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CS201x/B02

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Project 3 Write Up

## Identification of Mystery Abstract Data Types Using Time Cost Analysis

**Team Number:** 1015

### Results:

Structure Index	Structure Type
0	Binary Search Tree
1	Heap
2	Doubly Linked List
3	Hash Set
4	Binary Search Tree

### Raw Data:

Structure	N	T add	T remove	T contains	T removeLargest	T containsLargest
0	1	55	111	111	111	
0	2	55	139	138	166	
0	5	56	187	180	222	
0	10	56	229	218	277	
0	20	64	284	271	278	
0	30	67	305	291	333	
0	40	74	330	315	389	
0	50	86	342	335	333	
0	60	101	355	343	333	
0	70	104	374	358	333	
0	80	122	376	370	389	
0	90	123	381	375	389	
0	100	130	390	386	444	
0	200	247	435	439	444	
0	300	354	460	475	555	

0	400	409	486	492	556	
0	500	440	495	504	611	
0	600	454	511	525	555	
0	700	464	527	538	555	
0	800	480	533	546	556	
0	900	495	546	557	556	
0	1000	517	550	557	611	
0	2000	581	625	616	612	
0	3000	609	667	646	667	
0	4000	634	691	686	778	
0	5000	647	702	689	777	
0	6000	662	717	704	723	
0	7000	670	729	715	723	
0	8000	683	746	724	778	
0	9000	691	750	734	834	
0	10000	709	774	760	834	
1	1	85	170	85	170	85
1	2	101	213	125	170	85
1	5	107	391	253	340	85
1	10	110	615	466	341	85
1	20	104	1067	889	426	84
1	30	107	1470	1302	426	85
1	40	102	1922	1728	511	85
1	50	96	2401	2226	512	85
1	60	106	2706	2712	512	85
1	70	102	3153	2941	597	85
1	80	101	3598	3541	596	85
1	90	96	4127	3748	596	85
1	100	95	4465	4291	597	85
1	200	92	8679	8550	681	85
1	300	95	12695	13085	767	85
1	400	90	17429	17197	768	85
1	500	93	21515	21450	767	85
1	600	88	25924	25409	853	85
1	700	98	30959	30212	852	85
1	800	90	34204	34931	852	85
1	900	90	38796	39327	852	85
1	1000	99	43251	41976	853	85
1	2000	89	85177	86275	938	85
1	3000	110	131130	130522	1024	85
1	4000	87	164799	174926	1023	85
1	5000	96	218319	219048	1108	85
1	6000	89	253897	257491	1108	85

1	7000	124	307217	299078	1108	85
1	8000	87	340916	327947	1108	85
1	9000	89	377410	377575	1193	85
1	10000	91	436011	417331	1194	85
2	1	4	4	8	4	
2	2	4	5	9	8	
2	5	4	12	16	19	
2	10	4	22	26	39	
2	20	4	42	46	79	
2	30	4	60	67	119	
2	40	4	82	85	159	
2	50	4	103	107	199	
2	60	4	122	127	239	
2	70	4	145	144	279	
2	80	4	159	164	319	
2	90	4	188	182	359	
2	100	4	198	205	399	
2	200	4	417	406	799	
2	300	4	612	599	1199	
2	400	4	796	781	1599	
2	500	4	994	983	1998	
2	600	4	1244	1242	2398	
2	700	4	1415	1436	2799	
2	800	4	1622	1615	3198	
2	900	4	1787	1820	3599	
2	1000	4	1939	2005	3998	
2	2000	4	4010	3960	7994	
2	3000	4	6023	6031	11997	
2	4000	4	7890	8041	15993	
2	5000	4	9800	10052	19996	
2	6000	4	11568	12218	23993	
2	7000	4	14061	14397	27997	
2	8000	4	15926	16171	31999	
2	9000	4	17913	17991	35992	
2	10000	4	19613	19260	39992	
3	1	29	58	58	58	
3	2	29	58	58	58	
3	5	29	58	58	58	
3	10	29	58	58	58	
3	20	29	57	58	58	
3	30	30	57	58	58	
3	40	30	57	58	58	
3	50	30	57	58	58	

