

# ADAPTIVE GREEDY REJECTION SAMPLING

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# 1. IN COLLABORATION WITH



Lucas Theis

## **2. BACKGROUND**

## 2.1. CHANNEL SIMULATION

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How many bits does Alice need to send to Bob?

When common randomness  $S$  available (SFRL):

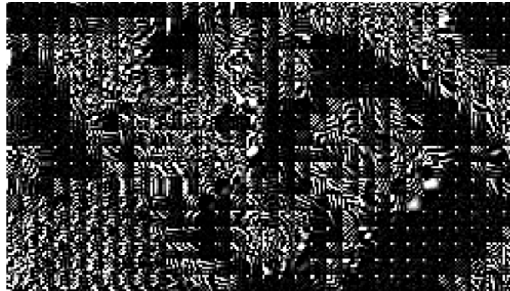
$$H[X | S] \leq I[X; Y] + \log(I[X; Y] + 1) + 4.$$

## **2.2. LOSSY COMPRESSION**

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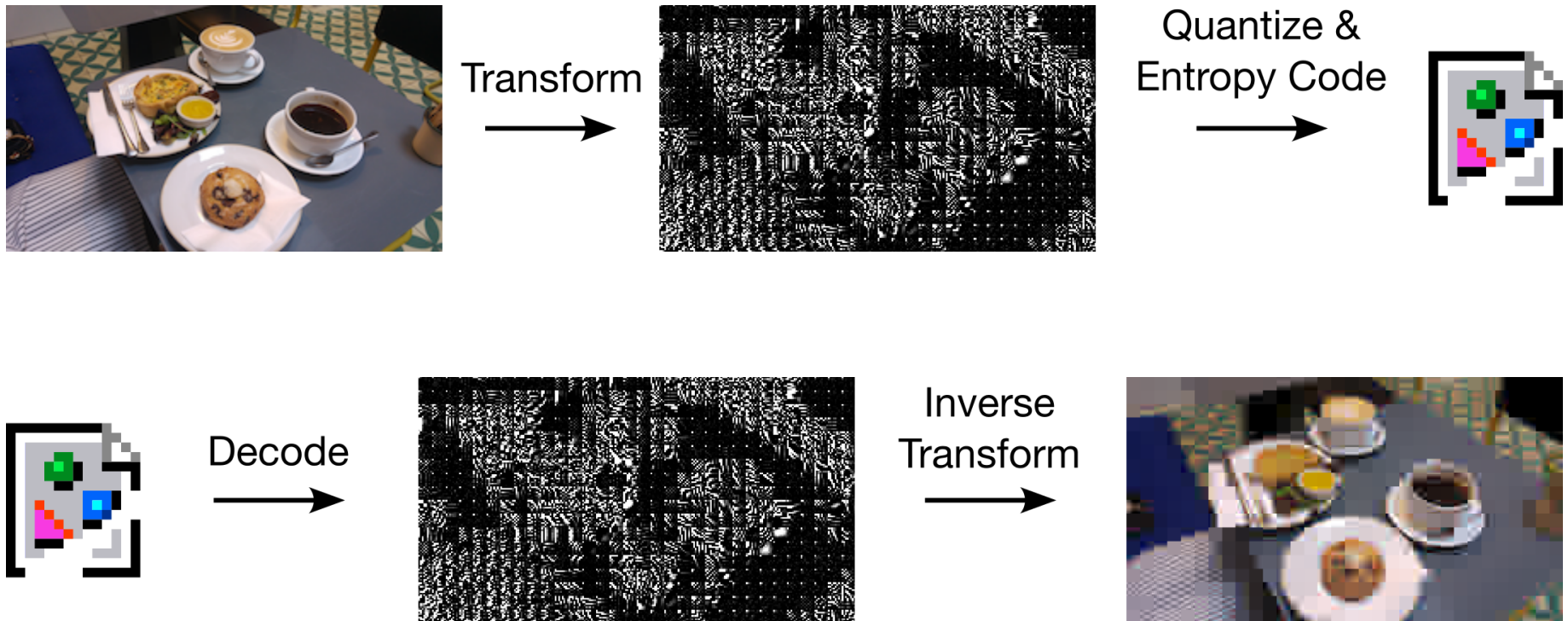
Transform  
→



Quantize &  
Entropy Code  
→



## 2.2. LOSSY COMPRESSION



# **3. GREEDY REJECTION SAMPLING**

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We extend it to Borel probability measures over Polish spaces.



