ICPC REFERENCE NOTEBOOK

Dominicoders

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MACROS

```
#define s(n)
                            scanf("%d",&n)
                            scanf("%c",&n)
#define sc(n)
#define sl(n)
                           scanf("%ld",&n)
                           scanf("%lld",&n)
#define sll(n)
                            scanf("%lf",&n)
#define sf(n)
#define ssp(n)
                            \operatorname{scanf}("\%[^{n}]\%*c",n)
#define prt(x)
                            printf("%d \mid n",x);
#define plt(x)
                            printf("%\underline{11d}\n",x);
#define INF
                            0x3f3f3f3f
#define EPS
                             1e-12
#define bitcount
                             __builtin_popcount
#define gcd
                            __gcd
#define forall(i,a,b)
                             for(int i=a;i<b;i++)
                               for(typeof((c).begin()) v = (c).begin(); v !=
#define foreach(v, c)
(c).end(); ++v)
#define all(a)
                            a.begin(), a.end()
                            a.rbegin(),a.rend()
#define rall(a)
#define in(a,b)
                            ((b).find(a) != (b).end())
#define pb
                           push back
#define fill(a,v)
                            memset(a, v, sizeof a)
#define sz(a)
                            ((int)(a.size()))
                            make_pair
#define mp
                             ((\operatorname{conj}(a)^*(b)).X)
#define dot(a,b)
#define cross(a,b)
                              ((conj(a)*(b)).imag())
#define normalize(v)
                               ((v)/length(v))
#define rotate(p,about,theta)
                                 ((p-about)*exp(point(0,theta))+about)
#define pointEqu(a,b)
                                (comp(a.X,b.X)==0 \&\& comp(a.Y,b.Y)==0)
#define maX(a,b)
                               ((a) > (b) ? (a) : (b))
                               ((a) < (b)?(a):(b))
#define miN(a,b)
#define checkbit(n,b)
                                ((n >> b) & 1)
```

```
x ## y
#define strjoin(x, y)
                             sort(all(a)); a.erase(unique(all(a)),a.end())
#define DREP(a)
#define INDEX(arr.ind)
                               (lower bound(all(arr),ind)-arr,begin())
typedef pair<int, int> ii;
tvpedef vector<ii>vii;
typedef vector<int> vi;
typedef stringstream ss;
typedef vector<string> vs;
typedef vector<double> vd;
typedef vector<vector<int> > vvi;
typedef long long ll;
typedef long double ld;
//g++ abc.cpp -o abc -DDEBUG
#ifdef DEBUG
       #define trace3(x,y,z) cerr<< FUNCTION <<":"<< LINE <<":
"#x" = "<<x<<" | "#y" = "<<y<<" | "#z" = "<<z<endl;
#else
       #define trace3(x,y,z)
#endif
```

FUNCTIONS

FastRead

```
for(; c>47LL && c<58LL ; c = getchar()) {
   x = (x << 1LL) + (x << 3LL) + c - 48LL;
  if(neg) x = -x;
  return x;
Power/GCD/LCM
ll power(ll a, ll b) {
   11 r = 1;
  while(b) {
     if(b & 1) r = r * a;
     a = a * a;
     b >>= 1;
  return r;
ll \gcd(ll a, ll b)
   return b==0?a:gcd(b,a%b);
ll lcm(ll a,ll b){
  return (a/gcd(a,b))*b;
nCr
long long C(int n, int r)
  if(r>n) return 0;
  if(r > n / 2) r = n - r; // because C(n, r) == C(n, n - r)
  long long ans = 1;
  int i:
  for(i = 1; i \le r; i++)
     ans *= n - r + i;
     ans = i;
  return ans;
```

```
Tobinary
int tobinary(bitset<8>x)
        string
       mystring=x.to_string<char,std::string::traits_type,std::string::allocator_typ
        e>();
       return atoi(mystring.c_str());
ab mod T
/* Iterative Function to calculate (x^n)%p in O(logy) */
int power(int x, unsigned int y, int p)
       int res = 1: // Initialize result
       x = x \% p; // Update x if it is more than or equal to p
       while (y > 0)
               // If y is odd, multiply x with result
               if (y & 1)
                        res = (res*x) \% p;
               // y must be even now
               y = y >> 1; // y = y/2
               x = (x*x) \% p;
        return res;
Modulo Arithmetic
#include<iostream>
#include<cstdio>
using namespace std;
int fast_pow(long long base, long long n,long long M)
  if(n==0)
    return 1;
```

```
if(n==1)
  return base:
  long long halfn=fast_pow(base,n/2,M);
  if(n\%2==0)
    return (halfn * halfn ) % M;
  else
    return ( ( ( halfn * halfn ) % M ) * base ) % M;
int findMMI fermat(int n,int M)
       return fast_pow(n,M-2,M);
int main()
  long long fact[100001];
  fact[0]=1;
  int i=1;
  int MOD=1000000007;
  while(i<=100000)
    fact[i]=(fact[i-1]*i)%MOD;
    i++;
  while(1)
    int n,r;
    printf("Enter n: ");
    scanf(" %d",&n);
    printf("Enter r: ");
    scanf(" %d",&r);
    long long numerator,denominator,mmi_denominator,ans;
    numerator=fact[n];
    denominator=(fact[r]*fact[n-r])%MOD;
    mmi denominator=findMMI fermat(denominator,MOD);
    ans=(numerator*mmi denominator)%MOD;
    printf("%lld\n",ans);
  return 0;
```

Base Conversion

```
int a,b; char sa[10000]; char sb[10000];
void rev(char s[]) {
        int l=strlen(s);
        for(int i=0; i<l-1-i; i++) swap(s[i],s[l-1-i]);
void multi(char s[], int k) {
        int i, c=0, d;
        for(i=0;s[i];i++)
                 d=(s[i]-'0')*k+c;
                c=d/b; d\%=b;
                 s[i]='0'+d;
        while(c)
                 s[i]='0'+(c\%b); i++;
                 c/=b:
        s[i]='\setminus 0';
void add(char s[], int k) {
        int i, c=k, d;
        for(i=0;s[i];i++)
                 d=(s[i]-'0')+c;
                 c=d/b; d\%=b;
                 72
                 s[i]='0'+d;
        while(c)
                 s[i]='0'+(c\%b); i++;
                 c/=b:
        s[i]='\setminus 0';
void trans(char s[]) {
        int i;
        for(i=0;s[i];i++)
```

```
char& c=s[i];
               if(c)='A' && c<='Z') c='0'+10+(c-'A');
               if(c)='a' && c<='z') c='0'+36+(c-'a');
void itrans(char s[]) {
       int i;
       for(i=0;s[i];i++)
               char& c=s[i]; int d=c-'0';
               if(d>=10 && d<=35) c='A'+(d-10);
               if(d>=36) c='a'+(d-36);
int main() {
       int q; cin>>q;
       int i,j;
       while(q)
               cin>>a>>b>>sa; sb[0]='0'; sb[1]='\0';
               cout<<a<<" "<<sa<<endl;
               trans(sa);
               for(i=0;sa[i];i++)
                       multi(sb, a);
                       add(sb, sa[i]-'0');
               rev(sb);
               itrans(sb);
               cout<<b<<" "<<sb<<endl;
               cout<<endl;
       return 0;
```

