

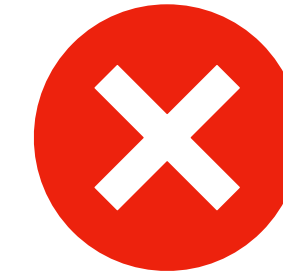
Week 8 — Sheet 7

Algorithms and Data Structures

13.11.2023 — Georg Hasebe





Debriefing of Submissions

```
A(n):  
  if memo[n] defined  
    return memo[n]  
  if n <= 4  
    return n  
  else  
    return A(n - 1) + A(n - 3) + 2A(n - 4)
```



```
A(n):  
  if memo[n] defined  
    return memo[n]  
  if n <= 4  
    return n  
  else  
    memo[n] <- A(n - 1) + A(n - 3) + 2A(n - 4)  
    return memo[n]
```



- Naming variable root can **decrease readability**, use `current_node` or `curr_node` or `node` etc. instead
- Initialization:
 - `Memo <- [1...n]`? 
 - `Memo[n] <- [1...n]`? 
 - `Memo[1...n]`? 
 - `Memo[1...n] <- [-1,...,-1]` etc. 
- Scope? Is `memo[n]` initialized globally? Or every recursion?

Suggestion: Read Solutions

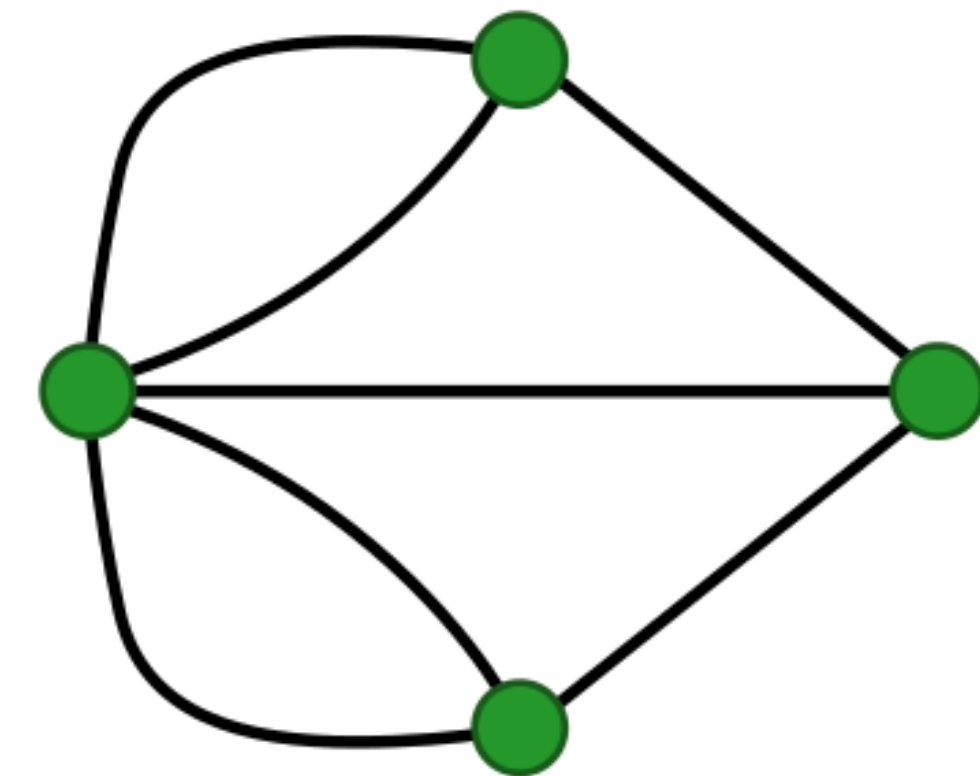
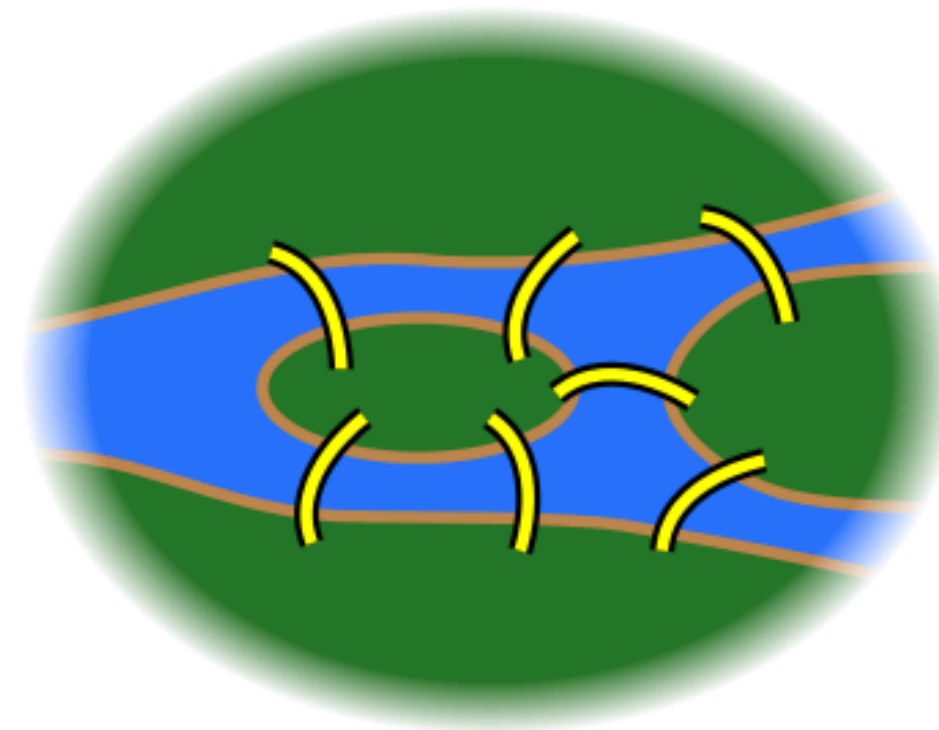
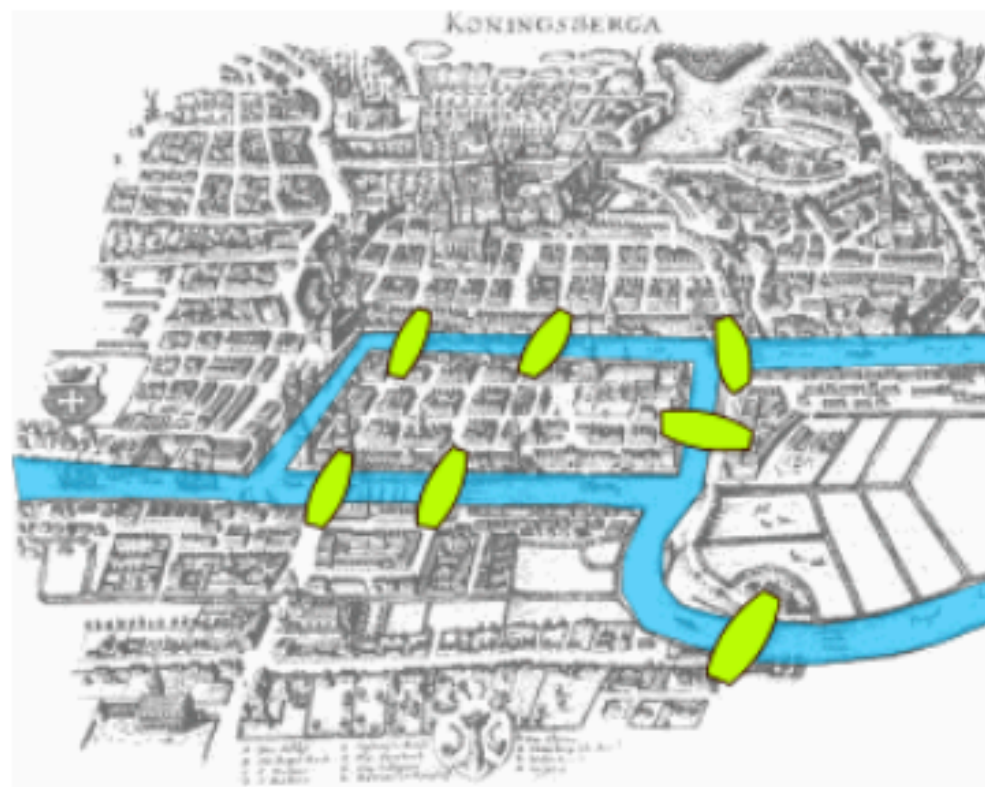
Exercise Sheet 7