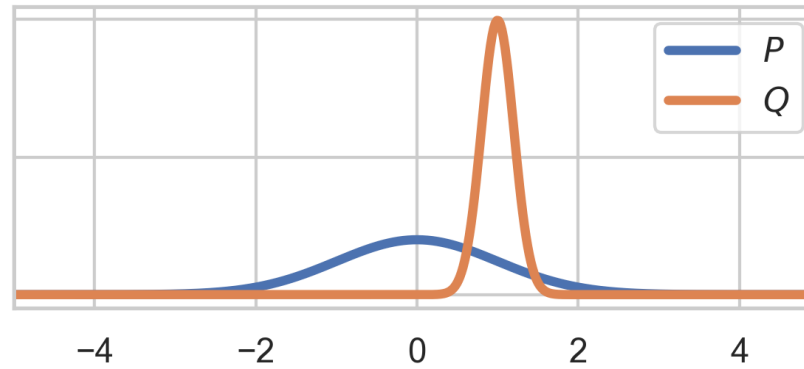
## 3.2. GRS EXAMPLE

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Target  $Q = \mathcal{N}(\mu, \sigma^2)$ , proposal  $P = \mathcal{N}(0, 1)$

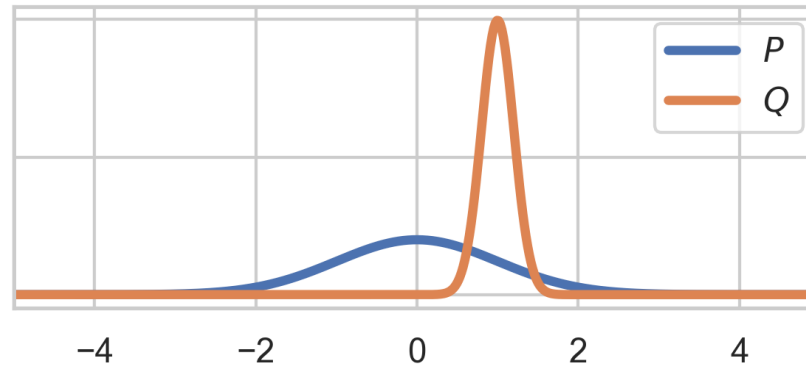
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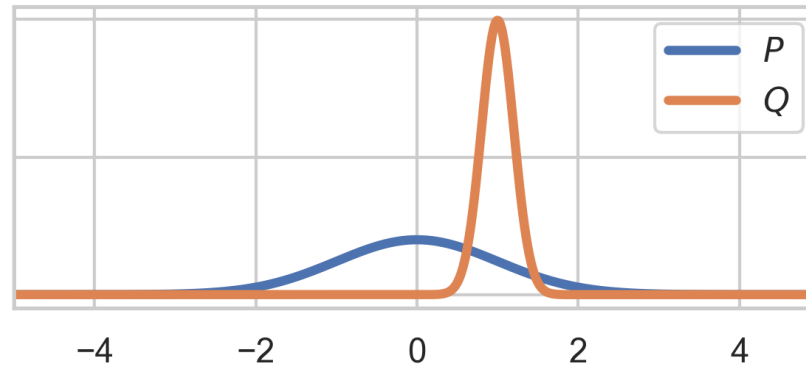
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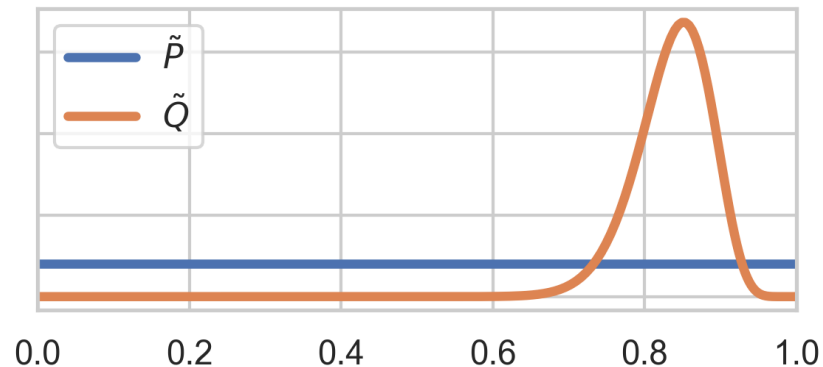
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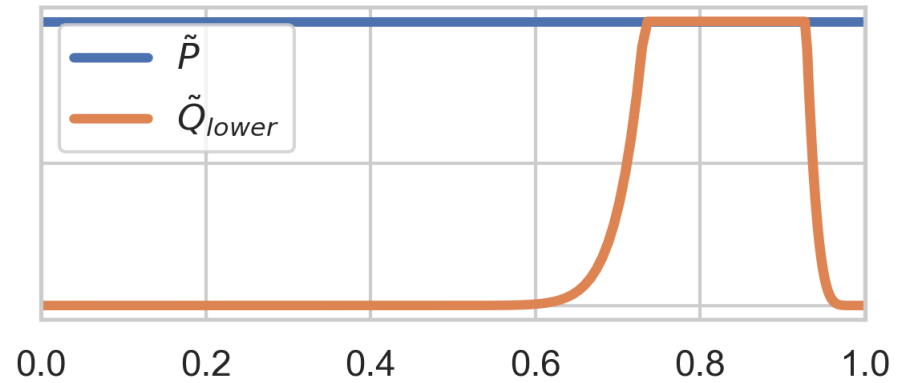


