

# "Zero-Knowledge Proofs: Privacy in Blockchain & Cloud Computing"

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## Abstract

Zero-Knowledge Proofs (ZKPs) have gained significant attention for their ability to ensure privacy in cryptographic systems by allowing one party (the prover) to prove to another party (the verifier) that they know a secret without revealing any information about the secret itself. This report investigates the application of ZKPs in blockchain and cloud computing, focusing on privacy-preserving protocols such as zk-SNARKs and zk-STARKs, which enable secure transactions and identity verification. A key aspect of the report is the implementation of a simple ZKP-based secure voting system, where voters cast their ballots without disclosing their choices. The system integrates ZoKrates for generating ZKPs, Solidity smart contracts for blockchain interaction, and a web interface for user interaction. The report elaborates on the ZKP logic used, including zk-SNARKs' trusted setup and proof generation, and details the challenges and advantages of implementing such a system in a real-world scenario. It further explores the potential impact of ZKPs on privacy, scalability, and security in blockchain-based applications and decentralized systems. Moreover, the report addresses the scalability and transparency advantages of zk-STARKs, as well as their post-quantum security features, making them a future-proof solution for blockchain and cloud computing. Finally, it highlights the ongoing research and the future trajectory of ZKPs in enhancing digital privacy and enabling secure collaborations.

**Keywords:** Zero-Knowledge Proofs, zk-SNARKs, zk-STARKs, Blockchain, Privacy-Preserving, Digital Identity, Secure Voting, Privacy-Enhancing Technologies, Cloud Computing, Post-Quantum Security, Decentralized Systems.

## 1 Introduction

Zero-Knowledge Proofs (ZKPs) are cryptographic techniques that allow a prover to demonstrate knowledge of a secret to a verifier without revealing the secret itself. ZKPs are essential for privacy-preserving protocols, addressing secure transactions, data verification, and identity assurance, particularly in decentralized systems like blockchain and cloud computing. They ensure that the verifier only learns the fact that the prover knows the secret, offering a robust solution to privacy concerns.

In blockchain and cloud environments, privacy is a critical challenge due to the need for transparency and verifiable operations

without exposing sensitive data. Traditional cryptographic methods like RSA provide security but can compromise privacy or efficiency. ZKPs offer an efficient alternative, enabling verification of transactions without disclosing the underlying data. This makes them ideal for private transactions, identity management, and voting systems, where confidentiality is paramount.

This report explores the application of ZKPs in enhancing privacy within blockchain and cloud systems, covering the principles of ZKPs and advanced protocols like zk-SNARKs and zk-STARKs. These protocols overcome the limitations of traditional cryptographic systems, offering succinct, non-interactive proofs that are both scalable and verifiable in milliseconds, making them suitable for privacy-preserving applications.

A key focus of this report is the implementation of a ZKP-based secure voting system as a case study. Using zk-SNARKs and the ZoKrates tool, along with Solidity smart contracts on the Ethereum blockchain, this system ensures that votes remain private, demonstrating the practical application of ZKPs for privacy in voting.

Additionally, this report traces the evolution of ZKPs from their origins in the 1980s with Goldwasser, Micali, and Rackoff's interactive proof systems to more advanced techniques like zk-SNARKs and zk-STARKs, which offer scalability and quantum-resistance.

The report also discusses the state-of-the-art tools and techniques in ZKPs, including Marlin, Groth16, ZoKrates, zk-STARKs, PlonK, and others, highlighting their use in privacy-preserving systems.

## 2 Research Methodology

This research employs Zero-Knowledge Proofs (ZKPs) for blockchain and cloud computing. Papers are selected based on specific criteria such as relevance to ZKPs, blockchain applications, citation counts, and recent advancements. These papers provide the foundation for the design and evaluation of the voting system in terms of security, privacy, scalability, and efficiency.

The selection criteria for the papers are outlined below:

The selected papers influenced the design of the ZKP-based system using zk-SNARKs for efficient proof generation and verification. Performance analysis ensures scalability for large-scale deployments.

Criteria	Description
ZKP Focus	zk-SNARKs, zk-STARKs, etc.
Blockchain	ZKP-based blockchain voting systems.
Source	Peer-reviewed journals/conferences.
Citations	High citation count for significance (> 100).
Privacy	Voter anonymity preservation.
Performance	Scalability and efficiency.
Advancements	Recent updates in ZKPs.