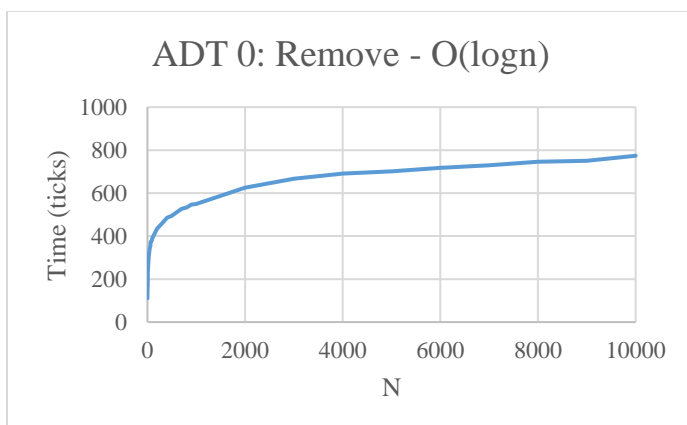
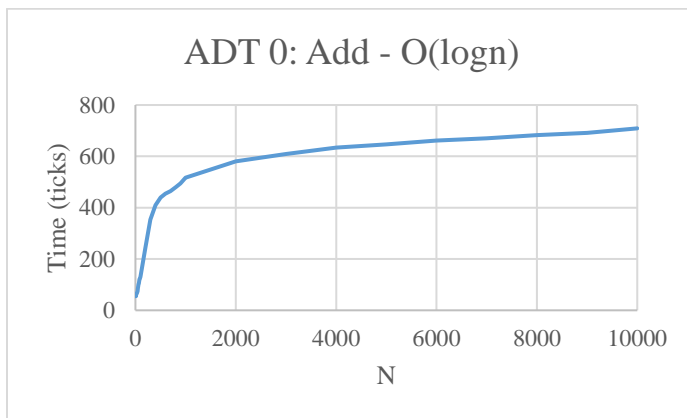
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3	90	31	55	58	58	
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3	200	35	52	57	58	
3	300	37	50	57	58	
3	400	39	48	56	58	
3	500	41	47	54	58	
3	600	41	46	54	58	
3	700	42	46	54	58	
3	800	43	45	52	58	
3	900	44	45	51	58	
3	1000	44	46	50	58	
3	2000	49	48	50	58	
3	3000	50	48	50	58	
3	4000	51	49	50	58	
3	5000	50	49	50	58	
3	6000	51	49	51	58	
3	7000	51	49	50	58	
3	8000	51	50	51	58	
3	9000	50	50	50	58	
3	10000	51	50	51	58	
4	1	57	115	115	115	
4	2	57	144	144	173	
4	5	58	195	188	231	
4	10	59	240	222	289	
4	20	70	286	268	346	
4	30	73	316	297	347	
4	40	77	344	330	347	
4	50	95	364	348	404	
4	60	103	375	364	462	
4	70	122	390	391	346	
4	80	118	392	378	520	
4	90	113	395	392	461	
4	100	136	407	403	462	
4	200	261	451	457	520	
4	300	365	478	492	520	
4	400	408	502	515	520	
4	500	458	523	528	577	
4	600	467	528	546	520	
4	700	501	553	569	635	
4	800	502	556	571	578	

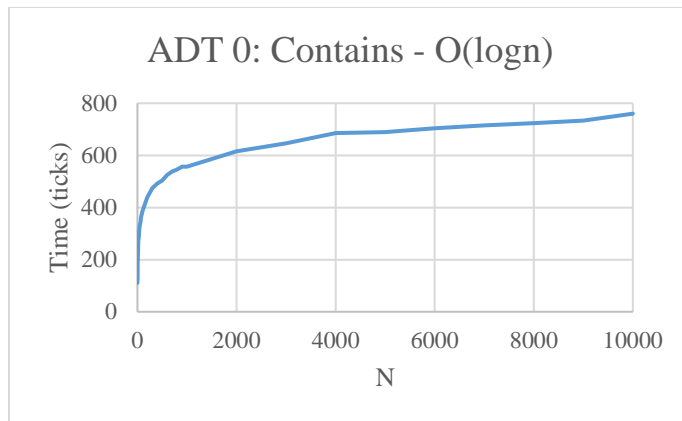
4	900	519	564	570	578	
4	1000	537	581	591	578	
4	2000	621	659	653	635	
4	3000	659	701	690	693	
4	4000	676	732	713	693	
4	5000	697	743	723	751	
4	6000	687	754	731	866	
4	7000	698	760	745	809	
4	8000	718	781	764	925	
4	9000	751	814	791	809	
4	10000	746	807	783	751	

### Conclusions (with Graphs):

In order to determine the time costs, we analyze the graphs to determine the slopes of the data. The  $O(1)$  time cost has a slope of zero, which is a straight horizontal line. The  $O(n)$  time cost has a slope of  $n$ , which is a straight diagonal line. The  $O(\log n)$  time cost has a slope of  $\log(n)$ , which is a line that begins almost completely vertical and continues curving until it is almost horizontal.

