

MAT344 Combinatorics

Graph Theory

Session 2

1. Let G be a simple regular graph on n vertices such that every vertex has degree 2, Prove that G is the union of disjoint cycles.

Note:

- A graph is said to be simple if it has no loops and every pair of vertices has at most one edge between them.
- A simple graph is said to be regular if every vertex has the same degree

1. Pick a random vertex v_1 from G , and let v_1 be connected to v_2 and v_k since the degree of all vertices is 2, **G is 2-regular**. Now, starting from v_1 keep going in a direction, to find all the vertices of the **current cycle**, since at some point there is no new vertex to go to, since the degree of the current vertex is 2, we have to go back to one of the vertices that we have already visited. Let that vertex be v_i , the only possibility for v_i is the initial vertex since it has degree 2 and we have only used one of its edges so far, otherwise if we go to some other vertex that we have visited than the initial one then the degree of that vertex will exceed 2, as we have already used both of its edges. Since all the vertices have degree 2, then the resulting walk is either, $v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_k \rightarrow v_1$, or $v_1 \rightarrow v_k \rightarrow v_{k-1} \rightarrow \dots \rightarrow v_2 \rightarrow v_1$. Both of which are the vertices of a **connected 2-regular subgraph, cycle**.

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1. Now, **take out the cycle** that we found, i.e. $v_1-v_2-\dots-v_k-v_1$, from graph G .
2. **Repeat** the process in 1 and 2 to **find and remove another cycle**.
3. **Stop** when there are **no more** vertices left in G to pick.
4. Repeat the steps above, we can find all the **disjoint** cycles of the graph G .

2. Let G be a simple graph on n vertices. We define the complement graph G^c of G to be the graph on the same vertex set of G in which the vertices x and y are connected by an edge in G^c if and only if x and y are not connected by an edge in G . Prove that for $n \geq 5$, either G or G^c has a cycle.

- Proof by contradiction.
- Assume that for $n \geq 5$, none of G and G^c have a cycle.
- If none of G and G^c have cycles, then each of them are **forests**, i.e. union of **disjoint** trees.
- Let G and G^c be n -vertex graphs.
- Then:
- $E(G) \leq n - 1$, and $E(G^c) \leq n - 1$,
- therefore; $E(G) + E(G^c) \leq 2n - 2$
- $E(G) + E(G^c) = E(K_n) = \binom{n}{2} = n(n - 1)/2$

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- Then,
- $E(G) + E(G^c) = n(n - 1)/2 \leq 2n - 2$
- $n^2 - n \leq 4n - 4$
- $n^2 - 5n + 4 \leq 0$
- $(n - 4)(n - 1) \leq 0$
- Then the possible values for $n = 1, 2, 3, 4$; however, $n \geq 5$, a Contradiction.
- **Q.E.D.**

3. Let G be a connected bipartite graph. Let v and w be two vertices of G . Let P_1 and P_2 be any (not necessarily vertex or edge disjoint) paths from v to w . Prove that the lengths (the number of edges) of P_1 and P_2 are either both even or both odd.

- Proof by contradiction:
- Assume one is even and one is odd, then walking from **v to w** using **p_1** and coming back from **w to v** using **p_2** , will make an **odd cycle**, which is a **contradiction**, since a bipartite graph has **no odd cycle**.
- **Q.E.D.**

4. Is there a bipartite graph G with 9 vertices v_1, v_2, \dots, v_9 such that their degrees are 3, 3, 3, 3, 3, 5, 6, 6, and 6, respectively?

Note:

- The total number of edges $= \text{sum of degrees} / 2 = 38 / 2 = 19$
- Since G is a **bipartite** graph, and it has **9 vertices**, let's consider different cases that G can exist:
 1. LHS = 1 RHS = 8 then, the maximum number of edges is **8**, a contradiction
 2. LHS = 2 RHS = 7 then, the maximum number of edges is **14**, a contradiction
 3. LHS = 3 RHS = 6 then, the maximum number of edges is **18**, a contradiction
 4. LHS = 4 RHS = 5 then, the maximum number of edges is **20**; However, there **cannot** exist a vertex of **degree 6**, a contradiction
- Therefore; none of the cases will satisfy the conditions for G ; Therefore, G does **not** exist.

Q.E.D.

1. Miklós Bóna. A Walk Through Combinatorics:An Introduction to Enumeration and Graph Theory (p. 206). World Scientific Publishing Company. Kindle Edition.
2. Miklós Bóna. A Walk Through Combinatorics:An Introduction to Enumeration and Graph Theory (p. 207). World Scientific Publishing Company. Kindle Edition.
3. Miklós Bóna. A Walk Through Combinatorics:An Introduction to Enumeration and Graph Theory (p. 209). World Scientific Publishing Company. Kindle Edition.