Outline: O(n log n) algorithm for The Longest Increasing Subseq

```
set<int> st;
set<int>::iterator it;
st.clear();
for(i=0; i<n; i++)
        st.insert(a[i]); it=st.find(a[i]);
        it++; if(it!=st.end()) st.erase(it);
cout<<st.size()<<endl;
Outline: O(nm) algorithm for the LCS with O(n) sapce
                int m[2][1000]; // instead of [1000][1000]
for(i=M; i>=0; i--)
        ii = i\&1:
        for(j=N; j>=0; j--)
                if(i==M \parallel j==N) { m[ii][j]=0; continue; }
                if(s1[i]==s2[j]) m[ii][j] = 1+m[1-ii][j+1];
                else m[ii][j] = max(m[ii][j+1], m[1-ii][j]);
cout << m[0][0]; // if you want m[x][y], write m[x&1][y];
Manacher's Algo O(2n)
// Transform S into T.
// For example, S = "abba", T = "^#a#b#b#a#$".
// ^ and $ signs are sentinels appended to each end to avoid bounds checking
string preProcess(string s) {
 int n = s.length();
 if (n == 0) return "^{"}";
 string ret = ^{"}^{"};
 for (int i = 0; i < n; i++)
  ret += "#" + s.substr(i, 1);
 ret += "#$";
 return ret;}
```

```
string longestPalindrome(string s) {
 string T = preProcess(s);
 int n = T.length();
 int P = \text{new int}[n];
 int C = 0, R = 0;
 for (int i = 1; i < n-1; i++) {
  int i mirror = 2*C-i; // equals to i' = C - (i-C)
  P[i] = (R > i) ? min(R-i, P[i mirror]) : 0;
  // Attempt to expand palindrome centered at i
  while (T[i + 1 + P[i]] == T[i - 1 - P[i]])
   P[i]++;
  // If palindrome centered at i expand past R,
  // adjust center based on expanded palindrome.
  if (i + P[i] > R) {
   C = i:
   R = i + P[i];
 // Find the maximum element in P.
 int maxLen = 0;
 int centerIndex = 0:
 for (int i = 1; i < n-1; i++) {
  if (P[i] > maxLen) {
   maxLen = P[i];
   centerIndex = i;
 delete[] P:
 return s.substr((centerIndex - 1 - maxLen)/2, maxLen);
```

