S_k V_k^{\mathsf{T}}$, where
 - $m{igstyle U}_k = [m{u}_1|\cdots|m{u}_k] \in \mathbb{R}^{n imes k}$, satisfies $m{U}^{\mathsf{T}}m{U} = m{I}$ $m{igstyle S}_k = \mathrm{diag}(\sigma_1,\ldots,\sigma_k) \in \mathbb{R}^{k imes k}$

 - $lackbox{V}_k = [m{v}_1|\cdots|m{v}_k] \in \mathbb{R}^{d \times r}$, satisfies $m{V}^{\mathsf{T}}m{V} = m{I}$
- ► **Theorem** (Schmidt/Eckart-Young):

$$\|m{A} - m{A}_k\|_F^2 = \min_{m{M}: \mathrm{rank}(m{M}) = k} \|m{A} - m{M}\|_F^2 = \sum_{i=k+1}^r \sigma_i^2$$

Encoder/decoder interpretation I

- lacktriangle Encoder: $oldsymbol{x}\mapsto oldsymbol{V}_k^{\scriptscriptstyle\mathsf{T}}oldsymbol{x}\in\mathbb{R}^k$
 - Encoding rows of A: $AV_k = U_k S_k$
- $lackbox{Decoder: } oldsymbol{z} \mapsto oldsymbol{V}_k oldsymbol{z} \in \mathbb{R}^d$
 - lacktriangle Decoding rows of $U_k S_k$: $U_k S_k V_k^{\mathsf{T}} = A_k$
- ► Same as k-dimensional PCA mapping!
 - $lackbox{ } A^{\mathsf{T}}A = VS^2V^{\mathsf{T}}$, so eigenvectors of $A^{\mathsf{T}}A$ are right singular vectors of A, non-zero eigenvalues are squares of the singular values
- ▶ PCA/SVD finds k-dimensional subspace of smallest sum of squared distances to data points.



5 / 15

6 / 15

Encoder/decoder interpretation II

▶ Example: OCR data, compare original image to decoding of k-dimensional PCA encoding $(k \in \{1, 10, 50, 200\})$











Figure 2: OCR digits, original and decoding of k-dimensional PCA encoding

Computation: power iteration

- ightharpoonup How to find v_1, \ldots, v_k ?
- ightharpoonup Easier problem: how to find v_1 ?
- ► Power iteration:
 - lacktriangle Start with any unit vector $oldsymbol{u} \in \mathbb{R}^d$
 - ▶ Repeat: Multiply u by A^TA to get new vector u'; renormalize back to unit length to replace u

Why does power iteration work?

- ► Ignore renormalization for now
- ► Then after t iterations, resulting vector is

$$\underbrace{(\boldsymbol{A}^{\scriptscriptstyle\mathsf{T}}\boldsymbol{A})\cdots(\boldsymbol{A}^{\scriptscriptstyle\mathsf{T}}\boldsymbol{A})}_{t \; \mathsf{times}}\boldsymbol{u} = \underbrace{(\boldsymbol{V}\boldsymbol{\Lambda}\boldsymbol{V}^{\scriptscriptstyle\mathsf{T}})\cdots(\boldsymbol{V}\boldsymbol{\Lambda}\boldsymbol{V}^{\scriptscriptstyle\mathsf{T}})}_{t \; \mathsf{times}}\boldsymbol{u} = \boldsymbol{V}\boldsymbol{\Lambda}^t\boldsymbol{V}^{\scriptscriptstyle\mathsf{T}}\boldsymbol{u}$$

▶ Look at Λ^t :

$$oldsymbol{\Lambda}^t = \lambda_1^t egin{bmatrix} 1 & & & & & & \\ & \left(rac{\lambda_2}{\lambda_1}
ight)^t & & & & \\ & & \ddots & & & \\ & & & \left(rac{\lambda_r}{\lambda_1}
ight)^t \end{bmatrix} pprox \lambda_1^t egin{bmatrix} 1 & & & & \\ & 0 & & & \\ & & \ddots & & \\ & & & 0 \end{bmatrix}$$

whenever $\lambda_2 < \lambda_1$ and t is large.

► Therefore

$$oldsymbol{V}oldsymbol{\Lambda}^toldsymbol{V}^{\scriptscriptstyle\mathsf{T}}pprox\lambda_1^toldsymbol{v}_1oldsymbol{v}_1^{\scriptscriptstyle\mathsf{T}}$$

and hence

$$m{V} m{\Lambda}^t m{V}^{\scriptscriptstyle\mathsf{T}} u pprox \lambda_1^t m{v}_1 m{v}_1^{\scriptscriptstyle\mathsf{T}} m{u} = (\lambda_1^t m{v}_1^{\scriptscriptstyle\mathsf{T}} m{u}) m{v}_1$$
 which renormalizes back to $m{v}_1$ (assuming $m{v}_1^{\scriptscriptstyle\mathsf{T}} m{u}
eq 0$).

