## **Extending Trace Estimation to Schatten Norm Estimation**

The methods outlined in this document follow sections 4.8 through 5.2, and 5.4 and 5.6 of Martinsson & Tropp (2021)<sup>1</sup>. Proofs and algorithms for these methods follow. Please find the accompanying code in <sup>2</sup>.

## The Frobenius norm and the Schatten 4-norm

Randomized trace estimators of the form given in the introduction can also be used to estimate some matrix norms, notably the Frobenius or Schatten 2-norm and the Schatten 4-norm, also referred to as the  $\ell_2$  and  $\ell_4$  norms.

Consider a rectangular matrix  $\mathbf{B} \in \mathbb{R}^{mxn}$ , accessed by a matrix-vector product  $\mathbf{u} \mapsto \mathbf{B}\mathbf{u}$ . Suppose that we extract test vectors from the standard normal distribution. Draw a standard normal matrix  $\mathbf{\Omega} \in \mathbb{R}^{nxk}$  with columns  $\mathbf{\omega}_i$ . Then construct the random variable:

$$\overline{X}_{k} := \frac{1}{k} \| \mathbf{B} \mathbf{\Omega} \|_{F}^{2} = \frac{1}{k} \sum_{i=1}^{k} \omega_{i}^{'} (\mathbf{B}' \mathbf{B}) \omega_{i} =: \frac{1}{k} \sum_{i=1}^{k} X_{i}$$

$$\tag{1}$$

To analyse  $\overline{X}_k$ , note that it is an instance of the randomized trace estimator in (1), where  $\mathbf{A} = \mathbf{B'B}$ .

Proposition 3:  $\overline{X}_k$  is an unbiased estimator for  $\|\mathbf{B}\|_F^2$ , with  $\operatorname{Var}(\overline{X}_k) = \frac{2}{k} \|\mathbf{B}\|_4^4$ .

Proof: 
$$\mathbb{E}[\overline{X}_k] = \mathbb{E}[\frac{1}{k}\sum_{i=1}^k \omega_i'(\mathbf{B'B})\omega_i] = \frac{1}{k}\mathbb{E}[\sum_{i=1}^k \omega_i'(\mathbf{B'B})\omega_i]$$
  $\frac{1}{k}$  can be factored out by linearity; 
$$= \frac{1}{k}\sum_{i=1}^k \mathbb{E}[\omega_i'(\mathbf{B'B})\omega_i]$$
 because  $X_i = \omega_i'(\mathbf{B'B})\omega_i$  are iid; 
$$= \frac{1}{k}\sum_{i=1}^k (\mathbf{B'B})\mathbb{E}[\omega_i'\omega_i]$$
 (B'B) can be factored out by linearity; 
$$= \frac{1}{k}k (\mathbf{B'B})\sum_{i=1}^k \mathbb{E}[\omega_i'\omega_i] = (\mathbf{B'B})\sum_{i=1}^k \mathbb{E}[\omega_i'\omega_i]$$
 rewrite (B'B) as: 
$$= (\sum_{l=1}^n b_{il}b_{lj})\sum_{i=1}^k \mathbb{E}[\omega_i'\omega_i] = \sum_{i=1}^n \sum_{j=1}^n b_{ij}^2 \sum_{i=1}^k \mathbb{E}[\omega_i'\omega_i]$$
 use the isotropic property of  $\omega_i$ : 
$$= \|\mathbf{B}\|_F^2$$

$$\operatorname{Var}(\overline{X}_{k}) = \operatorname{Var}(\frac{1}{k} \sum_{i=1}^{k} \omega_{i}'(\mathbf{B}'\mathbf{B})\omega_{i})$$

$$= \frac{1}{k^{2}} \operatorname{Var}(\sum_{i=1}^{k} \omega_{i}'(\mathbf{B}'\mathbf{B})\omega_{i})$$

$$= \frac{1}{k^{2}} \sum_{i=1}^{k} \operatorname{Var}(\omega_{i}'(\mathbf{B}'\mathbf{B})\omega_{i})$$

$$= \frac{1}{k^{2}} \sum_{i=1}^{k} \operatorname{Var}(\omega_{i}'(\mathbf{B}'\mathbf{B})\omega_{i})$$
because  $X_{i} = \omega_{i}'(\mathbf{B}'\mathbf{B})\omega_{i}$  are iid;

<sup>&</sup>lt;sup>1</sup> https://arxiv.org/abs/2002.01387v1, last accessed 4/30/21