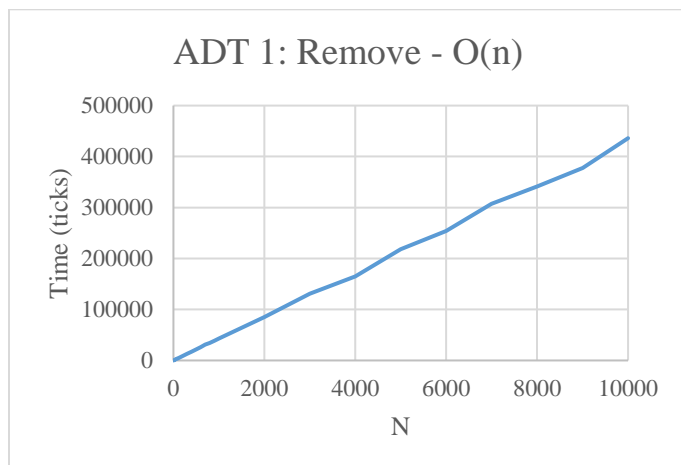
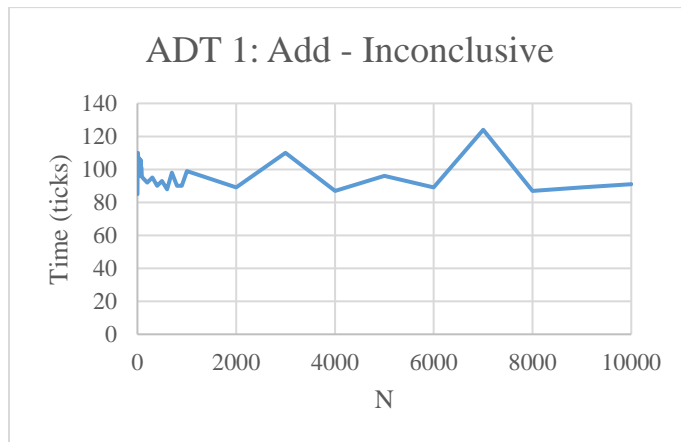
### Structure 0:

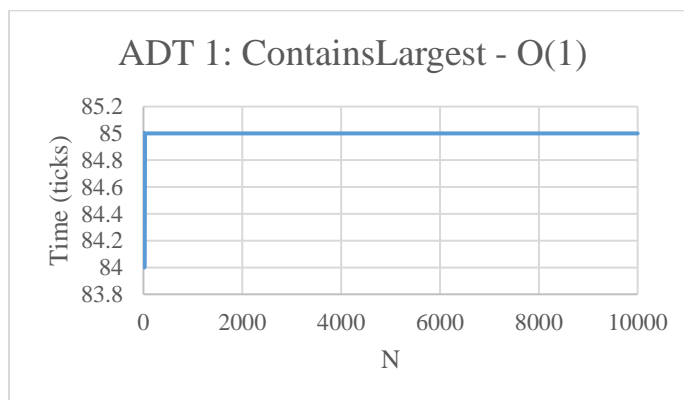
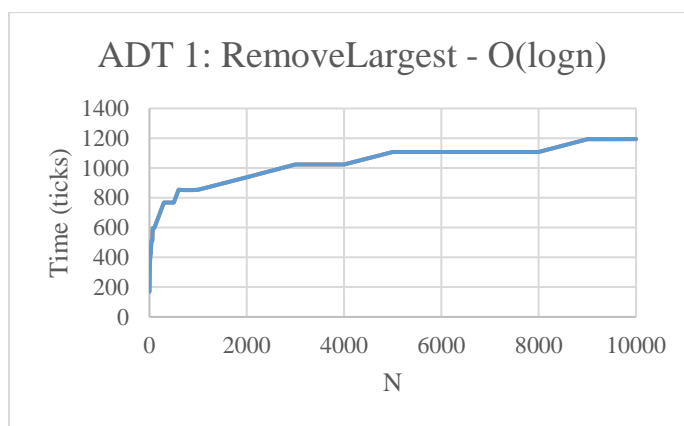
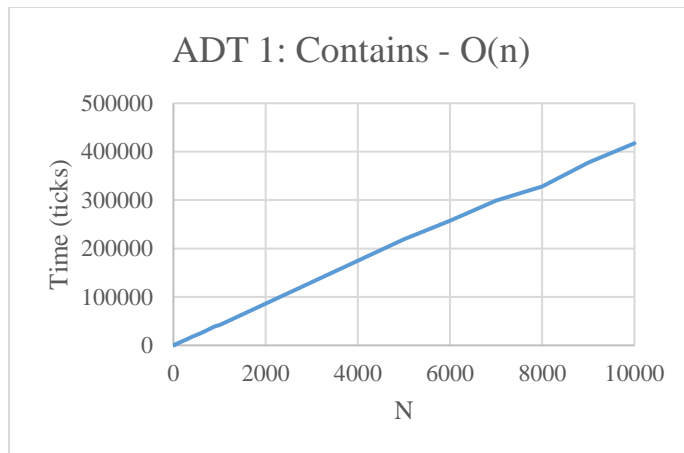




