Piotr Gajowniczek Instytut Telekomunikacji PW

Data Transmission



MPLS – Multiprotocol Label Switching

fundamentals and applications

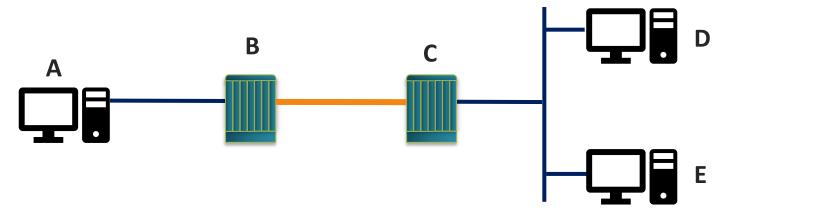
### MPLS – introduction

- RFC 3031, RFC 3032 MPLS architecture
  - initially invented in the '90s to improve the forwarding speed of IP routers and introduce additional functionality to IP control plane
- Nowadays, MPLS is very popular in the ISP core networks
  - network virtualization (VPNs)
  - resource and service management
  - traffic engineering, QOS
  - network resiliency
- Extension beyond the ISP core applications also exist
  - □ GMPLS (Generalized MPLS) single control plane extended to optical transport technologies
  - Seamless MPLS extension of MPLS towards the access network

# Layered model / switching vs routing

- Switching (L2)
  - findin the destination MAC address in
     MAC table (forwarding table, bridge table)
  - establish the outbound interface
  - send packet from input queue to output queue
  - packet is not modified in any way

- Routing (L3)
  - remove L2 header
  - □ lookup for next hop (forwarding table; routing table)
  - determine the correct destination L2 (MAC) address
  - rewrite MAC header
  - send packet to the output queue

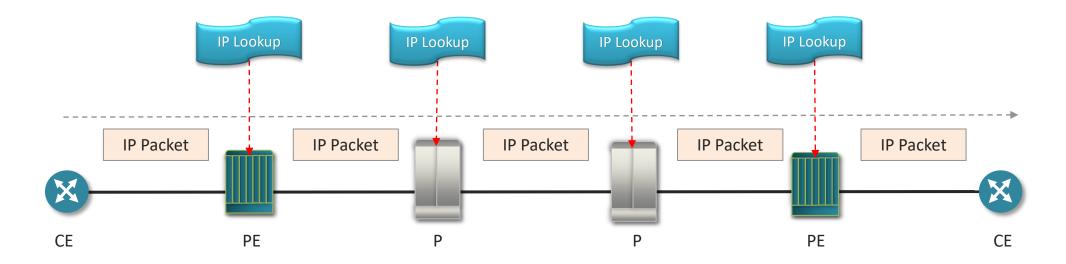


□ L3 switch?

# IP routing ...

### *IP* routing:

- data link layer frame validation
- network-layer protocol demultiplexing
- IP packet validation
- forwarding decision longest prefix match
- data link frame construction



## ... vs MPLS

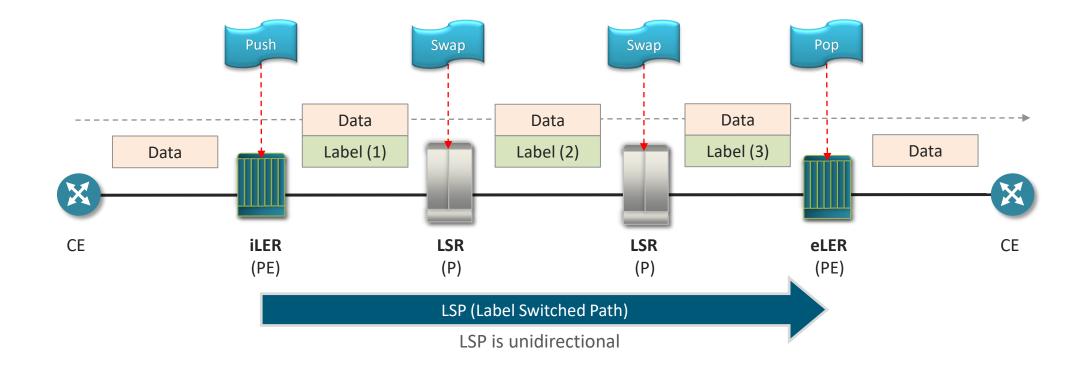
### *IP* routing:

- data link layer frame validation
- network-layer protocol demultiplexing
- IP packet validation
- forwarding decision longest prefix match
- data link frame construction



