

A Performance Comparison of Linux and a Lightweight Kernel

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ASCI Red Hardware

- 4640 compute nodes
 - Dual 333 MHz Pentium II Xeons
 - 256 MB RAM
- 800 MB/sec bi-directional network links
- 20 microsecond network latency
- 38x32x2 mesh topology
- Deployed in 1997









ASCI Red Storm

- 10,368 compute node processors (AMD Opterons @ 2.0 GHz)
- 10 TB of DDR memory @ 333MHz (1GB per node)
- Compute node topology:
 - $-27 \times 16 \times 24 (x, y, z)$
 - Mesh in x & y, torus in z
- 3 GB/s network bandwidth
- 5 microsecond network latency







A Lightweight Compute Node Operating System is a Fundamental Part of the Sandia Architecture

- It is essential for
 - Maximizing CPU resources
 - Reduce OS and runtime system overhead
 - Maximizing memory resources
 - Small memory footprint, large page support
 - Maximizing network resource
 - No virtual memory, physically contiguous address mapping
 - Increasing reliability
 - Small code base, reduced complexity
 - Deterministic performance
 - Repeatability
 - Scalability
 - OS resources should be independent of job size
- Others have realized these benefits
 - nCUBE (Vertex), Cray T3 (UNICOS/mk), IBM BG/L (HPK)





ASCI Red Compute Node Software

- Puma lightweight kernel
 - Follow-on to Sandia/UNM Operating System (SUNMOS)
 - Developed for 1024-node nCUBE-2 in 1993 by Sandia/UNM
 - Ported to 1800-node Intel Paragon in 1995 by Sandia/UNM
 - Ported to ASCI Red in 1996 by Intel and Sandia
 - Productized as "Cougar" by Intel







ASCI Red Software (cont'd)

Cougar

- Space-shared model (not time-shared)
- Exposes all resources to applications
- Consumes less than 1% of compute node memory
- Four different execution modes for managing dual processors
- Portals 2.0
 - High-performance message passing
 - Avoid buffering and memory copies
 - Supports multiple user-level libraries (MPI, Intel N/X, Vertex, etc.)







Cougar

Goals

- Target scientific and engineering applications on tightly coupled distributed memory architectures
- Scalable to tens of thousands of processors
- Fast message passing and execution
- Small memory footprint

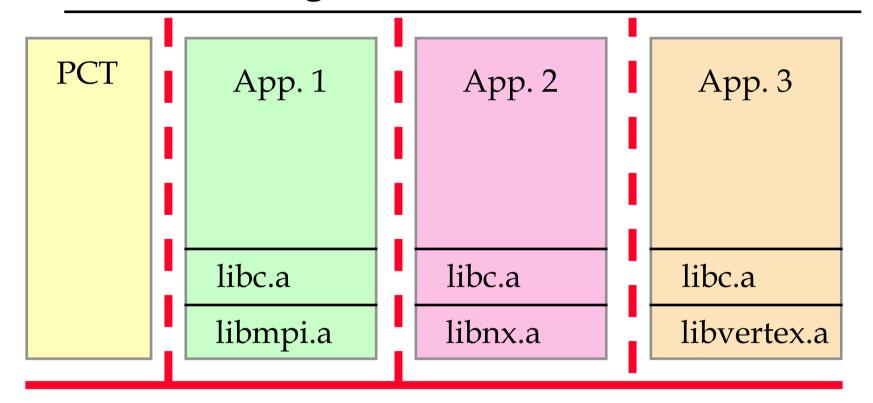
Approach

- Separate policy decision from policy enforcement
- Protect applications from each other
- Let user processes manage resources
- Get out of the way





Cougar General Structure



Q-Kernel: message passing, memory protection





Cougar Quintessential Kernel (QK)

- Policy enforcer
- Initializes hardware
- Handles interrupts and exceptions
- Maintains hardware virtual addressing (but no virtual memory)
- Small, static size
- Few, well defined entry points





Cougar Process Control Thread (PCT)

