

Fundamentals of Machine Learning

Yassine Laguel

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Maximal Inequality for bounded losses

Proposition : Upperbounding the max of bounded r.v.

Let X_1, \dots, X_n be n centered random variables ($\mathbb{E}[X_i] = 0$ for all i), such that $X_i \in [a, b]$ almost surely for all $i \in \{1, \dots, n\}$.

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This will be a common characteristic of our upcoming results.

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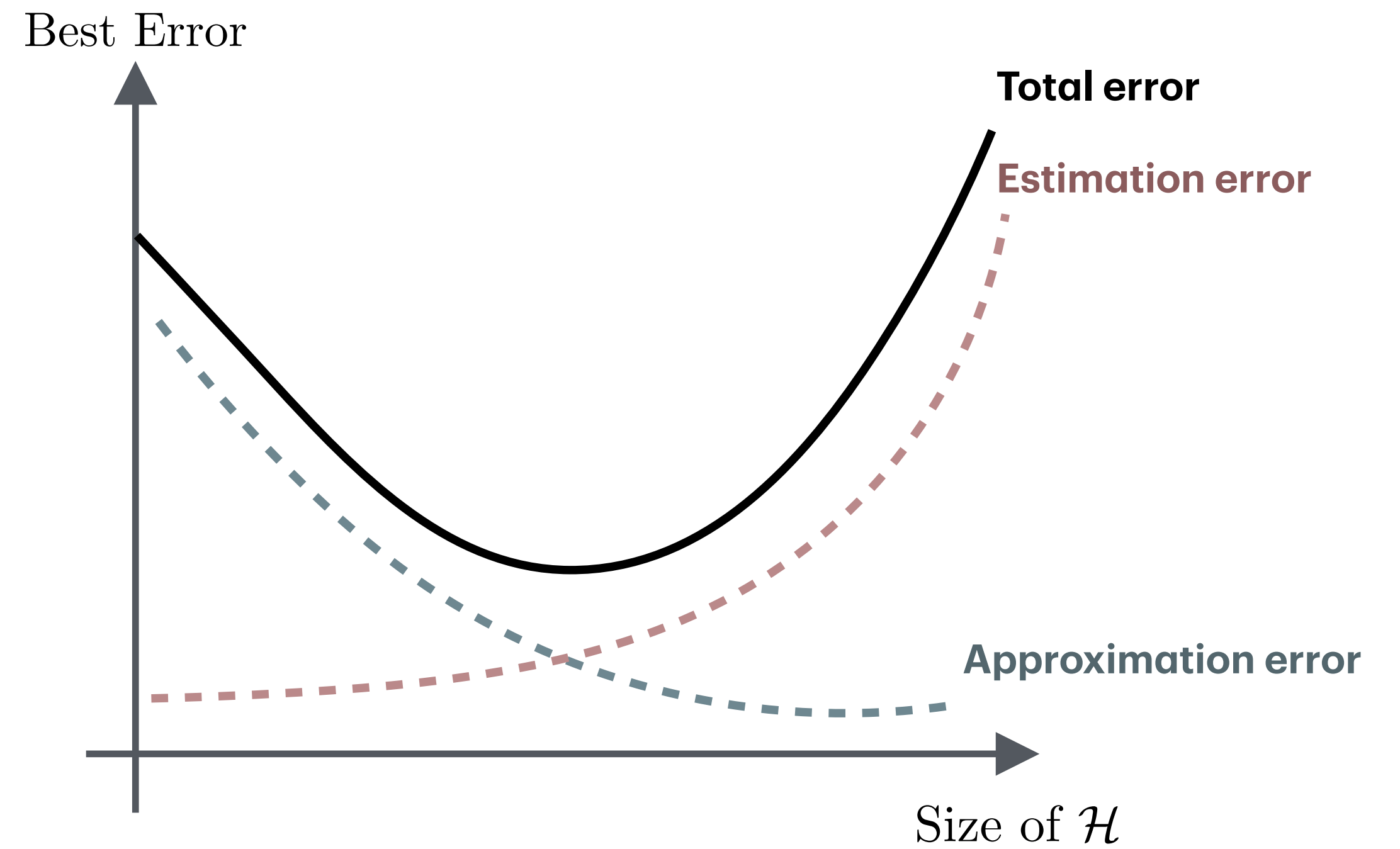
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- When $|\mathcal{H}|$ is infinite, this bound is not exploitable.

