

GAMS Tutorial

Systems and Computing Seminar

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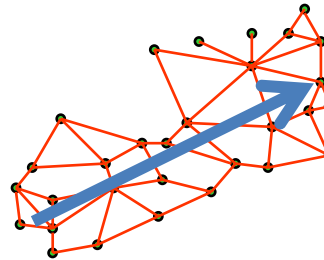
Generalities

What is GAMS?

- The General Algebraic Modeling System (GAMS) is a high-level modeling system for mathematical optimization. GAMS is designed for modeling and solving linear, nonlinear, and mixed-integer optimization problems.
- GAMS contains an integrated development environment (IDE) and is connected to a group of third-party optimization solvers. Among these solvers are BARON, COIN-OR solvers, CONOPT, CPLEX, DICOPT, GUROBI, MOSEK, SNOPT, SOLUM, and XPRESS.

General Purpose

Real-life problem



Mathematical Model

$$\text{Min}(d) = \sum_{(i,j) \in E} c_{ij} x_{ij}$$

S.A.

$$\sum_{(i,j) \in E} x_{ij} = 1, i = s$$

$$\sum_{(i,j) \in E} x_{ij} = 1, j = t$$

$$\sum_{(i,j) \in E} x_{ij} - \sum_{(j,i) \in E} x_{ji} = 0, i \neq s, t$$

$$x_{ij} \in \{0, 1\}$$

Implementation

```

Variables
x(i,j)      Indica si el enlace i j es utilizado o no en el SP
z           minimization ;
Binary Variable x;

Equations
Camino_Mas_Corto      Function Objetivo
nodo_origen(i)        nodo origen
nodo_destino(j)       nodo destino
nodo_intermedio       nodo intermedio;

Camino_Mas_Corto      .. z =e= sum((i,j), c(i,j) * x(i,j));

nodo_origen(i)$ord(i) = 1)      .. sum((j), x(i,j)) =e= 1;

nodo_destino(j)$ord(j) = 5)     .. sum((i), x(i,j)) =e= 1;

nodo_intermedio(i)$ord(i) ne 1 and ord(i) ne 5)
.. sum((j), x(i,j)) - sum((j), x(i,j)) =e= 0;

Model Transport /all/ ;
option mip=CPLEX
Solve transport using mip minimizing z;

Display x.l
Display z.l
  
```

Download

- <http://www.gams.com/>
- <http://www.gams.com/download/>

Solution Architecture

Solution Architecture

- File types:

- *.gms

- *.log

```

Variables
  x(i,j)      Indicates if the link i-j is selected or not.
  z           Objective function ;

Binary Variable x;

Equations
  objectiveFunction      objective function
  sourceNode(i)         source node
  destinationNode(j)    destination node
  intermediateNode       intermediate node;

objectiveFunction      .. z =e= sum((i,j), c(i,j) * x(i,j));
sourceNode(i)$ (ord(i) = 1)      .. sum((j), x(i,j)) =e= 1;
destinationNode(j)$ (ord(j) = 5) .. sum((i), x(i,j)) =e= 1;
intermediateNode(i)$ (ord(i) <> 1 and ord(i) ne 5) .. sum((j), x(i,j)) - sum((j), x(j,i)) =e= 0;

Model modell /all/ ;
option mip=CPLEX
Solve modell using mip minimizing z;

Display x.l;
Display z.l;
  
```

```

--- Job SP_Hops_v2.GMS Start 09/28/15 02:41:12 WEX-WEI 24.0.2 x86_64/MS Windows
GAMS Rev 240 Copyright (C) 1987-2013 GAMS Development. All rights reserved
Licensee: Ingenieria Quimica                      G120216:1205AP-WIN
          Universidad de los Andes                  DC9561
          License for teaching and research at degree granting institutions

--- Starting compilation
--- SP_Hops_v2.GMS (44) 3 Mb
--- Starting execution: elapsed 0:00:00.003
--- SP_Hops_v2.GMS (40) 4 Mb
--- Generating MIP model Transport
--- SP_Hops_v2.GMS (41) 4 Mb
--- 3 rows 26 columns 36 non-zeroes
--- 25 discrete-columns
--- Executing CPLEX: elapsed 0:00:00.026

IBM ILOG CPLEX Feb 14, 2013 24.0.2 WEX 38380.38394 WEI x86_64/MS Windows
Cplex 12.5.0.0

Reading data...
Starting Cplex...
Tried aggregator 2 times.
MIP Presolve eliminated 1 rows and 24 columns.
Aggregator did 2 substitutions.
All rows and columns eliminated.
Presolve time = 0.00 sec. (0.01 ticks)
MIP status(101): integer optimal solution
Cplex Time: 0.01sec (det. 0.02 ticks)
Fixing integer variables, and solving final LP...
Tried aggregator 1 time.
LP Presolve eliminated 3 rows and 26 columns.
  
```

Double-Click to Open File

Solution Architecture

- File types:
 - *.lst

```
----- 47 VARIABLE x.L Indicates if the link i-j is selected or not.
          n2          n5
n1        1.000
n2                1.000

----- 48 VARIABLE z.L = 2.000 Objective function

EXECUTION TIME = 0.000 SECONDS 3 Mb WEX240-240 Feb 14, 2013
```


Model Components

Example 1

- From 5 saleable goods, buy the ones that incur the minimal possible cost, taking into account our budget, that is, 10 monetary units.

Goods' values: 12, 5, 9, 6 y 4 respectively.

- Mathematical model:

value_i: parameter. Value of each article.

x_i: variable, where $x_i = \begin{cases} 1 & \text{if the article is bought} \\ 0 & \text{if not} \end{cases}$

$$value_1 * x_1 + value_2 * x_2 \dots value_5 * x_5$$

$$\min (value_1 * x_1 + value_2 * x_2 \dots value_5 * x_5)$$

$$value_1 * x_1 + value_2 * x_2 \dots value_5 * x_5 = BUDGET$$

Example 1

- Mathematical model:

$value_i$: parameter. Value of each article.

x_i : variable, where $x_i = \begin{cases} 1 & \text{if the article is bought} \\ 0 & \text{if not} \end{cases}$

$$\min (value_1 * x_1 + value_2 * x_2 \dots value_5 * x_5) \Rightarrow \min \sum_{i \in A} value_i * x_i$$

$$value_1 * x_1 + value_2 * x_2 \dots value_5 * x_5 = BUDGET$$

$$\sum_{i \in A} value_i * x_i = BUDGET$$

GAMS Implementation

- Mathematical model:

$value_i$: parameter. Value of each article.