Debriefing of Exercise Sheet 6

Theory Recap

Graph Theory

Graph Theory

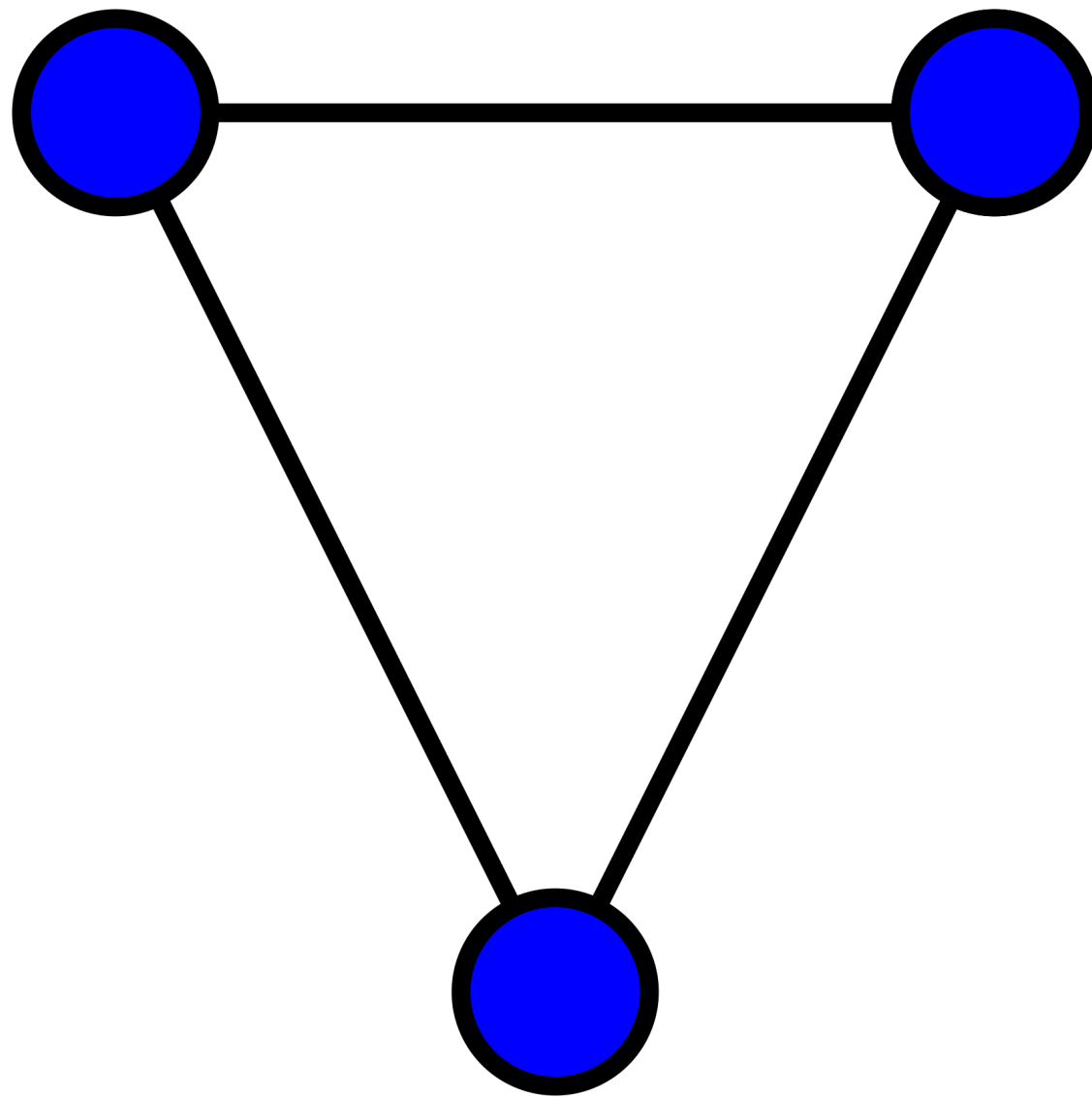


[1] https://upload.wikimedia.org/wikipedia/commons/6/60/Leonhard_Euler_2.jpg

[2] https://de.wikipedia.org/wiki/Königsberger_Brückenproblem

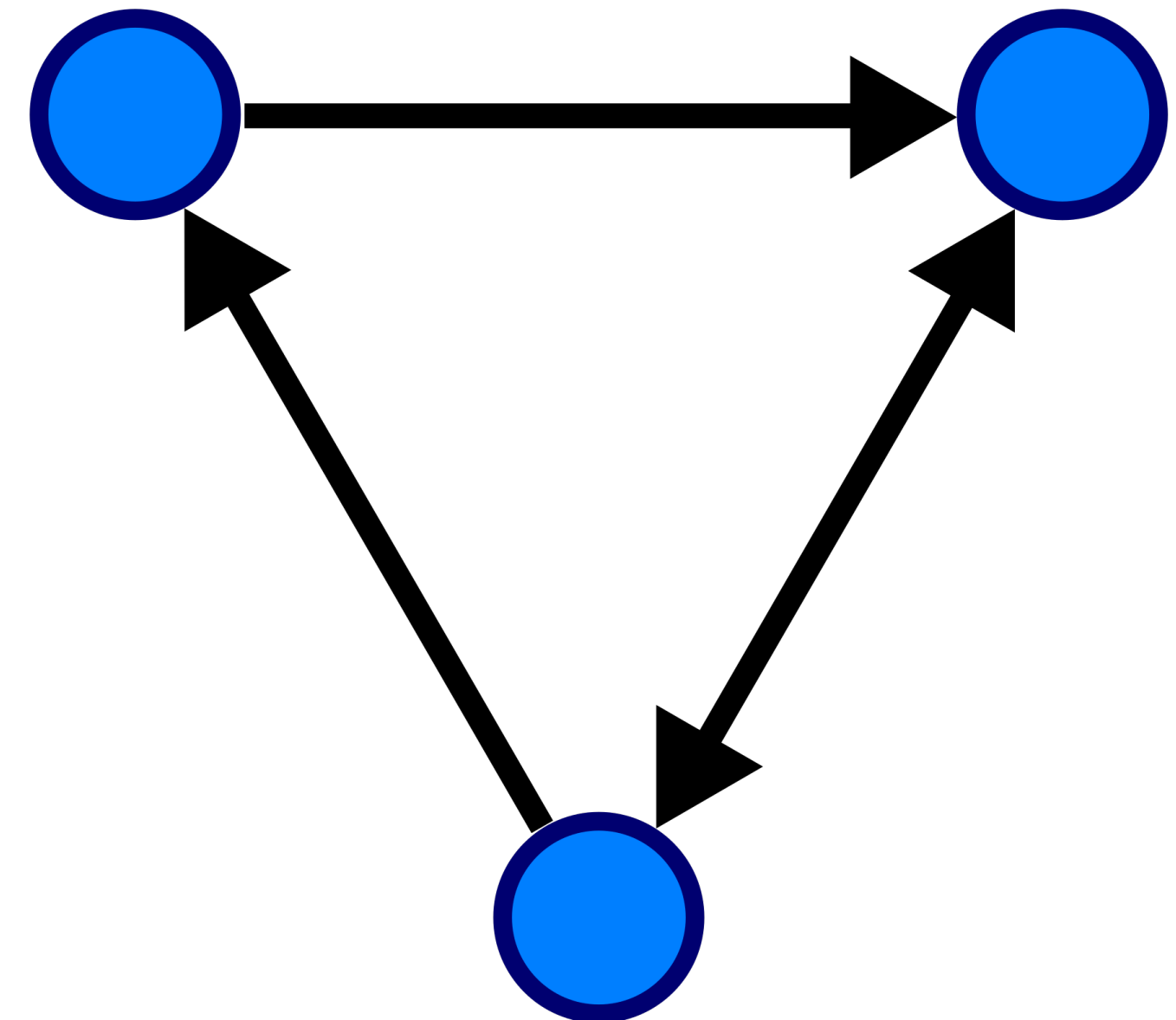
Graph Theory

- Undirected Graphs
- V a set of vertices
- $E \subseteq \{\{x, y\} \mid x, y \in V \text{ and } x \neq y\}$



https://en.wikipedia.org/wiki/Graph_theory#/media/File:Undirected.svg

- Directed Graphs
- V a set of vertices
- $E \subseteq \{(x, y) \mid (x, y) \in V^2 \text{ and } x \neq y\}$



https://en.wikipedia.org/wiki/Graph_theory#/media/File:Directed.svg

Eulerian trail (Eulerweg)

(Sometimes also Eulerian path in literature)

In **graph theory**, an **Eulerian trail** (or **Eulerian path**) is a **trail** in a finite graph that visits every **edge** exactly once (allowing for revisiting vertices).

A connected graph has an Euler trail. \Leftrightarrow Every vertex, except for at most 2, has even degree.

Eulerian circuit (Eulerkreis)

(Sometimes also Eulerian Cycle in literature)

An **Eulerian circuit** or **Eulerian cycle** is an Eulerian trail that starts and ends on the same **vertex**.