We classify the 0<sup>th</sup> data structure as a Binary Search Tree (BST). In the worst case, a BST should perform the add, remove, and contains methods all at an  $O(\log n)$  time. Graphing the data for the 0<sup>th</sup> data structure shows that the time cost of add is  $O(\log n)$ , the time cost of remove is  $O(\log n)$ , and the time cost of contains is  $O(\log n)$ . (See graphs.) Since this data corresponds with the expected time costs for a binary search tree, the 0<sup>th</sup> data structure is a binary search tree.

### Structure 1:

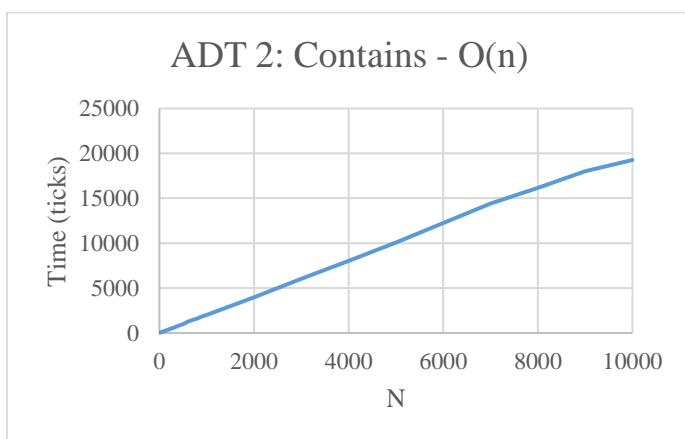
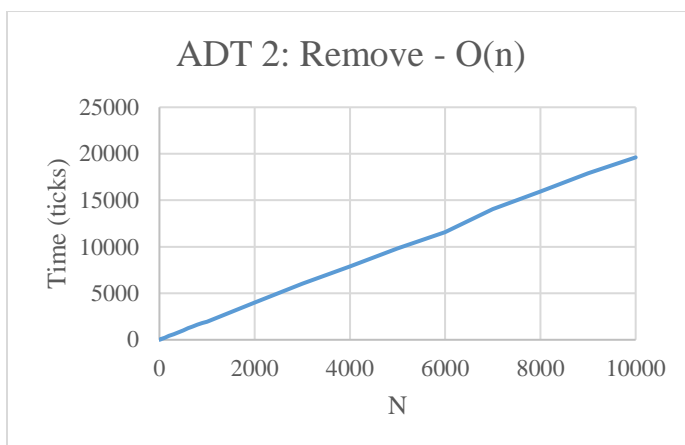
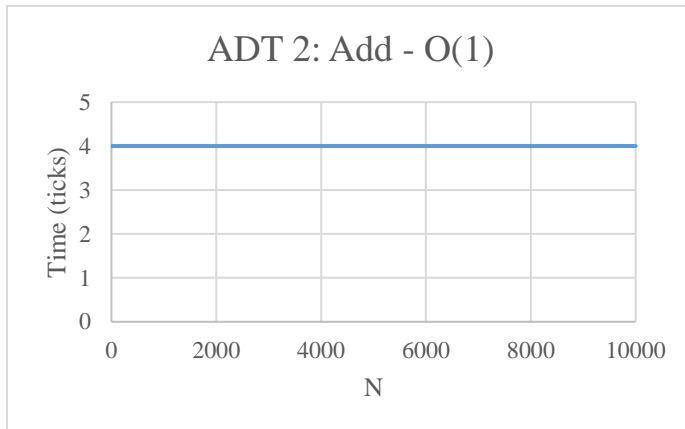