**Table 1: Paper Selection Criteria**

### 3 Literature Review

The concept of Zero-Knowledge Proofs (ZKPs) has revolutionized the field of cryptography by enabling a party (the prover) to convince another party (the verifier) of the validity of a statement without disclosing any information beyond the statement’s truth. This section reviews the key papers that have shaped the development of ZKPs from their theoretical introduction to their practical applications in blockchain and cloud computing.

#### 3.1 Early Foundations of Zero-Knowledge Proofs

The journey of Zero-Knowledge Proofs began in 1985 with the seminal work by Goldwasser, Micali, and Rackoff [1], which introduced interactive proof systems. They showed that a prover could convince a verifier of a statement’s truth without revealing anything about the statement itself. This concept was revolutionary as it demonstrated that cryptographic protocols could provide strong security guarantees without exposing any sensitive information. The work laid the foundation for many privacy-preserving cryptographic protocols.

A few years later, in 1989, Goldreich, Micali, and Wigderson [2] extended the concept of ZKPs to NP-complete problems, proving that every language in NP has an associated Zero-Knowledge Proof. This result significantly expanded the applicability of ZKPs, providing a theoretical framework for a broad range of secure applications, from private authentication to secure voting systems. Their work established the basis for future research into more efficient and scalable ZKPs.

#### 3.2 Advancements in Zero-Knowledge Proofs for Cryptography

In 2013, the development of zk-SNARKs (Succinct Non-Interactive Zero-Knowledge Arguments of Knowledge) by Ben-Sasson et al. [3] marked a major advancement in the field. zk-SNARKs are non-interactive proofs that can be verified in constant time, regardless of the size of the statement. This made ZKPs much more practical for use in decentralized systems like blockchain, where verifying large datasets quickly is crucial. zk-SNARKs allow for compact proofs, significantly improving the efficiency of cryptographic protocols in real-world applications.

In 2018, Ben-Sasson, Bentov, Yarkoni, Tromer, and Virza [4] introduced zk-STARKs (Zero-Knowledge Scalable Transparent Arguments of Knowledge), which further improved upon zk-SNARKs

by eliminating the need for a trusted setup. This was a major breakthrough, as zk-SNARKs required a trusted setup phase, which posed a potential security risk. zk-STARKs, by contrast, offer a transparent, scalable, and quantum-resistant approach to Zero-Knowledge Proofs. This makes zk-STARKs particularly relevant in the context of future-proof systems, including post-quantum cryptography.

#### 3.3 Zero-Knowledge Proofs in Blockchain and Cloud Applications

The application of ZKPs in blockchain systems has been particularly transformative, as it allows for privacy-preserving transactions without compromising the security or integrity of the blockchain. In 2020, Li, Zhang, and Wu [5] proposed a Zero-Knowledge-Proof-based digital identity management scheme on blockchain. Their system leverages the privacy-preserving capabilities of ZKPs to allow users to maintain control over their personal information while engaging in secure transactions on a blockchain. This application of ZKPs addresses the privacy concerns inherent in traditional digital identity management systems, where personal information is often centralized and vulnerable to breaches.

In 2024, Wang et al. [6] reviewed the advancements and challenges of using ZKPs in blockchain-based identity sharing. Their survey highlighted the potential of ZKPs to improve the privacy and scalability of decentralized identity systems, where users can share their identity in a secure and verifiable manner without disclosing sensitive information. They also discussed the challenges that remain, particularly with respect to the computational complexity of ZKPs in real-time applications. This paper outlined a roadmap for overcoming these challenges, emphasizing the need for more efficient protocols and better integration with existing blockchain infrastructures.

In summary, the literature on Zero-Knowledge Proofs illustrates their evolution from theoretical constructs to powerful tools in cryptographic applications, particularly in blockchain and cloud computing. These advancements in ZKPs are driving the future of privacy-preserving technologies and will continue to play a crucial role in the development of secure, decentralized systems.

Zero-Knowledge Proofs (ZKPs) can be broadly categorized based on the interaction between the prover and the verifier. These classifications influence how the proof is constructed and verified in different computational settings.

Each type is suited to specific applications depending on the communication model and privacy requirements. The two types of ZKPs are:

- (1) **Interactive ZKPs:** These require multiple rounds of communication between the prover and verifier. The prover convinces the verifier of the validity of a statement through a sequence of challenges and responses.
- (2) **Non-Interactive ZKPs (NIZKs):** These require only a single message from the prover to the verifier. Using techniques like the Fiat-Shamir heuristic, interaction is removed, making them ideal for blockchain and distributed environments.

