

**Problem background:**

Take smartphone market as an example, there is different models of smartphone ( $D = \{r_1, r_2, \dots, r_n\}$ ). Each model has different price, pixel, battery capacity, cooling capacity and etc. ( $r_i = [r_i^1, r_i^2, \dots, r_i^d]$ ).

Now brand  $P$  has  $m$  models ( $P = \{p_1, p_2, \dots, p_m\} \subset D$ ), and have user data set  $W = \{w_1, w_2, \dots, w_x\}$ . Each  $w_i = [w_i^1, w_i^2, \dots, w_i^d]$  present the preference weight vector of a single user.

The score of a model  $r_i$  respecting to a user  $w_i$  is the dot product  $r_i \cdot w_i$ . A model covers a user when the score of it ranks top  $k$  among  $D$ .

**Q:** How to introduce a new product  $p$  such that maximizes the cover ratio  $cp(p, P, k) = \frac{|\{w | \forall w \in W, \{P \cup \{p\}\} \cap TopK(w) \neq \emptyset\}|}{|W|}$  and the product creation cost  $C(p)$  is no greater than creation budget  $B$ ?

**Problem Definition:**

*Input:*

global model dataset  $D = \{r_1, r_2, \dots, r_m\}$ ,  $r_i = [r_i^1, r_i^2, \dots, r_i^d]$ ,

Another model dataset  $P = \{p_1, p_2, \dots, p_n\} \subset D$

Model creation budget  $B$ ,

a positive integer  $k$ ,

a creation cost function  $C(p_i) = \sum c^j p_i^j$ ,

a user preference tuple set  $W = \{w_1, w_2, \dots, w_x\}$ ,  $w_i = [w_i^1, w_i^2, \dots, w_i^d]$ ,

The score of a model  $r_i$  respecting to a user  $w_i$  is the dot product  $r_i \cdot w_i$

A model covers a user when the score of it ranks top  $k$  among  $D$ .

*Output:*

find a new product  $p$  such that satisfies the constrain  $C(p) \leq B$

and maximize  $P \cup \{p\}$ 's cover ratio  $cp(p, P, k) = \frac{|\{w | \forall w \in W, \{P \cup \{p\}\} \cap TopK(w) \neq \emptyset\}|}{|W|}$

**Hit:**

*The main difference between baseline and advance-1 is:*

1. advance-1 only consider weight vectors that not covered by  $P$
2. advance-1 only consider  $C(p) = B$  rather than  $C(p) \leq B$
3. advance-1 only consider the weight vectors  $w$  that  $w \cdot p = S_k$  intersects with  $C(p) = B$