- Runs in user space (but more privileged than apps)
- Customizable
 - Single-tasking or multi-tasking
 - Round robin or priority scheduling
 - High performance, debugging, or profiling version
- Changes behavior of OS without changing the kernel
- Policy Maker
 - Process loading and scheduling
 - Virtual address space management
 - Name Server
 - Fault handling







Cougar Processor Modes

- Chosen at job launch time
- Heater mode (proc 0)
 - QK/PCT and application process on system CPU
- Message co-processor mode (proc 1)
 - QK/PCT on system CPU
 - Application process on second CPU
- Compute co-processor mode (proc 2)
 - QK/PCT and application process on system CPU
 - Application co-routines on on second CPU
- Virtual node mode (proc 3)
 - QK/PCT and application process on system CPU
 - Second application process on second CPU







Research Goals

- Assess the performance and reliability of a lightweight kernel versus a traditional monolithic kernel
- Determine how to bring lightweight kernel advantages to general platforms







Current Approach

- Short-term
 - Compare Cougar and Linux on ASCI/Red hardware
- Beyond that
 - Figure out how best to leverage Linux or other open-source operating systems to achieve important characteristics of previous LWKs
 - Provide a basis for future OS research activities





Motivation for Linux/LWK Comparison

- No direct comparison of LWK versus full-service OS since SUNMOS versus OSF1/AD nearly ten years ago
- Much has changed (improved?) since
- A direct comparison between a LWK and Linux is important for providing insight into what is important
- Platform balance is important
- Need real numbers to show people like:
 <insert favorite skeptic here>







Linux on ASCI Red

- RedHat 7.2 Linux 2.4.18
- Adapted Linux bootloader and startup code to work with bootmesh protocol
- Service node receives Linux kernel via bootmesh and root filesystem from attached SCSI disk
- Compute nodes mount root filesystem from service node
- Sparse compute node services
 - sshd for remote access
 - Enough libraries for MPI jobs to run





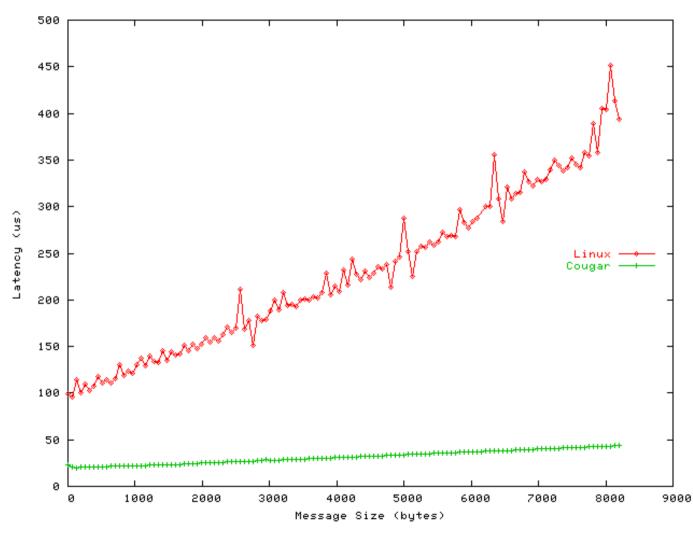
Linux IP Implementation for ASCI Red

- Implemented a Linux network driver for CNIC
 - Interrupt-driven ring buffer
 - Based on isa-skeleton.c
- Varying IP MTU from 4 KB (1 page) to 16 KB (4 pages) showed no noticeable difference in bandwidth
- Bandwidth is CPU limited
 - 45 MB/s for 333 Mhz processors
 - 32 MB/s for 200 MHz processors
- Custom raw device achieved 310 MB/s















But to be fair...