x_i : variable, where $x_i = \begin{cases} 1 & \text{if the article is bought} \\ 0 & \text{if not} \end{cases}$

$\min \sum_{i \in A} value_i * x_i \longrightarrow \text{Objective Function}$

$\sum_{i \in A} value_i * x_i = BUDGET \longrightarrow \text{Constraint}$

- GAMS:

```
*Model1

Set i    articles / a1, a2, a3, a4, a5 /;

Scalar BUDGET budget /10/;

Parameter value(i)    value of each article
                    / a1 12, a2 5, a3 9, a4 6, a5 4 /;

Variables
  x(i)            Indicates if the article is bought or not
  z              objective function;

Binary Variable x;

Equations
objectiveFunction      objective function
budgetConstraint       budget constraint;

objectiveFunction      ..      z =e= sum(i, value(i) * x(i));

budgetConstraint       ..      sum(i, value(i) * x(i)) =e= BUDGET;

Model Model1 /all/ ;

option mip=CPLEX
Solve Model1 using mip minimizing z

Display x.l;
Display z.l;
```

Model Components

```

*Model1

Set i   articles / a1, a2, a3, a4, a5 /;

Scalar BUDGET budget /10/;

Parameter value(i)    value of each article
                    /  a1 12, a2 5, a3 9, a4 6, a5 4  /;

Variables
  x(i)    Indicates if the article is bought or not
  z       objective function;

Binary Variable x;

Equations
  objectiveFunction    objective function
  budgetConstraint     budget constraint;

objectiveFunction      ..    z =e= sum(i, value(i) * x(i));
budgetConstraint       ..    sum(i, value(i) * x(i)) =e= BUDGET;

Model Model1 /all/ ;

option mip=CPLEX
Solve Model1 using mip minimizing z

Display x.l;
Display z.l;

```

→ Sets
 Parameters
 Variables
 Equations declaration
 Equations definition
 Solver configuration
 Results visualization in the *.lst file

Model Components

- *.lst file contents:
 - A copy of the mathematical model

```
1 *Modelo1
2
3 Set i    articles / a1, a2, a3, a4, a5 /;
4
5 Scalar BUDGET budget /10/;
6
7 Parameter value(i)    value of each article
8                        /  a1 12, a2 5, a3 9, a4 6, a5 4  /;
9
10 Variables
11   x(i)          if the article is bought
12   z              objective function;
13
14 Binary Variable x;
15
16 Equations
17 objectiveFunction          objective function
18 budgetConstraint           budget constraint;
19
20 objectiveFunction          ..      z =e= sum(i, value(i) * x(i));
21
22 budgetConstraint           ..      sum(i, value(i) * x(i)) =e= BUDGET;
23
24 Model Model1 /all/ ;
25
26 option mip=CPLEX
27 Solve Model1 using mip minimizing z
28
29 Display x.l;
30 Display z.l;
31
```

Model Components

- *.lst file contents:
 - Equations

```
---- objectiveFunction =E= objective function

objectiveFunction.. - 12*x(a1) - 5*x(a2) - 9*x(a3) - 6*x(a4) - 4*x(a5) + z =E=
    0 ; (LHS = 0)

---- budgetConstraint =E= budget constraint

budgetConstraint.. 12*x(a1) + 5*x(a2) + 9*x(a3) + 6*x(a4) + 4*x(a5) =E= 10 ;

    (LHS = 0, INFES = 10 ****)
```

Model Components

- *.lst file contents:
 - Solver

```
IBM ILOG CPLEX    Feb 14, 2013 24.0.2 WEX 38380.38394 WEI x86_64/MS Windows  
Cplex 12.5.0.0
```

```
MIP status(101): integer optimal solution  
Cplex Time: 0.00sec (det. 0.01 ticks)  
Fixing integer variables, and solving final LP...  
Fixed MIP status(1): optimal  
Cplex Time: 0.00sec (det. 0.00 ticks)  
Proven optimal solution.
```

```
MIP Solution:           10.000000    (0 iterations, 0 nodes)  
Final Solve:            10.000000    (0 iterations)
```


Model Components

- *.lst file contents:
 - Results

```

-----      29 VARIABLE x.L   if the article is bought
a4 1.000,      a5 1.000

-----      30 VARIABLE z.L           =      10.000  objective function

```

Sets

Sets

- Five elements:

```
Set i  articles / a1, a2, a3, a4, a5 /;
Set i  articles / a1*a5 /;
Set i  articles / 1*5 /;
```
- Create a copy set:

```
Set i  articles / 1*5 /;
alias(j,i);
```
- Parameterize the number of elements:

```
$Set NARTICLES 10

Set i  articles / a1*a%NARTICLES% /;

Set i  articles / 1*%NARTICLES% /;
```

- Description of elements:

```
Set i  articles
/a1  article 1
a2  article 2
a3  article 3
a4  article 4
a5  article 5
/;
```

Sets

- Subsets:

```
$Set TOTALARTICLES 10

Set k  articles / a1*a$TOTALARTICLES$ /;
Set i(k) five articles /a1*a5/;
Set j(k) four articles /a2*a5/;
```
- One-to-one Mapping: $A = \{(b, d), (a, c), (c, e)\}.$

```
set c countries
/ jamaica
haiti
guyana
brazil / ;

set p ports
/ kingston
s-domingo
georgetown
belem / ;

set ptoc(p, c) port to country relationship
/ kingston .jamaica
s-domingo .haiti
georgetown .guyana
belem .brazil /;

Display c;
Display p;
Display ptoc;
```

```
----      25 SET c  countries

jamaica,   haiti  ,    guyana ,    brazil

----      26 SET p  ports

kingston  ,    s-domingo ,    georgetown,    belem

----      27 SET ptoc  port to country relationship

                jamaica      haiti      guyana      brazil
kingston              YES
s-domingo              YES
georgetown              YES
belem                      YES
```

- Many-to-many Mapping:

```
set i / a, b /
j / c, d, e /
ij1(i,j) /a.c, a.d/
ij2(i,j) /a.c, b.c/
ij3(i,j) /a.c, b.c, a.d, b.d/ ;

Display i,j,ij1,ij2,ij3;
```

```
set i / a, b /
j / c, d, e /
ij1(i,j) /a.(c,d)/
ij2(i,j) /(a,b).c/
ij3(I,j) /(a,b).(c,d)/ ;
```

	----	36 SET ij1	
		c	d
a		YES	YES
	----	36 SET ij2	
		c	
a		YES	
b		YES	
	----	36 SET ij3	
		c	d
a		YES	YES
b		YES	YES

