

# Stacks and Queues

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# Stacks and Queues

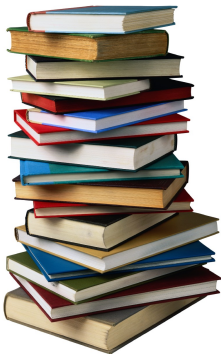


- **Stacks** are **Queues** are two popular easy-to-implement **abstract data types** (ADTs) used in algorithm design
- For an abstract data type, we **describe** the operations that can be executed on the data, but we do not define them
- For instance, the stack ADT can be implemented using a linked-list and an array as well

# Stack

## Definition

A stack is an ADT where items are inserted (**pushed**) and removed (**popped**) according to the **last-in, first-out** (LIFO) principle



# Operations on a stack



Let us denote a stack by  $S$ .

- ①  **$S.push(e)$** : **adds** element  $e$  to the top of  $S$
- ②  **$S.pop()$** : **removes** and **returns** the top element from  $S$ ; returns **null** if  $S$  is empty
- ③  **$S.top()$** : **returns** the top element from  $S$  but does not remove it
- ④  **$S.size()$** : **returns** the number of elements stored in  $S$  currently
- ⑤  **$S.isEmpty()$** : **returns** **true** if  $S$  is empty otherwise returns **false**

```
push(55)
```

55

`push(3)`

55, 3

`push(69)`

55, 3, 69

`push(-88)`

55, 3, 69, -88



`push(-17)`

55, 3, 69, -88, -17

`isEmpty()` returns false

55, 3, 69, -88, -17

`size()` returns 5

55, 3, 69, -88, -17

`pop()` removes and returns -17

55, 3, 69, -88

**pop() removes and returns -88**

55, 3, 69

`top()` returns 69

55, 3, 69

# Applications

- Can be used to implement the **FORWARD** and **BACK** buttons in a browser.
- Used by text editors to implement the **UNDO** operation. When you change something, the current state is pushed onto the stack before applying the new change. Hitting UNDO will simply revert the document to the previous state
- Method calls, recursion, and compiler design
- Used as an auxiliary data structure for many cornerstone algorithms

## An interface for stacks

```
public interface StackADT<E> {  
    int size();  
    boolean isEmpty();  
    void push(E e);  
    E pop();  
    E top();  
}
```

To use a stack, one should create a class by implementing this interface



## So how to implement a stack using the StackADT interface?

- Use an array (stack size must be known in advance, speedy in practice)  
See the class [ArrayStack](#)
- Use a linked-list (can grow arbitrarily, a bit slower than the array-based stack)  
See the class [LinkedStack](#)

Irrespective of the implementation, every one of these five stack operations, `size()`, `isEmpty()`, `push()`, `pop()`, `top()`, takes  $O(1)$  time

## An application: matching parentheses

- Given an algebraic expression, how to check if the parentheses, braces, and brackets in it are properly **balanced** or not
- Some **balanced** ones
  - $(a + b) + (66 * (s / t) - 12) * \{p + (z - [2 + 99.1 / (m - x)])\}$   
Sequence: ( ) ( ( ) ) { ( [ ( ) ] ) }
  - $30 + (40 + (9 * (6 - 10) * (1 - (17 + 19 + 99) - 1) / \{5 + (8 - [6 + 2 * (1 + (-99)) - 1] - 4) - 1\} - 1) - 1) * 88$   
Sequence: ( ( ( ) ( ( ) ) { ( [ ( ) ] ) } ) )
- Some **unbalanced** sequences
  - ) ( ( ) ) { ( [ ( ) ] ) }
  - ( { [ ] ) }
  - (

## How to solve it?

- Recall the three pairs we are interested in:

( ), { }, [ ]

- Here is a stack-based algorithm for this problem:
  - 1 Declare a **Character** stack **S**
  - 2 Scan the input algebraic expression from left to right
  - 3 If a left symbol (, {, or [ is encountered, **push** it onto **S**
  - 4 If a right symbol ), }, or ] is encountered, check the top element in **S**; if these two elements form a matching pair, then pop from **S** otherwise report **INVALID**
  - 5 At the end, when the string is fully scanned, if **S** is empty, then report **VALID**, else **INVALID**