KMP O(n)

```
#include <cstdio>
#include <cstring>
#include <time.h>
using namespace std;
#define MAX N 100010
char T[MAX N], P[MAX N]; // T = text, P = pattern
int b[MAX N], n, m; // b = back table, n = length of T, m = length of P
void naiveMatching() {
 for (int i = 0; i < n; i++) { // try all potential starting indices
  bool found = true:
  for (int j = 0; j < m && found; <math>j++) // use boolean flag `found'
   if (i + j \ge n \parallel P[i] != T[i + j]) // if mismatch found
     found = false; // abort this, shift starting index i by +1
  if (found) // if P[0 ... m - 1] == T[i ... i + m - 1]
   printf("P is found at index %d in T\n", i);
} }
void kmpPreprocess() { // call this before calling kmpSearch()
 int i = 0, j = -1; b[0] = -1; // starting values
 while (i < m) { // pre-process the pattern string P
  while (j \ge 0 \&\& P[i] != P[j]) j = b[j]; // if different, reset j using b
  i++; j++; // if same, advance both pointers
  b[i] = i; // observe i = 8, 9, 10, 11, 12 with i = 0, 1, 2, 3, 4
          // in the example of P = "SEVENTY SEVEN" above
void kmpSearch() { // this is similar as kmpPreprocess(), but on string T
 int i = 0, j = 0; // starting values
 while (i < n) { // search through string T
  while (i \ge 0 \&\& T[i] != P[i]) i = b[i]; // if different, reset i using b
  i++; j++; // if same, advance both pointers
  if (i == m) { // a match found when i == m
   printf("P is found at index %d in T \setminus n", i - j);
   i = b[i]; // prepare i for the next possible match
} } }
int main() {
```

```
strcpv(T, "I DO NOT LIKE SEVENTY SEV BUT SEVENTY SEVENTY
SEVEN");
 strcpy(P, "SEVENTY SEVEN");
 n = (int)strlen(T);
 m = (int)strlen(P);
 //if the end of line character is read too, uncomment the line below
 //T[n-1] = 0; n--; P[m-1] = 0; m--;
 printf("T = \frac{\% s}{n}", T);
 printf("P = \frac{1}{8} \frac{n}{n}, P);
 naiveMatching();
 printf("KMP\n");
 kmpPreprocess();
 kmpSearch();
 printf("String Library\n");
 char *pos = strstr(T, P);
 while (pos != NULL) {
  printf("P is found at index %d in T\n", pos - T);
  pos = strstr(pos + 1, P);
 return 0;
Subset Sum Problem O(sum*n)
// Returns true if there is a subset of set[] with sun equal to given sum
bool isSubsetSum(int set[], int n, int sum)
  // The value of subset[i][j] will be true if there is a
  // subset of set[0..j-1] with sum equal to i
  bool subset[sum+1][n+1];
  // If sum is 0, then answer is true
  for (int i = 0; i <= n; i++)
   subset[0][i] = true;
  // If sum is not 0 and set is empty, then answer is false
  for (int i = 1; i \le sum; i++)
   subset[i][0] = false;
```

```
// Fill the subset table in bottom up manner
   for (int i = 1; i \le sum; i++)
    for (int i = 1; i \le n; i++)
     subset[i][j] = subset[i][j-1];
     if (i \ge set[j-1])
       subset[i][j] = subset[i][j] \parallel
                     subset[i - set[i-1]][i-1];
  /* // uncomment this code to print table
   for (int i = 0; i \le sum; i++)
    for (int j = 0; j \le n; j++)
      printf ("%4d", subset[i][j]);
    printf("\n");
   } */
   return subset[sum][n];
int main()
 int set[] = \{3, 34, 4, 12, 5, 2\};
 int sum = 9:
 int n = sizeof(set)/sizeof(set[0]);
 if (isSubsetSum(set, n, sum) == true)
   printf("Found a subset with given sum");
   printf("No subset with given sum");
 return 0;
```

```
Prime Sieve of Eratosthenes
typedef long long ll;
typedef vector<int> vi;
typedef map<int, int> mii;
11 _sieve_size;
bitset<10000010> bs; // 10^7 should be enough for most cases
vi primes; // compact list of primes in form of vector<int>
// first part
void sieve(ll upperbound) {
                                  // create list of primes in [0..upperbound]
                                            // add 1 to include upperbound
 _sieve_size = upperbound + 1;
 bs.set();
                                        // set all bits to 1
 bs[0] = bs[1] = 0;
                                          // except index 0 and 1
 for (ll i = 2; i \le sieve\_size; i++) if (bs[i]) {
  // cross out multiples of i starting from i * i!
  for (ll j = i * i; j \le sieve_size; j += i) bs[j] = 0;
  primes.push back((int)i); // also add this vector containing list of primes
                              // call this method in main method
bool isPrime(11 N) {
                               // a good enough deterministic prime tester
 if (N <= _sieve_size) return bs[N];</pre>
                                                 // O(1) for small primes
 for (int i = 0; i < (int)primes.size(); i++)
  if (N % primes[i] == 0) return false;
                          // it takes longer time if N is a large prime!
 return true:
               // note: only work for N \le (last prime in vi "primes")^2
// second part
vi primeFactors(ll N) { // remember: vi is vector of integers, <u>ll</u> is long long
                       // vi `primes' (generated by sieve) is optional
 vi factors:
 11 PF idx = 0, PF = primes[PF idx]; // using PF = 2, 3, 4, ..., is also ok
 while (N != 1 \&\& (PF * PF \le N)) { // stop at sqrt(N), but N can get smaller
  while (N % PF == 0) { N /= PF; factors.push back(PF); } // remove this PF
                                              // only consider primes!
  PF = primes[++PF_idx];
 if (N!=1) factors.push_back(N); // special case if N is actually a prime
```

```
// if pf exceeds 32-bit integer, you have to change vi
 return factors;
// third part
11 numPF(11 N) {
11 \text{ PF\_idx} = 0, \text{ PF} = \text{primes}[\text{PF\_idx}], \text{ ans} = 0;
 while (N != 1 \&\& (PF * PF <= N))  {
  while (N \% PF == 0) \{ N /= PF; ans++; \}
  PF = primes[++PF idx];
 if (N != 1) ans++;
 return ans;
ll numDiffPF(ll N) {
If PF idx = 0, PF = primes[PF idx], ans = 0;
 while (N != 1 && (PF * PF <= N)) {
  if (N % PF == 0) ans++;
                                                // count this pf only once
  while (N % PF == 0) N \neq PF;
  PF = primes[++PF_idx];
 if (N != 1) ans++;
 return ans;
11 sumPF(11 N) {
 11 \text{ PF\_idx} = 0, \text{ PF} = \text{primes}[\text{PF\_idx}], \text{ ans} = 0;
 while (N != 1 && (PF * PF <= N)) {
  while (N % PF == 0) { N \neq PF; ans += PF; }
  PF = primes[++PF_idx];
 if (N != 1) ans += N;
 return ans:
11 numDiv(11 N) {
11 \text{ PF\_idx} = 0, \text{ PF} = \text{primes}[\text{PF\_idx}], \text{ ans } = 1;
                                                           // start from ans = 1
 while (N != 1 && (PF * PF <= N)) {
  11 power = 0;
                                                // count the power
```

```
while (N % PF == 0) { N /= PF; power++; }
  ans *= (power + 1);
                                          // according to the formula
  PF = primes[++PF idx];
                              // (last factor has pow = 1, we add 1 to it)
 if (N != 1) ans *= 2:
 return ans;
11 sumDiv(11 N) {
 ll PF idx = 0, PF = primes[PF idx], ans = 1;
                                                        // start from ans = 1
 while (N != 1 && (PF * PF <= N)) {
  11 power = 0;
  while (N % PF == 0) { N \neq PF; power++; }
  ans *= ((11)pow((double)PF, power + 1.0) - 1) / (PF - 1);