- Many-to-many Mapping:
 - Exercise:

<i>Construct</i>	<i>Result</i>
$(a,b).c.d$	
$(a,b).(c,d).e$	
$(a.1*3).c$	
$1*3. \quad 1*3. \quad 1*3$	

- Many-to-many Mapping:

- Exercise:

<i>Construct</i>	<i>Result</i>
$(a,b).c.d$	$a.c.d, b.c.d$
$(a,b).(c,d).e$	$a.c.e, b.c.e, a.d.e, b.d.e$
$(a.1*3).c$	$(a.1, a.2, a.3).c$ or $a.1.c, a.2.c, a.3.c$
$1*3. \quad 1*3. \quad 1*3$	$1.1.1, 1.1.2, 1.1.3, \dots, 3.3.3$

Data Entry

Parameters, Scalars &
Tables

Data Entry

- Scalars:


```
Scalar BUDGET budget /10/;

$Set BUDGET 10

Scalar BUDGET presupuesto /%BUDGET%/;
```
- Parameters:


```
Parameter value(i) value of each article
/ a1 12, a2 5, a3 9, a4 6, a5 4 /;
```

 - Parameter Data for Higher Dimensions:


```
parameter salaries(employee,manager,department)
/anderson .murphy .toy = 6000
hendry .smith .toy = 9000
hoffman .morgan .cosmetics = 8000 / ;
```
 - Exercise:


```
Set row / row1*row10 /
col / col1*col10 / ;

Parameter a(row, col)
/ (row1,row4) . col2*col7 12
row10 . col10 17
row1*row7 . col10 33 / ;
```

Data Entry

– Exercise:

```
Set row / row1*row10 /
    col / col1*col10 / ;

Parameter a(row, col)
    / (row1,row4) . col2*col7 12
    row10 . col10 17
    row1*row7 . col10 33 / ;
```

```
----- 64 PARAMETER a

           col2      col3      col4      col5      col6      col7
row1      12.000      12.000      12.000      12.000      12.000      12.000
row4      12.000      12.000      12.000      12.000      12.000      12.000

+         col10
row1      33.000
row2      33.000
row3      33.000
row4      33.000
row5      33.000
row6      33.000
row7      33.000
row10     17.000
```

Data Entry

- Tables:

$$\sum_{i \in P} \sum_{j \in B} C_{ij} X_{ij} \longrightarrow C_{ij} ?$$

Table c(i,j) link cost

	n1	n2	n3	n4	n5
n1	999	1	1	999	999
n2	999	999	999	999	1
n3	999	999	999	1	999
n4	999	999	999	999	1
n5	999	999	999	999	999;

---- 45 PARAMETER c link cost

	n1	n2	n3	n4	n5
n1	999.000	1.000	1.000	999.000	999.000
n2	999.000	999.000	999.000	999.000	1.000
n3	999.000	999.000	999.000	1.000	999.000
n4	999.000	999.000	999.000	999.000	1.000
n5	999.000	999.000	999.000	999.000	999.000

Data Entry

- Tables:
 - More than two dimensions?

```
table upgrade(strat,size,tech)
      small.tech1 small.tech2 medium.tech1 medium.tech2
strategy-1      .05      .05      .05      .05
strategy-2      .2      .2      .2      .2
strategy-3      .2      .2      .2      .2
strategy-4      .2      .2      .2      .2
```

```
Parameter upgrade(strat,size,tech);
upgrade('strategy-1','small','tech1')=0.05;
upgrade('strategy-1','small','tech2')=0.05;
upgrade('strategy-1','medium','tech1')=0.05;
upgrade('strategy-1','medium','tech2')=0.05;
```

Variables

Variables

- Declaration:

```
Variables
x(i)      Indicates if the article is bought or not
z         objective function;
```
- Variable type:

```
Binary Variable x;
```
- Variable types:

<i>Keyword</i>	<i>Default Lower Bound</i>	<i>Default Upper Bound</i>	<i>Description</i>
free (default)	-inf	+inf	No bounds on variable. Both bounds can be changed from the default values by the user
positive	0	+inf	No negative values are allowed for variable. The user can change the upper bound from the default value.
negative	-inf	0	No positive values are allowed for variables. The user can change the lower bound from the default value.
binary	0	1	Discrete variable that can only take values of 0 or 1
integer	0	100	Discrete variable that can only take integer values between the bounds. The user can change bounds from the default value.

Variables

```

Variables
  x(i)      Indicates if the article is bought or not
  z         objective function;

Binary Variable x;
.
.
.

Display x.1;

```

```

----      29 VARIABLE x.L   if the article is bought

a4 1.000,      a5 1.000

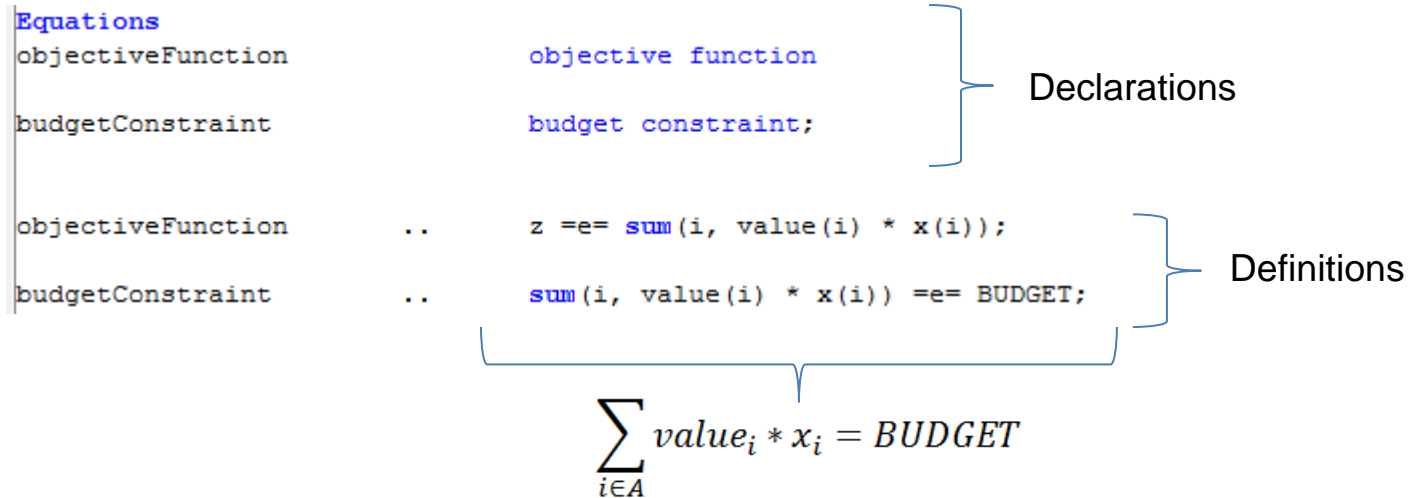
----      30 VARIABLE z.L           =      10.000  objective function