Year	Contribution
1985	ZKPs: Interactive proof systems
1989	ZKPs for NP: NP-complete problems
2013	zk-SNARKs: Succinct, non-interactive proofs
2018	zk-STARKs: Scalable, transparent proofs
2020	ZKP-based Digital Identity Management
2024	ZKP-based Privacy-Preserving Data Outsourcing

**Table 2: Key Contributions in Zero-Knowledge Proofs**

## 4 Privacy-Enhancing Techniques

Zero-Knowledge Proofs (ZKPs) are cryptographic protocols that allow a prover to demonstrate the knowledge of a solution to a statement without revealing the solution itself. Two prominent types of ZKPs are zk-SNARKs (Zero-Knowledge Succinct Non-Interactive Arguments of Knowledge) and zk-STARKs (Zero-Knowledge Scalable Transparent Arguments of Knowledge). Both offer privacy-preserving techniques, but differ in their cryptographic foundations, setup requirements, and scalability. To illustrate how these techniques work, we will use the example equation  $x^2 - 4x + 4 = 0$ , which simplifies to  $(x - 2)^2 = 0$ , with the solution  $x = 2$ .

### 4.1 zk-SNARKs (Zero-Knowledge Succinct Non-Interactive Arguments of Knowledge)

zk-SNARKs allow a prover to demonstrate that they know a solution to a mathematical problem without revealing the actual solution. In the context of the equation:

$$x^2 - 4x + 4 = 0 \quad \text{or} \quad (x - 2)^2 = 0$$

The solution to this equation is  $x = 2$ , which is the secret that the prover knows. Here's how zk-SNARKs would prove the prover knows this solution.

#### 4.1.1 zk-SNARKs Mathematical Workflow

- **Commitment to the Polynomial:** The prover constructs a commitment to the polynomial equation  $P(x) = x^2 - 4x + 4$ . This is done using a cryptographic function like elliptic curve cryptography (ECC), which binds the prover to the polynomial without revealing its coefficients or the solution.
- **Prover's Proof:** The prover generates a proof using cryptographic techniques (e.g., elliptic curve pairings) that they know the solution  $x = 2$ , without revealing the value of  $x$ . The proof consists of:
  - A *witness*, which is the secret value  $x = 2$ .
  - A *polynomial commitment* that proves the equation holds when  $x = 2$ .
- **Verification:** The verifier checks the proof by verifying the commitment's consistency with the equation, without knowing the actual value of  $x$ .
- **Succinctness:** zk-SNARKs are succinct, meaning the proof is very small and easy to verify. The verifier can check the proof without understanding the full polynomial or the actual solution.

#### 4.1.2 Mathematics of zk-SNARKs

The cryptographic setup for zk-SNARKs involves elliptic curve pairings. The prover generates a proof using the polynomial commitment and elliptic curve points. The verifier checks the proof by performing group operations on these elliptic curve points to verify the consistency of the proof with the original statement.

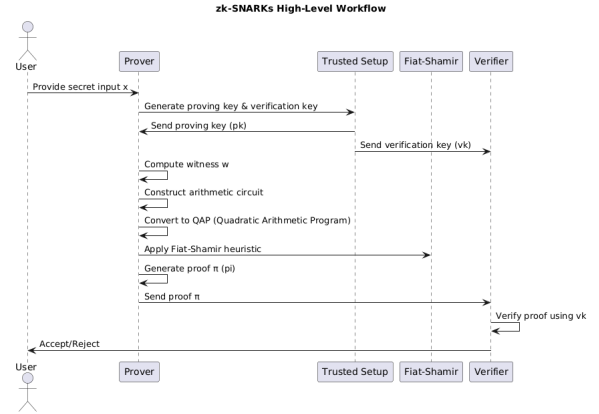
The key equation for verification is:

$$e(P, C) = e(A, P) \cdot e(B, Q)$$

Where:

- $P$  is the polynomial commitment.
- $C$  is the commitment to the prover's solution.
- $A, B$ , and  $Q$  are points on the elliptic curve.
- $e$  is the elliptic curve pairing function.

This equation ensures that the proof provided by the prover is consistent with the commitment, thus proving that the prover knows the solution without revealing it.