                                                                   // formula
  PF = primes[++PF_idx];
 if (N != 1) ans *= ((11)pow((double)N, 2.0) - 1) / (N - 1);
                                                                 // last one
 return ans:
11 EulerPhi(11 N) {
 11 \text{ PF\_idx} = 0, \text{ PF} = \text{primes}[\text{PF\_idx}], \text{ ans} = \text{N};
                                                        // start from ans = N
 while (N != 1 && (PF * PF <= N)) {
  if (N % PF == 0) ans -= ans / PF;
                                               // only count unique factor
  while (N % PF == 0) N \neq PF;
  PF = primes[++PF_idx];
 if (N != 1) ans -= ans / N;
                                                   // last factor
 return ans;
int main() {
 // first part: the Sieve of Eratosthenes
 sieve(10000000);
                                  // can go up to 10<sup>^{^{^{\prime}}}</sup>7 (need few seconds)
 printf("%d\n", isPrime(2147483647));
                                                        // 10-digits prime
                                                   // not a prime, 104729*1299709
 printf("%d\n", isPrime(136117223861LL));
 // second part: prime factors
 vi res = primeFactors(2147483647); // slowest, 2147483647 is a prime
 for (vi::iterator i = res.begin(); i != res.end(); i++) printf("> % d\n", *i);
```

```
res = primeFactors(136117223861LL); // slow, 2 large pfactors
104729*1299709
 for (vi::iterator i = res.begin(); i != res.end(); i++) printf("# \%d\n", *i);
 res = primeFactors(142391208960LL); // faster, 2^10*3^4*5*7^4*11*13
 for (vi::iterator i = res.begin(); i != res.end(); i++) printf("! %d\n", *i);
 //res = primeFactors((11)(1010189899 * 1010189899)); // "error"
 //for (vi::iterator i = res.begin(); i != res.end(); i++) printf("^{\wedge} %d^{\circ}, *i);
 // third part: prime factors variants
 printf("numPF(%d) = %lld\n", 50, numPF(50)); // 2^1 * 5^2 => 3
 printf("numDiffPF(%d) = %lld\n", 50, numDiffPF(50)); // 2^1 * 5^2 => 2
 printf("sumPF(%d) = %lld\n", 50, sumPF(50)); // 2^1 * 5^2 => 2 + 5 + 5 = 12
 printf("numDiv(%d) = %lld\n", 50, numDiv(50)); // 1, 2, 5, 10, 25, 50, 6 divisors
 printf("sumDiv(%d) = \%11d\n", 50, sumDiv(50)); // 1 + 2 + 5 + 10 + 25 + 50 = 93
 printf("EulerPhi(%d) = \%lld\n", 50, EulerPhi(50)); // 20 integers < 50 are
relatively prime with 50
 return 0;
Segment Tree
typedef vector<int> vi;
class SegmentTree {
                          // the segment tree is stored like a heap array
                        // recall that vi is: typedef vector<int> vi;
private: vi st, A;
 int n;
 int left (int p) { return p << 1; } // same as binary heap operations</pre>
 int right(int p) { return (p \ll 1) + 1; }
                                                  // O(n \log n)
 void build(int p, int L, int R) {
                              // as L == R, either one is fine
  if (L == R)
                                       // store the index
   st[p] = L;
                            // recursively compute the values
  else {
   build(left(p), L
                            (L + R) / 2;
   build(right(p), (L + R) / 2 + 1, R
   int p1 = st[left(p)], p2 = st[right(p)];
   st[p] = (A[p1] \le A[p2]) ? p1 : p2;
```

```
} }
int rmq(int p, int L, int R, int i, int j) {
                                                    // O(\log n)
 if (i > R || i < L) return -1; // current segment outside query range
 if (L >= i \&\& R <= j) return st[p];
                                               // inside query range
 // compute the min position in the left and right part of the interval
 int p1 = rmq(left(p), L
                                  , (L+R) / 2, i, j);
 int p2 = rmq(right(p), (L+R) / 2 + 1, R
 if (p1 == -1) return p2; // if we try to access segment outside query
                                             // same as above
 if (p2 == -1) return p1;
 return (A[p1] \le A[p2])? p1 : p2;
                                             // as as in build routine
int update_point(int p, int L, int R, int idx, int new_value) {
 // this update code is still preliminary, i == i
 // must be able to update range in the future!
 int i = idx, j = idx;
 // if the current interval does not intersect
 // the update interval, return this st node value!
 if (i > R \parallel j < L)
  return st[p];
 // if the current interval is included in the update range,
 // update that st[node]
 if (L == i \&\& R == i)  {
  A[i] = new value; // update the underlying array
  return st[p] = L; // this index
 // compute the minimum pition in the
 // left and right part of the interval
 int p1, p2;
 p1 = update point(left(p), L
                                       (L + R) / 2, idx, new value);
 p2 = update point(right(p), (L + R) / 2 + 1, R)
                                                       , idx, new value);
 // return the pition where the overall minimum is
 return st[p] = (A[p1] \le A[p2]) ? p1 : p2;
```

```
public:
 SegmentTree(const vi & A) {
  A = A; n = (int)A.size(); // copy content for local usage
  st.assign(4 * n, 0); // create large enough vector of zeroes
                            // recursive build
  build(1, 0, n - 1):
 int rmq(int i, int j) { return rmq(1, 0, n - 1, i, j); } // overloading
 int update point(int idx, int new value) {
  return update point(1, 0, n - 1, idx, new_value); }
};
int main() {
 int arr[] = { 18, 17, 13, 19, 15, 11, 20 };
                                              // the original array
 vi A(arr, arr + 7);
                         // copy the contents to a vector
 SegmentTree st(A);
                 idx 0, 1, 2, 3, 4, 5, 6\n");
 printf("
                 A is \{18,17,13,19,15,11,20\}\n");
 printf("
 printf("RMQ(1, 3) = \% d n", st.rmq(1, 3)); // answer = index 2
 printf("RMQ(4, 6) = %d\n", st.rmq(4, 6)); // answer = index 5
 printf("RMQ(3, 4) = \%d\n", st.rmq(3, 4)); // answer = index 4
 printf("RMQ(0, 0) = %d\n", st.rmq(0, 0)); // answer = index 0
 printf("RMO(0, 1) = \%d\n", st.rmg(0, 1)); // answer = index 1
 printf("RMQ(0, 6) = %d\n", st.rmq(0, 6)); // answer = index 5
 printf("
                 idx 0, 1, 2, 3, 4, 5, 6\n");
 printf("Now, modify A into {18,17,13,19,15,100,20}\n");
 st.update point(5, 100);
                               // update A[5] from 11 to 100
 printf("These values do not change\n");
 printf("RMQ(1, 3) = \% d n", st.rmq(1, 3));
                                                       // 2
 printf("RMQ(3, 4) = \% d n", st.rmq(3, 4));
                                                       // 4
 printf("RMQ(0, 0) = \% d n", st.rmq(0, 0));
                                                       // 0
 printf("RMQ(0, 1) = \% d n", st.rmq(0, 1));
                                                       // 1
 printf("These values change\n");
 printf("RMQ(0, 6) = \% d n", st.rmq(0, 6));
                                                     //5 -> 2
 printf("RMQ(4, 6) = \% d n", st.rmq(4, 6));
                                                     // 5->4
  printf("RMQ(4, 5) = \%d\n", st.rmq(4, 5));
                                                      // 5->4
 return 0;}
```