```

Equations

Equations

- Declaration and definition:



- Relational operators: =e=equal
 =g=greater than or equal to
 =l=less than or equal to

Equations

- Forall:

$$\min \sum_{i \in P} \sum_{j \in B} C_{ij} X_{ij}$$

$$\sum_{j \in B} X_{ij} \leq a_i \quad \forall i \in P$$

$$\sum_{i \in P} X_{ij} \geq b_j \quad \forall j \in B$$

Equations

```
cost define objective function
supply(i) observe supply limit at plant i
demand(j) satisfy demand at market j ;
```

```
cost .. z =e= sum((i,j), c(i,j)*x(i,j)) ;
supply(i) .. sum(j, x(i,j)) =l= a(i) ;
demand(j) .. sum(i, x(i,j)) =g= b(j) ;
```

Model and Solve Statements

```
Model Model1 /all/ ;  
Model Model1 /cost, supply, demand/ ;  
Model Model1 /cost, supply/ ;
```

→

```
cost .. z =e= sum((i,j), c(i,j)*x(i,j)) ;  
supply(i) .. sum(j, x(i,j)) =l= a(i) ;  
demand(j) .. sum(i, x(i,j)) =g= b(j) ;
```

```
option lp=CPLEX  
Solve Model1 using lp minimizing z;
```

```
Display x.l;  
Display z.l;
```

→ Solution for 'x' and 'z'.
'l' : level

Model and Solve Statements

- Problem type:

lp	for linear programming
qcp	for quadratic constraint programming
nlp	for nonlinear programming
dnlp	for nonlinear programming with discontinuous derivatives
mip	for mixed integer programming
rmip	for relaxed mixed integer programming
miqcp	for mixed integer quadratic constraint programming
rmiqcp	for relaxed mixed integer quadratic constraint programming
minlp	for mixed integer nonlinear programming
rminlp	for relaxed mixed integer nonlinear programming
mcp	for mixed complementarity problems
mpec	for mathematical programs with equilibrium constraints
rmpec	for relaxed mathematical program with equilibrium constraints
cns	for constrained nonlinear systems
emp	for extended mathematical programming

Model and Solve Statements

- Problem type:
 - File>Options>Solvers

Options

Editor | Execute | Output | Solvers | Licenses | Colors | File Extensions | Charts/GD

Project Defaults | Reset | Legend

Solver	License	CNS	DNLP	EMP	LP	MCP	MINLP	MIP	MIQCP	MPEC	NLP	QCP	RMINLP	RMIP	F
ALPHAEC	Demo						•		•						
AMPL	Full	-	-		-	-	-	-		-	-		-	-	
BARON	Demo		•		•		•	•	•		•	•	•	•	
BDMLP	Full				•			•						•	
BENCH	Full	-	-		-	-	-	-	-	-	-	-	-	-	
BONMIN	Full						•	•	•						
BONMINH	Demo						•	•	•						
CBC	Full				•			•						•	
CONOPT	Full	•	•		•						•	•	•	•	
CONVERT	Full	-	-		-	-	-	-	-	-	-	-	-	-	
COUENNE	Full		•				•		•		•	•	•		
CPLEX	Demo				•			X	•			•		•	
DE	Full			-											
DECIS	Demo			-											

OK Cancel

Solver Selection Legend

X Current selection

• Available selection

- Not available (Option statement only)

Conditionals

Conditionals

- a(b > 1.5) = 2 ;$
 - if $(b > 1.5)$, then $a = 2$
- Relational operators:
 - le, \leq less than or equal to
 - lt, $<$ strictly less than
 - eq, $=$ equal to
 - ne, \neq not equal to
 - ge, \geq greater than or equal to
 - gt, $>$ strictly greater than

Conditionals

- \$ operator in parameters:

```
b= sum(i$(a(i) ne 1), a(i)) ;
```

```
rho(i)$(sig(i) ne 0) = sig(i) - 1;
```

- It is equivalent to: $\text{rho}(i)$(\text{sig}(i)) = \text{sig}(i) - 1$

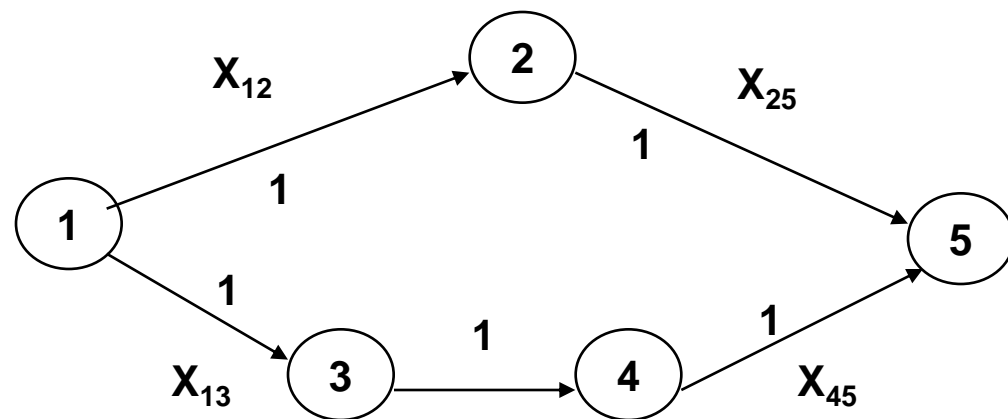
- \$ operator in constraints:

```
Eq2(i)$(cost(i)=5) .. x(i)=1=a(i) ;
```

```
Eq3(i)$(ord(i)<>1) .. x(i)=1=a(i) ;
```