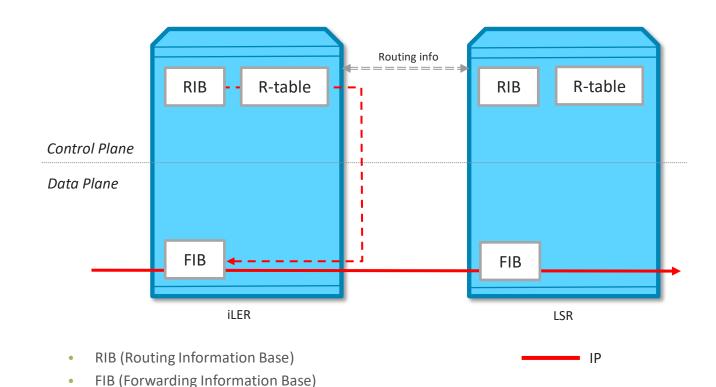
### **MPLS**:

Label Push, Swap & Pop



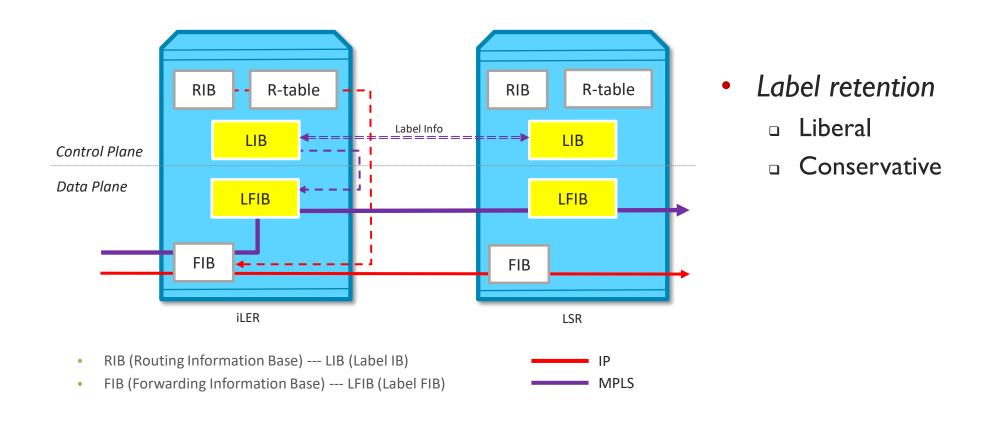
# IP control plane ...

- FEC (Forwarding Equivalence Class)
  - □ IP routing FEC = IP Prefix; FEC lookup done at each hop
  - MPLS other FEC criteria possible, FEC lookup only at an iLER

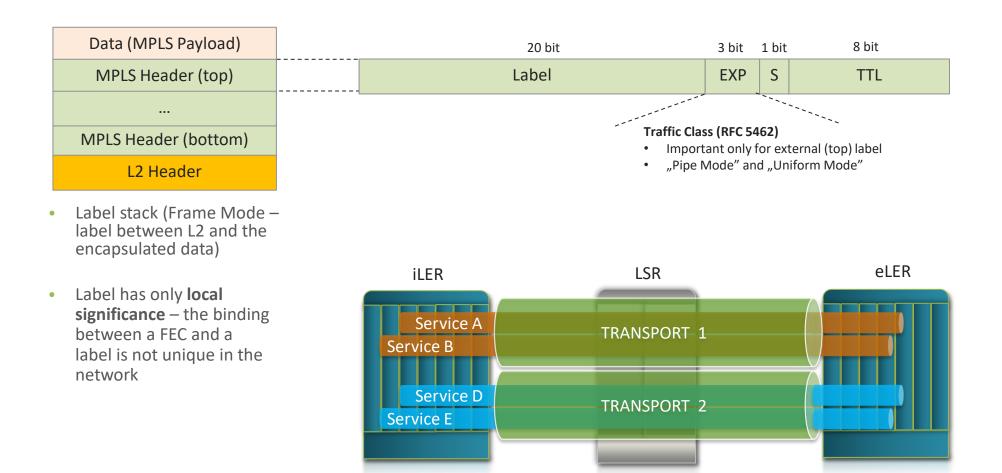


## ... vs MPLS control plane

- FEC (Forwarding Equivalence Class)
  - □ IP routing FEC = IP Prefix; FEC lookup done at each hop
  - MPLS other FEC criteria possible, FEC lookup only at an iLER

