

Transportation network with multiple lines

notes GG

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1 Formalism

1.1 The transportation network with multiple lines

Let $\mathcal{G} = (\mathcal{V}, \mathcal{E})$ be a simple, oriented, and connected graph representing a transportation network between $|\mathcal{V}| = n$ nodes, having $|\mathcal{E}| = m$ edges, and possessing p different transportation lines. Each node belongs to only one line, i.e. $\mathcal{V} = \bigcup_{k=1}^p \mathcal{V}_k$ and $\bigcap_{k=1}^p \mathcal{V}_k = \emptyset$, where \mathcal{V}_k represents the set of nodes in line k . The edge set \mathcal{E} , can also be decomposed with

$$\mathcal{E} = \mathcal{E}_W \cup \mathcal{E}_B, \quad \mathcal{E}_W := \bigcup_{k=1}^p \mathcal{E}_k, \quad \mathcal{E}_W \cap \mathcal{E}_B = \emptyset, \quad \mathcal{E}_k \cap \mathcal{E}_l = \emptyset, \quad \forall k, l. \quad (1)$$

where \mathcal{E}_k is the set of edges composing line k , \mathcal{E}_W the set containing all edges inside lines, and \mathcal{E}_B the set of transfer edges, connecting the different lines. The graph \mathcal{G} can be represented by its adjacency matrix $\mathbf{A} = (a_{ij})$, which can also be decomposed with

$$\mathbf{A} = \mathbf{A}_W + \mathbf{A}_B, \quad \mathbf{A}_W = \sum_{k=1}^p \mathbf{A}_k, \quad (2)$$

with $\mathbf{A}_k = (a_{ij}^k)$ are edges of line k , $\mathbf{A}_W = (a_{ij}^W)$ edges of inside all lines, and $\mathbf{A}_B = (a_{ij}^B)$ transfer edges. We suppose that there is an uniquely define route inside lines, i.e.

$$a_{i\bullet}^k \leq 1 \text{ and } a_{\bullet i}^k \leq 1, \quad \forall i, k. \quad (3)$$

where \bullet designates a summation over the replaced index.

1.2 The origin-destination matrix

The $(n \times n)$ *origin-destination matrix*, denoted by $\mathbf{N} = (n_{st})$, $n_{st} \geq 0$, $\forall s, t$, contains the flow (e.g. the number of passengers) entering the network in source node s and leaving it in target node t . We can denote its margins with

$$\boldsymbol{\sigma}_{\text{in}} := \mathbf{N} \mathbf{e}_n \quad (4)$$

$$\boldsymbol{\sigma}_{\text{out}} := \mathbf{N}^\top \mathbf{e}_n \quad (5)$$

where \mathbf{e}_n is the vector of ones of size n . The vector $\sigma_{\text{in}} = (\sigma_i^{\text{in}})$ is the *vector of flow entering the network* and $\sigma_{\text{out}} = (\sigma_i^{\text{out}})$ is the *vector of flow leaving the network*. We have

$$\sigma_{\bullet}^{\text{in}} = \sigma_{\bullet}^{\text{out}}. \quad (6)$$

Note that if only σ_{in} and σ_{out} are given, a flow matrix \mathbf{N} can be computed relatively to an *origin-destination affinity matrix* $\mathbf{S} = (s_{st})$, $0 \leq s_{st} \leq 1$, where $s_{st} = 1$ denote a perfect affinity and $s_{st} = 0$ no affinity, through

$$\mathbf{N} = \mathbf{Diag}(\mathbf{a})(\mathbf{S} + \epsilon)\mathbf{Diag}(\mathbf{b}), \quad (7)$$

where $\mathbf{Diag}(\cdot)$ denote the diagonal matrix obtained from a vector, ϵ a very small quantity, and vectors \mathbf{a} and \mathbf{b} are found through *proportional iterative fitting* algorithm in order to have margin constraints (4) and (5) respected for \mathbf{N} (a small ϵ has to be added to σ_{in} and σ_{out} if they possess null components).

1.3 The flow matrix

A flow on edges is represented by the $(n \times n)$ *flow matrix* $\mathbf{X} = (x_{ij})$, verifying

$$x_{ij} \geq 0, \quad \forall i, j, \quad (8)$$

$$a_{ij} = 0 \Rightarrow x_{ij} = 0, \quad \forall i, j, \quad (9)$$

$$x_{i\bullet} + \sigma_i^{\text{out}} = x_{\bullet i} + \sigma_i^{\text{in}}, \quad \forall i. \quad (10)$$

Again, we can decompose the flow matrix with

$$\mathbf{X} = \mathbf{X}_W + \mathbf{X}_B \quad \mathbf{X}_W := \sum_{k=1}^l \mathbf{X}_k \quad (11)$$

where \mathbf{X}_k represent the flow inside line k , \mathbf{X}_W is the flow inside all lines, and \mathbf{X}_B the flow between lines. This decomposition allows us to define the *vector of flow entering lines* $\rho_{\text{in}} = (\rho_i^{\text{in}})$ and the *vector of flow leaving lines* $\rho_{\text{out}} = (\rho_i^{\text{out}})$, with

$$\rho_{\text{in}} := \sigma_{\text{in}} + \mathbf{X}_B^{\top} \mathbf{e}_n, \quad (12)$$

$$\rho_{\text{out}} := \sigma_{\text{out}} + \mathbf{X}_B \mathbf{e}_n, \quad (13)$$

where \mathbf{e}_n is the vector of ones of size n . It is easy to see that we still have $\rho_{\bullet}^{\text{in}} = \rho_{\bullet}^{\text{out}}$.