## Sample run

| Current symbol scanned      | Stack (right is the stack top) | Action taken |
|-----------------------------|--------------------------------|--------------|
| ( ) ( ( ) ) { ( [ ( ) ] ) } | (                              | push         |
| ( ) ( ( ) ) { ( [ ( ) ] ) } |                                | pop          |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | (                              | push         |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | ( (                            | push         |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | (                              | pop          |
| ( ) ( ( ) ) { ( [ ( ) ] ) } |                                | pop          |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | {                              | push         |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | { (                            | push         |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | { ( [                          | push         |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | { ( [ (                        | push         |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | { ( [                          | pop          |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | { (                            | pop          |
| ( ) ( ( ) ) { ( [ ( ) ] ) } | {                              | pop          |
| ( ) ( ( ) ) { ( [ ( ) ] ) } |                                | pop          |

Result: **VALID**

## Another sample run

| Current symbol scanned | Stack (right is the stack top) | Action taken    |
|------------------------|--------------------------------|-----------------|
| ( { [ ] ) }            | (                              | push            |
| ( { [ ] ) }            | ( {                            | push            |
| ( { [ ] ) }            | ( { [                          | push            |
| ( { [ ] ) }            | ( {                            | pop             |
| ( { [ ] ) }            | ( {                            | <b>mismatch</b> |

Result: **INVALID**

# Code

```
import java.util.Scanner;
public class ExpressionChecker {
    public static void main(String[] args) {
        Scanner input = new Scanner(System.in);
        String expression = input.nextLine();
        LinkedStack<Character> S = new LinkedStack<>();
        int pos;

        for(pos = 0; pos < expression.length(); pos++) {
            char current = expression.charAt(pos);

            if( current == '(' || current == '{' || current == '[')        S.push(current);
            else if( current == ')' && !S.isEmpty() && S.top() == '(' )    S.pop();
            else if( current == ')' && !S.isEmpty() && S.top() != '(' )    break;
            else if( current == '}' && !S.isEmpty() && S.top() == '{' )    S.pop();
            else if( current == '}' && !S.isEmpty() && S.top() != '{' )    break;
            else if( current == ']' && !S.isEmpty() && S.top() == '[' )    S.pop();
            else if( current == ']' && !S.isEmpty() && S.top() != '[' )    break;
        }

        if( S.isEmpty() && pos == expression.length() )    System.out.println("VALID");
        else System.out.println("INVALID");

        input.close();
    }
}
```

# Food for thought

```
<html>
  <head>
    <title>Data Structure</title>
  </head>

  <body>
    <h1>
      <p>Stacks are fun...</p>
    </h1>

    <p>Queues are also fun...</p>
  </body>
</html>
```

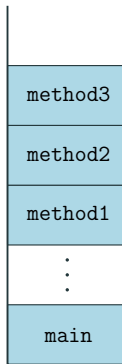


Write a program that can check if an HTML file is correctly formatted. This means every opening tag must have its closing tag at an appropriate location in the file.

# Method calls and use of stack

- Ever wondered how the control returns to the **caller** method after the **callee** method is done with its execution?

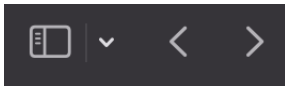
```
public char method1() {  
    // ...  
    method2();  
    counter++; // control returns here  
    i = j + 10;  
    //...  
}  
public double method2() {  
    // ...  
    s = p + s;  
    method3();  
    arr[q] = arr[q] / (10 * m); // control returns here  
    // ...  
}  
public void method3() {  
    // ...  
    s = p + s;  
    // ...  
    return;  
}
```



- A stack of **activation records** is used by the system; push a record when a method starts execution; pop its record when it is done



## Browser tabs: an application of stacks



- For every tab, we get to see the *go forward* and *go backward* buttons
- How to implement this button
- Use 2 **stacks**: one for the *go backward* button and one for the *go forward* button

