## ISLAMIC UNIVERSITY OF TECHNOLOGY (IUT) ORGANISATION OF ISLAMIC COOPERATION (OIC)

## Department of Computer Science and Engineering (CSE)

MID SEMESTER EXAMINATION

WINTER SEMESTER, 2021-2022

**DURATION: 1 HOUR 30 MINUTES** 

**FULL MARKS: 75** 

## CSE 4703: Theory of Computing

Programmable calculators are not allowed. Do not write anything on the question paper. Answer all 3 (three) questions. Marks of each question and corresponding CO and PO are written in the right margin with brackets.

a) Explain if each of the following assertion is correct or incorrect. 1.

 $2 \times 4 = 8$ 

(CO1)

- Deterministic Finite Automata (DFA) are strictly weaker classes of machines (PO2) i. than Non-deterministic Finite Automata (NFA), i.e., there exists a language that is accepted by an NFA and not by any DFA.
- $a^n b^m$  is a regular language, where the alphabet is a, b and  $n \ge 0$ ,  $m \ge 0$ . ii.
- The regular expressions  $(a^*b^*)^*$  and  $(a+b)^*$ , represent the same language, iii. where  $\Sigma = \{a, b\}^*$
- Considering the languages  $L_1 = \emptyset$  and  $L_2 = \{1\}$ , the strings are accepted iv. by the language  $L_1^{\circ} L_2^{*} + L_1^{*}$  is  $\{\varepsilon, 1\}$ .
- Prove that if L, M and N are any languages, then L(M + N) = LM + LN. (CO1) (PO1)
- c) Consider the following NFA shown in Figure 1. Convert this NFA into an equivalent 10 (CO1) DFA. Your answer should be the state diagram of a DFA. Your diagram should include (PO2) only the states that are reachable from the start state.

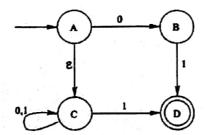
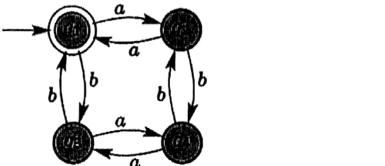


Figure 1: NFA for the question no. 1. (c)

- 2. a) Assume that the NFA,  $N_I$  recognizes language,  $L_I$ ; and  $N_2$  recognizes language,  $L_2$ . Construct the NFA N to recognize the language,  $L_1^{\circ}L_2$ . Use the *proof by construction* and also draw the *schematic diagram* for the new machine N. (CO1)
  - b) Convert the given DFA in Figure 2 over  $\Sigma = \{a, b\}$  into regular expression. (CO1)



(PO2)

Figure 2: DFA for question 2. (b)

- c) Design the DFA/NFA for the following regular expression R over  $\Sigma = \{a, b\}$ . 8  $R = (ab)^* + (a + ab)^*b^*(a + b)^* \tag{CO1}$ (PO2)
- 3. a) Design the regular expression for the following languages over the alphabet  $\Sigma = \{0, 1\}$ .

  i.  $L_1 = \{w \mid w \text{ consists of at least one '00' and '11' and not divisible by 2}\}$ .

  (CO1)
  - ii.  $L_2 = \{w \mid w \text{ starts and ends with } 1\}.$  (PO2)
  - iii.  $L_3 = \{w \mid in \ w, the \ number \ of \ '1' \ is \ exactly \ 3 \ and \ it \ must \ end \ with one \ or \ more \ '0's\}.$
  - b) Design a *Finite Stat Machine* that accepts strings in which the number '1' is congruent to  $1 \mod 3$  over the alphabet  $\Sigma = \{0, 1\}$ . (CO1) (PO2)
  - c) Apply the pumping lemma to prove that the language L over the alphabet  $\{a, b\}$  is not regular  $L = \{ww : w \in \{a, b\}^*.$  (CO1) (PO1)