Beyond the top eigenvector

- ightharpoonup To get $oldsymbol{v}_2,\ldots,oldsymbol{v}_k$:
 - ► Can <u>deflate</u> $A^{\mathsf{T}}A$: replace $A^{\mathsf{T}}A$ by $A^{\mathsf{T}}A \lambda_1 v_1 v_1^{\mathsf{T}}$; now v_2 is top eigenvector.
 - \triangleright This only works well for small k, since errors accumulate
- ▶ Alternative: to get some $U \in \mathbb{R}^{d \times k}$ with same (similar) range as V_k , use *subspace iteration*.

Application: Topic modeling

- ▶ Start with *n* documents, represent using "bag-of-words" count vectors
- lacktriangle Arrange in matrix $m{A} \in \mathbb{R}^{n \times d}$, where d is vocabulary size

	aardvark	abacus	abalone	
doc 1	3	0	0	
doc 2	7	0	4	
doc 3	2	4	0	
:	:	:	:	
		•	•	

ightharpoonup Rank k SVD provides an approximate factorization

$$A \approx BC$$

where $oldsymbol{B} \in \mathbb{R}^{n imes k}$ and $oldsymbol{C} \in \mathbb{R}^{k imes d}$

- ► This use of SVD is called *Latent Semantic Analysis (LSA)*
- ightharpoonup Interpret rows of C as "topics"
- ▶ $B_{i,t}$ is "weight" of document i on topic t
- ▶ If rows of C were probability distributions, could interpret as $C_{t,w}$ as probability that word w appears in topic t

► Start with ratings of movies given by users

Application: Matrix completion

- ▶ Arrange in a matrix $A \in \mathbb{R}^{n \times d}$, where $A_{i,j}$ is rating given by user i for movie j.
 - Netflix: n = 480000, d = 18000; on avg, each user rates 200 movies
 - ightharpoonup Most entries of A are unknown
- ▶ Idea: Approximate A with low-rank matrix, i.e., find

$$oldsymbol{U} = egin{bmatrix} \leftarrow & oldsymbol{u}_1^{\intercal} &
ightarrow \ dots & dots \ \leftarrow & oldsymbol{u}_n^{\intercal} &
ightarrow \end{bmatrix}, \qquad oldsymbol{M} = egin{bmatrix} \uparrow & & & \uparrow \ oldsymbol{m}_1 & \cdots & oldsymbol{m}_d \ \downarrow & & & \downarrow \end{bmatrix}$$

with goal of minimizing $\|oldsymbol{A} - oldsymbol{U} oldsymbol{M}\|_F^2$

- ▶ Note: If all entries of *A* were observed, we could do this with truncated SVD.
- ► Need to find a low-rank approximation without all of *A*: (low-rank) matrix completion

12 / 1

SGD for matrix completion

► Instead, try to minimize

$$\sum_{(i,j)\in\Omega}(A_{i,j}-\boldsymbol{u}_i^{\scriptscriptstyle\mathsf{T}}\boldsymbol{m}_j)^2$$

where Ω indexes the collection of observed entries of A.

- Not convex in (U, M); still use (stochastic) gradient descent anyway
- ▶ For epoch $p = 1, 2, \ldots$
 - ▶ For each $(i, j) \in \Omega$ in random order:

$$egin{aligned} oldsymbol{u}_i &\coloneqq oldsymbol{u}_i - 2\eta(oldsymbol{u}_i^{\mathsf{T}}oldsymbol{m}_j - A_{i,j})oldsymbol{m}_j \ oldsymbol{m}_j &\coloneqq oldsymbol{m}_j - 2\eta(oldsymbol{u}_i^{\mathsf{T}}oldsymbol{m}_j - A_{i,j})oldsymbol{u}_i \end{aligned}$$

(Could also switch or randomize order of updates)

Can also use regularization and other objectives / optimization algorithms

Feature representations from matrix completion

- ▶ Small data set (n=6040 users, d=3952 movies, $|\Omega|=800000$ ratings)
- lacktriangle Fit U and M by optimizing (regularized) least squares objective
- lacktriangle Are $m_1,\ldots,m_d\in\mathbb{R}^k$ useful feature vectors for movies?
- ► Some nearest-neighbor pairs $(m_j, NN(m_j))$:
 - ► Toy Story (1995), Toy Story 2 (1999)
 - ► Sense and Sensibility (1995), Emma (1996)
 - ► Heat (1995), Carlito's Way (1993)
 - ► The Crow (1994), Blade (1998)
 - ► Forrest Gump (1994), Dances with Wolves (1990)
 - Mrs. Doubtfire (1993), The Bodyguard (1992)

10 / 10