**Figure 1: High-level Working of zk-SNARKs**

### 4.2 zk-STARKs (Zero-Knowledge Scalable Transparent Arguments of Knowledge)

zk-STARKs offer a more scalable and transparent approach to Zero-Knowledge Proofs compared to zk-SNARKs. Unlike zk-SNARKs, zk-STARKs use cryptographic hash functions instead of elliptic curve cryptography, making them more transparent and scalable. Let's explain zk-STARKs using the same quadratic equation  $x^2 - 4x + 4 = 0$ .

#### 4.2.1 zk-STARKs Mathematical Workflow

- **Commitment to the Equation:** In zk-STARKs, the prover generates a commitment to the equation  $P(x) = x^2 - 4x + 4$  using a cryptographic hash function (like a Merkle hash tree). This commitment binds the prover to the equation and the secret solution  $x = 2$ , ensuring no manipulation of the equation.
- **Interactive Proof (Optional):** While zk-SNARKs are non-interactive, zk-STARKs can involve a series of challenges

and responses between the prover and verifier. This interaction is optional but allows zk-STARKs to be used for more complex statements involving large datasets.

- **Transparency:** zk-STARKs provide transparency because all information required to verify the proof (such as the hash commitments) is publicly available. This eliminates the need for a trusted setup phase, making zk-STARKs inherently more secure in decentralized environments.
- **Verification:** The verifier checks the proof by performing a series of hash computations to ensure that the prover's commitment to the polynomial is consistent with the solution, without needing to know the actual solution.

#### 4.2.2 Mathematics of zk-STARKs

In zk-STARKs, the proof verification involves checking whether the prover's commitment matches the expected values derived from hash functions. This typically involves Merkle trees, which are hash trees that allow efficient verification of large datasets.

The verifier checks the following equation, which is based on hash function consistency:

$$H(P) = H(Q) \cdot H(R)$$

Where:

- $H$  is the cryptographic hash function.
- $P$ ,  $Q$ , and  $R$  are hash values derived from the prover's polynomial and the solution commitment.

Since zk-STARKs rely on hash functions instead of elliptic curve pairings, they provide a more scalable and transparent system for proving knowledge of the solution, without needing a trusted setup phase.

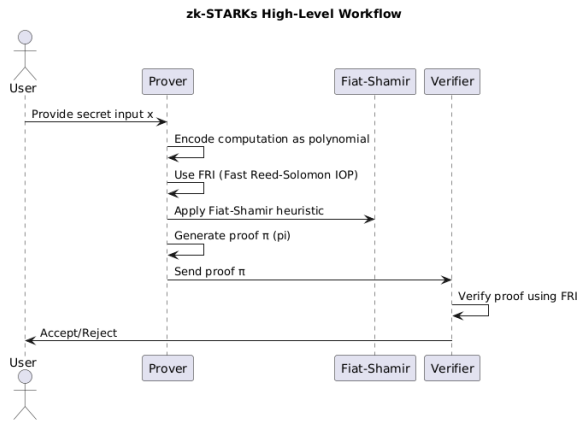


Figure 2: High-level Working of zk-STARKs

## 5 Problems

While zk-SNARKs and zk-STARKs offer significant privacy enhancements, there are still challenges. One of the key issues is the computational overhead involved in generating and verifying proofs, especially for complex statements or large datasets. Additionally, the trusted setup in zk-SNARKs can be a potential vulnerability, as

compromising the setup can undermine the entire system's security. zk-STARKs address some of these issues by removing the need for a trusted setup, but they come with larger proof sizes, which may make them less practical for certain applications.

## 6 New Development

Recent developments have focused on improving the scalability, efficiency, and transparency of ZKP systems. zk-STARKs have become increasingly popular for large-scale applications, such as blockchain and cloud computing, due to their scalability and transparency. In contrast, zk-SNARKs continue to be favored for applications requiring small, succinct proofs with quick verification times, such as privacy-preserving cryptocurrencies. Researchers are also exploring hybrid approaches that combine the strengths of both zk-SNARKs and zk-STARKs to address the limitations of each technique.

## 7 State-of-the-Art Tools and Techniques

This section provides a list of state-of-the-art tools and techniques used for zk-SNARKs and zk-STARKs, segregated into their respective categories.