*The main difference between advance-1 and advance-2 is:*

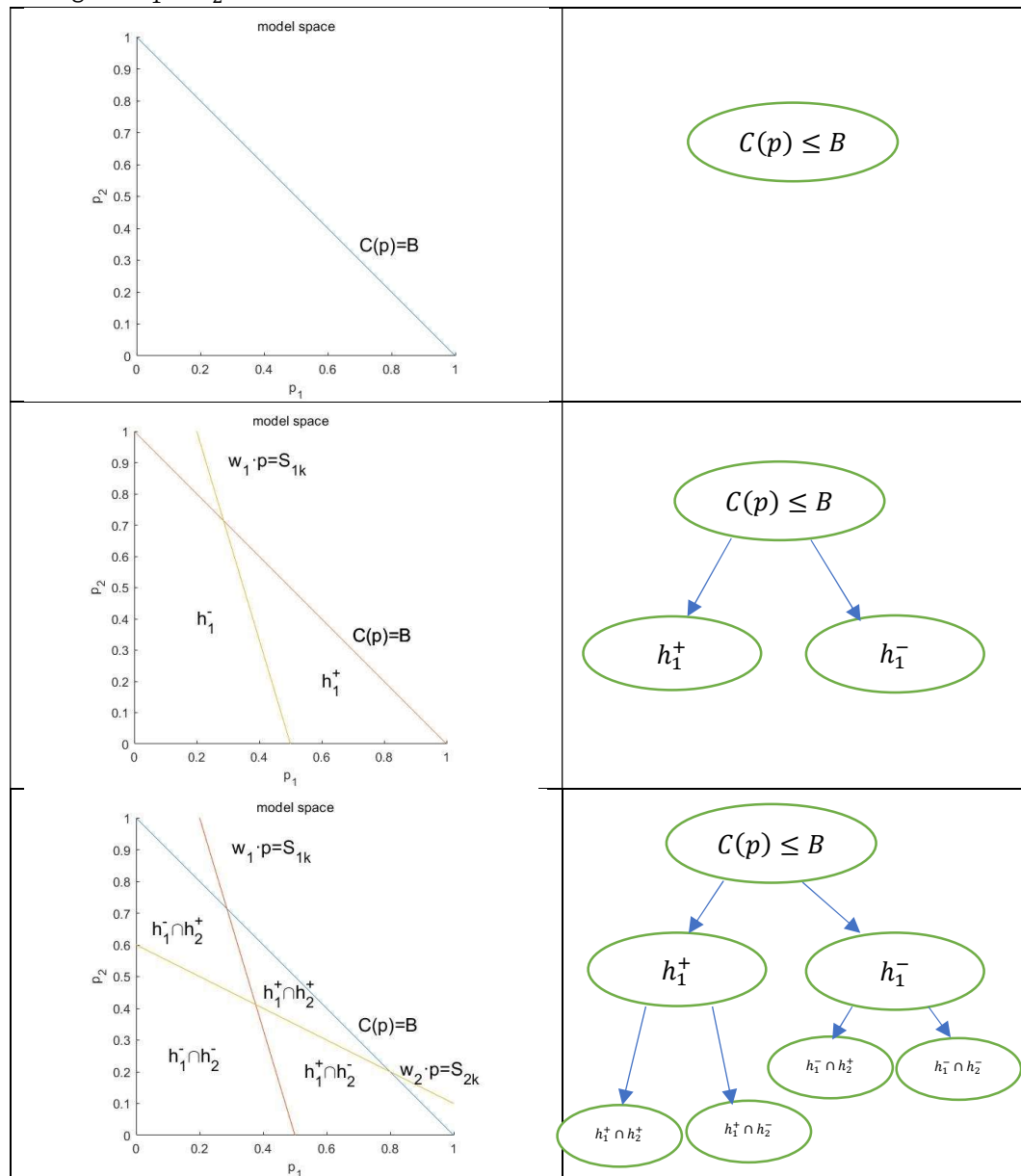
1. advance-2 considers convex hull to reduce the range of  $C(p) = B$

### Baseline solution example

Take  $d = 2$  as example,

where  $S_{ik}$  is the score to rank top  $k$  respecting to  $w_i$

1. In model space, only the half-space  $C(p) \leq B$  is valid
2. If we want to cover  $w_1$ , the score of the new model  $p$  should be greater than  $S_{1k}$   
The line  $w_1 \cdot p = S_{1k}$  divides the space  $C(p) \leq B$  into 2 parts.  
One is  $h_1^+$ , where any model in it could cover  $w_1$ ;  
Another is  $h_1^-$ , where any model in it couldn't cover  $w_1$ .
3. If we want to cover  $w_2$ , the score of the new model  $p$  should be greater than  $S_{2k}$   
The line  $w_2 \cdot p = S_{2k}$  divides the space  $h_1^+$  into 2 parts.  
One is  $h_1^+ \cap h_2^+$ , where any model in it could cover  $w_1, w_2$ ;  
Another is  $h_1^+ \cap h_2^-$ , where any model in it could only cover  $w_1$ .  
The line  $w_2 \cdot p = S_{2k}$  divides the space  $h_1^-$  into 2 parts.  
One is  $h_1^- \cap h_2^+$ , where any model in it could only cover  $w_2$ ;  
Another is  $h_1^- \cap h_2^-$ , where any model in it couldn't only any user.
4. Return region  $h_1^+ \cap h_2^+$



**Baseline solution ( $kSPR$ ):**

1. For each  $w_i$  in  $W$ , based on  $w_i \cdot p \leq S_{ik}$  using cell tree divide model space  $C(p) \leq B$ , where  $S_{ik}$  is the score to rank top  $k$  respecting to  $w_i$
2. count each divided region's user cover number
3. Return the region that satisfies  $C(p) \leq B$  and with the greatest cover count

Time complexity:  $O(n^d)$ , where  $n$  is the capacity of  $D$ ,  $d$  is the dimension of a model.