Fenwick Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
        while(x \le N) {
                tree[x] += v;
                x += (x \& -x);
// get cumulative sum up to and including x
int get(int x) {
        int res = 0;
        while(x) {
                res += tree[x];
                x = (x \& -x);
        return res;
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
        int idx = 0, mask = N;
        while(mask && idx < N) {
                int t = idx + mask:
                if(x \ge tree[t]) 
                        idx = t;
                        x = tree[t];
                mask >>= 1;
        return idx:
```

String Trie

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <stdbool.h>
#define ARRAY_SIZE(a) sizeof(a)/sizeof(a[0])
// Alphabet size (# of symbols)
#define ALPHABET_SIZE (26)
// Converts key current character into index
// use only 'a' through 'z' and lower case
#define CHAR TO INDEX(c) ((int)c - (int)'a')
// trie node
struct TrieNode
       struct TrieNode *children[ALPHABET_SIZE];
       // isLeaf is true if the node represents
       // end of a word
       bool isLeaf:
};
// Returns new trie node (initialized to NULLs)
struct TrieNode *getNode(void)
       struct TrieNode *pNode = NULL;
        pNode = (struct TrieNode *)malloc(sizeof(struct TrieNode));
       if (pNode)
                int i;
               pNode->isLeaf = false;
               for (i = 0; i < ALPHABET_SIZE; i++)
                       pNode->children[i] = NULL;
       return pNode;
// If not present, inserts key into trie
// If the key is prefix of trie node, just marks leaf node
void insert(struct TrieNode *root, const char *key)
        int level;
```

```
int length = strlen(key);
        int index:
        struct TrieNode *pCrawl = root;
        for (level = 0; level < length; level++)
                index = CHAR_TO_INDEX(key[level]);
                if (!pCrawl->children[index])
                        pCrawl->children[index] = getNode();
                pCrawl = pCrawl->children[index];
        // mark last node as leaf
        pCrawl->isLeaf = true;
// Returns true if key presents in trie, else false
bool search(struct TrieNode *root, const char *key)
        int level:
        int length = strlen(key);
        int index:
        struct TrieNode *pCrawl = root;
        for (level = 0; level < length; level++)
                index = CHAR_TO_INDEX(key[level]);
                if (!pCrawl->children[index])
                        return false:
                pCrawl = pCrawl->children[index];
        return (pCrawl != NULL && pCrawl->isLeaf);
int main()
        // Input keys (use only 'a' through 'z' and lower case)
        char keys[][8] = {"the", "a", "there", "answer", "any",
                                        "by", "bye", "their"};
        char output[][32] = {"Not present in trie", "Present in trie"};
        struct TrieNode *root = getNode();
        // Construct trie
        int i;
```

```
for (i = 0; i < ARRAY\_SIZE(keys); i++)
                insert(root, keys[i]);
       // Search for different keys
printf("%s --- %s\n", "the", output[search(root, "the")]);
printf("%s --- %s\n", "these", output[search(root, "these")]);
printf("%s --- %s\n", "their", output[search(root, "their")]);
printf("%s --- %s\n", "thaw", output[search(root, "thaw")]);
return 0;
Dijkstra's Algo
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
// Running time: O(|E| \log |V|)
const int INF = 20000000000;
typedef pair<int, int> PII;
int main() {
        int N, s, t;
        scanf("%d%d%d", &N, &s, &t);
        vector<vector<PII>> edges(N);
        for (int i = 0; i < N; i++) {
                int M;
                scanf("%d", &M);
                for (int j = 0; j < M; j++) {
                        int vertex, dist;
                        scanf("%d%d", &vertex, &dist);
                        edges[i].push_back(make_pair(dist, vertex)); // note order
of arguments here
        // use priority queue in which top element has the "smallest" priority
        priority_queue<PII, vector<PII>, greater<PII> > Q;
        vector<int> dist(N, INF), dad(N, -1);
        Q.push(make pair(0, s));
        dist[s] = 0;
        while (!Q.empty()) {
                PII p = Q.top();
                Q.pop();
                int here = p.second;
                if (here == t) break;
```

```
if (dist[here] != p.first) continue;
                for (vector<PII>::iterator it = edges[here].begin(); it !=
edges[here].end(); it++)
        if (dist[here] + it->first < dist[it->second])
                 dist[it->second] = dist[here] + it->first;
                dad[it->second] = here;
                Q.push(make pair(dist[it->second], it->second));
        printf("%d\n", dist[t]);
        if (dist[t] < INF)
                for (int i = t; i != -1; i = dad[i])
                         printf("%d%c", i, (i == s ? • f \setminus n \cdot f : • f \cdot f);
        return 0;
Sample input:
504
21231
22445
3 1 4 3 3 4 1
20123
21521
Expected:
5
4230
Prim & Kruskal MST
typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector<ii>vii;
// Union-Find Disjoint Sets Library written in OOP manner, using both path
```