1. Supervised Learning Setting
2. Estimation vs Approximation
3. Maximal inequalities
4. Rademacher Complexity



Rademacher Complexity of a set

- Rademacher variables and complexity of a set

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Definition : Rademacher random variables

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The Rademacher complexity of a set $T \subset \mathbb{R}^n$, denoted $\text{Rad}(T)$, is defined as:

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Proposition : Homotetie and translation

Let $T \subset \mathbb{R}^n, v \in \mathbb{R}^n, c \in \mathbb{R}$, and define $cT + v = \{ct + v, t \in T\}$.

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For any sets $T, T' \subset \mathbb{R}^d$, we have

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Proposition : Talagrand's lemma

Let $f_1, \dots, f_n : \mathbb{R} \rightarrow \mathbb{R}$ be n G -lipchitz functions.

Let $T \subset \mathbb{R}^n$ and define $(f_1, \dots, f_n) \circ T$ as

$$(f_1, \dots, f_n) \circ T = \{(f_1(t_1), \dots, f_n(t_n)), t \in T\}.$$

Then,

$$\text{Rad}((f_1, \dots, f_n) \circ T) \leq G \text{Rad}(T).$$

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$$\text{Rad}(\text{conv}(T)) = \text{Rad}(T).$$

Proof : on the black board.

Proposition : Massart's Lemma - Complexity for finite Sets

Let $T \subset \mathbb{R}^n$ be finite.

Then,

$$\text{Rad}(T) \leq \max_{t \in T} \|t\|_2 \frac{\sqrt{2 \log(|T|)}}{n}.$$

Proof : on the black board.

Remark :

- The two previous propositions allow to upper bound the Rademacher complexity of the containing polyhedra.



Proposition : Talagrand's lemma

Let $f_1, \dots, f_n : \mathbb{R} \rightarrow \mathbb{R}$ be n G -lipchitz functions.

Let $T \subset \mathbb{R}^n$ and define $(f_1, \dots, f_n) \circ T$ as

$$(f_1, \dots, f_n) \circ T = \{(f_1(t_1), \dots, f_n(t_n)), t \in T\}.$$

Then,

$$\text{Rad}((f, \dots, f_n) \circ T) \leq G \text{Rad}(T).$$

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More properties

Proposition : Convex hull

Let $T \subset \mathbb{R}^n, v \in \mathbb{R}^n, c \in \mathbb{R}$, and define $cT + v = \{ct + v, t \in T\}$.

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Remark :

- This lemma will later be used to establish statistical bounds when the loss ℓ satisfies a Lipschitz condition.

Rademacher complexity of a hypothesis class

■ Formal definition & Symmetrization Lemma

Definition : Rademacher complexity of a class of functions

Let $(\Omega_1, \dots, \Omega_n)$ be n i.i.d. Rademacher variables.

We define the empirical Rademacher complexity of the hypothesis class \mathcal{H} , denoted $\widetilde{\text{Rad}}(\mathcal{H})$, as

$$\widetilde{\text{Rad}}(\mathcal{H}) = \mathbb{E} \left[\sup_{f \in \mathcal{H}} \frac{1}{n} \sum_{i=1}^N \Omega_i \ell(f(X_i), Y_i) \mid (X_1, Y_1), \dots, (X_n, Y_n) \right],$$

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Proposition : Symmetrization lemma

Under the Empirical Risk Minimization Framework, we have :

$$\mathbb{E} \left[\sup_{f \in \mathcal{H}} \{r(f) - R(f)\} \right] \leq 2 \text{Rad}(\mathcal{H}).$$

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■ Applications in the regression setting

Proposition : Complexity for Lipschitz Losses

Under the Empirical Risk Minimization Framework, if $\hat{y} \mapsto \ell(\hat{y}, y)$ is G -lipschitz, then

$$\widetilde{\text{Rad}}(\mathcal{H}) \leq G \mathbb{E} \left[\sup_{f \in \mathcal{H}} \frac{1}{n} \sum_{i=1}^n \Omega_i f(X_i) \mid X_1, \dots, X_n \right].$$

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Remarks :

- Hence, assuming assuming a Lipschitz error function allows us to focus on the complexity spanned by the set of estimators rather than the composition $\ell(f(X), Y)$.