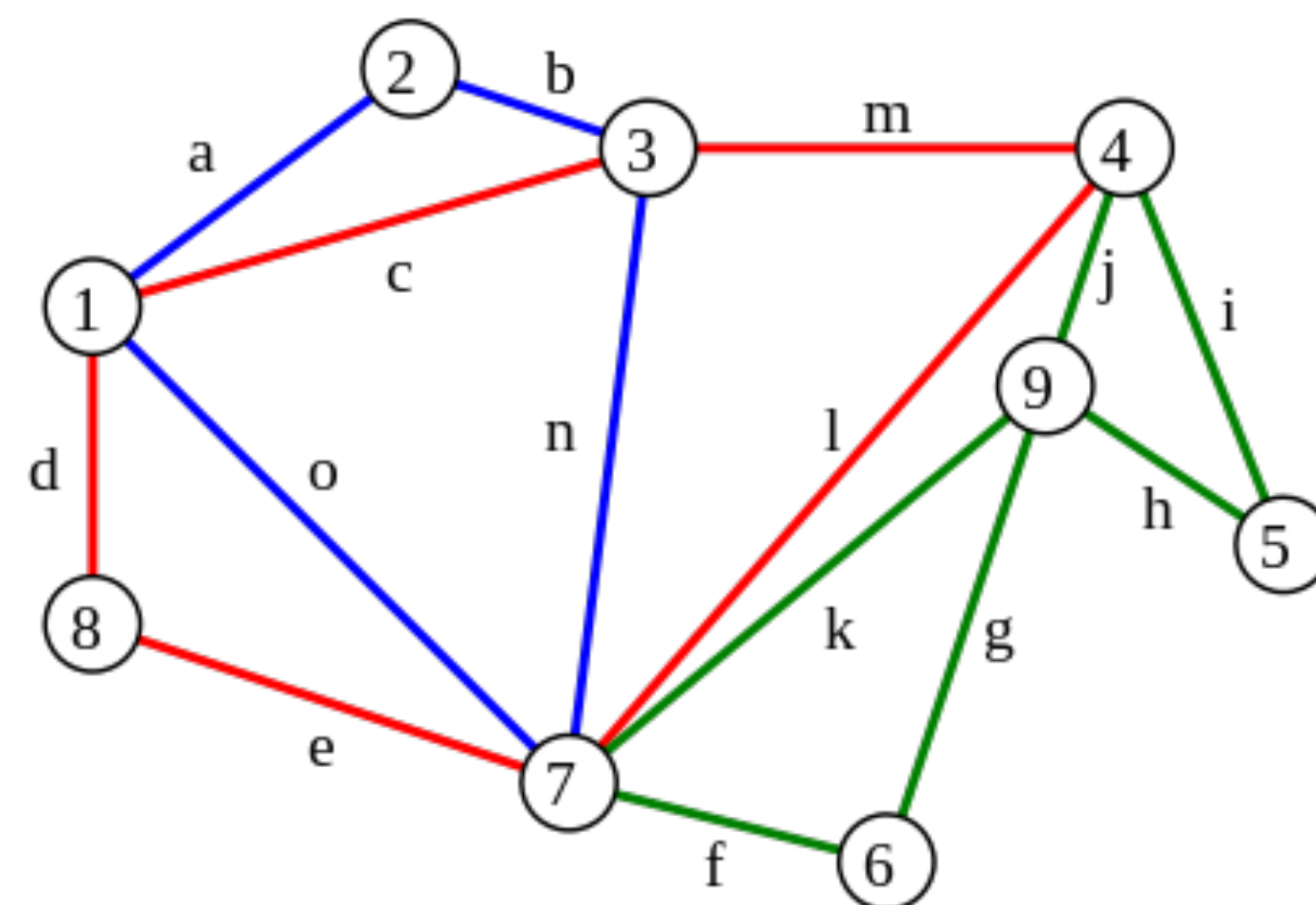
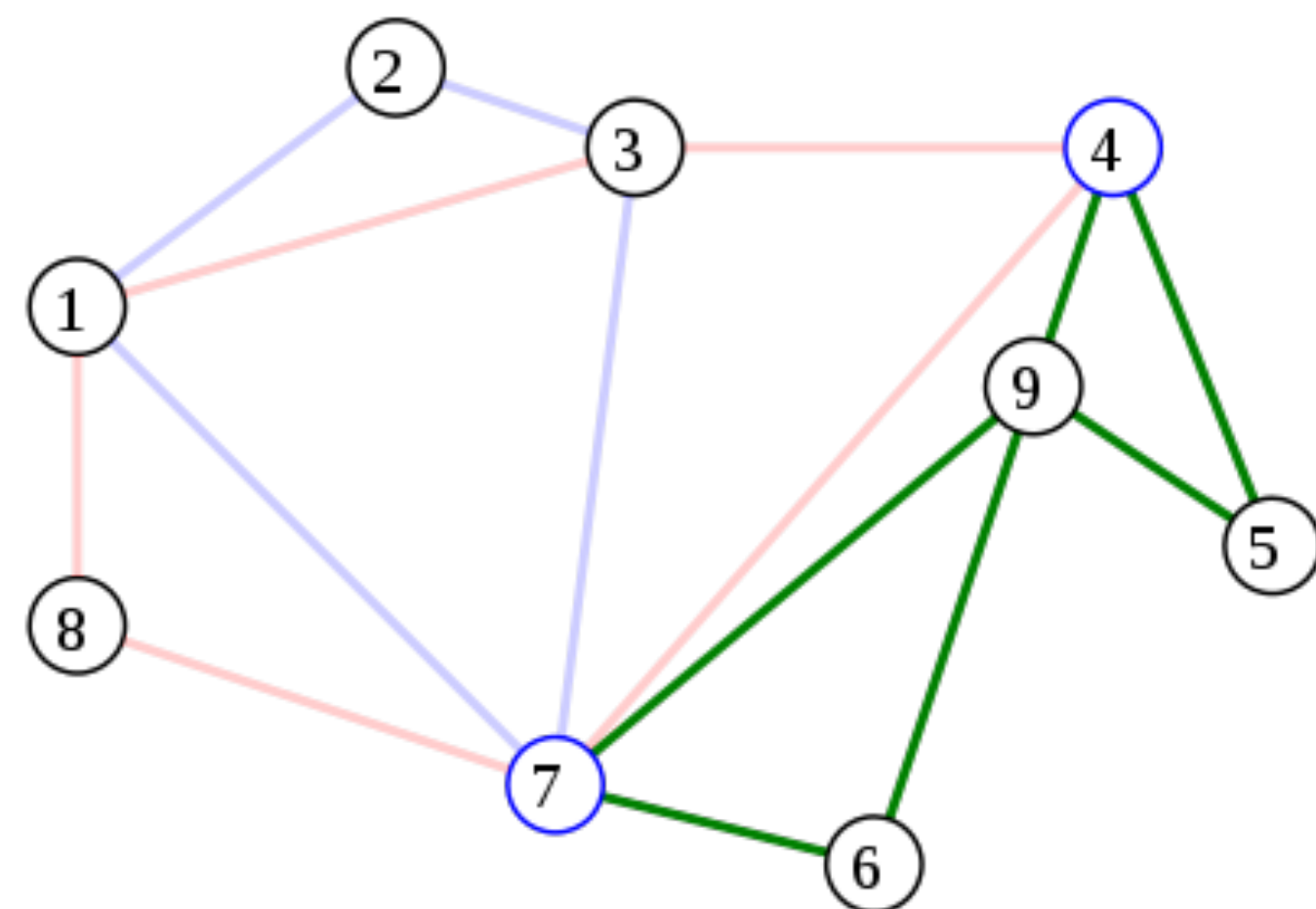