Conditionals

- Example 2:



$$\min \sum_{i \in N} \sum_{j \in N} x_{ij}$$

Subject to:

$$\sum_{j \in N} x_{ij} = 1, i = s$$

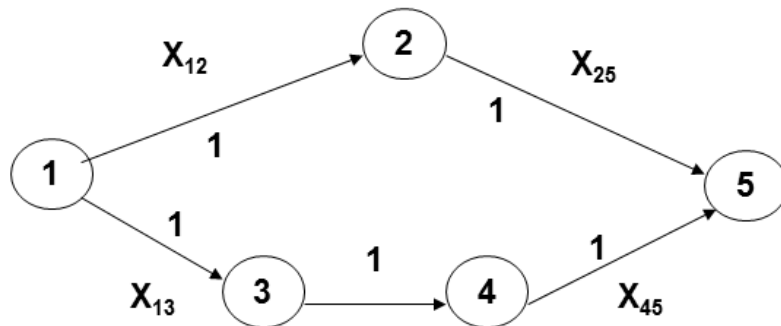
$$\sum_{i \in N} x_{ij} = 1, j = d$$

$$\sum_{j \in N} x_{ij} - \sum_{j \in N} x_{ji} = 0, i \neq s, d$$

$$x_{ij} \in \mathbb{Z}, 0/1$$

Conditionals

- Example 2:



$s=1$

$d=5$

```

Sets
  i  network nodes / n1, n2, n3, n4, n5 /

alias(j,i);

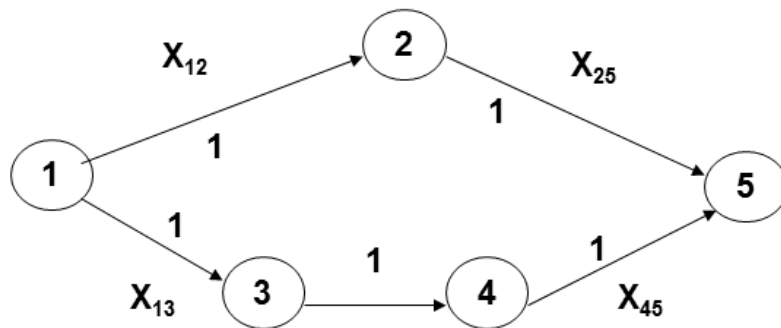
Table c(i,j) link cost
      n1      n2      n3      n4      n5
n1      999      1      1      999      999
n2      999      999      999      999      1
n3      999      999      999      1      999
n4      999      999      999      999      1
n5      999      999      999      999      999;

Variables
  x(i,j)      Indicates if the link i-j is selected or not.
  z           Objective function ;

Binary Variable x;
  
```

Conditionals

- Example 2:



$s=1$
 $d=5$

Equations

objectiveFunction
 sourceNode
 destinationNode
 intermediateNode2
 intermediateNode3
 intermediateNode4

objective function
 source node
 destination node
 intermediate node2
 intermediate node3
 intermediate node4;

objectiveFunction .. z =e= sum((i,j), c(i,j) * x(i,j));

sourceNode .. sum((j), x('n1',j)) =e= 1;

destinationNode .. sum((i), x(i,'n5')) =e= 1;

intermediateNode2 .. sum((i,j), x(i,'n2')) - sum((i,j), x('n2',j)) =e= 0;

intermediateNode3 .. sum((i,j), x(i,'n3')) - sum((i,j), x('n3',j)) =e= 0;

intermediateNode4 .. sum((i,j), x(i,'n4')) - sum((i,j), x('n4',j)) =e= 0;

Model Model1 /all/ ;

option mip=CPLEX

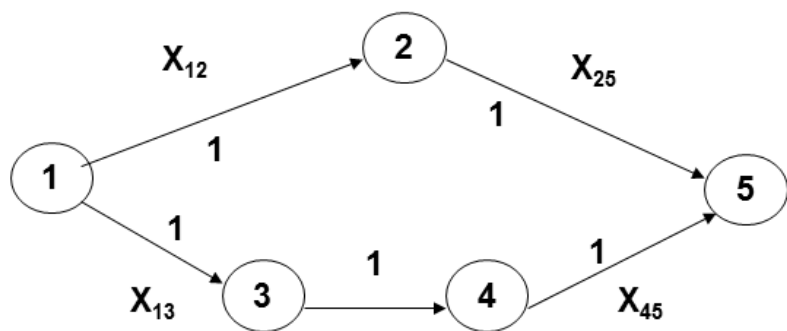
Solve Model1 using mip minimizing z;

Display x.l

Display z.l

Conditionals

- Example 2: generic version



s=1
d=5

```

Equations
objectiveFunction      objective function
sourceNode(i)          source node
destinationNode(j)     destination node
intermediateNode       intermediate node;

objectiveFunction      .. z =e= sum((i,j), c(i,j) * x(i,j));

sourceNode(i)$ (ord(i) = 1)      .. sum((j), x(i,j)) =e= 1;

destinationNode(j)$ (ord(j) = 5) .. sum((i), x(i,j)) =e= 1;

intermediateNode(i)$ (ord(i) <> 1 and ord(i) ne 5) .. sum((j), x(i,j)) - sum((j), x(j,i)) =e= 0;

Model model1 /all/ ;
option mip=CPLEX
Solve model1 using mip minimizing z;

Display x.l, c;
Display z.l;
  
```

```

---- 45 VARIABLE x.L Indicates if the link i-j is selected or not.

      n2      n5

n1      1.000
n2              1.000

---- 46 VARIABLE z.L = 2.000 Objective function
  
```

Programming Flow Control Features

Programming Flow Control Features

- Loop:

```
set t / 1985*1990 /  
parameter pop(t) / 1985 3456 /  
growth(t) / 1985 25.3, 1986 27.3, 1987 26.2  
          1988 27.1, 1989 26.6, 1990 26.6 /;  
loop(t,  
      pop(t+1) = pop(t) + growth(t) ) ;
```

- Conditional loop:

```
loop(i$(curacc > reltol),  
    value(i+1) = 0.5*(value(i) + target/value(i));  
    sqrtval = value(i+1);  
    curacc = abs (value(i+1)-value(i))/(1+abs(value(i+1)))  
);
```

- One cannot make declarations or define equations inside a loop statement.
- It is illegal to modify any controlling set inside the body of the loop.

Programming Flow Control Features

- If-elseif-else:

```
if (f <= 0,
    p(i) = -1 ;
    q(j) = -1 ;
elseif ((f > 0) and (f < 1)),
    p(i) = p(i)**2 ;
    q(j) = q(j)**2 ;
else
    p(i) = p(i)**3 ;
    q(j) = q(j)**3 ;
) ;
```

- One cannot make declarations or define equations inside an if statement.

Programming Flow Control Features

- while:

```
scalar count ; count = 1 ;
scalar globmin ; globmin = inf ;
while((count le 1000),
    a(j) = uniform(0,1) ;
    solve m1 using nlp minimizing obj ;
    if (obj.l le globmin,
        globmin = obj.l ;
    ) ;
    count = count+1 ;
) ;
```

- One cannot make declarations or define equations inside a while statement.

