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Name:	Student ID:	
Due Date: 8 Nov 2018, 11:59pm.		
Submit answers on eDimension in pdf	format. Submission without student information will No	OT
be marked! Any questions regarding t	ne homework can be directed to the TA through email (c	on-
tact information on eDimension).		
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Week 6

Question 1

- 1. Consider a hash table of size $m=\sqrt{n}$. Then under the uniform hashing assumption, as $n\to\infty$, the number of empty slots tends to:
- a) a constant
- b) zero
- c) infinity
- **2.** If m = cn; c > 0 then as $n \to \infty$, the number of empty slots tends to:
- a) a constant
- b) zero
- c) infinity
- **3.** If $m=n^a; a>1$ then as $n\to\infty$, the number of empty slots tends to:
- a) a constant
- b) zero
- c) infinity

Question 2

Consider a hash table of size 1000, and n = 1500 keys to be hashed uniformly.

- **1.** The expected number of empty slots is about:
- a) 150
- b) 220
- c) 350
- **2.** The expected number of collisions is about:
- a) 100
- b) 230
- c) 720
- **3.** For a random slot, the average number of keys that hash on that slot is:
- a) 0.5
- b) 2
- c) 1.5

Question 3

- **1.** If the loading factor is greater 1, we expect:
- a) all slots to have at least one item
- b) most slots to have at least one item
- c) it can never happen that more than half of the slots are empty
- **2.** If the loading factor is smaller 1, we expect:
- a) most slots to be empty
- b) very few slots to have one or more items
- c) very few slots to have two or more items

Question 4

Given the following adjacency lists:

```
adj(s) = [a, c, d],

adj(a) = [],

adj(c) = [e, b],

adj(b) = [d],

adj(d) = [c],

adj(e) = [s].
```

- **1.** Starting with node *s*, list the visited node order for:
- a) Breadth First Search
- b) Depth First Search

- 2. The adjacency list representation is better in terms of space than the matrix representation for graphs that are not dense (i.e., the number of edges m = O(n), where n is the number of nodes) because
- a) it is simpler to represent the edges from each node as a linked list
- b) an adjacency matrix representation does not explicitly name the set of edges
- c) when the graph is not dense, the number of edges is m = O(n) and the space complexity in bits when using the adjacency list representation is $O(n \log n)$ instead of $O(n^2)$
- d) for non-dense graphs both representations are equivalent since the complexity is proportional to the number of nodes