### Advance solution version-1 example:

Take  $d = 2$  as example,

where  $S_{ik}$  is the score to rank top  $k$  respecting to  $w_i$

1. discard all weight vectors such that have been cover by  $P$ :

$$W = W - \{w | \forall w \in W, P \cap \text{TopK}(w) \neq \emptyset\}$$

2. In model space, only consider the space  $C(p) = B$

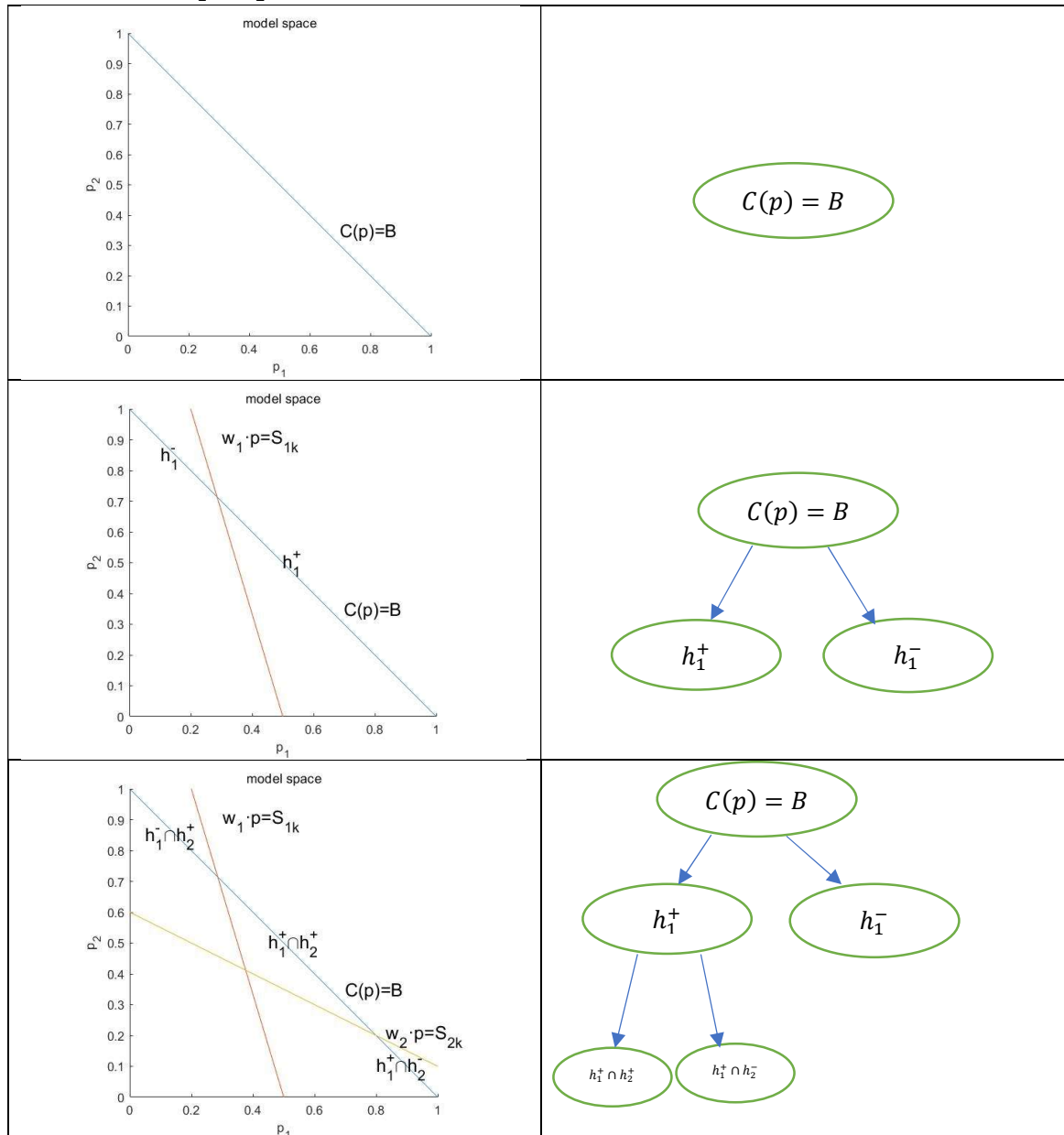
3. The line  $w_1 \cdot p = S_{1k}$  divides the line  $C(p) = B$  into 2 part.