Programming Flow Control Features

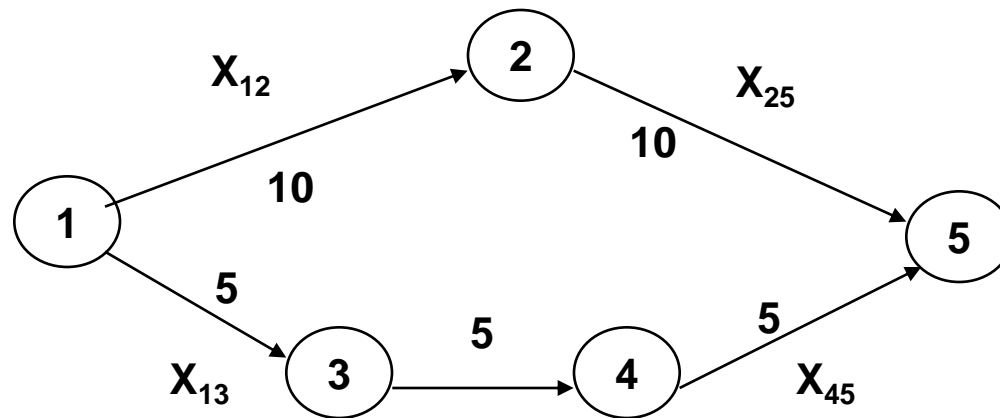
- for:

```
scalar i ;
scalar globmin ; globmin = inf ;
for (i = 1 to 1000,
    a(j) = uniform(0,1) ;
    solve m1 using nlp minimizing obj ;
    if (obj.l le globmin,
        globmin = obj.l ;
        globinit(j) = x.l(j) ;
    );
);
```

- One cannot make declarations or define equations inside a for statement.

Multiobjective Optimization

- Concept: Example
 - Function 1: minimize hops
 - Function 2: minimize cost



Multiobjective Optimization

- Theoretical basis:

Optimize [minimize/maximize]

$$F(X) = \{f_1(X), f_2(X), \dots, f_n(X)\}$$

subject to

$$H(X) = 0$$

$$G(X) \geq 0$$

Multiobjective Optimization

- Weighted Sum Method:

Optimize [minimize/maximize]

$$F'(X) = \sum_{i=1}^n r_i * f_i(X)$$

subject to

$$H(X) = 0$$

$$G(X) \geq 0$$

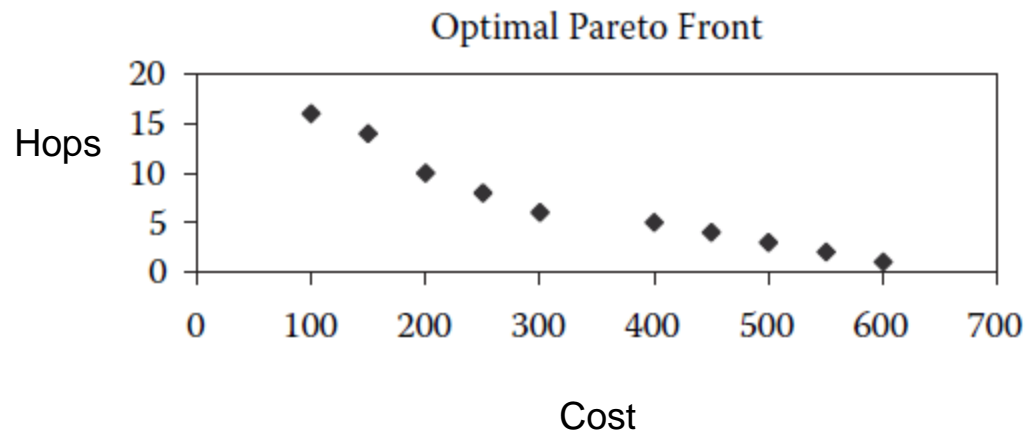
$$0 \leq r_i \leq 1, i = \{1, \dots, n\}$$