Tool/Technique	Category
Marlin	zk-SNARKs, zk-STARKs
Groth16	zk-SNARKs
Bellman	zk-SNARKs
ZoKrates	zk-SNARKs
zk-STARKs (Native)	zk-STARKs
PlonK	zk-SNARKs, zk-STARKs
STARKy	zk-STARKs
Winterfell	zk-STARKs

### Tools Descriptions:

- **Marlin:** A framework for constructing both zk-SNARKs and zk-STARKs with high efficiency and modularity.
- **Groth16:** A widely used zk-SNARK construction, particularly for its compact proof size and efficient verification.
- **Bellman:** A Rust library for building zk-SNARKs, often used as a foundation for other cryptographic applications.
- **ZoKrates:** A toolbox for zk-SNARKs focused on blockchain and smart contracts, providing a complete end-to-end solution.
- **zk-STARKs (Native):** A native implementation of zk-STARKs, focusing on scalability and transparency.
- **PlonK:** A versatile proving system used for both zk-SNARKs and zk-STARKs, optimized for scalability and post-quantum security.
- **STARKy:** A library for zk-STARKs that is designed to optimize the creation of transparent proofs.
- **Winterfell:** A framework for zk-STARKs optimized for cloud and decentralized applications, focusing on scalability.

## 8 Implementation

This section outlines the implementation of a zero-knowledge proof (ZKP) voting system designed to ensure voter privacy on the Ganache blockchain. The goal was to create a secure, verifiable

voting prototype where users cast private votes (0 or 1) without revealing their choices. The system integrates ZoKrates for generating ZKPs, Solidity smart contracts for blockchain logic, and a web interface for user interaction.

The implemented system consists of four primary components:

- A **ZoKrates script** (`vote.zok`) to define the voting constraint.
- A **Verifier.sol** smart contract to validate the proofs on-chain.
- A **ZKP Voting.sol** contract to tally votes.
- A **web frontend** (HTML, JavaScript, CSS) that connects to the Ganache blockchain via Web3.js.

All the source files, including ZoKrates scripts and smart contract code, are attached separately for clarity. This prototype demonstrates privacy-preserving voting, prevents double-voting, and ensures transparency through offline proof generation and on-chain verification.

## 8.1 Overview

The voting system leverages zero-knowledge proofs (ZKPs) to ensure voter privacy on the Ganache blockchain, which is a local Ethereum testnet used for testing purposes. The primary goal was to build a secure, verifiable voting system where users can cast votes (either 0 or 1) without revealing their exact choice.

The three core components of the system are:

- **ZoKrates:** Used for generating zk-SNARKs for the zero-knowledge proofs.
- **Solidity smart contracts:** The backbone of the blockchain logic for vote tallying and verification.
- **Web interface:** Allows users to interact with the blockchain in a seamless manner.

The ZoKrates script (`'vote.zok'`) defines the voting constraint, ensuring that votes can only be 0 or 1. The `'Verifier.sol'` smart contract checks the validity of ZKPs on the blockchain. The `'ZKP Voting.sol'` contract is responsible for recording votes and tallying the results. The web interface allows users to submit their votes, which are validated and tallied on the blockchain.

## 8.2 ZKP Logic with ZoKrates

The core of the ZKP logic is implemented in a ZoKrates script, `'vote.zok'`, included in the attached files. This script specifies the voting constraint that a vote must be a private 32-bit unsigned integer, which can only be either 0 or 1. The constraint is enforced by the following line in the script:

```
assert(vote == 0 || vote == 1);
```

The vote is then returned publicly for tallying, ensuring that only the vote's validity (0 or 1) is revealed, while maintaining privacy.

### 8.2.1 Mathematics Behind the ZKP Logic

To ensure that the vote is either 0 or 1, ZoKrates compiles the script into an arithmetic circuit. This circuit represents the constraint in mathematical form, transforming it into equations that can be verified cryptographically. The critical equation for the ZKP in this case is:

$$(vote - 0) \times (vote - 1) = 0$$

This equation holds true only if the vote is either 0 or 1, because:

- If  $vote = 0$ :  $(0 - 0) \times (0 - 1) = 0 \times -1 = 0$
- If  $vote = 1$ :  $(1 - 0) \times (1 - 1) = 1 \times 0 = 0$

This equation ensures that the vote must either be 0 or 1, providing a simple yet effective method to guarantee the correctness of the vote without revealing it.