One is  $h_1^+$ , where could cover  $w_1$ ; Another is  $h_1^-$ , where couldn't cover  $w_1$ .

4. The line  $w_2 \cdot p = S_{2k}$  divides the space  $h_1^+$  into 2 part.

One is  $h_1^+ \cap h_2^+$ , could cover  $w_1, w_2$ ; Another is  $h_1^+ \cap h_2^-$ , where could only cover  $w_1$ .

5. Return line  $h_1^+ \cap h_2^+$



**Advance solution version1:**

1. Preprocessing *I*, discard all user that *P* already covers:  

$$W = W - \{w | \forall w \in W, P \cap TopK(w) \neq \emptyset\}$$
2. we only consider  $C(p) = B$
3. Preprocessing *II*, discard all  $w$  such that  $w \cdot p = S_{ik}$  doesn't intersect with  $C(p) = B$   

$$W = W - \{w | \forall w \in W, \{p | w \cdot p = S_k\} \cap \{p | C(p) = B\} = \emptyset\}$$
4. For each  $w_i$  in  $W$ , based on  $w_i \cdot p \leq S_{ik}$  using cell tree divide model space  $C(p) = B$ , where  $S_{ik}$  is the score to rank top  $k$  respecting to  $w_i$
5. count each divided region's user cover number
6. Return the region that satisfies  $C(p) = B$  and with the greatest cover count

Time complexity:  $O(n_{new}^{d-1})$ ,

where  $n_{new}$  is the number of  $w$  such that not cover by  $P$  and  $w \cdot p = S_k$  intersect with  $C(p) = B$ ;

$d$  is the dimension of a model.

### Advance solution version-2 example:

Suppose  $d = 2$  and  $\{p_1, p_2, p_3\} \subset P$

1. Preprocessing *I*, discard all user that  $P$  already covers:

$$W = W - \{w | \forall w \in W, P \cap \text{TopK}(w) \neq \emptyset\}$$

2. we only consider  $C(p) = B$

3. Find the biggest convex hull of point set  $P$  (name it as  $B_{CH}$ )

4. Preprocessing *II*, on  $C(p) = B$ , discard the part which is in convex hull (*see proof 1*)

$$A = \{p | C(p) = B\}$$

$$A = A - B_{CH}$$

5. Preprocessing *III*, discard all  $w$  such that  $w \cdot p = S_{ik}$  doesn't intersect with  $C(p) = B$

$$W = W - \{w | \forall w \in W, \{p | w \cdot p = S_k\} \cap \{p | C(p) = B\} = \emptyset\}$$