# Code

```
public class BrowserTab {
    String currentPage = "";
    private final LinkedStack<String> backPages = new LinkedStack<>(),
    forwardPages = new LinkedStack<>();
    public static String getTimestamp(){ return "[" + new Date() + "] "; }

    public BrowserTab() { System.out.println(getTimestamp() + "New tab opened."); }

    public void typeAndGoNewSite(String newPage) {
        if( !currentPage.isEmpty() ) // some page is being viewed currently
            backPages.push(currentPage); // send the current page to the previousPages DLL

        currentPage = newPage;
        System.out.println(getTimestamp() + "Currently viewing: " + currentlyViewing());
    }
    // contd. on the next slide
}
```

# Code

```
public class BrowserTab {
    public void clickOnGoBackButton() {
        if( backPages.isEmpty() ) {
            System.out.println(getTimeStamp() + "Back button is greyed out (unavailable).");
            return;
        }
        else {
            System.out.print(getTimeStamp() + "Back button is clicked. ");
            forwardPages.push(currentPage);
            currentPage = backPages.top();
            backPages.pop();
        }
        System.out.println("Currently viewing: " + currentlyViewing());
    }

    public void clickOnGoForward() {
        if( forwardPages.isEmpty() ) {
            System.out.println(getTimeStamp() + "Forward button is greyed out (unavailable).");
            return;
        }
        else {
            backPages.push(currentPage);
            currentPage = forwardPages.pop();
        }
        System.out.println(getTimeStamp() + "Forward button is clicked. Currently viewing: " + currentlyViewing());
    } // contd. on the next slide
}
```

# Code

```
public class BrowserTab {  
  
    public String currentlyViewing() { return currentPage; }  
  
    public String toString() {  
        StringBuilder prettyString = new StringBuilder();  
  
        for (String site : previousPages)  
            prettyString.append(site).append(" ");  
  
        prettyString.append(" ***").append(currentPage).append(" *** ");  
  
        for (String site : nextPages)  
            prettyString.append(site).append(" ");  
  
        return prettyString.toString();  
    }  
}
```

For an usage refer to the class **TestBrowserTab**

# Queue

## Definition

A queue is an ADT where items are inserted (**enqueued**) and removed (**dequeued**) according to the **first-in, first-out** (FIFO) principle



# Operations on a queue

Let us denote a queue by Q.



- 1 **Q.enqueue(e)**: **adds** an element e to the back of Q
- 2 **Q.dequeue()**: **removes** and **returns** the first element from Q; returns **null** if Q is empty
- 3 **Q.first()**: **returns** the first element of Q, without removing it; returns **null** if Q is empty
- 4 **Q.size()**: **returns** the number of elements stored in Q currently
- 5 **Q.isEmpty()**: **returns** **true** if Q is empty, otherwise returns **false**

enqueue(55)

55

55, 3



enqueue(69)

55, 3, 69

enqueue(-88)

55, 3, 69, -88

`enqueue(-17)`

55, 3, 69, -88, -17

`isEmpty()` returns false

55, 3, 69, -88, -17

`size()` returns 5

55, 3, 69, -88, -17

**dequeue() removes and returns 55**

3, 69, -88, -17

`dequeue()` removes and returns 3

69, -88, -17

`first()` returns 69

69, -88, -17



# Applications

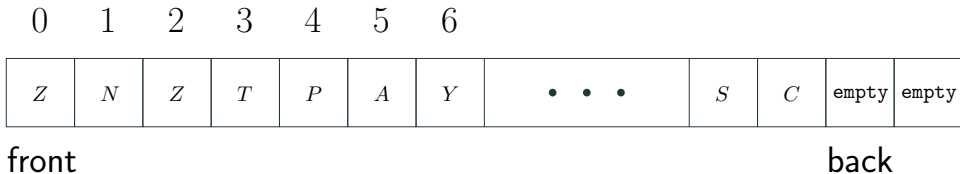
- Simulation of real-world queues (airlines, ticket counter, etc.)
- Graph algorithms
- Resource sharing in multi-user systems
- Operating systems
- Computer networks
- ...