- Implemented a Portals 3.2 CNIC driver in Linux
 - 46 μs latency, 280 MB/s
- Not quite perfectly fair because Portals 3.2 vs. 2.0
- ...but as close as we can get







NPB 2.4 - CG

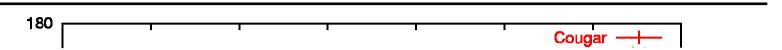
Cougar —







NPB 2.4 - IS

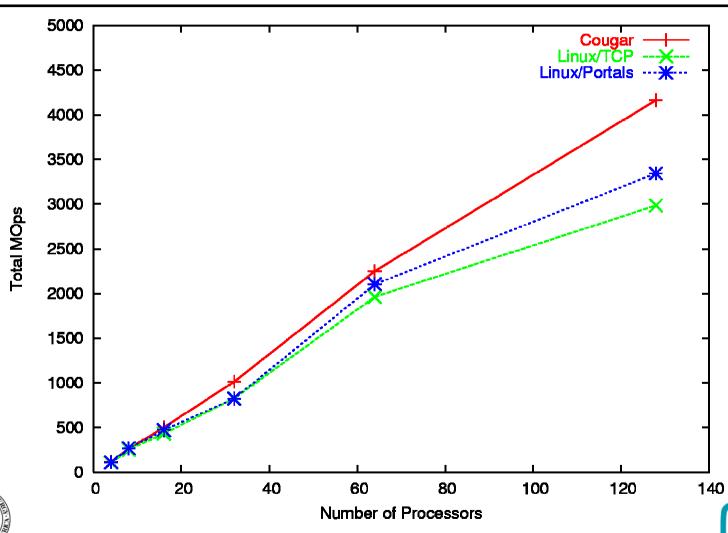








NPB 2.4 - MG







CTH Family of Codes

- Models complex multi-dimensional, multi-material problems characterized by large deformations and/or strong shocks
- Uses two-step, second-order accurate finitedifference Eulerian solution
- Material models for equations of state, strength, fracture, porosity, and high explosives
- Impact, penetration, perforation, shock compression, high explosive initiation and detonation problems







CTH Steps

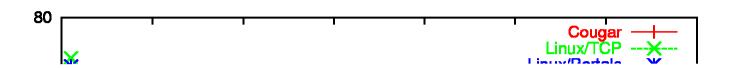
- Read initial restart file, one file per node
- Simulate shock wave physics
 - Many nearest-neighbor communications, a few global reductions per time step
- Write results to restart, history, and viz files
- Performance measured in grind time
 - Time to compute all calculations on a single cell for a single time step







CTH Performance

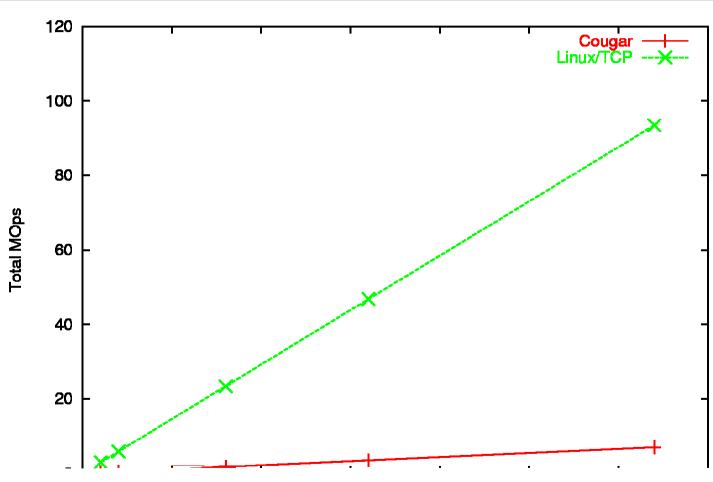








NPB 2.4 - EP









Issues

- Compilers and runtime
 - Cougar numbers are from (old) PGI compilers
 - Linux numbers are from (new) Intel compilers
- Determinism
 - No variability in Cougar execution times
 - Even on a loaded machine
 - Significant (>5%) variability in Linux execution times
- Level of effort
 - Maintaining LWK may be equivalent to maintaining a Linux driver







Conclusions

- Finally have a real apples-to-apples comparison (albeit granny smith to red delicious)
- Numerous issues make fair comparison hard (should be perfectly fair on Red Storm, but that is still a long time away)
- Definite evidence of limitation of Linux at scale (still investigating)
- Definite advantages of LWK for some apps (IS)







Future Work

- Linux 2.6
 - Large page support
- Cougar
 - Provide a modern set of compilers/libraries
- Broader range of applications







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