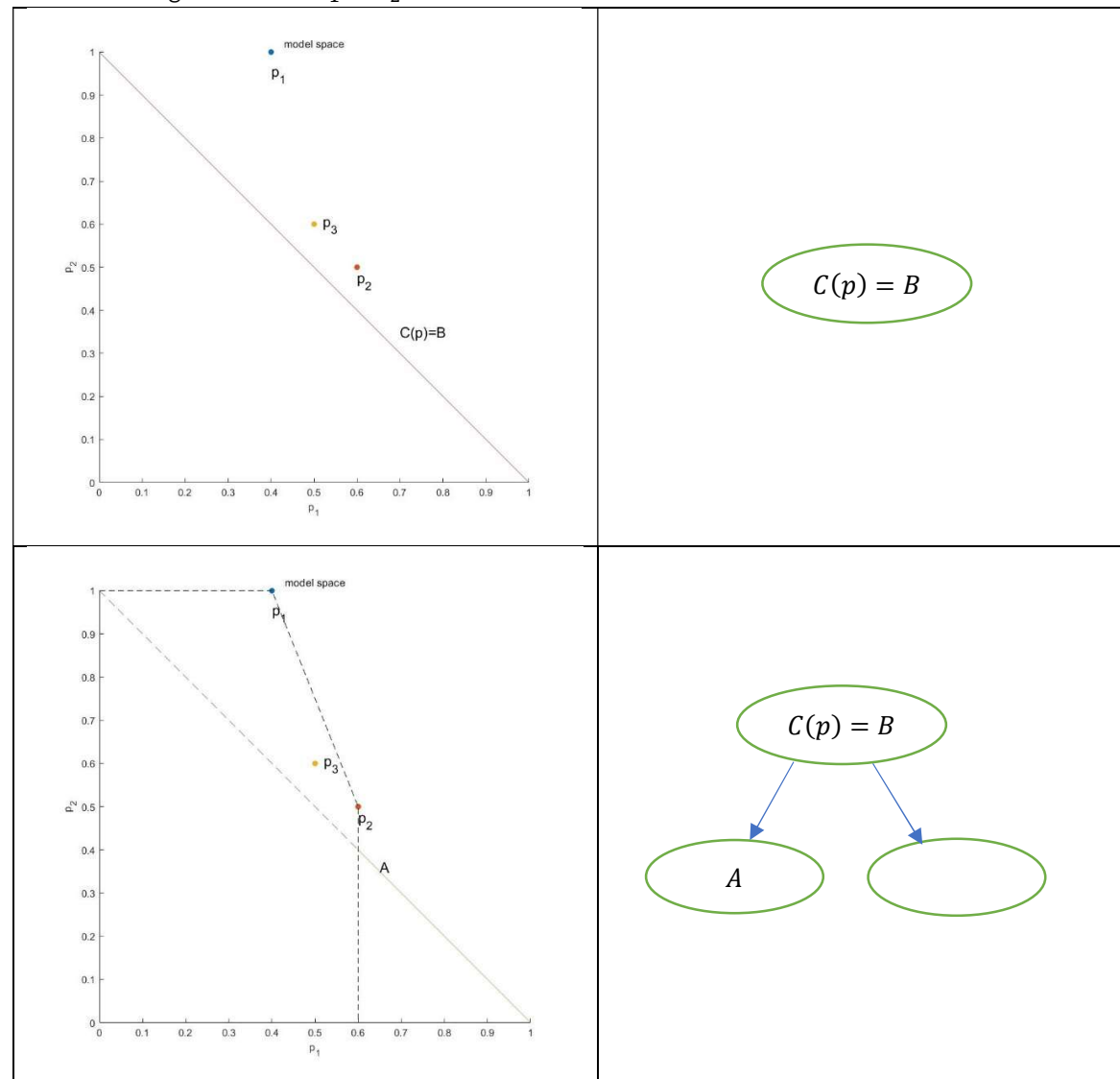
As shown below, we only need the solid part of  $C(p) = B$ , which is  $A$