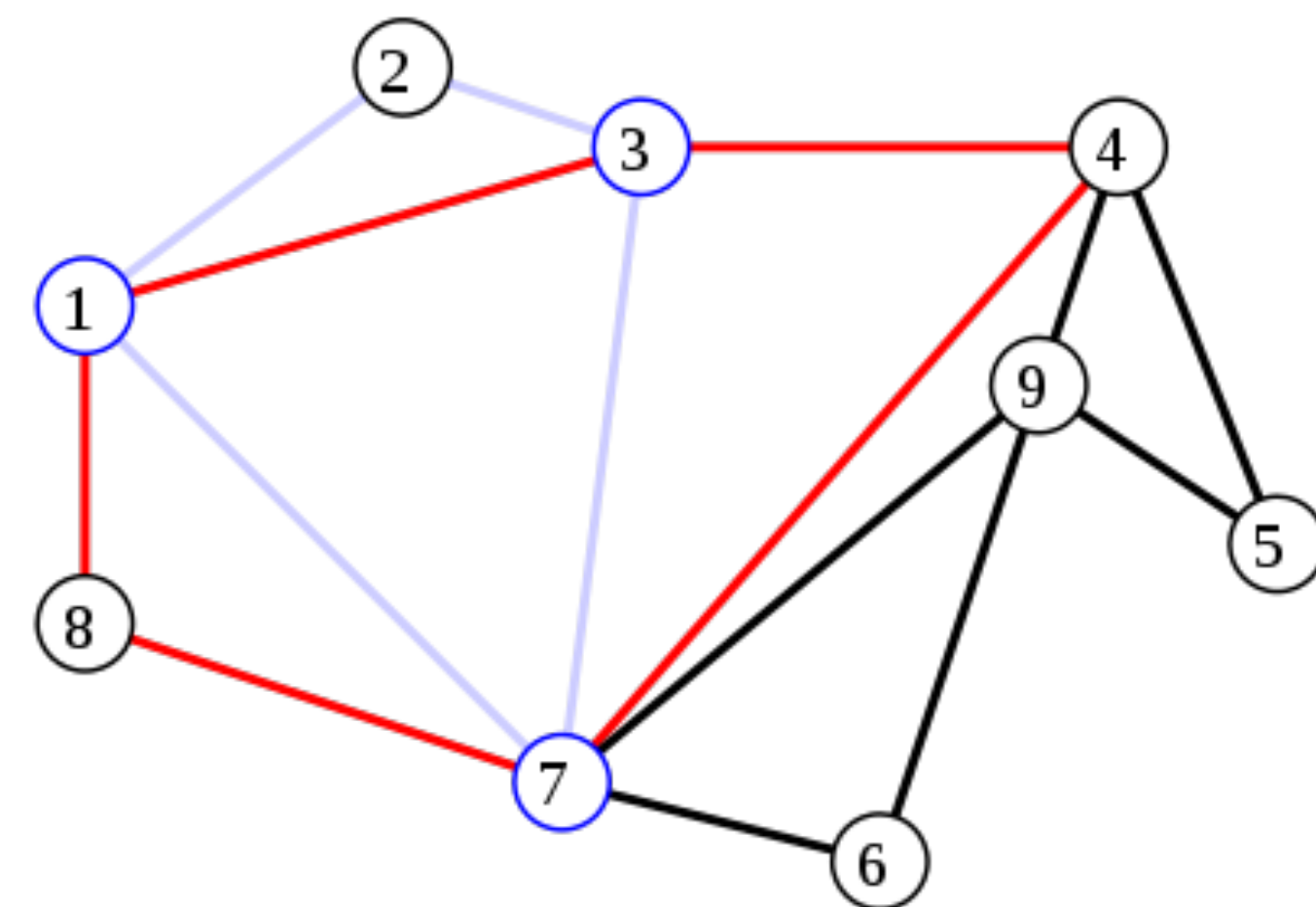
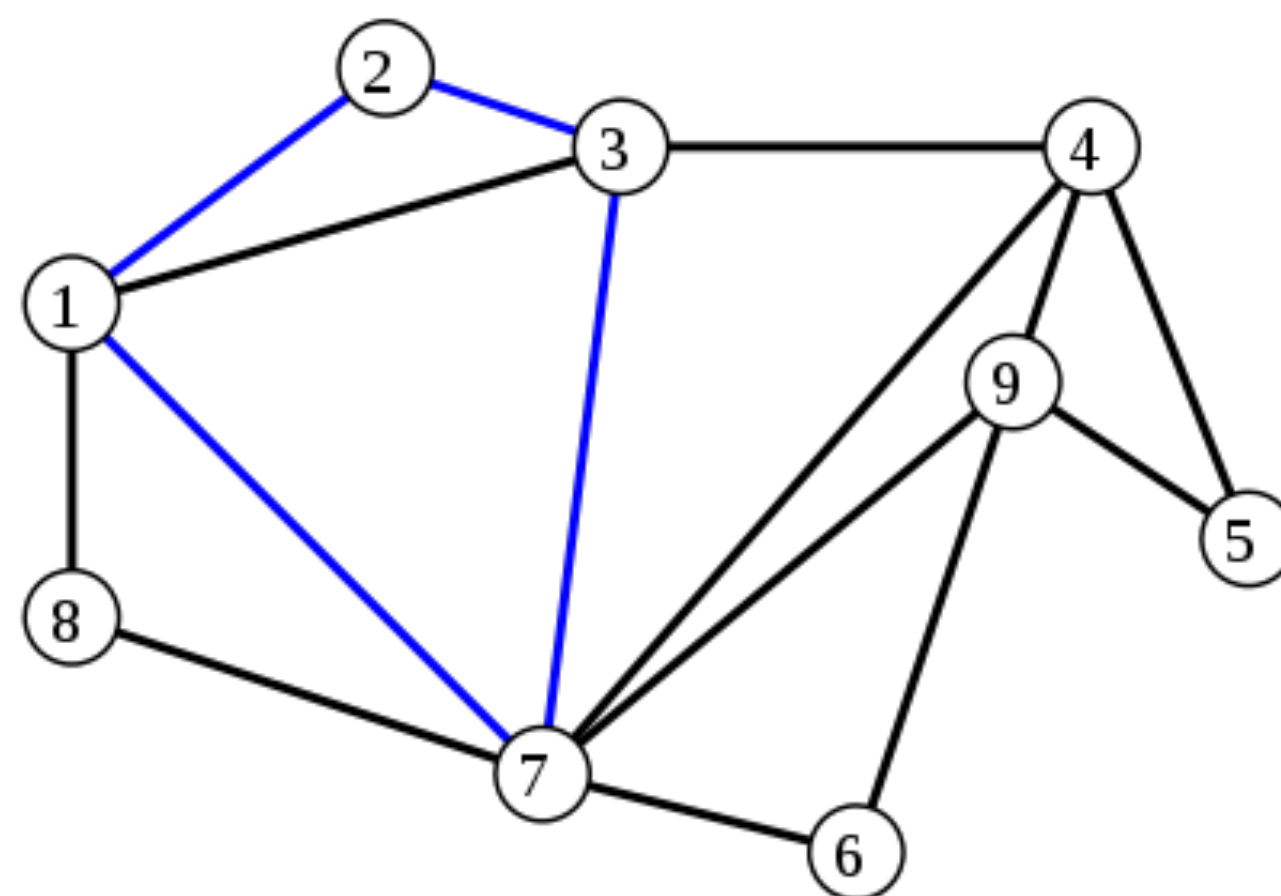
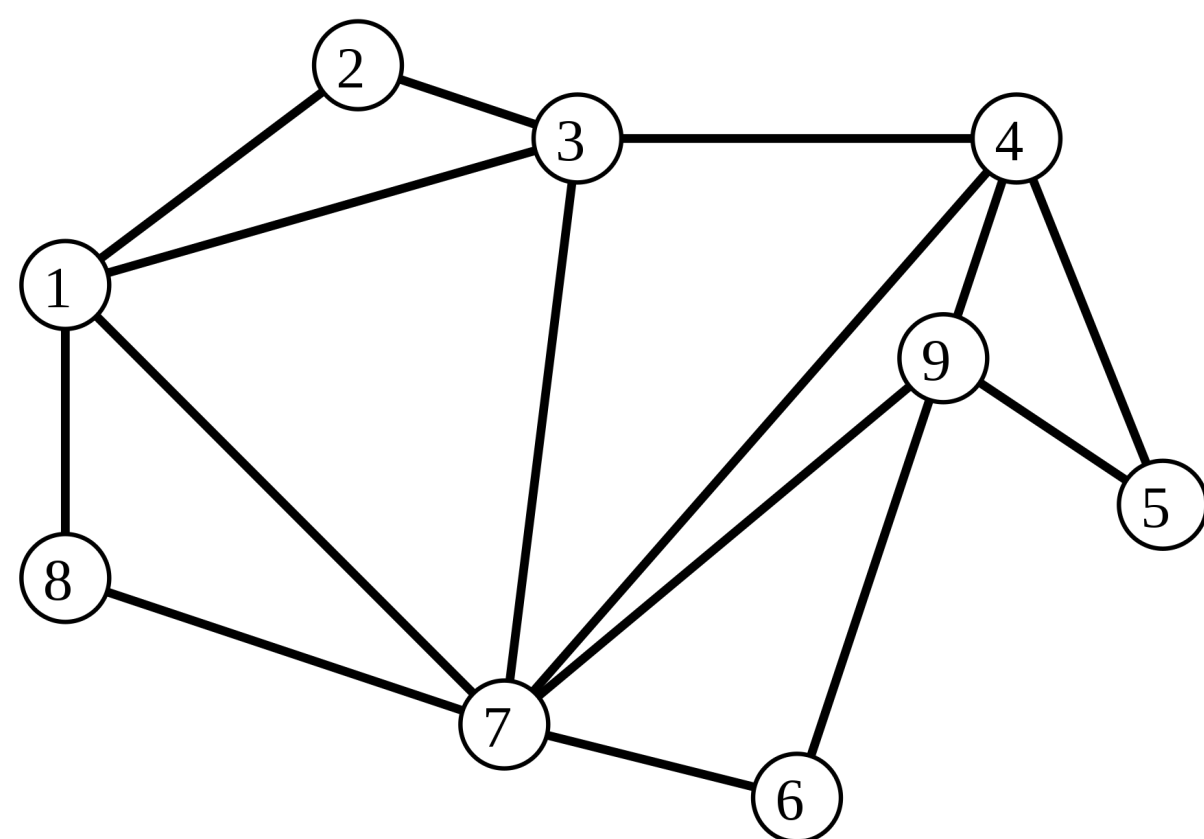
A connected graph has an Euler cycle. \Leftrightarrow Every vertex has even degree.