// OOP style

compression and union by rank heuristics

class UnionFind {

```
private:
 vi p, rank, setSize;
                                   // remember: vi is vector<int>
 int numSets:
public:
 UnionFind(int N) {
  setSize.assign(N, 1); numSets = N; rank.assign(N, 0);
  p.assign(N, 0); for (int i = 0; i < N; i++) p[i] = i;
 int findSet(int i) { return (p[i] == i) ? i : (p[i] = findSet(p[i])); }
 bool isSameSet(int i, int j) { return findSet(i) == findSet(j); }
 void unionSet(int i, int j) {
  if (!isSameSet(i, j)) { numSets--;
  int x = findSet(i), y = findSet(j);
  // rank is used to keep the tree short
  if (rank[x] > rank[y]) \{ p[y] = x; setSize[x] += setSize[y]; \}
                  \{p[x] = y; setSize[y] += setSize[x];
  else
                  if (rank[x] == rank[y]) rank[y]++; } } }
 int numDisjointSets() { return numSets; }
 int sizeOfSet(int i) { return setSize[findSet(i)]; }
};
vector<vii>> AdjList;
                              // global boolean flag to avoid cycle
vi taken;
                              // priority queue to help choose shorter edges
priority_queue<ii>> pq;
void process(int vtx) { // so, we use -ve sign to reverse the sort order
 taken[vtx] = 1;
 for (int j = 0; j < (int)AdjList[vtx].size(); <math>j++) {
  ii v = AdjList[vtx][i];
  if (!taken[v.first]) pq.push(ii(-v.second, -v.first));
                       // sort by (inc) weight then by (inc) id
} }
int main() {
 int V, E, u, v, w;
 // Graph in Figure 4.10 left, format: list of weighted edges
 // This example shows another form of reading graph input
 57
 0 1 4
 024
 036
 046
```

```
1 2 2
 238
 349
 freopen("in 03.txt", "r", stdin);
 scanf("%d %d", &V, &E);
 // Kruskal's algorithm merged with Prim's algorithm
 AdjList.assign(V, vii());
 vector< pair<int, ii>> EdgeList; // (weight, two vertices) of the edge
 for (int i = 0; i < E; i++) {
  scanf("%d %d %d", &u, &v, &w); // read the triple: (u, v, w)
  EdgeList.push back(make pair(w, ii(u, v))); // (w, u, v)
  AdjList[u].push_back(ii(v, w));
  AdjList[v].push_back(ii(u, w));
 sort(EdgeList.begin(), EdgeList.end()); // sort by edge weight O(E log E)
        // note: pair object has built-in comparison function
 int mst cost = 0;
 UnionFind UF(V);
                                // all V are disjoint sets initially
 for (int i = 0; i < E; i++) {
                                   // for each edge, O(E)
  pair<int, ii> front = EdgeList[i];
  if (!UF.isSameSet(front.second.first, front.second.second)) { // check
   mst_cost += front.first; // add the weight of e to MST
   UF.unionSet(front.second.first, front.second.second); // link them
                  // note: the runtime cost of UFDS is very light
 // note: the number of disjoint sets must eventually be 1 for a valid MST
 printf("MST cost = %d (Kruskal's)\n", mst cost);
// inside int main() --- assume the graph is stored in AdjList, pq is empty
                             // no vertex is taken at the beginning
 taken.assign(V, 0);
 process(0); // take vertex 0 and process all edges incident to vertex 0
 mst cost = 0:
 while (!pq.empty()) { // repeat until V vertices (E=V-1 edges) are taken
  ii front = pq.top(); pq.pop();
  u = -front.second, w = -front.first; // negate the id and weight again
  if (!taken[u]) // we have not connected this vertex yet
   mst_cost += w, process(u); // take u, process all edges incident to u
                     // each edge is in pq only once!
```

```
printf("MST cost = %d (Prim's)\n", mst_cost);
 return 0;
DFS Depth First Search
class Graph
  int V: // No. of vertices
  list<int> *adj;// Pointer to an array containing adj lists
  void DFSUtil(int v, bool visited[]); // A func used by DFS
public:
  Graph(int V); // Constructor
  void addEdge(int v, int w);// function to add an edge to graph
  void DFS(); // prints DFS traversal of the complete graph
};
Graph::Graph(int V)
  this->V = V:
  adj = new list<int>[V];
void Graph::addEdge(int v, int w)
  adj[v].push_back(w); // Add w to v's list.
void Graph::DFSUtil(int v, bool visited[])
  visited[v] = true;
  cout << v << " ";
  // Recur for all the vertices adjacent to this vertex
  list<int>::iterator i;
  for(i = adj[v].begin(); i != adj[v].end(); ++i)
     if(!visited[*i])
       DFSUtil(*i, visited);
// The function to do DFS traversal. It uses recursive DFSUtil()
void Graph::DFS()
```

```
// Mark all the vertices as not visited
  bool *visited = new bool[V];
  for (int i = 0; i < V; i++)
     visited[i] = false;
  // Call the recursive helper function to print DFS traversal
  // starting from all vertices one by one
  for (int i = 0; i < V; i++)
     if (visited[i] == false)
       DFSUtil(i, visited);
int main()
  Graph g(4);
  g.addEdge(0, 1);
  g.addEdge(0, 2);
  g.addEdge(1, 2);
  g.addEdge(2, 0);
  g.addEdge(2, 3);
  g.addEdge(3, 3);
  cout << "Following is Depth First Traversal\n";</pre>
  g.DFS();
  return 0;
BFS Breadth First Search
int V, E, a, b, s;
vector<vii> AdjList;
                // addition: the predecessor/parent vector
vi p;
void printPath(int u) { // simple function to extract information from `vi p'
 if (u == s) { printf("%d", u); return; }
 printPath(p[u]); // recursive call: to make the output format: s -> ... -> t
 printf(" %d", u); }
int main() {
```

```
// format: list of unweighted edges
 // This example shows another form of reading graph input
 13 16
 01 12 23 04 15 26 37 56
 48 89 510 611 712 910 1011 1112
 freopen("in_04.txt", "r", stdin);
 scanf("%d %d", &V, &E);
 AdjList.assign(V, vii()); // assign blank vectors of pair<int, int>s to AdjList
 for (int i = 0; i < E; i++) {
  scanf("%d %d", &a, &b);
  AdjList[a].push back(ii(b, 0));
  AdjList[b].push_back(ii(a, 0));
 // as an example, we start from this source
  s = 5;
// BFS routine
// inside int main() -- we do not use recursion, thus we do not need to create
separate function!
 vi dist(V, 1000000000); dist[s] = 0; // distance to source is 0 (default)
 queue<int> q; q.push(s);
                                   // start from source
 p.assign(V, -1); // to store parent information (p must be a global variable!)
 int layer = -1; // for our output printing purpose addition of
 bool isBipartite = true; //one boolean flag, initially true
 while (!q.empty()) {
  int u = q.front(); q.pop();
  if (dist[u] != layer) printf("\nLayer %d: ", dist[u]);
  layer = dist[u];
  printf("visit %d, ", u);
  for (int j = 0; j < (int)AdjList[u].size(); <math>j++) {
   ii v = AdjList[u][i];
                           // for each neighbors of u
   if (dist[v.first] == 10000000000) {
    dist[v.first] = dist[u] + 1; // v unvisited + reachable
     p[v.first] = u; // addition: the parent of vertex v->first is u
    q.push(v.first);
                             // enqueue v for next step
   else if ((dist[v.first] \% 2) == (dist[u] \% 2))// same parity
    isBipartite = false;
```

```
} }
 printf("\nShortest path: ");
 printPath(7), printf("\n");
 printf("isBipartite? %d\n", isBipartite);
 return 0;
Lowest Common Ancestor LCA
const int max nodes, log max nodes;
int num nodes, log num nodes, root;
vector<int> children[max nodes]; // children[i] contains the children of node i
int A[max nodes][log max nodes+1]; // A[i][i] is the 2<sup>i</sup>-th ancestor of node i, or
-1 if that
ancestor does not exist
int L[max_nodes]; // L[i] is the distance between node i and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
        if(n==0)
                 return -1;
         int p = 0;
        if (n >= 1 << 16) \{ n >= 16; p += 16; \}
        if (n >= 1 << 8) \{ n >>= 8; p += 8; \}
        if (n >= 1 << 4) \{ n >>= 4; p += 4; \}
        if (n >= 1 << 2) \{ n >>= 2; p += 2; \}
        if (n >= 1 << 1) \{ p += 1; \}
        return p;
void DFS(int i, int l)
        L[i] = 1;
        for(int j = 0; j < children[i].size(); <math>j++)
                 DFS(children[i][j], l+1);
int LCA(int p, int q)
        // ensure node p is at least as deep as node q
        if(\underline{L}[p] < \underline{L}[q])
```

```
swap(p, q);
        // "binary search" for the ancestor of node p situated on the same level as q
        for(int i = log num nodes; i >= 0; i--)
                if(L[p] - (1 << i) >= L[q])
                        p = A[p][i];
        if(p == q)
                return p;
       // "binary search" for the LCA
        for(int i = log num nodes; i >= 0; i--)
                if(A[p][i] != -1 && A[p][i] != A[q][i])
                        p = A[p][i];
                        q = A[q][i];
        return A[p][0];
int main(int argc,char* argv[])
        // read num nodes, the total number of nodes
       log_num_nodes=lb(num_nodes);
        for(int i = 0; i < num nodes; i++)
                int p;
        // read p, the parent of node i or -1 if node i is the root
                A[i][0] = p;
                if(p != -1)
                        children[p].push back(i);
                else
                        root = i;
        // precompute A using dynamic programming
        for(int j = 1; j <= log_num_nodes; j++)
                for(int i = 0; i < num nodes; i++)
                        if(A[i][i-1]!=-1)
                                A[i][i] = A[A[i][i-1]][i-1];
                        else
                                A[i][j] = -1;
        // precompute L
        DFS(root, 0);
        return 0;
```

