Quarter 1 Final, Data Structures

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Note: This is the work I did within the class time. I will also be sos_submitting a better version later on.

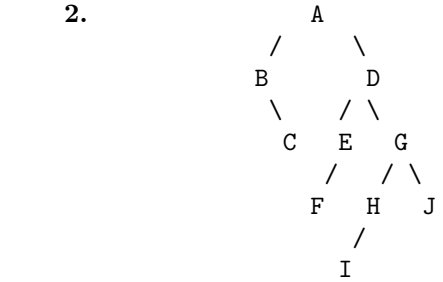
1. *(a)*

```
abcde
-f-g-
--h--
-i-j-
klmno
compact = [a,b,c,d,e,f,g,h,i,j,k,l,m,n,o]
(which is each row concatenated)
```

(b)

```
get(row, col, list)
  if(row == 1)
    return index(list, col)
  else if(row == 5)
    return index(list, 10+col)
  else if(row == col)
    return index(list, row + col + 1)
  else if(row == (6 - col))
    return index(list, 4 + row)
  else
    error("Invalid index")
```

2.




```
For some reason I wrote this algo in haskell last night.
```

```
Tree a = Node a (Tree a, Tree a) | Nil
rtrav [] [] = Nil
rtrav (p:ps) in = let i = find p in
                      lin = take i in
                      rin = drop (i+1) in
                      lpre = take (length lin) ps
                      rpre = drop (length lin) ps
                  in Tree p ((rtrav lpre lin),(rtrav rpre rin))
```

3. *(a)*

Index	Key
0	38
1	14
2	11
3	42
4	
5	
6	
7	7
8	73
9	8
10	22
11	34
12	25

(b) Not too sure about linear probe deletion. This is my best guess: collided keys migrate back towards their correct bucket.

Index	Key
0	38
1	14
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5	
6	
7	7
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10	
11	11
12	25

4. *(a)*

```
ElementType computeS(TreeNode t) {
  if(t == NULL) return -1 /* base case */
  else return 1 + max(computeS(t->left), computeS(t->right))
}
/* Leaves return 0. This is similar logic to the book's height balancing
```

```
for AVL trees. */
```

(b) $\mathcal{O}(N)$

Or more closely: $\mathcal{O}(N + ceil(log(N)) * 2)$ (to account for two extra calls at each leaf).

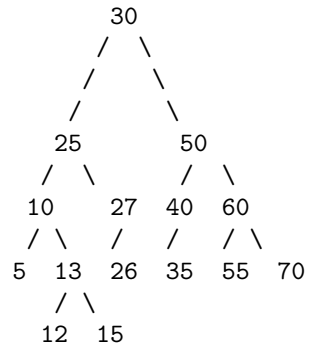
Actually, that's wrong. I guess I'll leave it for the idea. Running out of time!

5. *(a)*

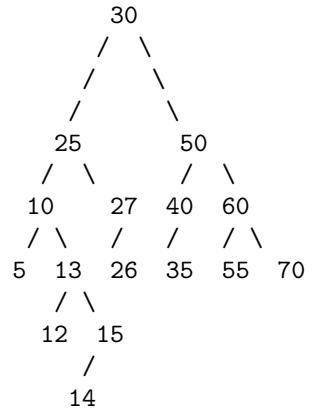
```
I'll do positive = more left
         negative = more right
          +1
         /  \
        /    \
       /      \
     +1        0
    /  \      /  \
  -1   +1   +1   0
  / \  /    /   / \
 0  +1 0   0   0   0
    /
   0
```

(b)

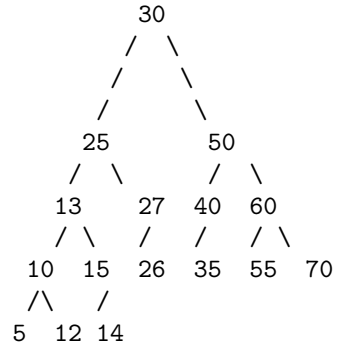
```
Ahhhhh!! Ten minutes left!!!

1. perform insertion

       30
      /  \
     /    \
    /      \
   25       50
  /  \      /  \
 10   27  40  60
 / \  /   /   / \
5  15 26  35  55  70
   /
  12
   \
    13

2. We find as we percolate up recursively that 15 is +2.
   So we do an LR rotation
```



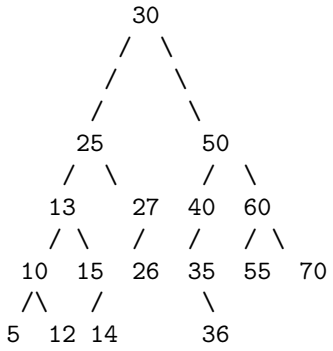
3. We continue to percolate and find that all is balanced.
   Then we insert 14



4. We find that 10 is -2
   We do a RR rotation



5. All is balanced.
   We then insert 36.



6. 40 is imbalanced
   We do an LR

And so on. Ran out of time