$$\sum_{i=1}^n r_i = 1$$

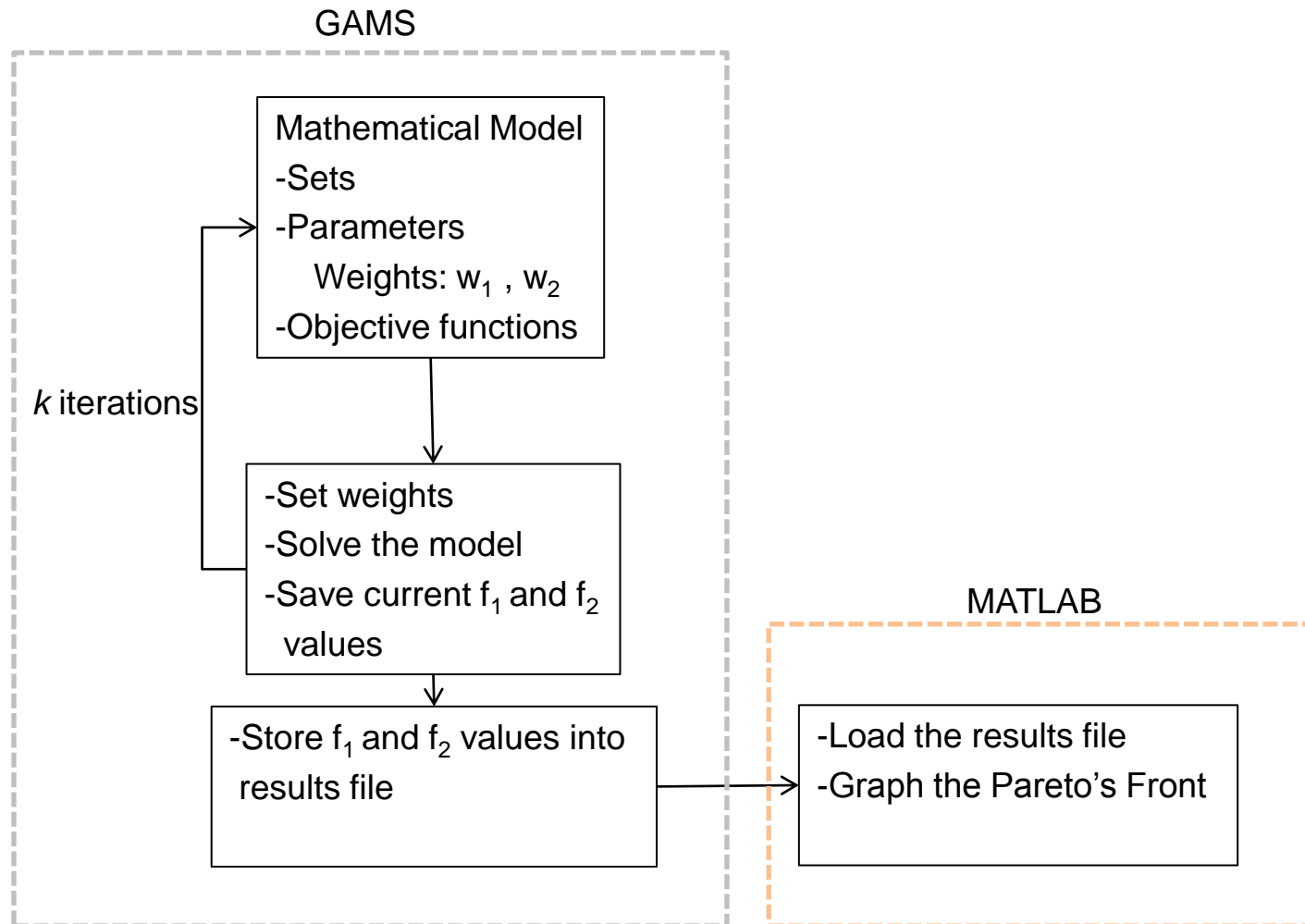
Multiobjective Optimization

- Weighted Sum Method:

$$F(X) = r_1 \cdot f_1(X) + r_2 \cdot f_2(X)$$



Multiobjective Optimization



Multiobjective Optimization

- Example 3:

Suppose there are two types of packets in a network: without priority and with priority. The network has three source nodes and four destination nodes. From the source nodes it is required to send 60, 80 and 50 packets without priority, respectively, and 20, 20 and 30 packets with priority respectively. At the destination nodes are requested 50, 90, 40 and 10 packets without priority respectively, and 10, 20, 10 and 30 packets with priority, respectively. The sending costs from source nodes to destination nodes are presented in the following table:

Source node	Destination 1	Destination 2	Destination 3	Destination 4
1	10	9	10	11
2	9	10	11	10
3	11	9	10	10

In addition, it is necessary to take into account the sending delay from source nodes to destination nodes:

Source node	Destination 1	Destination 2	Destination 3	Destination 4
1	12	14	10	11
2	11	8	7	13
3	6	11	4	15

Propose a multiobjective optimization model for minimizing the sending cost and the sending delay, finding the Pareto front using the Weighted Sum method.

Multiobjective Optimization

• Example 3:

Suppose there are two types of packets in a network: without priority and with priority. The network has three source nodes and four destination nodes. From the source nodes it is required to send 60, 80 and 50 packets without priority, respectively, and 20, 20 and 30 packets with priority respectively. At the destination nodes are requested 50, 90, 40 and 10 packets without priority respectively, and 10, 20, 10 and 30 packets with priority, respectively. The sending costs from source nodes to destination nodes are presented in the following table:

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3	6	11	4	15

Propose a multiobjective optimization model for minimizing the sending cost and the sending delay, finding the Pareto front using the Weighted Sum method.

Sets:

i : set of packet types

j : set of source nodes

k : set of destination nodes

Parameters:

c_{jk}	d_1	d_2	d_3	d_4
s_1	10	9	10	11
s_2	9	10	11	10
s_3	11	9	10	10

t_{jk}	d_1	d_2	d_3	d_4
s_1	12	14	10	11
s_2	11	8	7	13
s_3	6	11	4	15

inv_{ij}	s_1	s_2	s_3
p_1	60	80	50
p_2	20	20	30

$demand_{ik}$	d_1	d_2	d_3	d_4
p_1	50	90	40	10
p_2	10	20	10	30

Multiobjective Optimization

• Example 3:

Suppose there are two types of packets in a network: without priority and with priority. The network has three source nodes and four destination nodes. From the source nodes it is required to send 60, 80 and 50 packets without priority, respectively, and 20, 20 and 30 packets with priority respectively. At the destination nodes are requested 50, 90, 40 and 10 packets without priority respectively, and 10, 20, 10 and 30 packets with priority, respectively. The sending costs from source nodes to destination nodes are presented in the following table:

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3	11	9	10	10

In addition, it is necessary to take into account the sending delay from source nodes to destination nodes:

Source node	Destination 1	Destination 2	Destination 3	Destination 4
1	12	14	10	11
2	11	8	7	13
3	6	11	4	15

Propose a multiobjective optimization model for minimizing the sending cost and the sending delay, finding the Pareto front using the Weighted Sum method.

Variables: $x_{ijk} \in Re^+$

$$f_1 = \sum_{ijk} c_{jk} * x_{ijk} \quad f_2 = \sum_{ijk} t_{jk} * x_{ijk}$$

Objective function:

$$\min (w_1 * f_1 + w_2 * f_2)$$

$$where w_1 + w_2 = 1$$

Constraints:

$$\sum_k x_{ijk} \leq inv_{ij} \quad \forall i, j$$

$$\sum_j x_{ijk} = dem_{ik} \quad \forall i, k$$

Multiojective Optimization

• Example 3:

Sets:

- i : set of packet types
- j : set of source nodes
- k : set of destination nodes

Definition of weights:

Parameters:

c_{jk}	d_1	d_2	d_3	d_4
s_1	10	9	10	11
s_2	9	10	11	10
s_3	11	9	10	10

t_{jk}	d_1	d_2	d_3	d_4
s_1	12	14	10	11
s_2	11	8	7	13
s_3	6	11	4	15

inv_{ij}	s_1	s_2	s_3
p_1	60	80	50
p_2	20	20	30

dem_{ik}	d_1	d_2	d_3	d_4
p_1	50	90	40	10
p_2	10	20	10	30

```

Set i   packet types / p1, p2 /;
Set j   source nodes / s1, s2, s3 /;
Set k   destination nodes / d1, d2, d3, d4 /;

set iter iterations /it1*it11/;
scalar w1 weight 1 / 0 /;
scalar w2 weight 2 / 0 /;

parameter w1_vec(iter) w1 values
           /it1 1, it2 0.9, it3 0.8, it4 0.7, it5 0.6, it6 0.5,
           it7 0.4, it8 0.3, it9 0.2, it10 0.1, it11 0/;
parameter w2_vec(iter) w2 values;

Table c(j,k) sending cost
           d1      d2      d3      d4
s1         10      9      10      11
s2          9     10     11      10
s3         11      9     10     10;

Table t(j,k) sending delay
           d1      d2      d3      d4
s1         12     14     10     11
s2         11      8      7     13
s3          6     11      4     15;

Table inv(i,j) inventory
           s1      s2      s3
p1         60     80     50
p2         20     20     30;

Table dem(i,k) demand
           d1      d2      d3      d4
p1         50     90     40     10
p2         10     20     10     30;