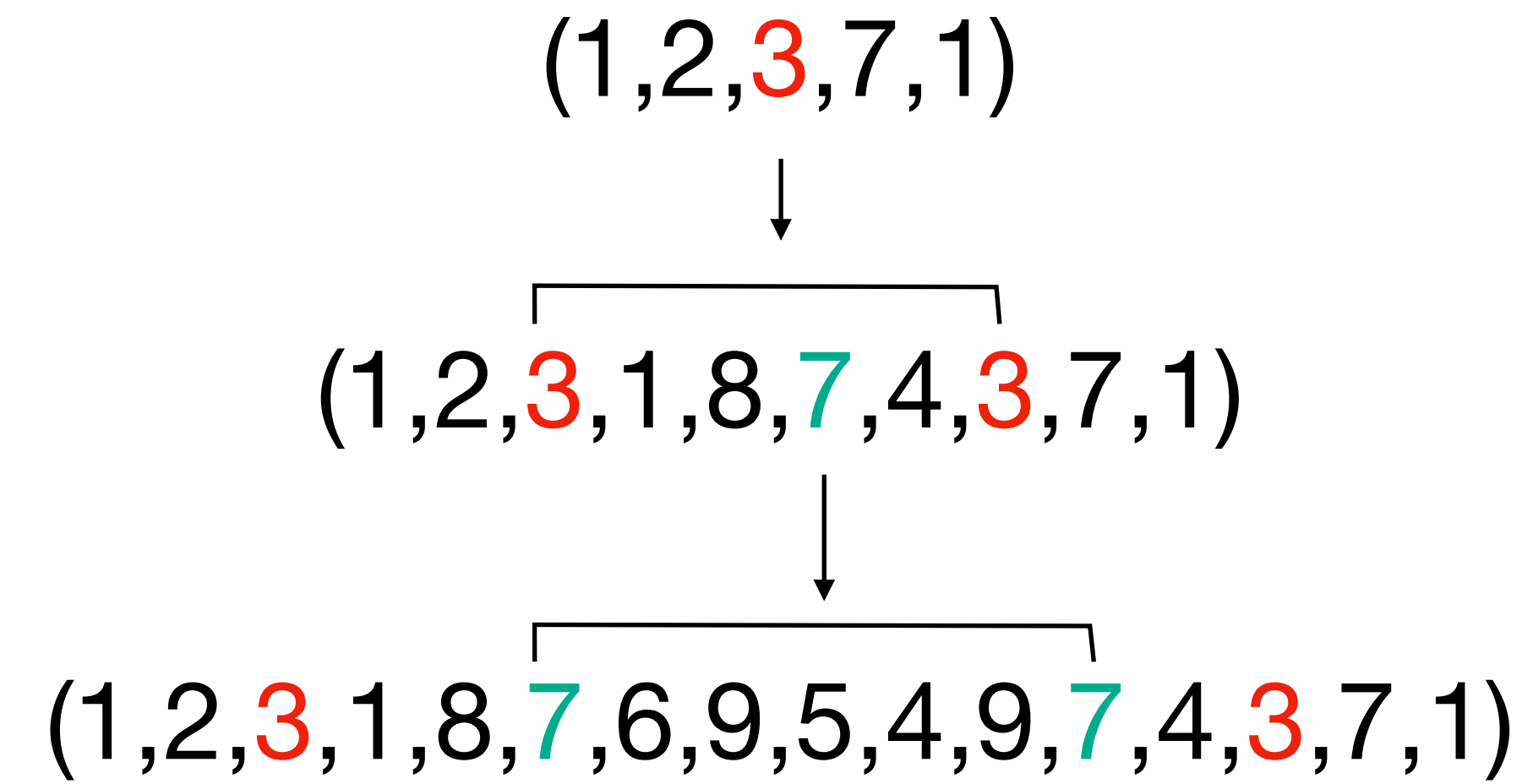
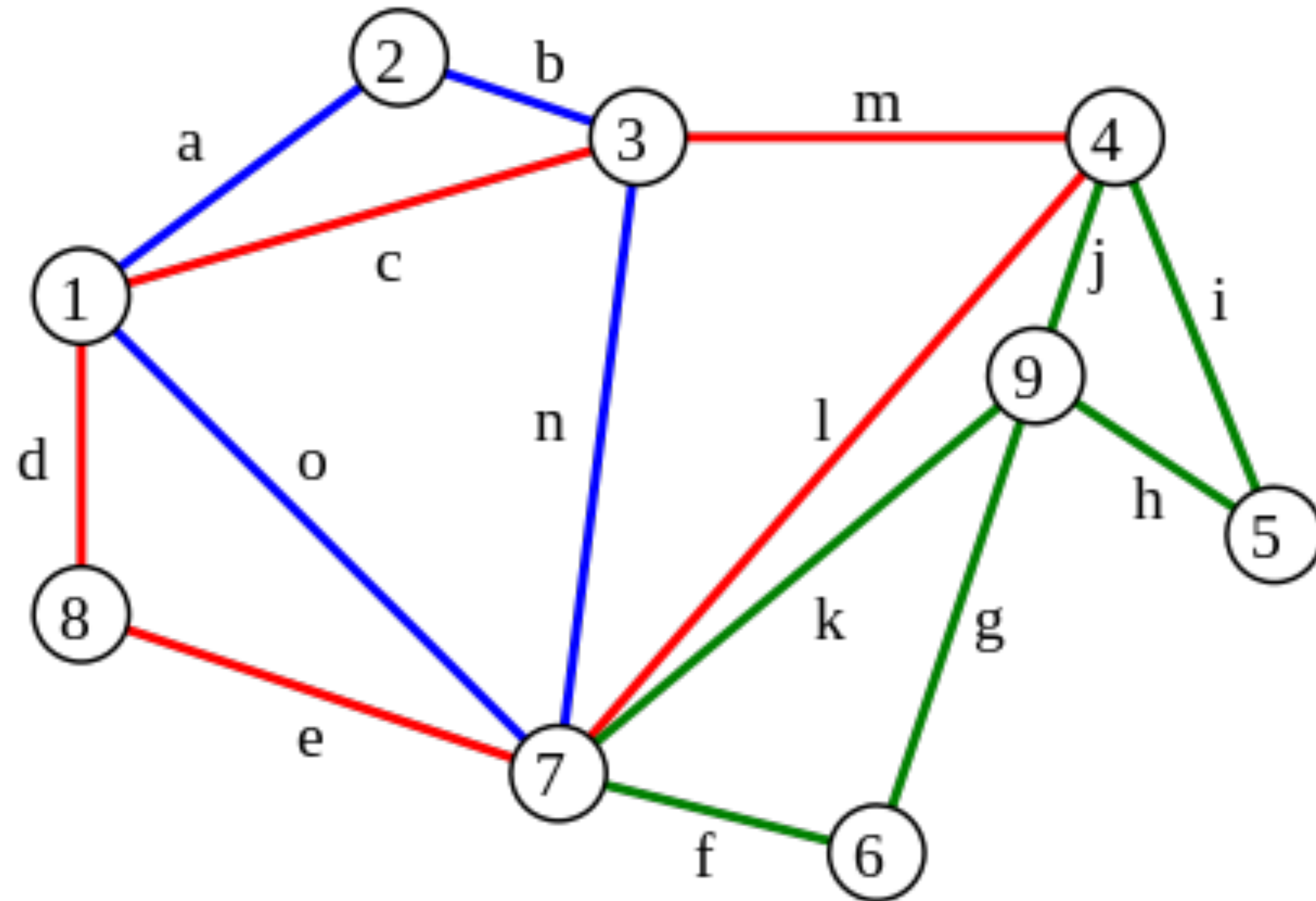
Hierholzer's Algorithm in $O(|E|)$