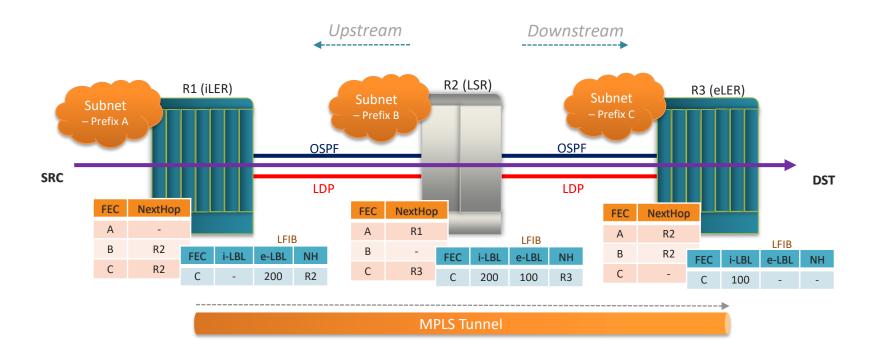


## MPLS – labels and tunnels



LSR handles only transport tunnels

## MPLS – tunel set-up



- MPLS tunnel set-up requires:
  - FEC (IP prefixes)reachability = OSPF
  - distribution of label
     mappings between routers

- Label distribution protocols
  - LDP (Label Distribution Protocol) – "Downstream Unsolicited"
  - RSVP-TE (Resource Reservation Protocol) – "Downstream on Demand"

- RI Request(FEC C)
- R3 Response(FEC C, 100)
- R2 Response(FEC C, 200)

# MPLS – label distribution protocols

### transport tunnels

- LDP (Label Distribution Protocol)
  - TCP based
  - tunnels built based on IGP (full-mesh)
  - simple configuration
  - automatic creation of tunnels
  - no traffic engineering
  - convergence time depends on IGP
  - label distribution in "downstream unsolicited" approach

- RSVPTE (Resource Reservation Protocol with Traffic Engineering)
  - explicit tunnels (not following IGP paths)
  - additional constraints (administrative and TE-related) for advanced path calculation
  - bandwidth reservation for LSP (CAC)
  - advanced LSP protection against failures
  - □ label distribution in "downstream on demand" approach

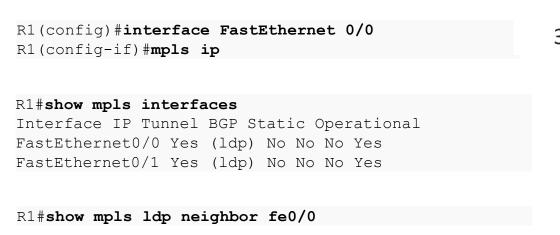
#### service tunnels

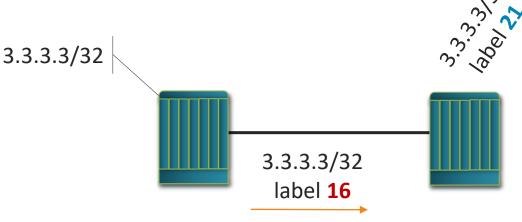
- T-LDP (Targeted LDP) RFC 4447
  - multi-hop LDP
  - for L2 services
  - for end-to-end tunnels between PE routers

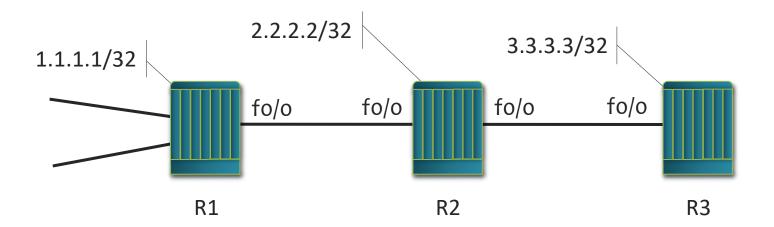
- MP-BGP (Multi-Protocol BGP) RFC 4364
  - for L3 services

# LDP (Label Distribution Protocol)

- To establish LDP adjacency:
  - router sends UDP multicast (224.0.0.2:646) hello packets to discover other neighbors
  - next, builds the neighbor adjacency (on loopback addresses) using a TCP connection for exchanging label information (each router has a unique ID called the LSR ID; much like in OSPF).
- LDP creates a label binding for each prefix by default and sends them to neighbors (downstream unsolicited) Label Mapping Advertisement, Label Withdraw
- LDP configuration







