Lock-Free Ring Buffer for Market Data

Low Latency Data Pipeline Team

April 7, 2025

Outline

Introduction

Ring Buffer Design

Testing

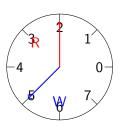
Applications & Conclusions

Why Lock-Free Data Structures?

- Challenge: Modern market data systems process millions of updates per second
- ► **Problem**: Traditional mutex-based synchronization creates bottlenecks
- ► **Solution**: Lock-free data structures provide:
 - No mutex acquisition/release overhead
 - No thread blocking
 - Lower and more predictable latency
 - Better scalability under contention
- ► **Application**: Ideal for producer-consumer patterns in market data processing

Ring Buffer Fundamentals

- Fixed-size circular buffer
- ► Two atomic indices:
 - write_idx_ Next position to write
 - read_idx_ Next position to
 read
- Empty when read_idx_ ==
 write_idx_
- ► Full when (write_idx_ + 1) % Size == read_idx_
- One slot always reserved for empty detection



Implementation Highlights

```
template <typename T, size_t Size>
class LockFreeRingBuffer {
public:
 bool TryPush(const T& item) {
    const size_t current_write = write_idx_.load(std::
   memory order relaxed);
   const size_t next_write = (current_write + 1) % Size
   ;
    if (next_write == read_idx_.load(std::
   memory_order_acquire))
      return false; // Buffer full
   buffer_[current_write] = item;
    write_idx_.store(next_write, std::
   memory_order_release);
   return true;
 bool TryPop(T* output) { /* Similar implementation */
```

Key Technical Features

- ► Lock-free algorithm: Using atomic variables
- Memory ordering: Careful use of memory ordering constraints
 - memory_order_relaxed for initial loads
 - memory_order_acquire for checking conditions
 - memory_order_release for committing updates
- ► Cache line alignment: Preventing false sharing between indices
 - Producer and consumer threads operate on different cache lines
 - ▶ alignas (64) ensures indices are on separate cache lines
- Non-blocking behavior: TryPush/TryPop never block, return boolean success

Basic Functionality Test

► Tests fundamental properties with realistic market data structures:

```
struct MarketTick {
   int64_t timestamp_ns;
   std::string symbol;
   double price;
   double quantity;
   char side; // 'B' for buy, 'S' for sell

   // For testing equality in our assertions
   bool operator==(const MarketTick(*@&@*) other)
   const;
};
```

- Verifies:
 - Empty buffer detection
 - Push until full
 - ► Pop in FIFO order
 - Alternating push/pop sequences

Market Data Pipeline Test

- Simulates realistic market data processing:
 - Producer thread: generates synthetic market tick data
 - Consumer thread: processes ticks and measures latency
- Test parameters:
 - ▶ 5,000 market ticks
 - ▶ 1024-slot buffer
 - Alternating BTC/ETH symbols
 - Realistic price movements
- Measures end-to-end latency from tick creation to processing

Performance Results

- ➤ **Test machine**: Modern x86_64 system
- ► Compiler: GCC 13.3.0
- **► C++ standard**: C++23
- ▶ Build type: Release (-O3)

- ► Minimum latency: 215 ns
- ► Median latency: 26,305 ns
- **▶ 99th percentile**: 3,284,727 ns
- ► Maximum latency: 4,554,931 ns

Applications in Market Data Systems

- ▶ Feed handlers: Buffering incoming market data packets
- ▶ Parsing pipeline: Moving raw data to normalization stage
- Order book updates: Efficiently queuing price updates
- Strategy components: Communicating signals between modules
- ▶ **Risk checks**: Queuing pre-trade validations
- ▶ **Logging**: Non-blocking capture of events for later analysis

Conclusions & Next Steps

Achievements:

- Created a fully lock-free, high-performance ring buffer
- Demonstrated functionality with real market data structures
- Measured sub-microsecond minimum latency
- Successfully processed thousands of events with predictable performance

Potential improvements:

- ► Multi-producer/multi-consumer variants
- Batched operations for higher throughput
- Memory reclamation for dynamically allocated elements
- Integration with hardware acceleration (FPGA, GPU)

Thank You

Questions?