

Throughout this note, ^{connected}

M : cpt d -dim. submfld of \mathbb{R}^n

(boundary and boundary corners are permitted).

Def

For $x, y \in M$, define

$$d_M(x, y) := \inf \{ L(\gamma) \mid \gamma: \text{admissible from } x \text{ to } y \}$$

Def

$X = \{x_i\} \subseteq M$, a finite set.

G : a graph on X .

Define, $\forall x, y \in X$,

$$d_G(x, y) := \min_P (\|x_0 - x_1\| + \dots + \|x_{p-1} - x_p\|) \text{ and}$$

$$d_S(x, y) := \min_P (d_M(x_0, x_1) + \dots + d_M(x_{p-1}, x_p)),$$

where $P = (x_0, x_1, \dots, x_p)$ is a path on G

w/ $x_0 = x$ and $x_p = y$.

Goal: We want to prove $d_G \approx d_M$.

Approach: $d_G \approx d_S$ and $d_S \approx d_M$.

Prop 1

$d_M(x, y) \leq d_S(x, y)$ and

$d_G(x, y) \leq d_S(x, y)$, $\forall x, y \in X$.

(Pf)

The first is obvious.

The second comes from the fact that straight lines are shortest in \mathbb{R}^n . \oplus

Thm 2 ($d_M \approx d_S$).

$\varepsilon, \delta > 0$ w/ $4\delta < \varepsilon$. Suppose

① G contains all edges xy w/ $d_M(x, y) \leq \varepsilon$

② $\forall m \in M, \exists x_i \in X$ s.t. $d_M(x_i, m) \leq \delta$.

Then, $\forall x, y \in X$,

$$d_M(x, y) \leq d_S(x, y) \leq (1 + 4\delta/\varepsilon) d_M(x, y). \quad \begin{matrix} \text{ISO} \\ \text{PI} \end{matrix}$$

Rank:

Condition ② is called the δ -sampling condition.

(Pf)

" \leq " of LHS follows from Prop 1.

Now consider " \leq " on the RHS.

It suffices to prove the following claim.

claim: \forall piecewise C^∞ curve γ from x to y ,
w/ $l = L(\gamma)$, \exists path $P = (x_0, \dots, x_p)$ on G
from x to y s.t. $d_M(x_0, x_1) + \dots + d_M(x_{p-1}, x_p) \leq$
(then the result follows from taking $(1 + 4\delta/\varepsilon)l$
inf. among all such γ).

(Pf of claim).

case 1 $l \leq \varepsilon - 2\delta$. (Then $d_M(x, y) \leq l \leq \varepsilon$).

By condition ①, $xy \in G$.

$$\Rightarrow d_S(x, y) \leq d_M(x, y) \leq l \leq (1 + 4\delta/\varepsilon) l.$$

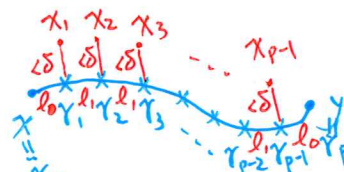
(x, y) is
a path in G .

case 2 $l > \varepsilon - 2\delta$.

By division algorithm,

$$l = (\varepsilon - 2\delta)q + r, \text{ w/ } q \in \mathbb{N} \cup \{0\} \text{ and } r \in [\varepsilon - 2\delta, 2(\varepsilon - 2\delta)].$$

$$= l_0 + (\underbrace{l_1 + \dots + l_1}_{q = p-2}) + l_0, \text{ w/ } \underbrace{l_1 = \varepsilon - 2\delta}_{\text{and}} \text{ and } \underbrace{l_0 = r/2 \in [\frac{\varepsilon - 2\delta}{2}, \varepsilon - 2\delta]}.$$



Divide γ according to this length subdivision as in the picture. (Denote the pts by x_1, \dots, x_{p-1}).

By condition ②, for $1 \leq i \leq p-1$, $\exists x_i \in X$

s.t. $d_M(x_i, x_{i+1}) \leq \delta$.

Then $d_M(x_i, x_{i+1}) \leq \delta + l_i + \delta = \varepsilon = l_i \cdot \frac{\varepsilon}{\varepsilon - 2\delta}$.

$d_M(x, x_1) \leq l_0 + \delta$ and $d_M(x_{p-1}, y) \leq \delta + l_0$.

Notice that $\frac{1}{2} \leq \frac{l_0}{\varepsilon - 2\delta} \Rightarrow \delta \leq l_0 (\frac{2\delta}{\varepsilon - 2\delta})$

$$\Rightarrow l_0 + \delta \leq l_0 \left(\frac{\varepsilon}{\varepsilon - 2\delta} \right).$$

$$\text{Thus } d_M(x, x_1) \leq l_0 \cdot \left(\frac{\varepsilon}{\varepsilon - 2\delta} \right) \text{ and}$$

$$d_M(x_{p-1}, y) \leq l_0 \cdot \left(\frac{\varepsilon}{\varepsilon - 2\delta} \right).$$

Hence,

$$d_M(x, x_1) + d_M(x_1, x_2) + \dots + d_M(x_{p-1}, y)$$

$$\leq \left(\frac{\varepsilon}{\varepsilon - 2\delta} \right) (l_0 + l_1 + \dots + l_{p-1} + l_0) \leq \frac{\varepsilon}{\varepsilon - 2\delta} \cdot l.$$

$$\text{Notice that } \frac{1}{1-t} < 1+2t, \forall t \in (0, 1/2).$$

$$\frac{\varepsilon}{\varepsilon - 2\delta} = \frac{1}{1 - (2\delta/\varepsilon)}.$$

$$2\delta/\varepsilon \in (0, 1/2) \text{ (due to } 4\delta < \varepsilon).$$

$$\text{Thus } \frac{\varepsilon}{\varepsilon - 2\delta} \cdot l < (1 + 2 \cdot (2\delta/\varepsilon)) \cdot l = (1 + \frac{4\delta}{\varepsilon}) l$$

Therefore, the claim follows. (#)

Def

The minimum radius of curvature $r_0 = r_0(M)$ is

$$\text{defined by } \frac{1}{r_0} = \max_{\gamma, t} \{ \|\ddot{\gamma}(t)\| \}, \text{ where}$$

γ varies over all unit speed geod. and

t is in the domain of γ .