## An interface for queues

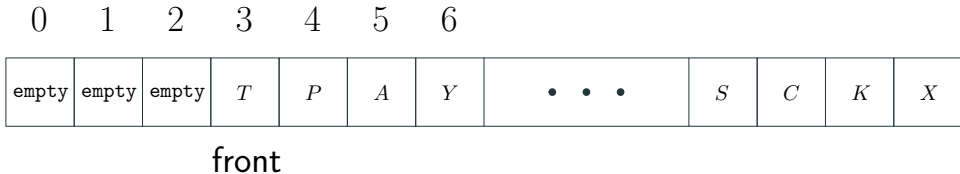
```
public interface QueueADT<E> {  
    void enqueue(E e);  
    E dequeue();  
    int size();  
    boolean isEmpty();  
    E first();  
}
```

To use a queue, one should create a class by implementing this interface

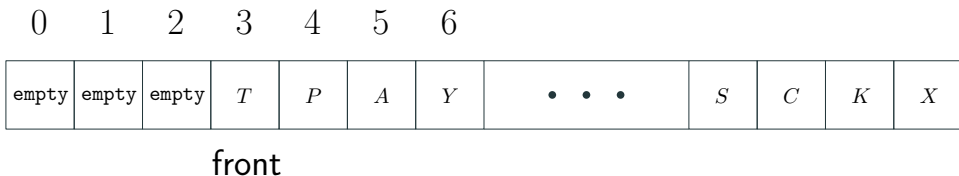
## Implementing using array



After 3 dequeues and 2 enqueues the situation is: we have space but we cannot enqueue anymore!

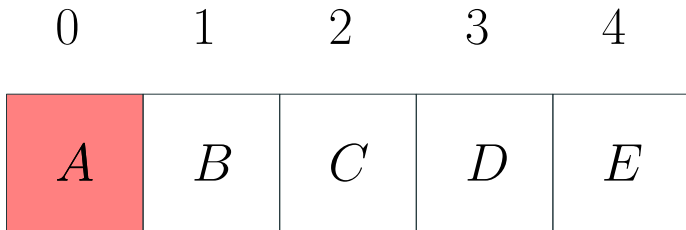


## What to do then?



- Trivial solution: shift the content to the left and add new stuff
- Downside? This is very expensive! Takes  $O(n)$  time
- We need to implement enqueue faster than this
- Solution: wrap around

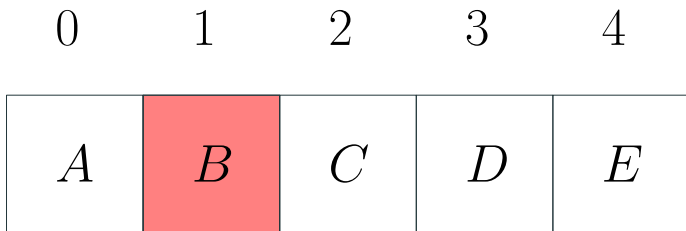
## Using the mod operator for wrapping around



$$i = 0$$

$$\text{Next index: } (i + 1) \bmod 5 = 1 \bmod 5 = 1$$

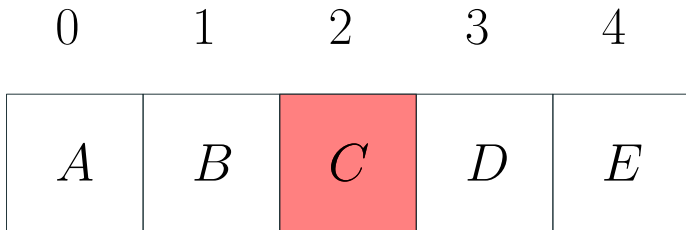
## Using the mod operator for wrapping around



$$i = 1$$

$$\text{Next index: } (i + 1) \bmod 5 = 2 \bmod 5 = 2$$

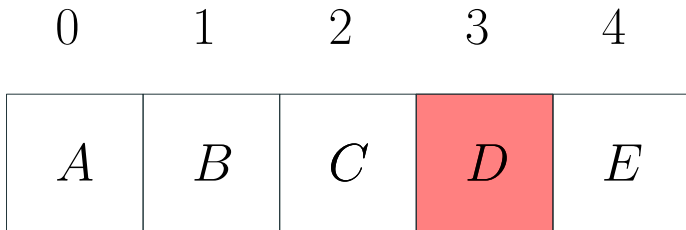
## Using the mod operator for wrapping around