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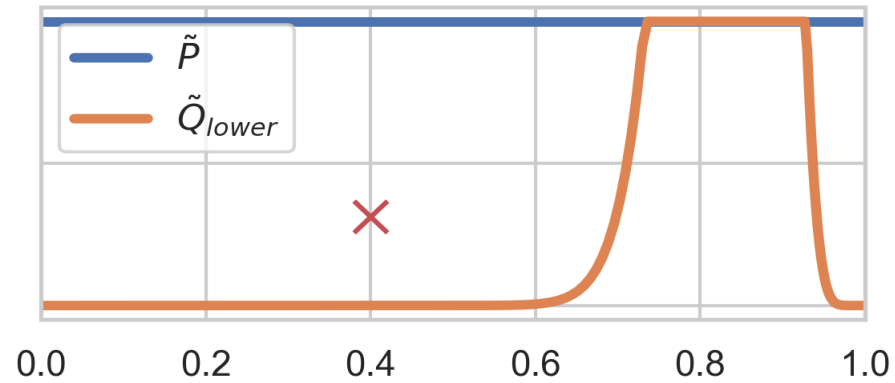


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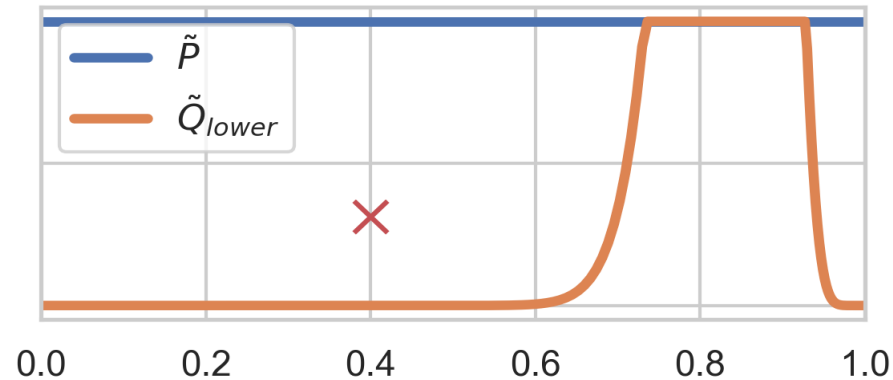


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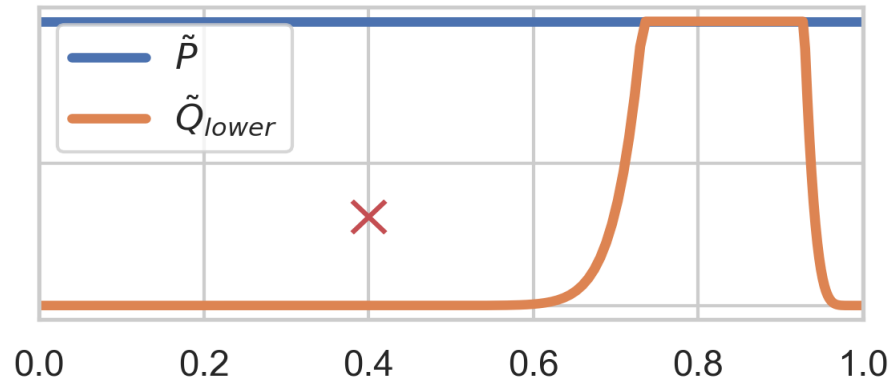


