

## Chapter 2.6

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Jan 28, 2009

I had planned to put the entries in my documents and then cut and paste them into xfig. But xfig did not have cut and paste. So, I'm not doing a figure for the CST. Instead, each path is represented as a sequence.

- 1

$$F(\exists x)(\varphi(x) \vee \psi(x)) \leftrightarrow (\exists x)\varphi(x)\psi(x) \vee (\exists x)\psi(x)$$

$$T(\exists x)(\varphi(x) \vee \psi(x))$$

$$F(\exists x)\varphi(x) \vee (\exists x)\psi(x)$$

$$T(\exists x)(\varphi(x) \vee \psi(x))$$

$T\varphi(c_1)$  8a: a new constant

$$F(\exists x)\varphi(x) \vee (\exists x)\psi(x)$$

$$F(\exists x)\varphi(x)$$

$$F(\exists x)\psi(x)$$

$F(\exists x)\varphi(x)$  8b: any constant

$$F\varphi(c_1)$$

\*

$$F(\exists x)(\varphi(x) \vee \psi(x))$$

$$T(\exists x)\varphi(x) \vee (\exists x)\psi(x)$$

$$F(\exists x)(\varphi(x) \vee \psi(x))$$

$F(\varphi(c_1) \vee \psi(c_1))$  any constant

$$F\varphi(c_1)$$

$$F\psi(c_1)$$

$$T(\exists x)\varphi(x) \vee (\exists x)\psi(x)$$

this is a branch using atomic tableau 4a

$$T(\exists x)\varphi(x)$$

$T\varphi(c_2)$  new constant, unify with any above  $c_1 = c_2$

\*

$$T(\exists x)\psi(x)$$

$T\psi(c_3)$  new constant, unify with any above  $c_1 = c_3$

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- 12a

$$F\forall x\exists y\neg\forall z\varphi(x, y, z) \leftrightarrow \forall x\exists y\exists z\neg\varphi(x, y, z)$$

this is branch using atomic tableau 6b

$$\begin{aligned} & T\forall x\exists y\neg\forall z\varphi(x, y, z) \\ & F\forall x\exists y\exists z\neg\varphi(x, y, z) \\ & T\forall x\exists y\neg\forall z\varphi(x, y, z) \\ & T\exists y\neg\forall z\varphi(c_1, y, z) \text{ any } c_1 \\ & T\neg\forall z\varphi(c_1, c_2, z) \text{ new } c_2 \\ & F\forall z\varphi(c_1, c_2, z) \\ & F\varphi(c_1, c_2, c_3) \text{ new } c_3 \\ & F\forall x\exists y\exists z\neg\varphi(x, y, z) \\ & F\exists y\exists z\neg\varphi(c_4, y, z) \text{ new } c_4 \\ & F\exists z\neg\varphi(c_4, c_5, z) \text{ any } c_5 \\ & F\neg\varphi(c_4, c_5, c_6) \text{ any } c_6 \\ & T\varphi(c_4, c_5, c_6) \\ & c_1 \text{ any } , c_4 \text{ new} \\ & c_2 \text{ new } , c_5 \text{ any} \\ & c_3 \text{ new } , c_6 \text{ any} \\ & * \end{aligned}$$

$$\begin{aligned} & F\forall x\exists y\neg\forall z\varphi(x, y, z) \\ & T\forall x\exists y\exists z\neg\varphi(x, y, z) \\ & F\forall x\exists y\neg\forall z\varphi(x, y, z) \\ & F\exists y\neg\forall z\varphi(c_1, y, z) \text{ new } c_1 \\ & F\neg\forall z\varphi(c_1, c_2, z) \text{ any } c_2 \\ & T\forall z\varphi(c_1, c_2, z) \\ & T\varphi(c_1, c_2, c_3) \text{ any } c_3 \\ & T\forall x\exists y\exists z\neg\varphi(x, y, z) \\ & T\exists y\exists z\neg\varphi(c_4, y, z) \text{ any } c_4 \\ & T\exists z\neg\varphi(c_4, c_5, z) \text{ new } c_5 \\ & T\neg\varphi(c_4, c_5, c_6) \text{ new } c_6 \\ & F\varphi(c_4, c_5, c_6) \\ & c_1 \text{ new } , c_4 \text{ any} \\ & c_5 \text{ new } , c_2 \text{ any} \\ & c_6 \text{ new } , c_3 \text{ any } * \end{aligned}$$

- 15 This is an infinite tableau, because the entry  $T(\forall x)R(x, x)$  will have an  $i$ 'th occurrence. For every  $i$ 'th occurrence, we will reduce (finish) the  $i$ 'th occurrence by introducing an  $i + 1$ 'th occurrence, as well as a new individual(for any ground term). Using rule 2 from definition 6.7 on p 116, the  $i$ 'th occurrence is finished when we insert the entry with the ground term as well as the  $i + 1$ 'th entry. So we have an infinite tableau, meaning that we have an infinite sequence of finite tableau. If there are entries of the form  $T\forall$  or  $F\exists$ , then each time we reduce such an entry, we introduce a new unreduced entry. However, for any unreduced entry, we will reduce it in a finite number of steps. So we can consider

a tableau finished (if infinite) if the only unreduced entries remaining are ones that are of the above form.