$$i = 2$$

Next index:  $(i + 1) \bmod 5 = 3 \bmod 5 = 3$

## Using the mod operator for wrapping around

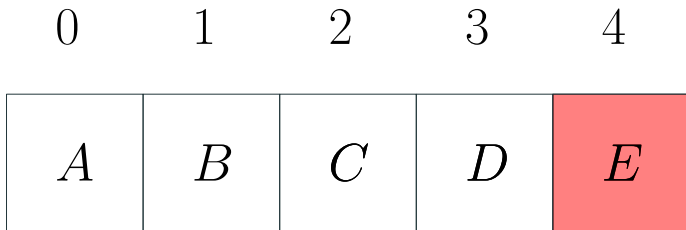


$$i = 3$$

$$\text{Next index: } (i + 1) \bmod 5 = 4 \bmod 5 = 4$$



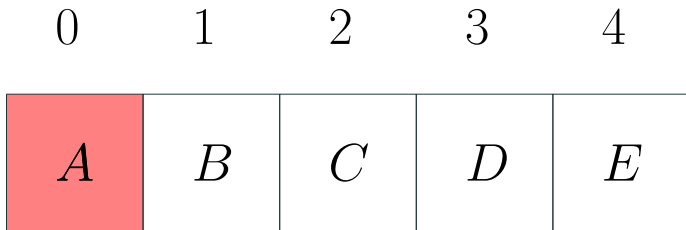
## Using the mod operator for wrapping around



$$i = 4$$

$$\text{Next index: } (i + 1) \bmod 5 = 5 \bmod 5 = 0$$

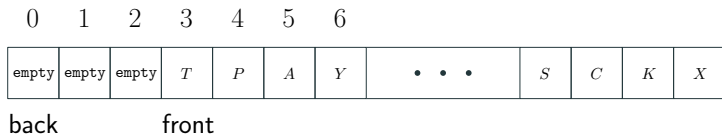
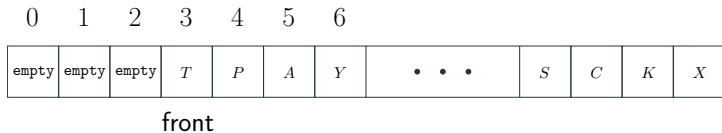
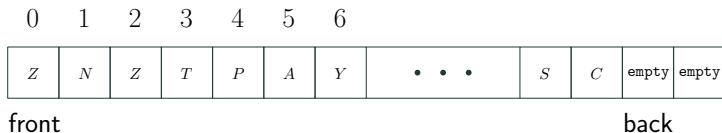
## Using the mod operator for wrapping around



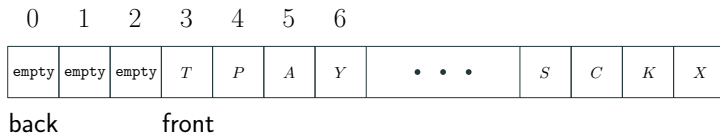
$$i = 0$$

$$\text{Next index: } (i + 1) \bmod 5 = 1 \bmod 5 = 1$$

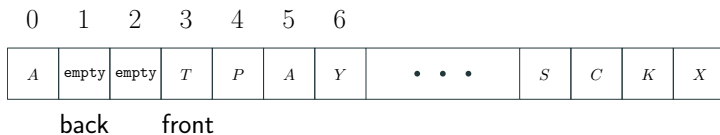
# Using this wrap-around approach for array-based queues



## Using this wrap-around approach for array-based queues



After `Q.enqueue('A');`, we get



## So how to implement a queue using the QueueADT interface?

- Use an array (queue size must be known in advance, speedy in practice)  
See the class [ArrayQueue](#)
- Use a linked-list (can grow arbitrarily, a bit slower than the array-based queue)  
See the class [LinkedList](#)