Tricks in cmath

```
// when the number is too large. use powl instead of pow.
// will provide you more accuracy.
powl(a, b)
(int)round(p, (1.0/n)) // nth root of p
```

Initialize array with predefined value

```
// for 1d array, use STL fill_n or fill to
initialize array fill(a, a+size_of_a,
value)
fill_n(a, size_of_a, value)
//for 2d array, if want to fill in 0 or -1
memset(a, 0, sizeof(a));
// otherwise, use a loop of fill or fill_n
through every a[i] fill(a[i], a[i]+size_of_ai,
value) // from 0 to number of row.
```

<u>Java</u>

StringBuilder Use of Functions

```
import java.lang.StringBuilder;
public class String_builder {
        public static void main(String[] args) {
                // Create a new StringBuilder.
                StringBuilder builder1 = new StringBuilder();
                // Loop and append values.
                for (int i = 0; i < 5; i++) {
                        builder1.append("abc ");
                // Convert to string.
                String result = builder1.toString();
                System.out.println(result);
                // " INSERT"
                StringBuilder builder2 = new StringBuilder("abc");
                // Insert this substring at position 2.
                builder2.insert(2, "xyz");
                System.out.println(builder2);//abxyzc
                // INDEX-OF
                StringBuilder builder3 = new StringBuilder("abc");
                // Try to find this substring.
                int result1 = builder3.indexOf("bc");
                System.out.println(result1);// 1
                // This substring does not exist.
                int result2 = builder3.indexOf("de");
                System.out.println(result2);// -1
                // DELETE
```

```
StringBuilder builder4 = new StringBuilder("carrot");
               // Delete characters from index 2 to index 5.
               builder4.delete(2, 5);
               System.out.println(builder4);// cat
                // REPLACE
               StringBuilder b = new StringBuilder("abc");
               // Replace second character with "xyz".
               b.replace(1, 2, "xyz");
               System.out.println(b);// axyzc
                // SUBSTRING
               StringBuilder builder = new StringBuilder();
               builder.append("Forest");
               String firstTwo = builder.substring(0, 2);
               System.out.println(firstTwo);// Fo
                // REVERSE
               StringBuilder builder5 = new StringBuilder();
               builder5.append("abc");
               builder5.reverse();
               System.out.println(builder5);
Decimal Formatter
import java.util.Scanner;
import java.util.Formatter;
public class readdouble {
       public static void main(String args[]){
       Scanner ob=new Scanner(System.in);
       float s = ob.nextFloat();
       Formatter fmt = new Formatter();
  fmt = new Formatter();
  fmt.format("%2.3f",s);
  System.out.println(fmt);}}
Fast Read
import java.io.BufferedReader;
import java.io.IOException;
import java.io.InputStreamReader;
```

```
import java.util.Scanner;
import java.util.StringTokenizer;
import java.lang.*;
public class fastread {
       static class FastReader {
               BufferedReader br;
                StringTokenizer st;
                public FastReader() {
                        br = new BufferedReader(new
InputStreamReader(System.in));
                String next() {
                        while (st == null || !st.hasMoreElements()) {
                                try {
                                        st = new StringTokenizer(br.readLine());
                                } catch (IOException e) {
                                        e.printStackTrace();
                       return st.nextToken();
               int nextInt() {
                       return Integer.parseInt(next());
               long nextLong() {
                        return Long.parseLong(next());
                double nextDouble() {
                        return Double.parseDouble(next());
                String nextLine() {
                        String str = "";
                        try {
                                str = br.readLine();
                        } catch (IOException e) {
                               e.printStackTrace();
                        return str;
```

```
String str = "123456789";
        public static void main(String[] args) {
                                                                                              BigInteger \underline{C} = A.add(new BigInteger(str));
                FastReader s = new FastReader():
                                                                                              int val = 123456789:
                int t = s.nextInt();
                                                                                              BigInteger C = A.add(BigIntger.valueOf(val));
                while (t-->0) {
                                                                                              //Extraction of value from BigInteger:
                                                                                              int x = A.intValue();// value should be in limit of int x
        }}}
Factorial Using BigInteger
                                                                                              long y = A.longValue(); // value should be in limit of long y
                                                                                              String z = A.toString();
        public static void calculateFactorial(int n) {
                                                                                              // Comparison
                BigInteger result = BigInteger.ONE;
                                                                                              if (a < b) \{ \}
                                                                                                                         // For primitive int
                for (int i=1; i<=n; i++) {
                                                                                              if (A.compareTo(B) < 0) {} // For BigInteger
                        result = result.multiply(BigInteger.valueOf(i));
                                                                                              // Equality
                                                                                              if (A.equals(B)) {} // A is equal to B }}
                System.out.println(n + "! = " + result); }}
                                                                                         Check Prime for BigIntegers
BigInteger Functions
                                                                                         import java.util.*;
                                                                                         import java.math.*;
import java.math.BigInteger;
                                                                                         class CheckPrimeTest
public class BigIntegerDemo {
public static void main(String[] args) {
                                                                                            //Function to check and return prime numbers
                                                                                            static boolean checkPrime(long n)
    BigInteger b1 = new BigInteger("987654321987654321000000000");
    BigInteger b2 = new BigInteger("987654321987654321000000000");
                                                                                              // Converting long to BigInteger
                                                                                              BigInteger b = new BigInteger(String.valueOf(n));
    BigInteger product = b1.multiply(b2);
                                                                                              return b.isProbablePrime(1);
    BigInteger division = b1.divide(b2);
                                                                                              // returns whether prime or not
    System.out.println("product = " + product);
                                                                                              return Long.parseLong(b.nextProbablePrime().toString());
    System.out.println("division = " + division);
                                                                                              // returns next probable prime, long
    int a, b;
    BigInteger A, B;
                                                                                            // Driver method
    a = 54;
                                                                                            public static void main (String[] args) throws java.lang.Exception
    b = 23;
    A = BigInteger.valueOf(54);
    B = BigInteger.valueOf(37);
                                                                                              long n = 13;
                                                                                             System.out.println(checkPrime(n));
    A = new BigInteger("54");
    B = new BigInteger("123456789123456789");
    A = BigInteger.ONE;
    int c = a + b;
                                                                                         Stack Queue
                                                                                         class ch2_04_stack_queue {
    BigInteger C = A.add(B); // Other similar function are subtract() multiply(),
                                                                                          public static void main(String[] args) {
divide(), remainder(), mod()
                                                                                            Stack<Character> s = new Stack<Character>();
```

```
// Queue is abstract, must be instantiated with LinkedList
  // (special case for Java Queue)
  Queue<Character> q = new LinkedList<Character>();
  Deque<Character> d = new LinkedList<Character>();
  System.out.println(s.isEmpty());
                                    // currently s is empty, true
  System.out.println("======");
  s.push('a'); //push b,c
  // stack is LIFO, thus the content of s is currently like this:
                                             // output 'c'
System.out.println(s.peek());
                                       // pop topmost
  s.pop();
  System.out.println(s.peek());
                                               // output 'b'
  while (!s.isEmpty()) {
                                 // stack s still has 2 more items
   q.offer(s.peek()); // enqueue 'b', and then 'a' (the method name in Java Queue
for push/enqueue operation is 'offer')
   s.pop(); }
  q.offer('z');
                                    // add one more item
  System.out.println(q.peek());
                                                // prints 'b'
  // in Java, it is harder to see the back of the queue...
  // output 'b', 'a', then 'z' (until queue is empty), according to the insertion order
above
  System.out.println("=======");
  while (!q.isEmpty()) {
   System.out.printf("%c\n", q.peek());
                                          // take the front first
   q.poll();
                   // before popping (dequeue-ing) it
  System.out.println("======"):
  d.addLast('a');
  d.addLast('b');
  d.addLast('c');
  System.out.printf("%c - %c\n", d.getFirst(), d.getLast()); // prints 'a - c'
  d.addFirst('d');
  System.out.printf("%c - %c\n", d.getFirst(), d.getLast()); // prints 'd - c'
  d.pollLast();
  System.out.printf("%c - %c\n", d.getFirst(), d.getLast()); // prints 'd - b'
  d.pollFirst();
```

System.out.printf("%c - %c\n", d.getFirst(), d.getLast());}} // prints 'a - b'