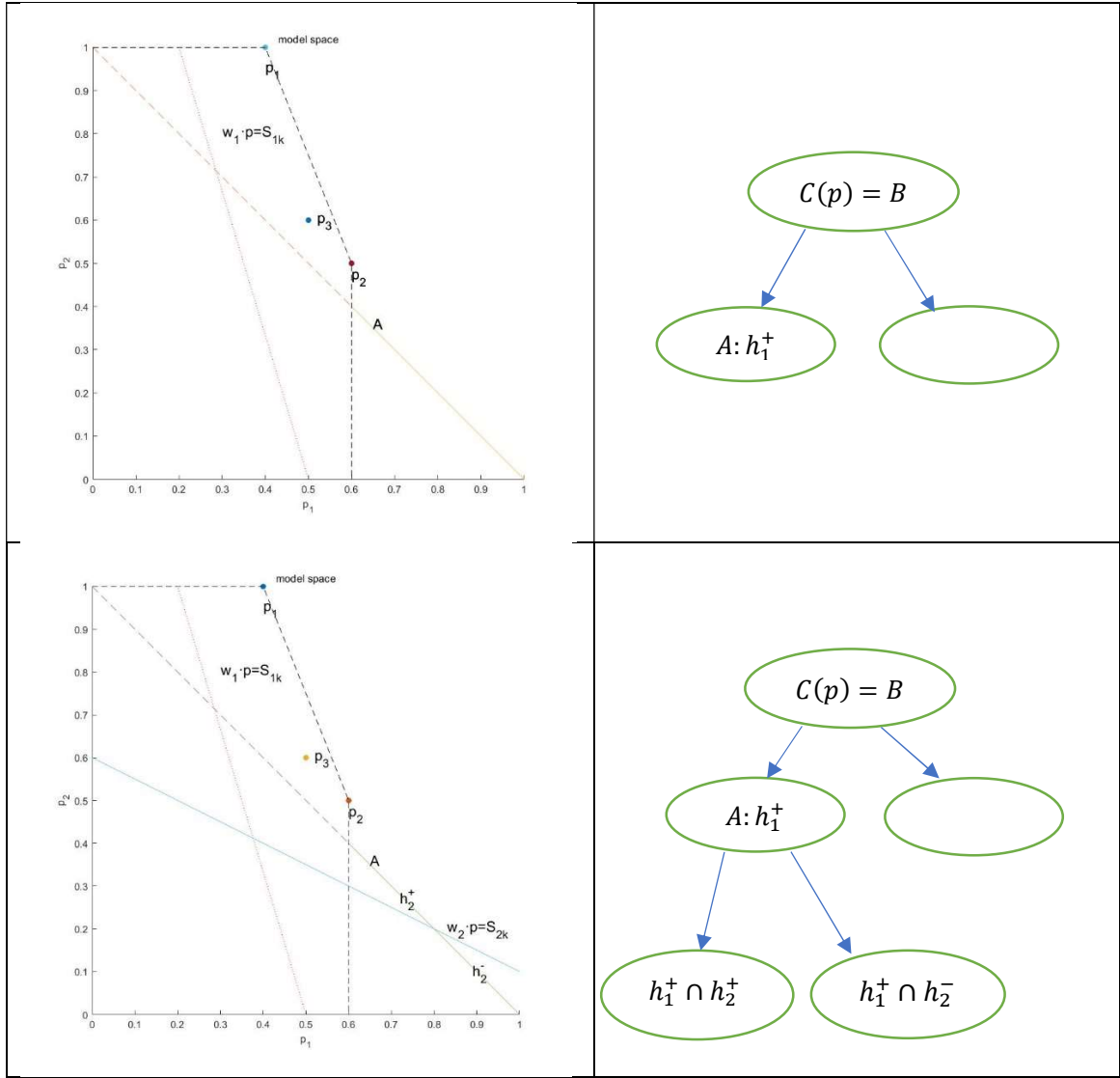
6. From line  $w_1 \cdot p = S_{1k}$ , we could know all of  $A$  is in  $h_1^+$ , which means  $A$  could cover  $w_1$

7. The line  $w_2 \cdot p = S_{2k}$  divides the line  $A$  into 2 part.

One is  $h_1^+ \cap h_2^+$ , could cover  $w_2$ ; Another is  $h_1^+ \cap h_2^-$ , where couldn't cover  $w_2$ .

8. Return segmentation  $h_1^+ \cap h_2^+$





#### Advance solution version-2:

1. Preprocessing *I*, discard all user that  $P$  already covers:  

$$W = W - \{w | \forall w \in W, P \cap TopK(w) \neq \emptyset\}$$
2. we only consider  $C(p) = B$
3. Find the biggest convex hull of point set  $P$  (name it as  $B_{CH}$ )
4. Preprocessing *II*, on  $C(p) = B$ , discard the part which is in convex hull (**see proof 1**)  

$$A = \{p | C(p) = B\}$$

$$A = A - B_{CH}$$
5. Preprocessing *III*, discard all  $w$  such that  $w \cdot p = S_{ik}$  doesn't intersect with  $C(p) = B$   

$$W = W - \{w | \forall w \in W, \{p | w \cdot p = S_k\} \cap \{p | C(p) = B\} = \emptyset\}$$
6. For each  $w_i$  in  $W$ , based on  $w_i \cdot p \leq S_{ik}$  using cell tree divide model space  $A$ , where  $S_{ik}$  is the score to rank top  $k$  respecting to  $w_i$
7. count each divided region's user cover number
8. Return the region that with the greatest cover count

Time complexity:  $O(n_{new}^{d-1})$ ,

where  $n_{new}$  is the number of  $w$  such that not cover by  $P$  and  $w \cdot p = S_k$  intersect with  $A$ ;  
 $d$  is the dimension of a model;

Compare with version1, version2's  $n_{new}$  is no more than version1's  $n_{new}$ .