After compiling the ZoKrates script, the following steps are executed:

- (1) **Compilation:** The ZoKrates script is compiled using the `zokrates compile` command, which transforms the script into an arithmetic circuit.
- (2) **Setup:** The trusted setup is performed using `zokrates setup`, which generates a proving key and a verification key.
- (3) **Witness Computation:** For each possible vote (0 or 1), ZoKrates computes a "witness" using: `zokrates compute-witness -a <value>`. The witness is a set of intermediate values that satisfy the circuit for that specific input (either vote 0 or vote 1).
- (4) **Proof Generation:** A cryptographic proof is generated using the command `zokrates generate-proof`, which creates the proof components (a, b, c) stored in a JSON file.
- (5) **Verification Key:** The verification key, which is used to validate the proofs on the blockchain, is exported with the command `zokrates export-verifier`.

## 8.3 Smart Contracts on Ganache

The blockchain layer of the voting system consists of two main Solidity smart contracts deployed on Ganache:

- **Verifier.sol:** This contract is generated by ZoKrates and is responsible for verifying the ZKP proofs on-chain.
- **ZKP Voting.sol:** A custom Solidity contract to manage the voting process. This contract interacts with the `Verifier.sol` to validate proofs before recording the votes.

### 8.3.1 Verifier.sol Contract

The `Verifier.sol` contract is responsible for validating the ZKPs. It uses elliptic curve pairing functions, implemented in the `Pairing.sol` library, to verify the cryptographic proofs. The verification process checks that the proof corresponds to the correct verification key and public input.

The code within the contract utilizes the following verification logic:

```
function verifyTx(Proof memory proof, uint[2] memory input)
public view returns (bool) {
    return verify(proof, input);
}
```

This function takes the proof and input (vote) and checks the validity against the verification key, ensuring that the vote is valid without revealing the actual vote.

### 8.3.2 ZKPVoting.sol Contract

This contract integrates with the `Verifier.sol` contract and provides the logic for voting. Key functions in this contract include:

- `registerVoter`: Marks a voter as eligible.
- `castVote`: Takes the proof components (a, b, c) and the public input (either 0 or 1) and calls the `Verifier.sol` contract to validate the proof before recording the vote.
- `getVoteCounts`: Returns the total count of votes for each option (0 or 1).
- `getWinner`: Determines the winner by comparing the vote counts.

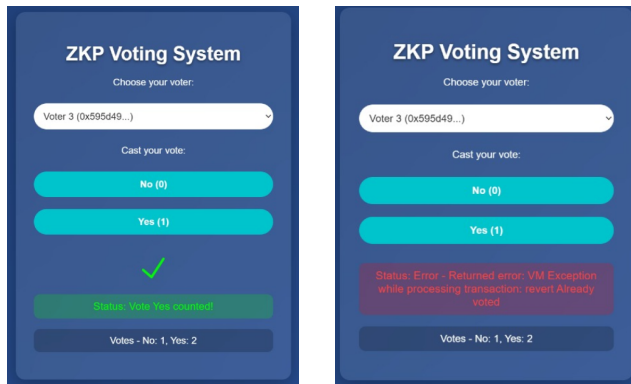
Mappings are used in the contract to track voters and prevent double-voting.

### 8.4 Web Interface

The web interface allows users to interact with the voting system. It is implemented using HTML, JavaScript, and CSS. The HTML file provides a dropdown menu listing Ganache accounts, buttons to vote "No" (0) or "Yes" (1), and visual feedback with status messages, vote counts, and animations.

The JavaScript file uses Web3.js to interact with the smart contracts on the blockchain. The key functionality includes:

- **Registering a Voter**: The user is registered as a voter when they interact with the system for the first time.
- **Casting a Vote**: When the user casts their vote, the proof (a, b, c) and the public input (either 0 or 1) are submitted to the `ZKPVoting.sol` contract via Web3.js.
- **Displaying Results**: The frontend displays the updated vote counts by calling the `getVoteCounts` function from the smart contract.



(a) Web Interface - Casting Vote (b) Error if Double-voting

Figure 3: Comparison of Web Interfaces to Cast Vote

### 8.5 Workflow and Integration

The voting system operates through a seamless workflow:

- (1) The user selects a vote (0 or 1) from the web interface.
- (2) Precomputed ZoKrates proofs for each vote (0 or 1) are stored in the `script.js` file.

- (3) The proof and public input are submitted to the blockchain via Web3.js.
- (4) The `castVote` function in `ZKPVoting.sol` validates the proof using `Verifier.sol`.
- (5) If the proof is valid, the vote is recorded, and the vote count is updated.
- (6) The web interface displays the updated results in real-time.