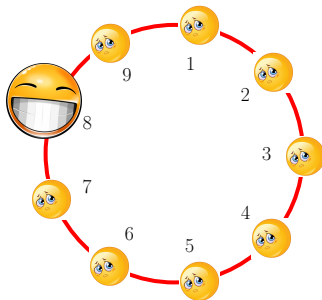
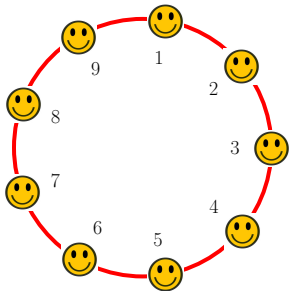
Irrespective of the implementation, every one of these five queue operations, `size()`, `isEmpty()`, `enqueue()`, `dequeue()`, `first()`, takes  $O(1)$  time

# An application: the Josephus problem

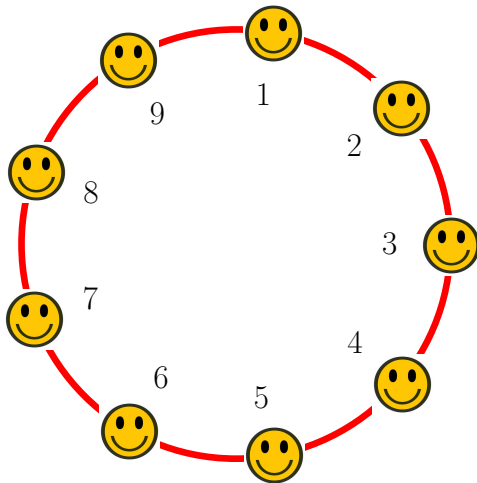
## The problem

A group of  $n$  people agree to play the following fun game. They arrange themselves on a circle (at positions numbered from 1 to  $n$ ) and proceed around the circle clockwise, eliminating every  $m$ th person until only one person is left.

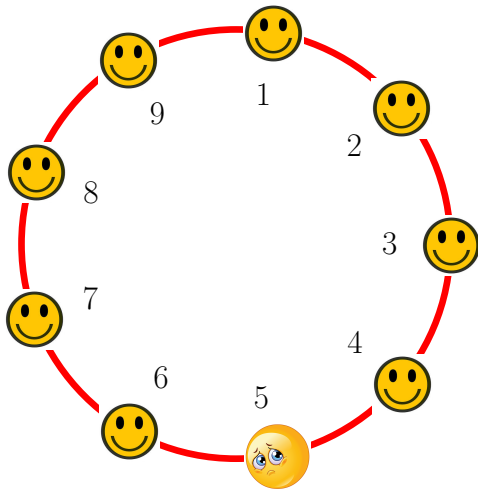
The last remaining person wins the game. How to figure out the winning position?