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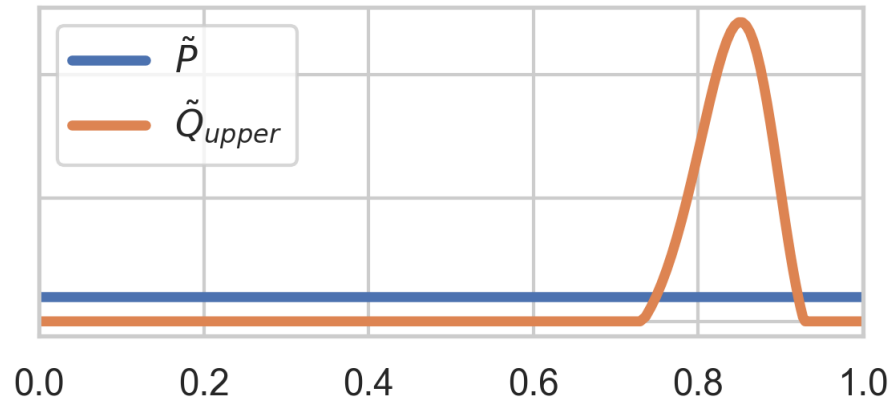


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Transmit  $K$  using a Zeta coding distribution:

$$\zeta(k \mid \alpha) \propto k^{-\alpha}$$

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When  $Q \leftarrow P_{X|Y}$  and  $P \leftarrow P_X$ , average over  $Y$ :

$$H[K] \leq I[X; Y] + \log(I[X; Y] + 1) + 4$$

## **3.6. GRS ANALYSIS CONT'D**

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$$\mathbb{E}[K + 1] = 2^{D_\infty[Q \parallel P]}$$

where

$$D_\infty[Q \parallel P] = \log \left( \text{ess sup}_{x \in \Omega} \left\{ \frac{dQ}{dP}(x) \right\} \right)$$

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Now, we find:

$$\mathbb{E}[K + 1] = \infty$$



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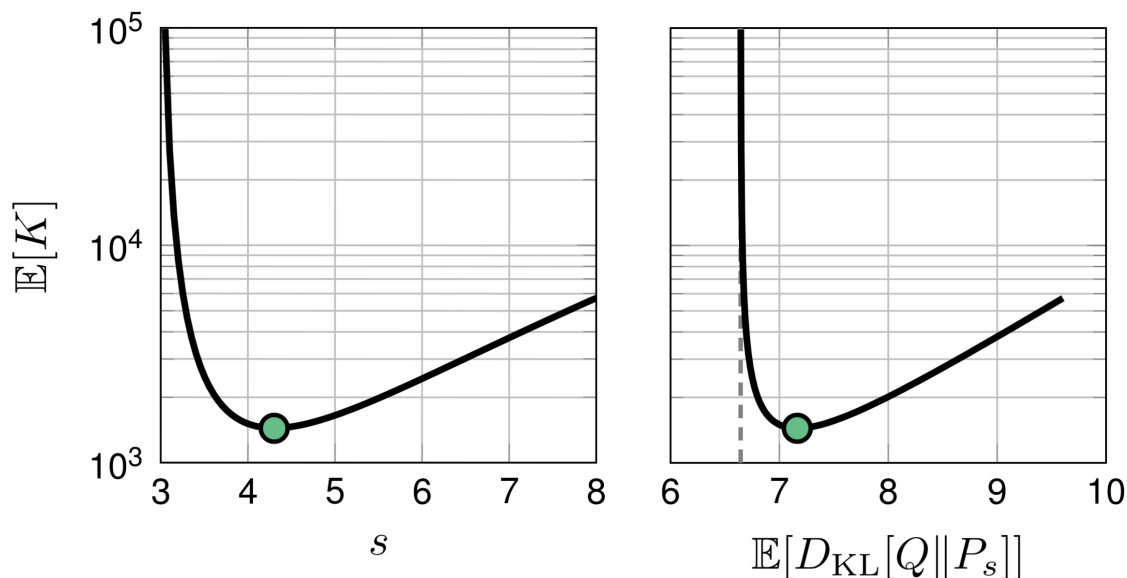
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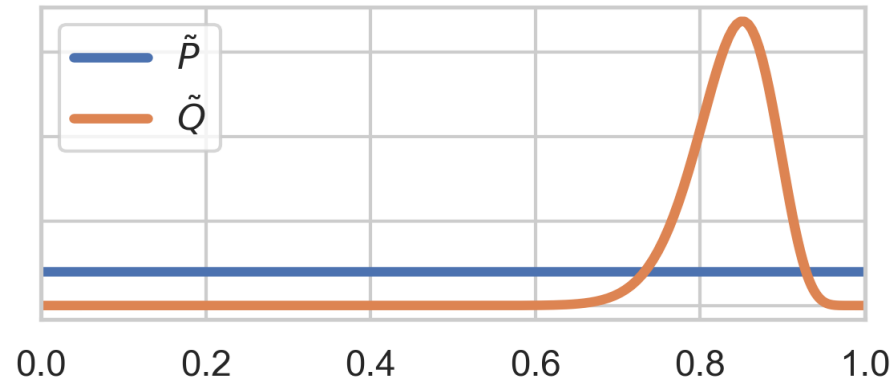
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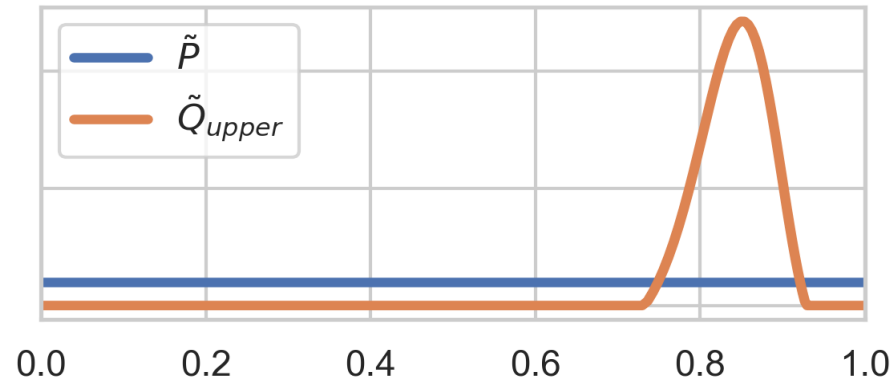
# **4. ADAPTIVE GREEDY REJECTION SAMPLING**