This workflow, combining offline proof generation, on-chain verification, and user interaction, creates a privacy-preserving, transparent, and secure voting system.

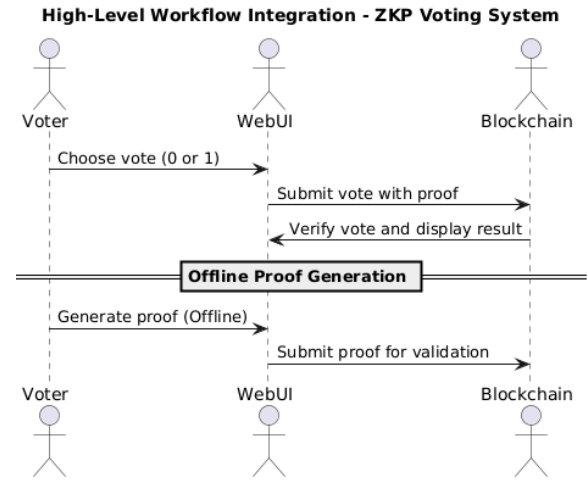


Figure 4: High-level Workflow and Integration

## 9 Thoughts and Remarks

Throughout the implementation of this zero-knowledge proof (ZKP) voting system, I gained valuable insights into the complexities and capabilities of blockchain-based privacy-preserving solutions. Below are my thoughts and remarks on the project's challenges, successes, and potential future improvements.

### 9.1 Challenges Faced

One of the primary challenges in implementing this system was understanding and correctly applying the cryptographic concepts behind zk-SNARKs. ZoKrates, while powerful, has a steep learning curve, especially in terms of its trusted setup and the intricacies of generating cryptographic proofs and verification keys. Ensuring the trusted setup's security was crucial since the integrity of the entire voting system depends on the correctness of the proof generation process.

Another challenge was integrating ZoKrates with Solidity. Although ZoKrates provides a convenient way to generate ZKPs, translating the generated proof data and verification code into Solidity and ensuring compatibility with the Ethereum blockchain required careful handling. The need for elliptic curve pairings and compatibility with Ganache also added complexity, as some of the cryptographic functions needed to be manually adjusted.

## 9.2 Successes and Achievements

On the positive side, the integration of ZoKrates with Ethereum smart contracts provided a functional privacy-preserving voting system. The system ensures that votes remain private while still being verifiable, maintaining transparency and preventing fraud or double voting. By generating offline proofs and validating them on-chain, the system effectively combines the privacy benefits of ZKPs with the immutability and transparency of blockchain.

The web interface provided a user-friendly method for interacting with the blockchain. By using Web3.js, voters could seamlessly cast their votes, which were then securely validated and recorded on the blockchain. This feature demonstrated the practical viability of using ZKPs in a real-world application, which I believe is a promising use case for future privacy-preserving decentralized applications.

## 9.3 Future Improvements and Considerations

Looking forward, there are several areas where this voting system could be improved or extended:

- **Scalability:** Currently, the system works well for a small number of voters, but in real-world scenarios with large-scale elections, performance optimizations would be necessary. This could include using more efficient proof generation schemes or offloading some computations off-chain.
- **Security Enhancements:** While the use of zk-SNARKs provides strong privacy guarantees, additional security measures such as ensuring voter anonymity and preventing attacks like Sybil attacks could further enhance the system's robustness.
- **User Experience:** The current web interface is functional, but it could benefit from additional features such as individual voter login registration and authentication, better error handling to improve overall usability.
- **Extended Use Cases:** The concept of using ZKPs for privacy-preserving voting could be extended beyond simple yes/no votes to more complex systems, such as ranked-choice voting or elections with multiple candidates, which could involve more advanced cryptographic protocols.

## 10 Conclusion

In this project, we successfully designed and implemented a privacy-preserving voting system using zero-knowledge proofs (ZKPs) on the Ganache blockchain. The system ensures voter privacy, prevents double-voting, and maintains transparency through the use of zk-SNARKs. While the prototype demonstrates the viability of privacy-preserving voting, several areas for improvement were identified, including scalability and security. Future work will focus on optimizing the system for larger-scale use cases and enhancing the user experience. This project highlights the potential of zero-knowledge proofs in creating secure and private decentralized applications. Additionally, the project emphasizes the need for further research into enhancing the efficiency of zk-SNARKs to handle real-time applications. By addressing these challenges, privacy-preserving systems can be widely adopted across multiple domains, including financial transactions and digital identity management.

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