**Example:**  $n = 9, m = 5$

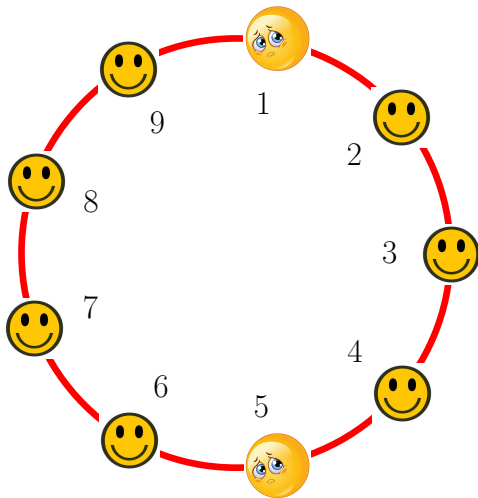


**Example:**  $n = 9, m = 5$

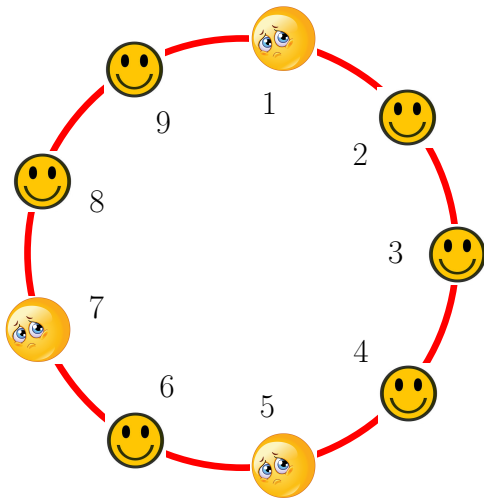




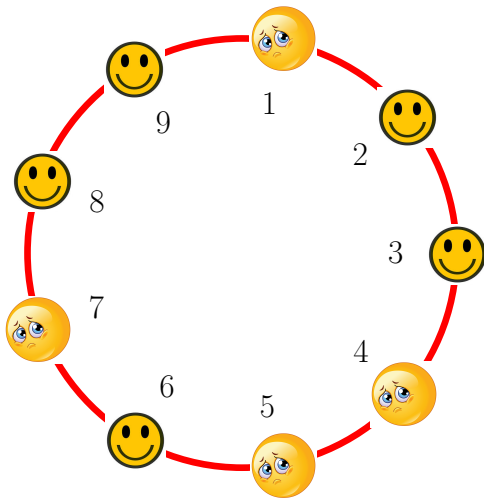
**Example:**  $n = 9, m = 5$



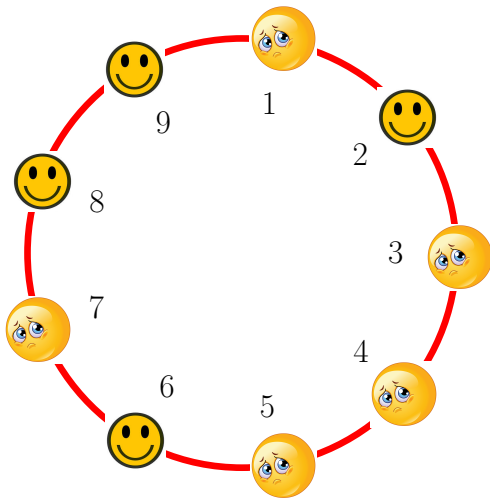
**Example:**  $n = 9, m = 5$



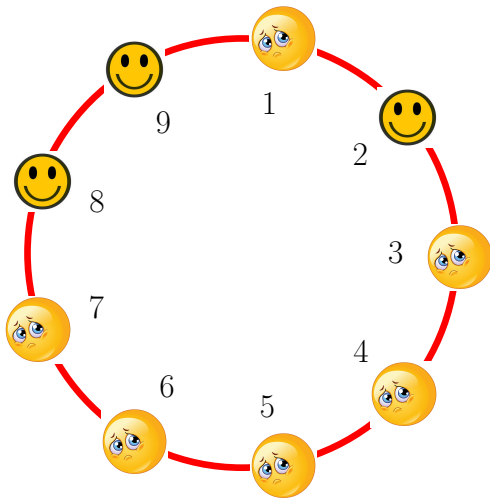
**Example:**  $n = 9, m = 5$



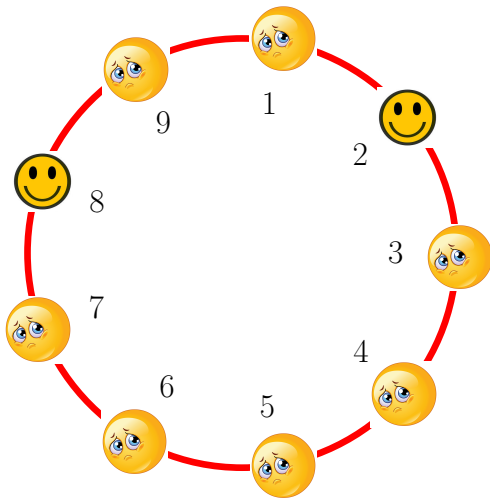
**Example:**  $n = 9, m = 5$



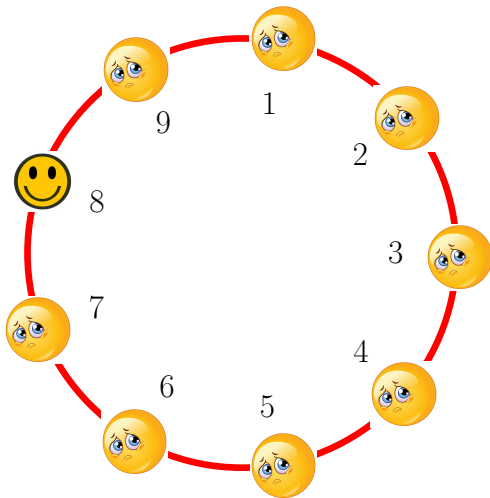
**Example:**  $n = 9, m = 5$



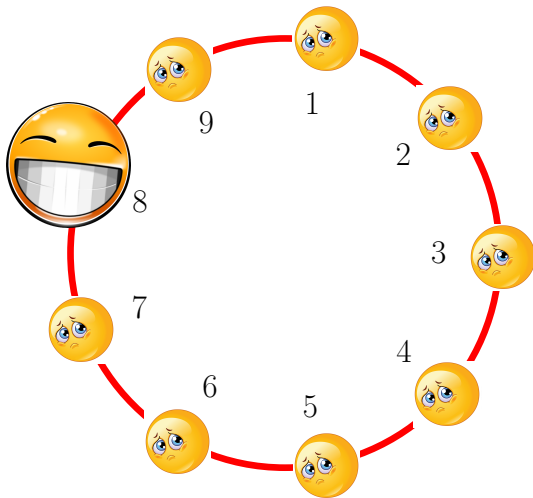
**Example:**  $n = 9, m = 5$



**Example:**  $n = 9, m = 5$



**Example:**  $n = 9, m = 5$





# Code

```
import java.util.Scanner;

public class JosephusSolver {
    public static void main(String[] args) {
        int n = 9, m = 5;

        LinkedList<Integer> Q = new LinkedList<>();

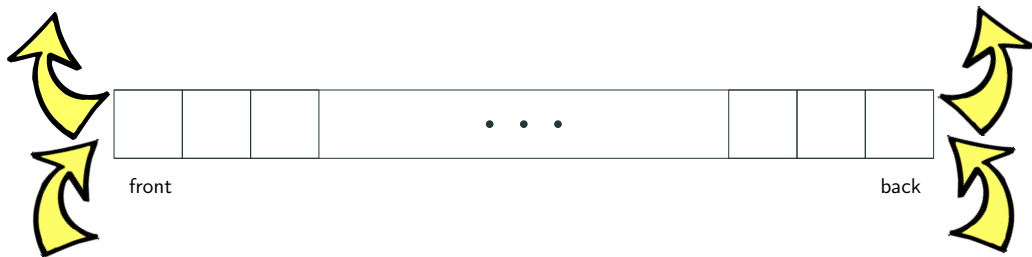
        for (int i = 1; i <= n; i++) // populate the queue with the integers 1,2,...,n
            Q.enqueue(i);

        while ( ( Q.size() > 1 ) ) {
            for (int i = 0; i < m-1; i++) // send the first m-1 elements to the back of the queue
                Q.enqueue(Q.dequeue());

            System.out.println("Eliminating player: " + Q.dequeue() + " ");
        }

        System.out.print("\nWinning position: " + Q.first()); // the sole integer in the queue is the winner
    }
}
```

## Double ended queues: Deque (pronounced as DECK)



Supports insertion and removal at both the front and the back

Deque can be used as a stack or a queue

# Deque ADT

```
public interface DequeADT<E> {  
    int size();  
    boolean isEmpty();  
    E first();  
    E last();  
    void addFirst(E e);  
    void addLast(E e);  
    E removeFirst();  
    E removeLast();  
}
```

## An easy way to implement a deque

Doubly linked-list (using a singly linked-list will make the `removeLast` operation run in  $O(n)$  time)

👉 Refer to the class [LinkedDeque](#) for an implementation

## Time complexity

Just like stacks and queues, all operations on deques run in  $O(1)$  time

## Suggested exercise

Create an array-based iterable deque class

## Stacks and queues from **Chapter 5**

<https://opensa-server.cs.vt.edu/OpenDSA/Books/CS3/html/index.html>