LIB (Label Information Base)

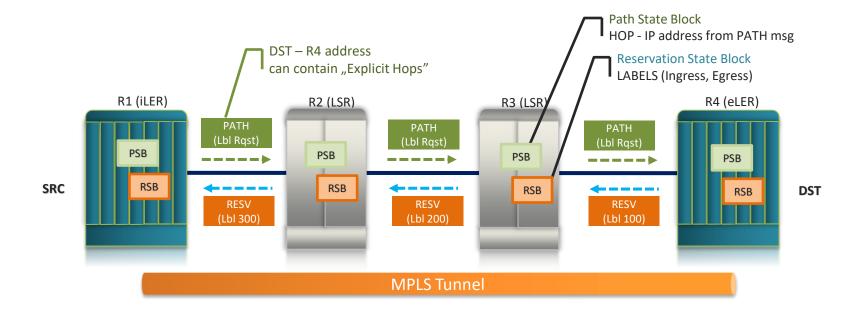
```
R1#show mpls ldp bindings
lib entry: 2.2.2.2/32, rev 7
local binding: label: 16
remote binding: lsr: 2.2.2.2:0, label: imp-null
lib entry: 3.3.3.3/32, rev 9
local binding: label: 17
remote binding: lsr: 2.2.2.2:0, label: 21
```

## LFIB (Label Forwarding Information Base)

R1#show mpls forwarding-table					
Loca	l Outgoing	Prefix	Bytes Label	Outgoing	Next Hop
Label Label		or Tunnel	Switched	interface	
		Id			
16	Pop Label	2.2.2.2/32	0	F0/0	10.1.1.1
17	21	3.3.3.3/32	0	F0/0	10.1.1.1

• implicit null label – penultimate hop popping (PHP)

# MPLS – the RSVP protocol



- RFC 3209 RSVP as LDP
- RSVPTE:
  - LSP definition
  - path calculation "outside" IGP metrics ("link colors", bandwidth etc.)
  - tunel protection (Secondary Paths, Fast Reroute)
  - resource reservation (CAC)

- MPLS tunnel can be composed of many paths (LSP-Paths).
  - one "primary" path and seven "secondary"
  - only one active at a time

- Other message types:
  - PATH Tear: (downstream), RESVTear: (upstream)
  - PATH Error, RESV Error:
  - Hello (RSVP heartbeat)
  - Summary Refresh (for less signalling)

# RSVP Traffic Engineering: path calculation

- Source routing
- Path Option Explicit
  - manual configuration at source router
  - high signalling overhead
- Path Option Dynamic = APC (Advanced Path Calculation)
  - CSPF (Constrained Shortest Path First)
  - additional criteria
    - bandwidth reservation state
    - administrative criteria (link colors)
    - hop limit
    - TE metric
    - Explicit route (,,strict hops", ,,loose hops")
    - Shared Link Groups

### reservations are made in the Control Plane

- actual bandwidth usage in the Data
   Plane is not considered
- requires relevant QoS solutions in the Data Plane

### resiliency

allows automated creation of backup
 paths and detours (Fast Reroute) that
 are disjoint with the primary path

## RSVP-TE – how to force tunner route?

## Signalling

- information about the route is conveyed in the RSVP PATH message in an ERO (Explicit Route Object)
- ERO is updated in each intermediate router

#### Bandwidth reservation

- CSPF algorithm calculates a path with the required amount of unreserved bandwidth using data from TED database at source router
- downstream:
  - reservation request is signaled in RSVP PATH message
  - each router checks bandwidth availability on outgoing link (CAC)
- upstream:
  - bandwidth is reserved in each router on path (RSVP RESV message)
  - Unreserved Bandwidth updated and advertised

#### Least-Fill Bandwidth Reservation rule

- if CSPF has found multiple paths with the same metric
- relevant QOS policies in the Data Plane are required

## RSVP – TE variables

the need for additional constraints and link state data has to be reflected in routing protocol

- OSPF-TE (OSPF Traffic Engineering)
  - RFC 2370:The OSPF Opaque LSA Option
  - Opaque LSAs deneric mechanism for OSPF extensions