We classify the 1<sup>st</sup> data structure as a Heap. In the worst case, a heap should perform add in  $O(\log n)$  time, remove in  $O(n)$  time, and contains in  $O(n)$ . According to the time cost graphs for these functions, the first data structure performs remove in  $O(n)$  time and contains in  $O(n)$  time. However, the time for the add method is inconclusive. (See graphs.) Thus, more testing is required. In addition to the aforementioned functions, a heap should remove the largest element in  $O(\log n)$  time and should find the largest element in  $O(1)$  time. The graphs for remove largest element and contains largest element display  $O(\log n)$  and  $O(1)$  time, respectively, which also coincides with the heap structure. (See below.) Since four of the five of the first structure's time

costs conclusively match the expected time costs for a heap structure, we can conclude that the 1<sup>st</sup> data structure is a heap.

## Structure 2:

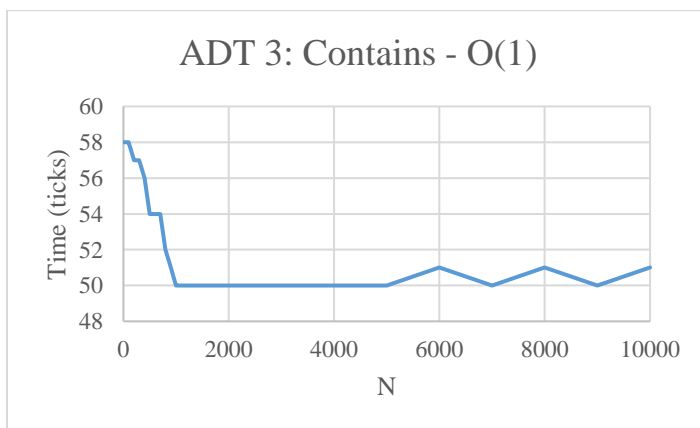
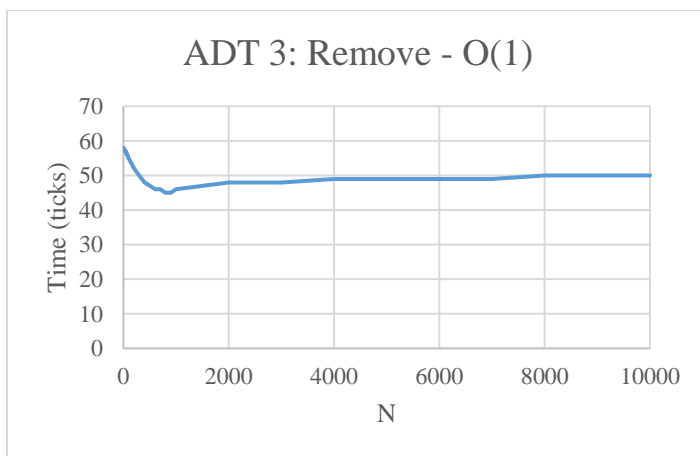
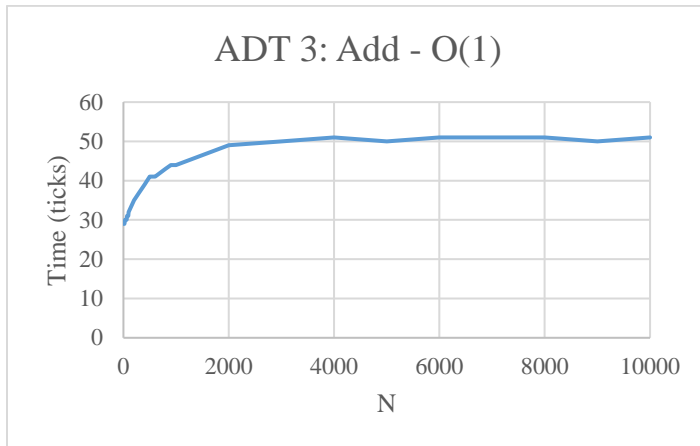


We classify the 2<sup>nd</sup> data structure as a Doubly Linked List. In the worst case, a doubly linked list should perform add in  $O(1)$  time, remove in  $O(n)$  time, and contains in  $O(n)$  time. The graphs of the add, remove, and contains functions for the 2<sup>nd</sup> show a time cost of  $O(1)$ ,  $O(n)$ , and  $O(n)$ ,



respectively. Since the time costs of the 2<sup>nd</sup> data structure match the expected time costs of the doubly linked list, we can conclude that the 2<sup>nd</sup> data structure is a doubly linked list.