KMP

```
import java.util.*;
class ch6 02 kmp {
 char[] T, P; // T = text, P = pattern
 int n, m; // n = length of T, m = length of P
 int [] b; // b = back table
 void naiveMatching() {
  for (int i = 0; i < n; i++) { // try all potential starting indices
   Boolean found = true;
   for (int j = 0; j < m && found; <math>j++) // use boolean flag `found'
     if (i + j \ge n || P[j] != T[i + j]) // if mismatch found
      found = false; // abort this, shift starting index i by +1
   if (found) // if P[0 ... m - 1] == T[i ... i + m - 1]
     System.out.printf("P is found at index %d in T\n", i);
void kmpPreprocess() { // call this before calling kmpSearch()
  int i = 0, j = -1; b[0] = -1; // starting values
  while (i < m) { // pre-process the pattern string P
   while (i \ge 0 \&\& P[i] != P[i]) i = b[i]; // if different, reset i using b
   i++; j++; // if same, advance both pointers
   b[i] = i; // observe i = 8, 9, 10, 11, 12 with i = 0, 1, 2, 3, 4
          // in the example of P = "SEVENTY SEVEN" above
 void kmpSearch() { // this is similar as kmpPreprocess(), but on string T
  int i = 0, j = 0; // starting values
  while (i < n) { // search through string T
   while (i \ge 0 \&\& T[i] != P[i]) i = b[i]; // if different, reset i using b
   i++; j++; // if same, advance both pointers
   if (i == m) { // a match found when i == m
     System.out.printf("P is found at index %d in T\n", i - j);
    i = b[i]; // prepare i for the next possible match
 } } }
 void run() {
  String Tstr = "I DO NOT LIKE SEVENTY SEV BUT SEVENTY SEVENTY
SEVEN":
  String Pstr = "SEVENTY SEVEN";
  T = new String(Tstr).toCharArray();
  P = new String(Pstr).toCharArray();
  n = T.length;
  m = P.length;
```

```
System.out.println(T);
  System.out.println(P);
  System.out.println();
  System.out.printf("Naive Mathing\n");
  naiveMatching();
  System.out.println();
  System.out.printf("KMP\n");
  b = new int[100010];
  kmpPreprocess();
  kmpSearch();
  System.out.println();
  System.out.printf("String Library\n");
  int pos = Tstr.indexOf(Pstr);
  while (pos != -1) {
   System.out.printf("P is found at index %d in T\n", pos);
   pos = Tstr.indexOf(Pstr, pos + 1);
  System.out.println();
 public static void main(String[] args)
  new ch6_02_kmp().run();
Binomial coefficient
#define MAXN 100 // largest n or m
long binomial_coefficient(n,m) // compute n choose m
int n,m;
{int i,j;
long bc[MAXN][MAXN];
    for (i=0; i<=n; i++) bc[i][0] = 1;
    for (j=0; j<=n; j++) bc[j][j] = 1;
```

```
for (i=1; i \le n; i++)
         for (j=1; j<i; j++)
        bc[i][j] = bc[i-1][j-1] + bc[i-1][j];
  return bc[n][m];
Euler's totient function
        the positive integers less than or equal to n that are relatively prime to n.
int phi (int n)
int result = n;
for (int i=2; i*i <= n; ++i) if (n \% i==0)
while(n \%i==0) n /= i;
result -= result / i;
if (n > 1)
result -= result / n; return result;
```

```
0/1 Knapsack problems
#include<iostream>
using namespace std;
int f[1000] = \{0\};
int n=0, m=0;
int main(void)
\{cin >> n >> m;
for (int i=1;i<=n;i++)
int price=0, value=0;
 cin >> price >> value;
  for (int j=m;j>=price;j--)
       if (f[j-price]+value>f[j])
         f[j]=f[j-price]+value;
  cout \ll f[m] \ll endl;
  return 0;
Longest common subsequence (LCS)
int dp[1001][1001];
int lcs(const string &s, const string &t)
  int m = s.size(), n = t.size();
  if (m == 0 || n == 0) return 0;
```

```
for (int i=0; i<=m; ++i)
    dp[i][0] = 0;
  for (int j=1; j <=n; ++j)
    dp[0][j] = 0;
  for (int i=0; i < m; ++i)
    for (int j=0; j< n; ++j)
       if (s[i] == t[j])
              dp[i+1][j+1] = dp[i][j]+1;
            else
              dp[i+1][j+1] =
  \max(dp[i+1][j], dp[i][j+1]);
  return dp[m][n];
Maxmium Matrix
int a[150][150]=\{0\};
int c[200]=\{0\};
int maxarray(int n)
int b=0, sum=-100000000;
for (int i=1;i<=n;i++)
if (b>0) b+=c[i];
```

```
else b=c[i];
if (b>sum) sum=b;
return sum;
int maxmatrix(int n)
int sum=-100000000, max=0;
for (int i=1; i <=n; i++)
for (int j=1; j <=n; j++)
c[j]=0;
for (int j=i;j <=n;j++)
for (int k=1;k<=n;k++)
c[k]+=a[j][k];
max=maxarray(n);
if (max>sum) sum=max;
return sum;
int main(void)
int n=0;
```

```
cin >> n;
for (int i=1;i<=n;i++)
for (int j=1;j<=n;j++)
cin >> a[i][j];
cout << maxmatrix(n);
return 0;
}</pre>
```

Count number of ways to partition a set into k subsets

$$S(n, k) = k*S(n-1, k) + S(n-1, k-1)$$

Flood fill algorithm

```
//component(i) denotes the
//component that node i is in
void flood_fill(new_component)
do
num_visited = 0
for all nodes i
if component(i) = -2
num_visited = num_visited + 1
component(i) = new_component
for all neighbors j of node i
if component(j) = nil
component(j) = -2
until num_visited = 0
void find_components()
```

```
num\_components = 0
                                                                                          /*
for all nodes i
                                                                                                  for undirected graph a[x][0]++; a[y][a[y][0]]=x; w[y][x]=z;
component(node i) = nil
                                                                                          //
                                                                                          */
for all nodes i
if component(node i) is nil
num\_components = num\_components + 1
                                                                                          int s=0, e=0;
component(i) = -2
                                                                                          cin >> s >> e; // s: start, e: end
flood_fill(component num_components)
                                                                                          SPFA(s);
                                                                                          cout \ll d[e] \ll endl;
SPFA — shortest path
                                                                                          return 0;
int q[3001] = \{0\}; // queue for node
int d[1001]=\{0\}; // record shortest path from start to ith node bool f[1001]=\{0\};
                                                                                          void SPFA(int v0)
int a[1001][1001]={0}; // adjacency list
int w[1001][1001]={0}; // adjacency matrix
                                                                                          int t,h,u,v;
                                                                                          for (int i=0;i<1001;i++) d[i]=INT_MAX;
int main(void)
                                                                                          for (int i=0;i<1001;i++) f[i]=false;
int n=0, m=0;
                                                                                          d[v0]=0;
cin >> n >> m;
                                                                                          h=0; t=1; q[1]=v0; f[v0]=true;
for (int i=1; i <= m; i++)
                                                                                          while (h!=t)
int x=0, y=0, z=0;
                                                                                           h++;
                                                                                          if (h>3000) h=1;
cin >> x >> y >> z; // node x to node y has weight z
a[x][0]++;
                                                                                          u=q[h];
a[x][a[x][0]]=y;
                                                                                          for (int j=1; j \le a[u][0]; j++)
w[x][y]=z;
```

```
v=a[u][j]
if (d[u]+w[u][v]< d[v]) // change to > if calculating longest path
d[v]=d[u]+w[u][v];
if (!f[v])
t++;
if (t>3000) t=1;
q[t]=v;
f[v]=true;
f[u]=false;
Floyd-Warshall algorithm – shortest path of all pairs
//map[i][j]=infinity at start void floyd()
for (int k=1; k<=n; k++)
for (int i=1; i<=n; i++)
for (int j=1; j<=n; j++)
if (i!=j && j!=k && i!=k)
if (map[i][k]+map[k][j] < map[i][j])
map[i][j]=map[i][k]+map[k][j];
```

Strings

```
string str("The Geeks for Geeks");
// find() returns position to first
// occurrence of substring "Geeks"
// Prints 4
cout << "First occurrence of \"Geeks\" starts from : ";</pre>
cout << str.find("Geeks") << endl;</pre>
string str="We think in generalities, but we live in
details.";
string str2 = str.substr (3,5);
                                    // "think"
size t pos = str.find("live");
                                    // position of "live"
in str
string str3 = str.substr (pos);  // get from "live" to
the end
//OUTPUT: think live in details.
string str1 ("green apple");
string str2 ("red apple");
if (str1.compare(str2) != 0)
    cout << str1 << " is not " << str2 << '\n';</pre>
if (strl.compare(6,5,"apple") == 0)
    cout << "still, " << str1 << " is an apple\n";</pre>
if (str2.compare(str2.size()-5,5,"apple") == 0)
    cout << "and " << str2 << " is also an apple\n";</pre>
if (str1.compare(6,5,str2,4,5) == 0)
    cout << "therefore, both are apples\n";</pre>
/*OUTPUT: green apple is not red apple
still, green apple is an apple
and red apple is also an apple
therefore, both are apples*/
```