#### **OSPFTE**

- routers create additional database TED (Traffic Engineering Database) for storing additional link attributes distributed by Opaque LSAs (Type 10)
- Opaque LSA Flooding activated when:
  - link state (up/down), link configuration of bandwidth reservation state changes
  - periodically (as in IGP)
- Opaque LSA Type 10 contains Link TLV object, used to advertise information about links handled by RSVP-enabled routers:
  - link type, link ID
  - IP addresses of interfaces on both sides of the link
  - TE metrics
  - maximum bandwidth
  - maximum reservable bandwidth (per LSP priority)
  - unreserved bandwidth (100 = 100%, overbooking possible)
  - administrative group
  - Shared Risk Link Group (SRLG)

# MPLS – priorities and preemption

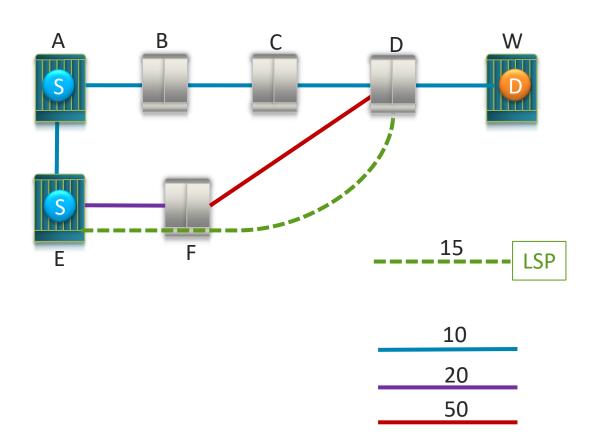
- LSP Soft Preemption
  - higher priority LSPs can preempt lower priority paths
  - priorities work in conjunction with knowledge of the Unreserved Bandwidth parameter – current values are advertised by OSPF TE for each priority level
- Setup and Hold priorities (0 to 7, lower value = higher priority)
  - LSP A can preempt LSP B if Setup
     Priority(A) < Hold Priority(B)</li>
  - LSP priorities are signaled in RSVP PATH message, in SESSION\_ATTRIBUTE object

- RSVP Preemption-Timer & LSP Retry-Timer
  - preemption by MBB (Make Before Break)
  - CSPF tries to find another route for preempted LSP
    - periodically (Retry-Timer)
    - preemption (status = down) after time defined in Preemption-Timer (unless a new route was found earlier)

# Using MPLS tunnels in IP routing

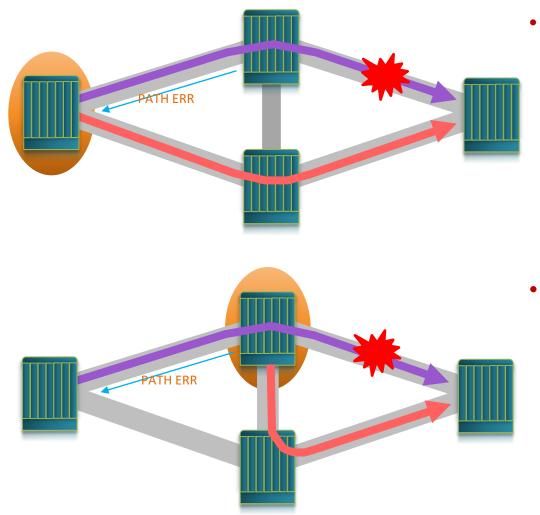
## Static mapping

- LSR is configured to send packets through the tunnel
- scalability issues
- Dynamic mapping (done by IGP)
  - LSP is treated as an interface (tunnel) with an associated metric
  - metric can influence routing decisions (choice between CSPF and SPF)
  - two cases:
    - IGP can use tunnel for SPF calculation only in an iLER (autoroute announce)
    - tunnel may be signaled in LSA for use by other routers to calculate shortest paths (forwarding adjacency)



# MPLS – failure resiliency

- factors influencing quality of protection
  - □ avg. failure detection speed
    - OSPF Hello (30 s)
    - RSVP Hello (9 s)
    - Bidirectional Forwarding Detection (<|s)</li>
      - » "IP level heartbeat"
  - speed of failure advertising
  - service restoration time (switchover speed)



### Secondary LSP

- switchover at source router
- switchover time depends on PATH ERR message delivery to source router
- max. 7 standby (Secondary) paths
  - Hot Standby
  - Non-Standby

#### Fast Reroute

- local switchover (<50 ms)</li>
- PATH ERR conveys only the information about failure
- detour paths are calculated automatically (CSPF)
- protects against node and link failures
- protection types:
  - One-to-One Backup (Detour)
  - Facility Backup (Bypass Tunnel)