1.4 Shortest-paths flow

Let \mathcal{P}_{st} be the set of *admissible* shortest-paths between s and t on \mathcal{G} . We can denote by $P_{st}(i, j)$ the probability of having edge $(i, j) \in \wp$ when drawing a path \wp from \mathcal{P}_{st} . We have

$$P_{st}(i, j) := \frac{1}{|\mathcal{P}_{st}|} \sum_{\wp \in \mathcal{P}_{st}} \delta((i, j) \in \wp), \quad (14)$$

where $\delta(\cdot)$ designate the indicator function. Note that if there is an unique shortest-path between node s and t , noted \wp_{st} , we have $P_{st}(i, j) = \delta((i, j) \in \wp_{st})$.

If we are given an origin-destination matrix $\mathbf{N} = (n_{st})$, we can compute the *shortest-path flow matrix*, noted $\mathbf{X}_{\text{sp}} = (x_{ij}^{\text{sp}})$, with

$$x_{ij}^{\text{sp}} = \sum_{st} P_{st}(i, j) n_{st}. \quad (15)$$

This matrix contains the flow on each edge if we suppose that the flow follows shortest-paths from origin to destination.

We can rewrite equation (15) by defining the $(n^2 \times n^2)$ *shortest-path - edge matrix* $\mathbf{P} = (p_{\alpha\beta})$ with

$$p_{\alpha\beta} = \begin{cases} P_{st}(i, j) & \text{if } \alpha = t + n(s - 1) \text{ and } \beta = j + n(i - 1), \\ 0 & \text{otherwise.} \end{cases} \quad (16)$$

Then (15) writes

$$\mathbf{vec}(\mathbf{X}_{\text{sp}}) = \mathbf{P}^\top \mathbf{vec}(\mathbf{N}), \quad (17)$$

where $\mathbf{vec}(\cdot)$ denotes the vectorization function of a matrix, obtained by stacking matrix columns on top of one another.

From equation (15), we see that

$$\frac{\partial x_{ij}^{\text{sp}}}{\partial n_{st}} = P_{st}(i, j), \quad (18)$$

which equal to 1 if there is a unique shortest-path between s and t and (i, j) belongs to this path.

1.5 Problem definition

The problem: We suppose that we know the flow entering and leaving each line, i.e. ρ_{in} and ρ_{out} and we want to find origin-destination trajectories n_{st} .

By setting the problem like that, we easily see that it is ill posed. Several solutions exists (e.g. units remain on the same line and follows a first-in/first-out scheme) and we need to add some hypotheses to restrain it.

Hypothesis 1: Trajectories follow shortest-paths from origin s to destination t .

Hypothesis 2: The number of trajectories n_{st} should be as close as possible to $\rho_s^{\text{in}} \rho_t^{\text{out}} s_{st}$, where $S = (s_{st})$ is a given affinity matrix between origin and destination nodes.

1.6 Algorithm

Set $\sigma_{\text{in}}^{(1)} = \rho_{\text{in}}$, $\sigma_{\text{out}}^{(1)} = \rho_{\text{out}}$, and $\mathbf{S}^{(1)} = \mathbf{S}$. Until convergence, do:

1. Compute $\mathbf{N}^{(i)} = \text{Diag}(\mathbf{a}^{(i)})(\mathbf{S}^{(i)} + \epsilon)\text{Diag}(\mathbf{b}^{(i)})$ with iterative fitting, such that $\mathbf{N}^{(i)} \mathbf{e}_n = \sigma_{\text{in}}^{(i)} + \epsilon$ and $(\mathbf{N}^{(i)})^\top \mathbf{e}_n = \sigma_{\text{out}}^{(i)} + \epsilon$.
2. Compute the associated shortest-path flow matrix $\mathbf{X}^{(i)}$ with (15).
3. Compute $\sigma_{\text{in}}^{(i+1)} = \rho_{\text{in}} - (\mathbf{X}_{\text{B}}^{(i)})^\top \mathbf{e}_n$ and $\sigma_{\text{out}}^{(i+1)} = \rho_{\text{out}} - \mathbf{X}_{\text{B}}^{(i)} \mathbf{e}_n$.
4. Set negative components of $\sigma_{\text{in}}^{(i+1)}$ and $\sigma_{\text{out}}^{(i+1)}$ to 0.
5. Update $\mathbf{S}^{(i+1)}$ according to the flow with is unable to pass.