

## Logic for Applications

(Ch 12) P 144: 1

(Ch 13) P 151: 1, 5

(Ch 14) P 157: 2,3

**Ch 12****1.**

(a)

 $S0 = \{P(x,y), P(y,f(z))\}$  $d1 = \{x,y\}$  $s1 = \{x/y\}$  $S1 = \{P(y,y), P(y,f(z))\}$  $d2 = \{y, f(z)\}$  $s2 = \{f(z), y\}$  $S2 = \{P(y,f(z))\}$ 

(b)

 $S0 = \{P(a,y,f(y)), P(z,z,u)\}$  $d1 = \{a,z\}$  $s1 = \{a/z\}$  $S1 = \{P(z,y,f(y)), P(z,z,u)\}$  $d2 = \{y,z\}$  $s2 = \{y/z\}$  $S2 = \{P(z,z,f(y)), P(z,z,u)\}$  $d3 = \{f(y), u\}$  $s3 = \{u/f(y)\}$  $S3 = \{P(z,z,f(y))\}$ 

And so on.

**Ch 13****1.**

(a)

 $\{P(x,y), P(y,z)\}, \{\sim P(u, f(u))\}$  $z/f(u), y/u$  $\{P(x,u), P(u,f(u))\}, \{\sim P(u, f(u))\}$ unify  $P(u, f(u))$  $\{P(x,u)\}$ 

(b)

 $\{P(x,x), \sim R(x,f(x))\}, \{R(x,y), Q(y,z)\}$  $y/f(x)$  $\{P(x,x), \sim R(x,f(x))\}, \{R(x,f(x)), Q(f(x),z)\}$ unify  $R(x,f(x))$  $\{P(x,x)\}, \{Q(f(x),z)\}$ 

And so on.

**5. v, ii, iv, i**

Did these in a pretty much random order; I'm sure the steps can be much shortened.

1. resolve (v) and (ii) to get:

$$(vi) = \{\sim Q(h(b),v)\}$$

2. resolve (vi) and (iv) to get:

$$(vii) = \{P(a,u,f(h(u))), H(u,a), Q(h(b),b)\}$$

3. resolve (vii) and (v) to get:

$$(viii) = \{P(a,v,f(h(v))), Q(h(b),b)\}$$

4. resolve (viii) and (i) to get:

$$(ix) = \{P(a,v,f(h(v))), P(a,x,f(h(b))), P(a, f(h(b)))\}$$

5. resolve (ix) and (iii) to get:

$$(x) = \{P(a,x,f(h(v))), P(a, f(h(v))), H(x,a)\}$$

We can then resolve (x) with (iii) repeatedly until our only terms are the predicates  $H(a,b)$ , and then resolve that with (v) repeatedly. I don't think I need to type it out.

**Ch 14****2.**

By 13.6 and 13.10, we already have the soundness and completeness of resolution (that is not necessarily linear). For completeness, we must then show that if  $S$  is unsatisfiable, there is a linear resolution refutation of  $S$ , which we are given by theorem 14.4. For soundness, we assume that  $S$  produces box through linear resolution, and we must show that  $S$  must also be unsatisfiable. Suppose  $S$  is satisfiable. Then, by 13.6, box cannot be an element of the resolution of  $S$ , but it is, so this is the desired contradiction. (This probably isn't very correct, but I don't currently have time to wrestle with it more and will try to return to it later).

**3.**

This doesn't use the exact same syntactic rules from our book

i.

Ex.  $\text{Dragon}(x)$

ii.

Ex.  $\text{Dragon}(x) \ \& \ (\text{Sleeping}(x) \mid \text{Hunting}(x))$

iii.

Ex.  $\text{Dragon}(x) \ \& \ (\text{Hungry}(x) \rightarrow \sim \text{Sleeping}(x))$

iv.

Ex.  $\text{Dragon}(x) \ \& \ (\text{Tired}(x) \rightarrow \sim \text{Hunting}(x))$

(a).

Hunts in the forest (rules iii and ii)

(b).

Sleeping (rules iv and ii)