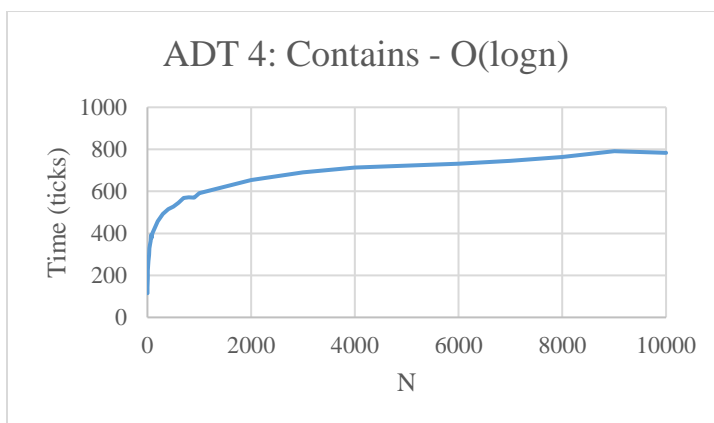
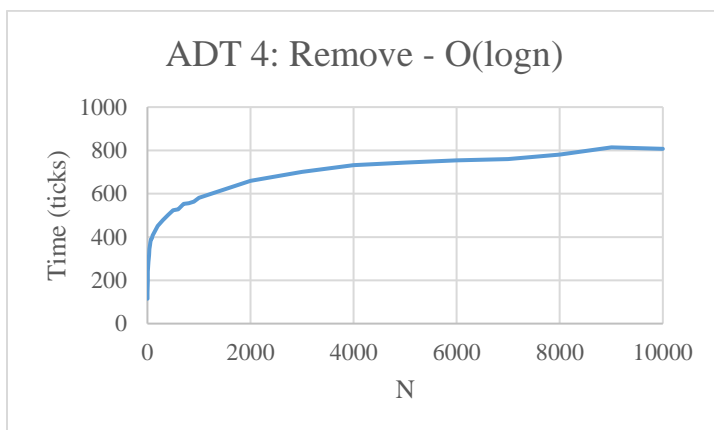
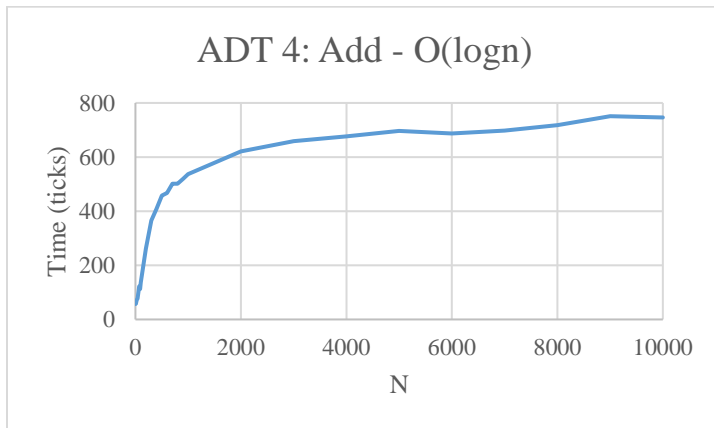
### Structure 3:



We classify the 3<sup>rd</sup> data structure as a Hash Set. In the average case, a hash set should perform add, remove, and contains in  $O(1)$  time. According to the graphs, the 3<sup>rd</sup> data structure performs

add in  $O(1)$  time, remove in  $O(1)$  time, and contains in  $O(1)$  time. These graphs look slightly out of the ordinary because there tends to be a curve in the beginning of the data. However, when zoomed out to the scale of the other data types, the curve is almost non-existent, and the end behavior is consistent with  $O(1)$  time. Therefore, we can conclude that the 3<sup>rd</sup> data structure performs the three methods in  $O(1)$  time, which matches the expected times for hash sets. Thus, the 3<sup>rd</sup> data structure is a hash set.

#### Structure 4:



We classify the 4<sup>th</sup> data structure as a Binary Search Tree. As stated in the conclusion for structure 0, a BST performs add, remove, and contains in  $O(\log n)$  time in the worst case. The graphs for the 4<sup>th</sup> data structure show that it performs the add, remove, and contains functions in  $O(\log n)$  time. Therefore, we can conclude that the 4<sup>th</sup> data structure is a Binary Search Tree.