# 4.1. GRS EXAMPLE REVISITED

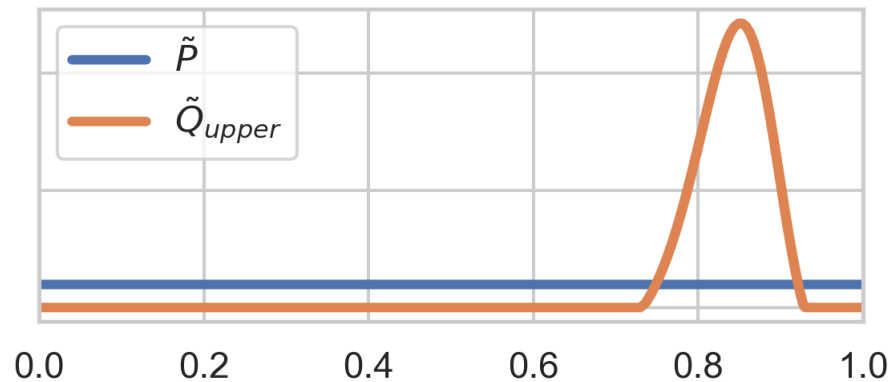
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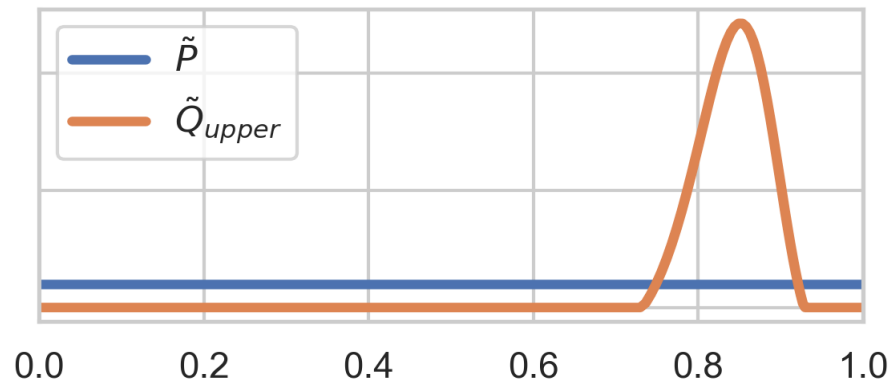
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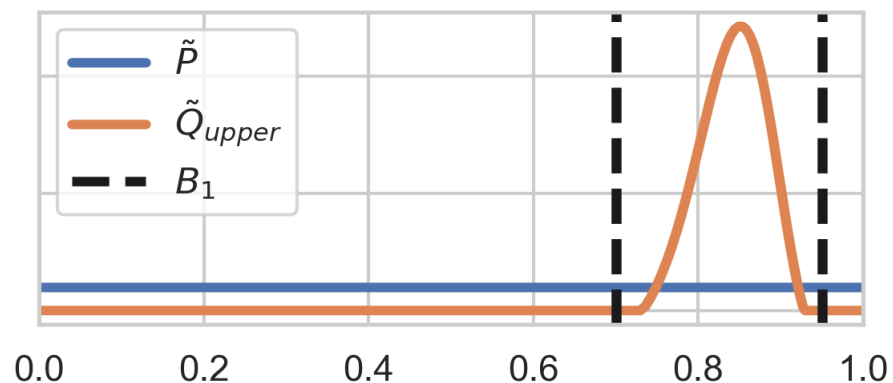
Most of sample space useless, *adapt proposal*:



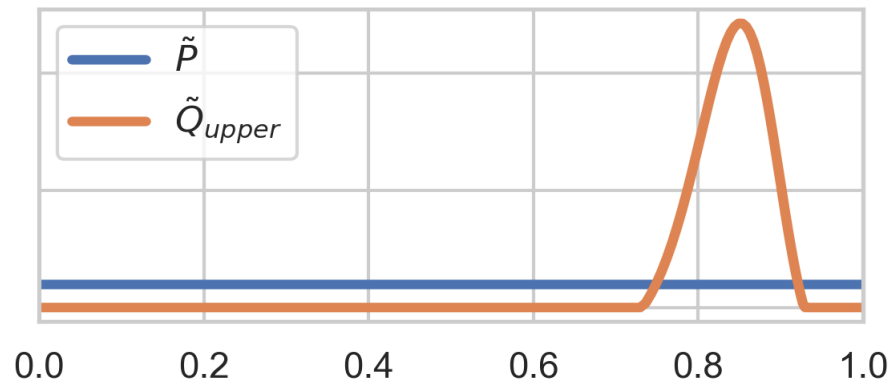
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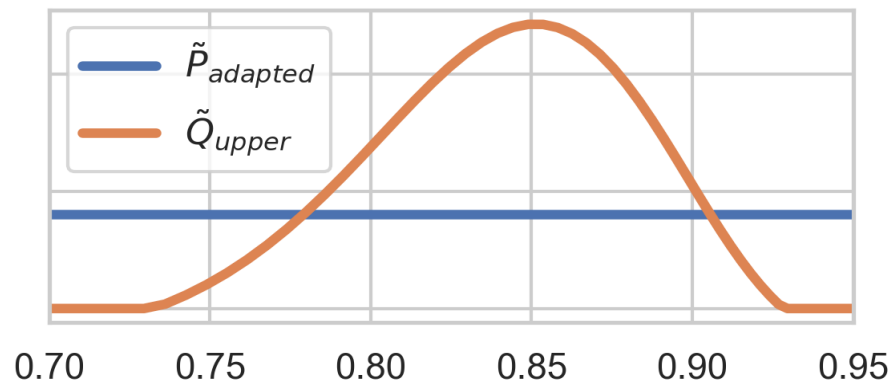
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**Theorem:**

$$H[K] \leq C + \log(C + 1) + 3.63,$$

where  $C = I[X; Y] + \mathbb{E}[\log P(B_{m_K})]$ .