```

Multiobjective Optimization

• Example 3:

Variables: $x_{ijk} \in \mathbb{R}^+$

$$f_1 = \sum_{ijk} c_{jk} * x_{ijk} \quad f_2 = \sum_{ijk} t_{jk} * x_{ijk}$$

Objective function:

$$\min (w_1 * f_1 + w_2 * f_2)$$

$$\text{where } w_1 + w_2 = 1$$

Constraints:

$$\sum_k x_{ijk} \leq \text{inv}_{ij} \quad \forall i, j$$

$$\sum_j x_{ijk} = \text{dem}_{ik} \quad \forall i, k$$

```

Variables
  x(i,j,k)    Amount of i type packets sent from the source node j
               to the destination node k.
  z            minimization
  f1           function 1
  f2           function 2;

Positive Variable x;

Equations
funObj                                Objective Function

invConstraint(i,j)                    inventory constraint

demConstraint(i,k)                    demand constraint

f1_value                                f1 value
f2_value                                f2 value;

f1_value                                ..    f1=e= sum((i,j,k), c(j,k) * x(i,j,k));

f2_value                                ..    f2=e= sum((i,j,k), t(j,k) * x(i,j,k));

funObj                                ..    z =e= w1*f1 + w2*f2;

invConstraint(i,j)                    ..    sum((k), x(i,j,k)) =l= inv(i,j);

demConstraint(i,k)                    ..    sum((j), x(i,j,k)) =e= dem(i,k);

Model Modell1 /all/ ;

```

Multiobjective Optimization

- Example 3:

z_res, f1_res, f2_res and x_res:

Parameters to store the values of z , $f1$, $f2$ and x at each iteration. These parameters are arrays to save a specified value at each iteration.

Loop statement:

Instruction for solving the problem at each iteration, according to certain values of w_1 and w_2 .

Display parameters:

Once all iterations have been performed, in the *.lst file are shown the results of these parameters.

Using the put writing instruction:

In the “result.dat” file are stored the values of the set “iter” and the arrays “f1_res” and “f2_res”.

```
parameter z_res(iter) "z results to store";
parameter f1_res(iter) "f1 results to store";
parameter f2_res(iter) "f2 results to store";
parameter x_res(i,j,k,iter) "x results to store";

loop (iter,
    w1=w1_vec(iter);
    w2=1 - w1_vec(iter);
    w2_vec(iter)=w2;

    option lp=CPLEX;
    Solve Model1 using lp minimizing z;
    z_res(iter)=z.l;
    f1_res(iter)=f1.l;
    f2_res(iter)=f2.l;
    x_res(i,j,k,iter)=x.l(i,j,k);
);

display z_res;
display f1_res;
display f2_res;
display w1_vec;
display w2_vec;
display x_res;

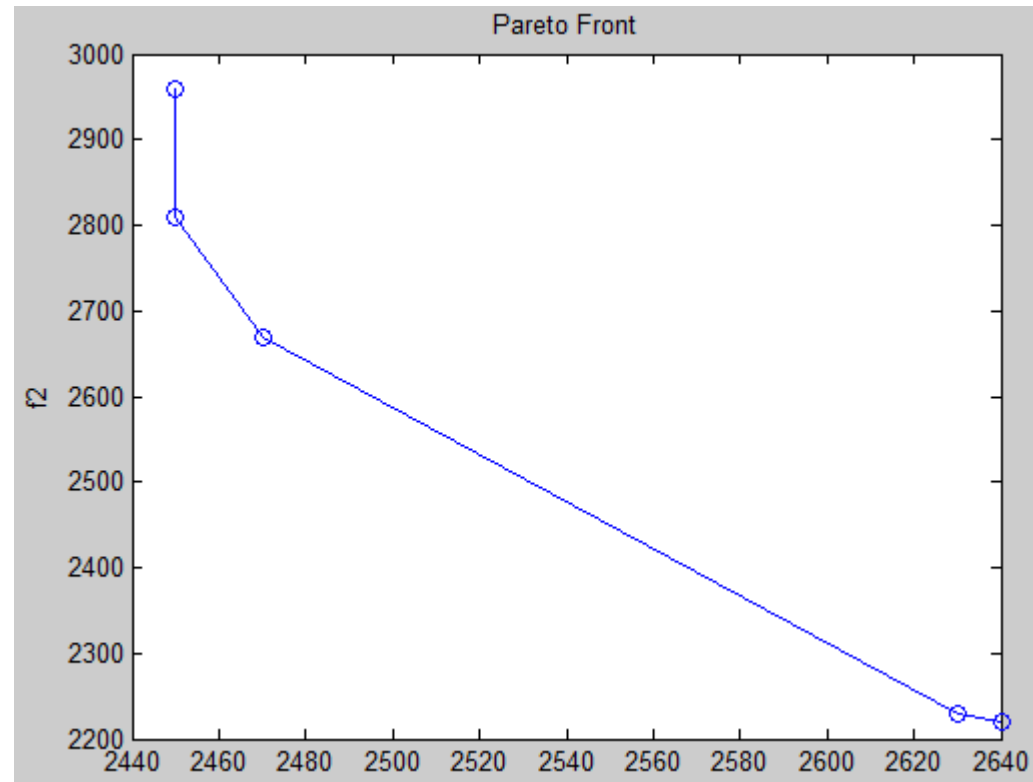
file GAMSresults /C:\DIRECTORY\results.dat/;
put GAMSresults;
loop(iter,
    put iter.tl, @5, f1_res(iter), @18, f2_res(iter) /

);
```

Multiojective Optimization

- Example 3:

```
1 - clc, clear all, close all
2
3 - [iter, f1, f2] = textread('results.dat', '%s %f %f', 20);
4
5 - figure
6 - plot(f1, f2, '-o')
7 - title('Pareto Front');
8 - xlabel('f1');
9 - ylabel('f2');
```



References

- https://en.wikipedia.org/wiki/General_Algebraic_Modeling_System
- <http://www.gams.com/>
- <http://www.gams.com/help/index.jsp>
- <http://www.gamsworld.org/performance/xpresslib/>

Thanks!

Questions?