

# **Fundamentals of Machine Learning**

Yassine Laguel

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**Proposition : Upperboudning the max of bounded r.v.**

Let  $X_1, \dots, X_n$  be  $n$  centered random variables ( $\mathbb{E}[X_i] = 0$  for all  $i$ ), such that  $X_i \in [a, b]$  almost surely for all  $i \in \{1, \dots, n\}$ .

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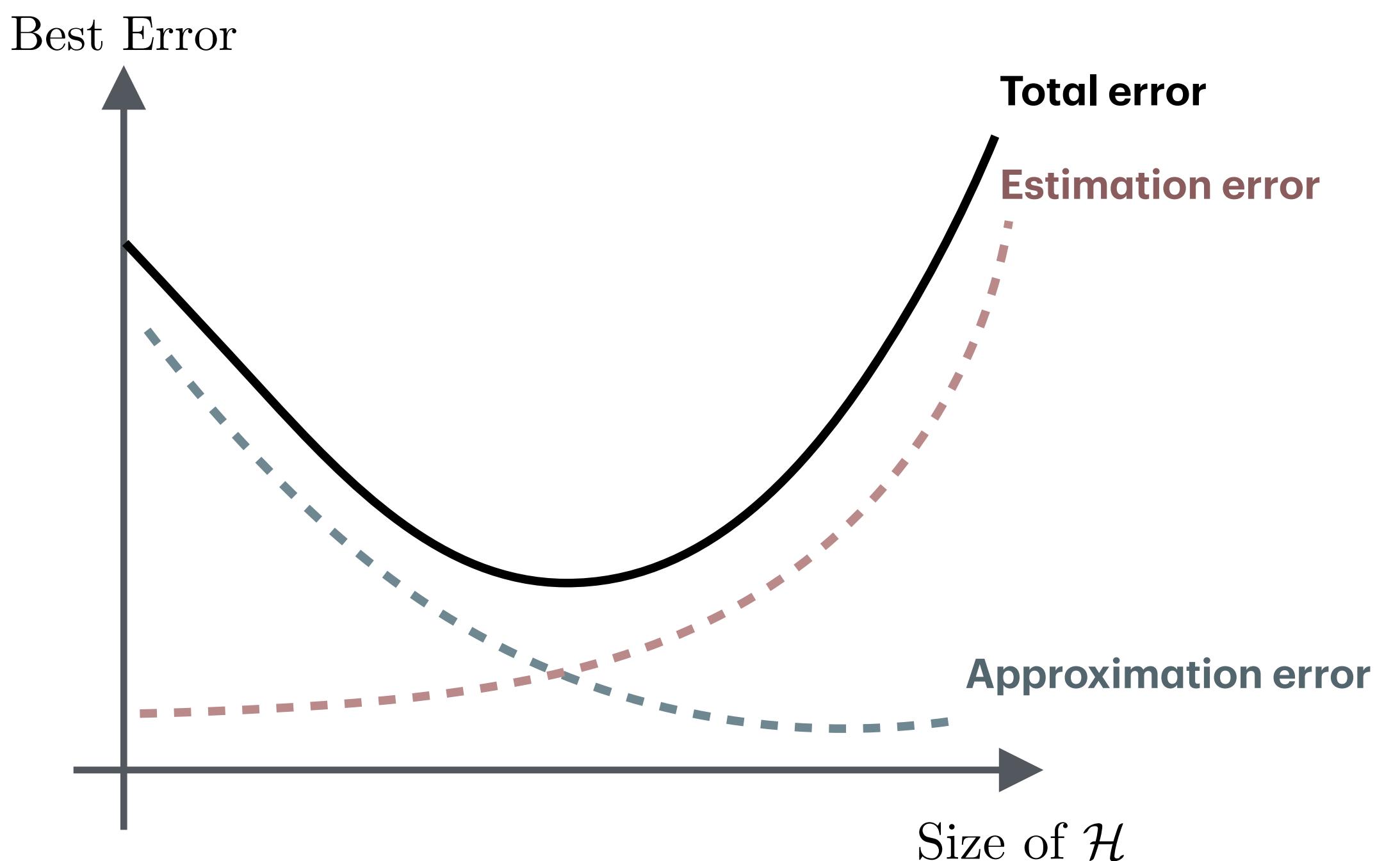
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- When  $|\mathcal{H}|$  is infinite, this bound is not exploitable.

1. Supervised Learning Setting
2. Estimation vs Approximation
3. Maximal inequalities
4. Rademacher Complexity



## **Rademacher Complexity of a set**

- Rademacher variables and complexity of a set

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## ■ Rademacher variables and complexity of a set

### Definition : Rademacher random variables

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such that

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### Proposition : Minkowski sum

For any sets  $T, T' \subset \mathbb{R}^d$ , we have

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## Proposition : Massart's Lemma - Complexity for finite Sets

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## Proposition : Convex hull

Let  $T \subset \mathbb{R}^n, v \in \mathbb{R}^n, c \in \mathbb{R}$ , and define  $cT + v = \{ct + v, t \in T\}$ .

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## Remark :

- This lemma will later be used to establish statistical bounds when the loss  $\ell$  satisfies a Lipschitz condition.

# Rademacher complexity of a hypothesis class

## ■ Formal definition & Symmetrization Lemma

### Definition : Rademacher complexity of a class of functions

Let  $(\Omega_1, \dots, \Omega_n)$  be  $n$  i.i.d. Rademacher variables.

We define the empirical Rademacher complexity of the hypothesis class  $\mathcal{H}$ , denoted  $\widetilde{\text{Rad}}(\mathcal{H})$ , as

$$\widetilde{\text{Rad}}(\mathcal{H}) = \mathbb{E} \left[ \sup_{f \in \mathcal{H}} \frac{1}{n} \sum_{i=1}^N \Omega_i \ell(f(X_i), Y_i) \mid (X_1, Y_1), \dots, (X_n, Y_n) \right],$$

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Under the Empirical Risk Minimization Framework, we have :

$$\mathbb{E} \left[ \sup_{f \in \mathcal{H}} \{r(f) - R(f)\} \right] \leq 2 \text{Rad}(\mathcal{H}).$$

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## ■ Applications in the regression setting

### Proposition : Complexity for Lipschitz Losses

Under the Empirical Risk Minimization Framework, if  $\hat{y} \mapsto \ell(\hat{y}, y)$  is  $G$ -lipschitz, then

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### Remarks :

- Hence, assuming a Lipschitz error function allows us to focus on the complexity spanned by the set of estimators rather than the composition

$$\ell(f(X), Y).$$