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$$\lfloor c + U \rfloor - U \stackrel{d}{=} c + U'$$

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$$\frac{\lfloor c + U \rfloor - U}{M} \sim \text{Unif} \left( \frac{c}{M} - \frac{1}{2M}, \frac{c}{M} + \frac{1}{2M} \right).$$

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**DQ + Bits-back:** Bits-back Quantization (BBQ)

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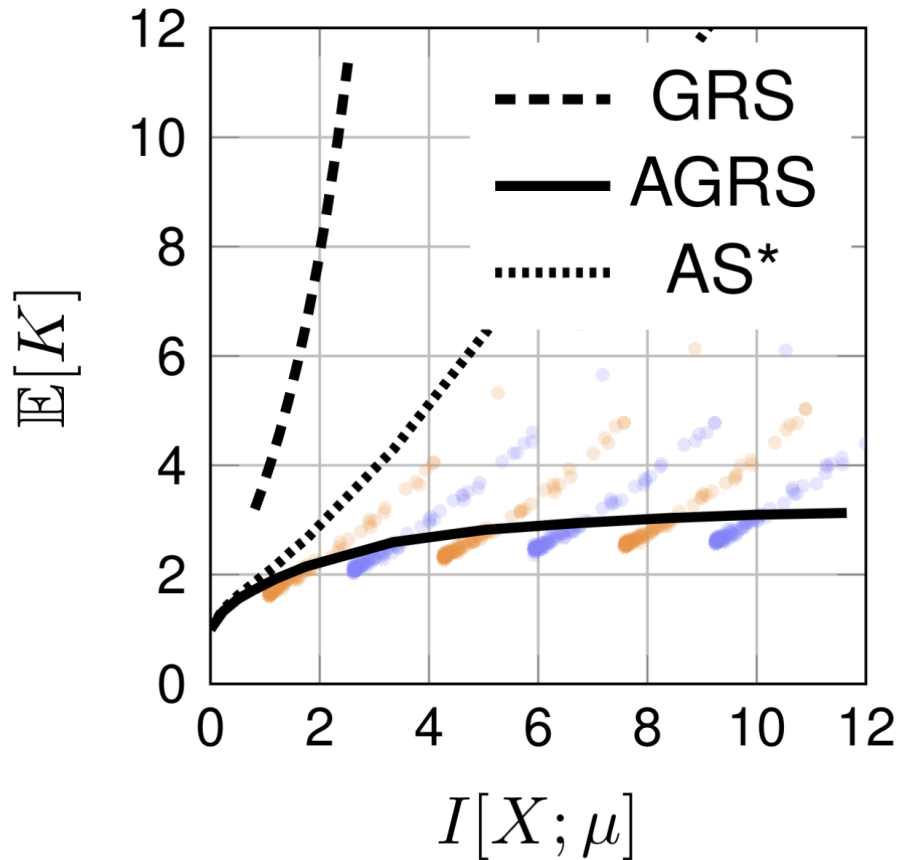
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Apply AGRS to our Gaussian example:



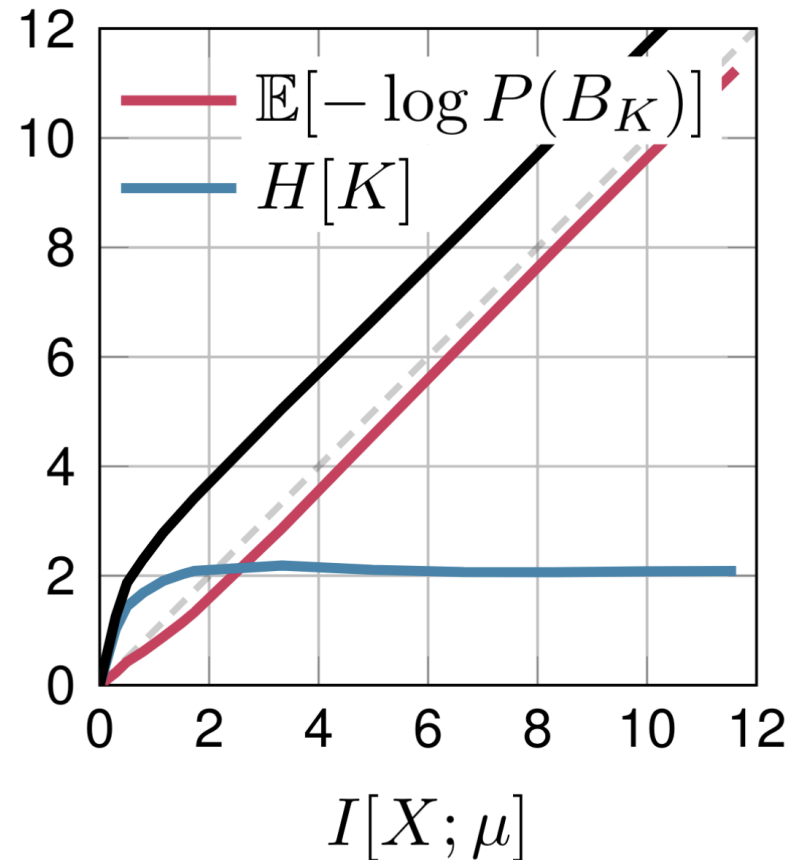
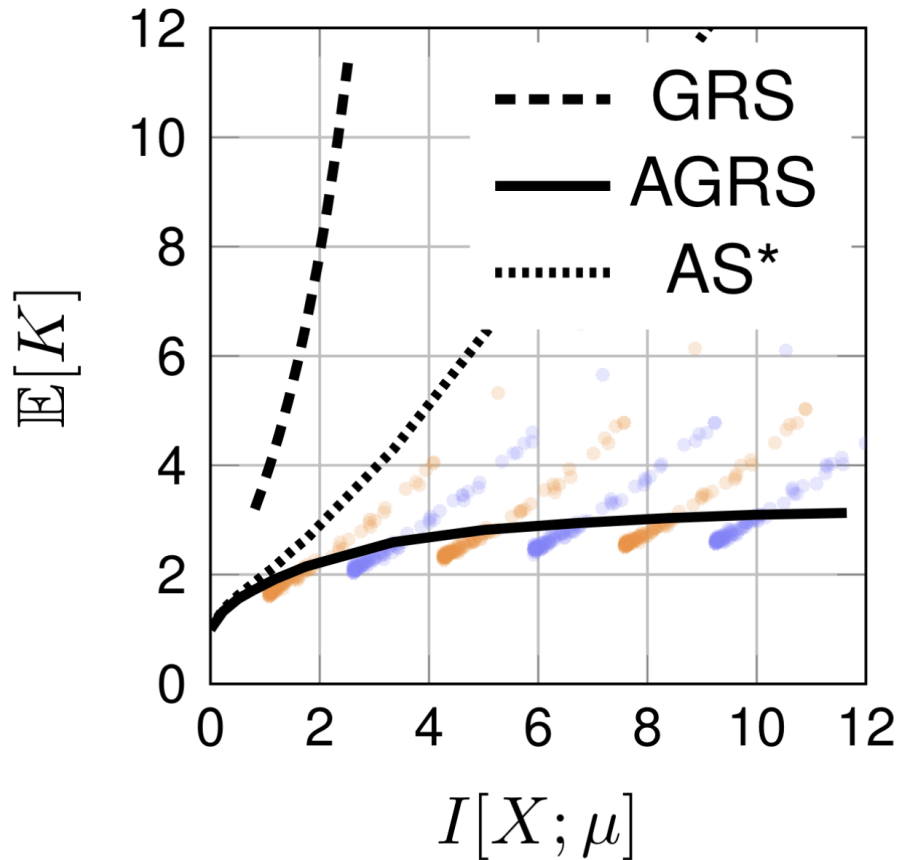
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3. Specialized algorithms for Gaussians?

# 6. REFERENCES

- P. Harsha, R. Jain, D. McAllester, and J. Radhakrishnan, “The communication complexity of correlation,” *IEEE Transactions on Information Theory*, vol. 56, no. 1, pp. 438–449, 2010.
- C. T. Li and A. El Gamal, “Strong functional representation lemma and applications to coding theorems,” *IEEE Transactions on Information Theory*, vol. 64, no. 11, pp. 6967–6978, 2018.
- F. “Greedy Poisson Rejection Sampling,” *arXiv preprint arXiv:2305.15313*, 2023.

# 7. BBQ