Rmk.: (Intuition)

Geod. in M "curl around less tightly" than circles of radii $r_0(M)$.

Def

The minimum branch separation $s_0 = s_0(M)$ is

$$s_0 := \sup \left\{ s \mid d_M(x, y) \leq \pi r_0 \right. \\ \left. \forall x, y \in M \text{ w/ } \|x - y\| \leq s \right\}$$

Rmk.:

Existence and positivity of r_0 and s_0 are guaranteed by compactness of M .

(For s_0 , we may need Lebesgue number lemma).

Now, we use r_0 and s_0 to give a technical lemma.

Lemma 3

γ : a geod. in M from x to y .

$$l := L(\gamma) \leq \pi r_0.$$

$$\text{Then } 2r_0 \sin(l/2r_0) \leq \|x - y\| \leq l.$$

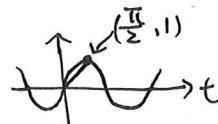
Rmk.:

By calculus, $\sin(t) \geq t - t^3/6, \forall t \geq 0$.

Thus, we can weaken the ineq. as:

$$\left(1 - \frac{l^2}{24r_0^2} \right) l \leq \|x - y\| \leq l.$$

Thus, for small l , $l \approx \|x - y\|$.



② For $l \leq \pi r_0$, $l/2r_0 \leq \pi/2$.

Treating $t := l/2r_0$, $t \in (0, \pi/2)$ and $\sin t \leq \frac{2}{\pi} t$.

$$\Rightarrow 2r_0 \sin(l/2r_0) = 2 \cdot \frac{l}{2t} \sin t \geq 2 \cdot \frac{l}{2t} \cdot \frac{2}{\pi} t$$

Thus, we may weaken the ineq. as $= (2/\pi) l$.

$$(2/\pi) l \leq \|x - y\| \leq l.$$

(Proof of L3 is postponed in the appendix).

Cor 4 (Lemma for $ds \approx d_G$)

$$\lambda > 0.$$

$$x_i, x_{i+1} \in M.$$

Assume (i) $\|x_i - x_{i+1}\| \leq s_0$,

(ii) $\|x_i - x_{i+1}\| \leq (2/\pi) r_0 \sqrt{24\lambda}$, and

(iii) \exists geod. on M of length $d_M(x_i, x_{i+1})$ connecting x_i and x_{i+1} .

Then

$$(1 - \lambda) d_M(x_i, x_{i+1}) \leq \|x_i - x_{i+1}\| \leq d_M(x_i, x_{i+1}).$$

<Pf>

By (i) and def of s_0 , $d_M(x_i, x_{i+1}) \leq \pi r_0$.

By (iii), we can apply L3.

By ② in the Rmk above, $(2/\pi) l \leq \|x_i - x_{i+1}\|$.

By (ii), we thus have $l \leq r_0 \sqrt{24\lambda}$.

$\Rightarrow 1 - \lambda \leq 1 - l^2/24r_0^2$. By Rmk ① above, we are done. (#)

Def
 M is called geodesically convex if every $x, y \in M$ can be connected by a geod. of length $d_M(x, y)$.

Main Thm A.

M : cpt subset of \mathbb{R}^n . (possibly w/ boundary)

$\mathcal{X} = \{x_i\}$: finite set of data pts in M .

G : a graph on \mathcal{X} .

$0 < \lambda_1, \lambda_2 < 1$.

$\varepsilon_{\min}, \varepsilon_{\max} > 0, \delta > 0$.

Assume

(i) G contains all edges xy w/ $\|x - y\| \leq \varepsilon_{\min}$.

(ii) All edges of G have length $\|x - y\| \leq \varepsilon_{\max}$.

(iii) \mathcal{X} satisfies δ -sampling condition.

(iv) M is geod. convex.

and

(v) $\varepsilon_{\max} < S_0$, where S_0 : minimum branch sep. of M .

(vi) $\varepsilon_{\max} \leq (2/\pi) \kappa_0 \sqrt{24\lambda_1}$,

where κ_0 : minimum radius of curvature.

(vii) $\delta \leq \lambda_2 \varepsilon_{\min} / 4$. (hence $4\delta < \varepsilon_{\min}$).

Then

$$(1 - \lambda_1) d_M(x, y) \leq d_G(x, y) \leq (1 + \lambda_2) d_M(x, y),$$

$$\forall x, y \in \mathcal{X}.$$

<PF>

Condition (i), (iii), (vii) are for applying Thm 2.

Condition (ii), (v), (vi), (iv) are for applying Cor 4.

Thus, $d_M \leq d_S \leq (1 + 4\delta/\varepsilon_{\min}) d_M \leq (1 + \lambda_2) d_M$

and $(1 - \lambda_1) d_S \leq d_G \leq d_S$.

Hence, $(1 - \lambda_1) d_M \leq (1 - \lambda_1) d_S \leq d_G \leq d_S \leq (1 + \lambda_2) d_M$.

(Thm 2)

(Cor 4)

(#)