

Definition 1. For a parameter n , Let $L \subset \mathbb{N}^+$ be the set of non-periodic words over the alphabet $\Sigma = \mathbb{N}$ that are bigger in Arabic (right-to-left) lexicographical order than all of their rotations. Let L_n be the set of all the words in L whose length divides n .

Definition 2. For a word w whose length is smaller or equal than n , let $f(w)$ be the transformation defined by successive applications of the following steps to w :

- f_1 : Increase the first letter of the word by one.
- f_2 : Pad with zeros on the left to get a word of length n .
- f_3 : Apply the substitution rules $u(vu)^+ \mapsto vu$ and then $w^+ \mapsto w$, with the longest possible u and the shortest possible w .

Definition 3. For a word w whose length is smaller or equal than n , let $b(w)$ be the transformation defined by successive applications of the following steps to w :

- b_1 : Expand w to uw^m where $m = \lfloor n/|w| \rfloor$ and u is the suffix of length $n - m|w|$ of w .
- b_2 : Remove leading zeros.
- b_3 : Decrease the first letter by one.

Observation 1. For any $w \in L_n$, $f(b(w)) = b(f(w)) = w$.

Observation 2. For any $w = 0^l w_{m-l} \cdots w_{m-1} w_m \in L_n$ where $w_{m-l} \neq 0$ and for each $0 \leq i < l$ we have $b(w)_{|b(w)|-i} = w_{m-i}$. If, in addition, $m \neq n$ we also have $b(w)_{|b(w)|-l} = w_{m-l}$.

Proposition 3. If we start with $w(0) = 0$ and generate a sequence of words by $w(i+1) = f(w(i))$, we get an enumeration of all the words in L whose length is smaller or equal to n .

Proof. This is a version of Duval's algorithm with a reversed order of the alphabet and a reversed order of letters in a word. \square

Definition 4. Let $f^*(w)$ be the first word in $f(w), f(f(w)), \dots$ whose length divides n and, similarly, let $b^*(w)$ be the first word in $b(w), b(b(w)), \dots$ whose length divides n .

Observation 4. Observation 2 is true also when b is replaced with b^* .

Proof. We only apply b again if the length is not n , so after the first application of b , that may change the first letter in the tail $w_{m-1} \cdots w_m$ when the length of w is n , the tail remains as it is. \square

Definition 5. Let $w(0), w(1), \dots$ be the sequence generated by starting with $w(0) = 0$ and then continuing ad infinitum by $w(i+1) = f^*(w(i))$ and let w^∞ be the infinite word obtained by concatenating all the words in this sequence.

Definition 6. For a word $w = w_1 \cdots w_{n-1} w_n$ let $R(w) = w_n w_1 \cdots w_{n-1}$ be the rotation of w to the right. Then, the nested invocation $R^m(w)$ is the m letter rotation to the right and its inverse $R^{-m}(w)$ is the m letter rotation to the left.

Proposition 5. For a given $w \in L_n$, let l be the number of trailing zeros (from the left) in w and let $\bar{w} = b_1(w)$. Then For all $0 \leq i \leq |w| - l - 1$ the word $R^i(\bar{w})$ comes $i + n - |w|$ letters before w as a sub-word of w^∞ .

Proof. If the length of w is n then $\bar{w} = w$ and, by Observation 4, the word that precedes w ends with the last $l - 1$ letters of w . In particular, the word $R^i(\bar{w})$ comes i letters before w as claimed. \square

Proposition 6. For a given $w \in L_n$, let l be the number of trailing zeros (from the left) in w and let $\bar{w} = b_1(w)$. Then For all $n - k \leq i \leq n - 1$ the word $R^i(\bar{w})$ comes $i - (n - |f_3(u)| \pmod{n})$ letters before the first $u \in \langle 0^{m-1}(\bar{w}_m + 1)\bar{w}_{m+1} \cdots \bar{w}_n \rangle_{m=i}^n$ such that $u \in L_n$ and $\max\{u_1, \dots, u_{n-1}\} \leq \max(w)$.

Proposition 7. The word w^∞ contains all the words in \mathbb{N}^n as subwords.

Proof. Any word of length n is a rotation of the expansion of a word in L_n . \square

Proposition 8. For any k the prefix $w_1^\infty \cdots w_k^\infty$ is an n -order de Bruijn sequences. Moreover, it is the reversed of the n -order prefer-max sequence on the alphabet $\langle 0, \dots, k - 1 \rangle$ (in this order).

Proof. Counting argument + arguing that if $|w| = n - 1$ and $\sigma_1 < \sigma_2$ then $w\sigma_1$ comes before $w\sigma_2$ as subwords of w^∞ . \square

Proposition 9. For $w \in \mathbb{N}^n$, let i be the minimal index such that $R^{-i}(w) \in L$ and let $\bar{w} = R^{-i}(w)$. Let $\bar{w}^+ = \bar{w}_{1..i}(\bar{w}_{i+1} + 1)\bar{w}_{(i+2)..n}$, i.e., the word obtained by increasing the $(i + 1)$ th letter of \bar{w} by one. Then, the function

$$\text{next}(w) = \begin{cases} f^*(f_3(w))_1 & \text{if } w \in L; \\ w_1 + 1 & \text{if } \bar{w}_{1..i} = 0^i \wedge \bar{w}^+ \in L \wedge \max(\bar{w}_{1..(n-1)}^+) \leq \max(w); \\ 0 & \text{if } \bar{w}_{1..i} = 0^i \wedge (\bar{w}^+ \notin L \vee \max(\bar{w}_{1..(n-1)}^+) > \max(w)); \\ w_1 & \text{otherwise.} \end{cases}$$

represents the mapping of a word w to the letter that follows the (one and only) occurrence of w as a subword of w^∞ .

Definition 7. Let $w(0) = 0, w(1) = f^*(w(0)), \dots, w(i) = f^*(w(i-1)), \dots$ be our enumeration of all the words in L_n . Let $w^{(i)} = w(0) \cdots w(i)$ be the concatenation of the first i words in this enumeration and let $u(j) = w_{j-n+1}^{(i)} \cdots w_j^{(i)}$ be the “window” of length n before the j th letter in $w^{(i)}$.

Proposition 10. *For a word $w = w(i)$ be the i th word in the above enumeration. Let l be the number of leading zeros in w and let $\hat{w} = R^{-l}(b_1(b^*(w)))$. Then, inserting the cycle $\langle R^{-s}(b_1(w)) \rangle_{s=0}^{n-1}$ after the word $\hat{w}_{1..(n-l)}0^l$ in $\langle u(j) \rangle_{j=0}^{i-1}$ yields the sequence $\langle u(j) \rangle_{j=0}^i$.*