**Proof 1:**

To proof:

if a model is covered by a convex hull and this model could cover weight vector  $w$ , then one of the models that form this convex hull can also cover  $w$ .

Proof:

Suppose point  $O$  is the original point; point  $C$  is inside convex hull;  $D$  is the crossover point between  $OC$  and convex hull; on convex hull,  $D$  is on the hyperplane  $H_1H_2 \dots H_d$ . Then

$$\overrightarrow{OD} = \lambda_1 \overrightarrow{OH_1} + \lambda_2 \overrightarrow{OH_2} + \dots + \lambda_{d-1} \overrightarrow{OH_{d-1}} + \lambda_d \overrightarrow{OH_d}, \quad \sum_{i=1}^d \lambda_i = 1; \forall i \in [1, d], \lambda_i \geq 0$$

Next,

$$w \cdot \overrightarrow{OD} = w \cdot (\lambda_1 \overrightarrow{OH_1} + \lambda_2 \overrightarrow{OH_2} + \dots + \lambda_{d-1} \overrightarrow{OH_{d-1}} + \lambda_d \overrightarrow{OH_d})$$

Without loss of generality, suppose  $\forall i \in [1, d], w \cdot \overrightarrow{OH_1} \geq w \cdot \overrightarrow{OH_i}$

Then

$$\begin{aligned} & w \cdot \overrightarrow{OC} \\ & < w \cdot \overrightarrow{OD} \\ & = w \cdot (\lambda_1 \overrightarrow{OH_1} + \lambda_2 \overrightarrow{OH_2} + \dots + \lambda_{d-1} \overrightarrow{OH_{d-1}} + \lambda_d \overrightarrow{OH_d}) \\ & \leq w \cdot (\lambda_1 \overrightarrow{OH_1} + \lambda_2 \overrightarrow{OH_1} + \dots + \lambda_{d-1} \overrightarrow{OH_1} + \lambda_d \overrightarrow{OH_1}) \\ & = w \cdot \overrightarrow{OH_1} \end{aligned}$$

That is

$$w \cdot \overrightarrow{OC} < w \cdot \overrightarrow{OH_1}$$

Proof ends.

**Proof 1's corresponding original lemma in  $k$  – Hit query:**

*Lemma 4. Let  $p$  be a tuple in  $D$  which is inside  $\text{Conv}(D)$  but is not on the surface of  $\text{Conv}(D)$ . Then,  $HP(p) = 0$ .*

*Proof of Lemma 4:*

*Firstly, we show that for any point  $p' \in D$  and any (linear) function  $f \in \mathcal{L}$  associated with its weight vector  $w$ ,  $f(p')$  is equal to the length of the projection of the vector from the origin to point  $p'$  on vector  $w$ . This can be shown as follows. We denote the vector from the origin to point  $p'$  by  $\overrightarrow{op'}$ . Let  $\theta$  be the angle between vector  $w$  and vector  $\overrightarrow{op'}$ .  $f(p') = w \cdot p = \|w\| \|\overrightarrow{op'}\| \cos \theta$ . Since the norm of the weight vector  $w$  is 1 (by definition) (i.e.,  $\|w\| = 1$ ), we have  $f(p') = \|\overrightarrow{op'}\| \cos \theta$ . Note that  $\|\overrightarrow{op'}\| \cos \theta$  is equal to the length of the projection of vector  $\overrightarrow{op'}$  on vector  $w$ , which completes the proof.*

*Secondly, we show that  $HP(p) = 0$ . For any weight vector  $w$  where its corresponding utility function  $f$  is in  $\mathcal{L}$ , there exists a point  $p'$  on the surface of  $\text{Conv}(D)$  such that the length of the projection of the vector from the origin to point  $p'$  on vector  $w$  is greater than the length of the projection of the vector from the origin to point  $p$  on vector  $w$ . This means that  $f(p') > f(p)$ . In other words, for any utility function  $f$  in  $\mathcal{L}$ , there exists a point  $p'$  on the surface of  $\text{Conv}(D)$  such that  $f(p') > f(p)$ . This implies that  $HP(p) = 0$ .*