Voraussetzung: Sei $G=(V,E)$ ein zusammenhängender Graph, der nur Knoten mit geradem Grad aufweist.

1. Wähle einen beliebigen Knoten v_0 des Graphen und konstruiere von v_0 ausgehend einen Unterkreis K in G , der keine Kante in G zweimal durchläuft.
2. Wenn K ein Eulerkreis ist, breche ab. Andernfalls:
3. Vernachlässige nun alle Kanten des Unterkreises K .
4. Am ersten Eckpunkt von K , dessen Grad größer 0 ist, lässt man nun einen weiteren Unterkreis K' entstehen, der keine Kante in K durchläuft und keine Kante in G zweimal enthält.
5. Füge in K den zweiten Kreis K' ein, indem der Startpunkt von K' durch alle Punkte von K' in der richtigen Reihenfolge ersetzt wird.
6. Nenne jetzt den so erhaltenen Kreis K und fahre bei Schritt 2 fort.

Hierholzer's Algorithm Example





$$E = (1, 2, 3, 1, 8, 7, 6, 9, 5, 4, 9, 7, 4, 3, 7, 1)$$

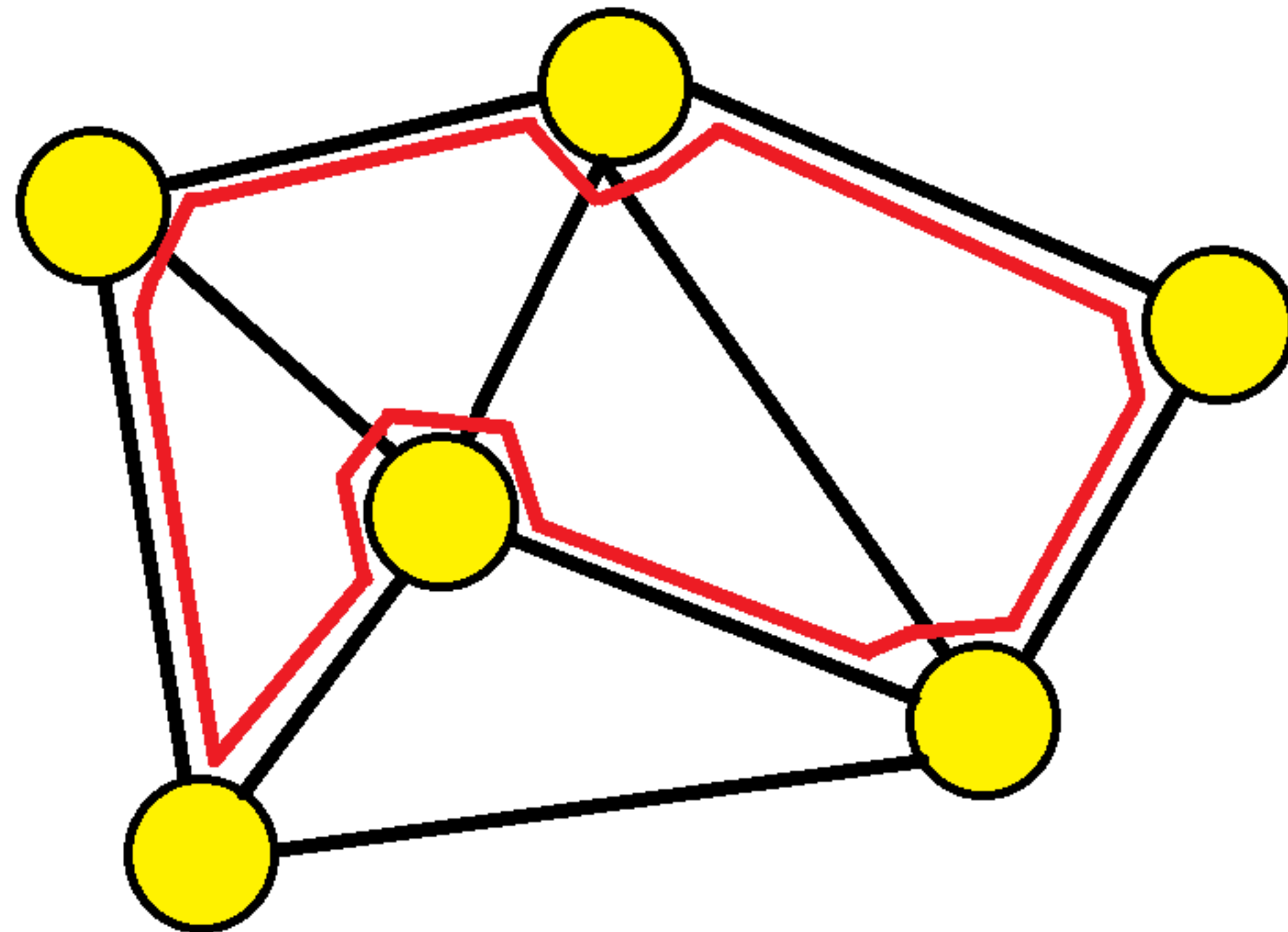
$$C_{\text{blau}} = (1, 2, 3, 7, 1)$$

$$C_{\text{red}} = (3, 1, 8, 7, 4, 3)$$

$$C_{\text{green}} = (7, 6, 9, 5, 4, 9, 7)$$

Hamiltonian path and Hamiltonian cycle (circuit)

In the [mathematical](#) field of [graph theory](#), a **Hamiltonian path** (or **traceable path**) is a [path](#) in an undirected or directed graph that visits each [vertex](#) exactly once. A **Hamiltonian cycle** (or **Hamiltonian circuit**) is a [cycle](#) that visits each vertex exactly once.



Hamiltonian path and Hamiltonian cycle (circuit)

**Determining if such paths or
cycles exist is NP-Complete
(very hard).**

Graph Terminology

Sadly sometimes ambiguous and confusing

- Julian Steinmann's Graph Terminology Cheatsheet
 - <https://exams.vis.ethz.ch/user/jsteinmann/document/graph-terminology-cheatsheet>
- Wikipedia's Glossary of Graph Theory
 - https://en.wikipedia.org/wiki/Glossary_of_graph_theory#trail

Kahoot