# Project 1 (Water Pitcher)

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### Overview

The Water Pitcher problem is:

- You have several finite-capacity pitchers (3 liters, 5 liters, etc.) plus one infinite pitcher.
- Your goal is to measure out a specific target quantity of water into the infinite pitcher.
- Each move (pouring or filling) has a cost of 1 step.
- You want the **shortest** sequence of steps that results in the infinite pitcher holding the target amount.

## File Structure

- main.py
  - parse\_file: Reads capacities and target from file.
  - can\_measure: Checks if target is possible (via GCD).
  - heuristic: Lower bound for  $A^*$  = ceil(remaining / smallest\_pitcher) (never overestimates).
  - fill / pour: Generate next states.
  - a\_star: Runs A\* to find shortest solution or -1 if impossible.

### · test.py

 Uses unittest to validate parsing, measuring, heuristic correctness, pouring/filling operations, and final A\* results.

# Tests ensure that:

- The input file is correctly parsed.
- Feasibility of measuring the target with given containers is accurately determined.
- The heuristic function provides reasonable estimates.
- Filling and pouring operations produce correct new states.
- The overall  $\mathbf{A}^*$  search algorithm correctly identifies both unsolvable and trivial cases.

### A\* Overview

 $A^*$  expands states in order of their priority f = g + h, where:

•  $\mathbf{g}$  = the number of steps taken so far.

•  $\mathbf{h} = \text{the heuristic (our lower bound on how many more steps are needed)}$ .

# Why the Lower Bound Is Crucial

```
def heuristic(state, goal, sizes):
needed = goal - state[-1]
if not needed:
    return 0
finite = sizes[:-1]
return math.inf if not finite else ceil(abs(needed) / min(finite))
```

- The heuristic here in the code snippet, doesn't *overestimate* the remaining cost, so it is **admissible**.
- An admissible heuristic guarantees that  $A^*$  always finds the optimal (shortest) path. Without a valid lower bound, it would just be a uniform-cost search, losing  $A^*$ 's efficiency advantage.

# Run program:

```
python main.py input1.txt
```

### Run test:

```
python -m unittest test.py
```