<sup>&</sup>lt;sup>2</sup> https://github.com/ghostpress/comp-stats-sims/tree/final-project/final-project/trace-estim

$$= \frac{1}{k^2} \sum_{i=1}^{K} (\mathbf{B'B}) \operatorname{Var}(\omega_i^{'}\omega_i^{'})(\mathbf{B'B})'$$

$$= \frac{1}{k^2} \sum_{i=1}^{k} (\mathbf{B'B}) \mathbb{E}[(\omega_i^{'}\omega_i^{'} - \mathbb{E}[\omega_i^{'}\omega_i^{'}])(\omega_i^{'}\omega_i^{'} - \mathbb{E}[\omega_i^{'}\omega_i^{'}])'](\mathbf{B'B})$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} \mathbb{E}[(\omega_i^{'}\omega_i^{'} - \mathbf{I})(\omega_i^{'}\omega_i^{'} - \mathbf{I})']$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} \mathbb{E}[(\omega_i^{'}\omega_i^{'} - \mathbf{I})(\omega_i^{'}\omega_i^{'} - \mathbf{I})]$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} \mathbb{E}[(\omega_i^{'}\omega_i^{'}\omega_i^{'}\omega_i^{'} - \omega_i^{'}\omega_i^{'} - \omega_i^{'}\omega_i^{'} - \mathbf{I}]$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} (\mathbb{E}[\omega_i^{'}\omega_i^{'}\omega_i^{'}\omega_i^{'}] - 2\mathbb{E}[\omega_i^{'}\omega_i^{'}] - \mathbf{I})$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} (\mathbb{E}[\omega_i^{'}\omega_i^{'}]\mathbb{E}[\omega_i^{'}\omega_i^{'}] - \mathbf{I})$$
because  $\omega$  are independent;
$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} (\mathbf{I} - \mathbf{3I})$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} (\mathbf{I} - \mathbf{3I})$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) \sum_{i=1}^{k} 2\mathbf{I}$$

$$= \frac{1}{k^2} (\mathbf{B'B})(\mathbf{B'B}) 2k$$

$$= \frac{2}{k} (\mathbf{B'B})(\mathbf{B'B})$$

$$= \frac{2}{k} \| \mathbf{B} \|_4^4$$

From these results we can also clearly see that rescaling  $Var(\overline{X}_k)$  by  $\frac{k}{2}$  gives an unbiased estimate for  $\|\mathbf{B}\|_4^4$ , the Schatten 4-norm of  $\mathbf{B}$ .

 $\overline{X}_k$  is computed by simulating nk standard normal variables, taking k matrix-vector products with  $\mathbf{B}$ , and performing O(kn) additional arithmetic. Therefore the total runtime of this method is O(2kn + knm).

## Schatten p-norm Estimation by Sampling

We now extend the discussion of approximating Schatten 2- and 4-norms by sampling to the 2p-norm for each  $p \in \mathbb{N}$ . Consider the general matrix  $\mathbf{B} \in \mathbb{R}^{mxn}$ , accessed via the matrix-vector product  $\mathbf{u} \mapsto \mathbf{B}\mathbf{u}$ . For a sample size k, let  $\mathbf{\Omega} \in \mathbb{R}^{nxk}$  be a (random) test matrix that does not depend on  $\mathbf{B}$ . For a natural number  $p \geq 3$ , the problem is to estimate the Schatten 2p-norm  $\|\mathbf{B}\|_{2p}$  from the sample matrix  $\mathbf{Y} = \mathbf{B}\mathbf{\Omega}$  such that  $\mathbf{Y}$  is an unbiased estimator of the norm. Moreover, the methods used should ideally be less expensive than the  $O(\min\{mn^2, nm^2\})$  cost of computing the Singular Value Decomposition (SVD) of  $\mathbf{B}$ , as in the classical algorithm.

The authors remark that the sample size k needed to estimate  $\|\mathbf{B}\|_{2p}$  up to a fixed constant factor with 75% probability is unfortunately  $k \ge \min\{m, n\}^{1-2/p}$ . In other words, the sample size must grow polynomially with the dimensions of the matrix for p > 2. Nevertheless, as shown below, the algorithm used is still faster than the SVD method.

Assume that the random test matrix  $\Omega \in \mathbb{R}^{nxk}$  has isotropic columns  $\omega_i$  that are iid. Form the sample matrix  $\mathbf{Y} = \mathbf{B}\Omega$ . Abbreviate  $\mathbf{A} = \mathbf{B'B}$  and  $\mathbf{X} = \mathbf{Y'Y}$ . Then:

$$(\mathbf{X})_{ij} = (\mathbf{Y}'\mathbf{Y})_{ij} = \omega_i^{'} \mathbf{A} \omega_i$$
 (2)

For any natural numbers that satisfy  $1 \le i_1, ..., i_p \le k$ ,

$$(\mathbf{X})_{i_1 i_2} (\mathbf{X})_{i_2 i_3} \dots (\mathbf{X})_{i_p i_1} = trace(\omega_{i_1} \omega_{i_1}^{'} \mathbf{A} \dots \omega_{i_p} \omega_{i_p}^{'} \mathbf{A})$$

Under the assumption that  $i_1, ..., i_p$  are distinct, by independence and isotropy the expectation becomes:

$$\mathbb{E}[(\mathbf{X})_{i_1 i_2} (\mathbf{X})_{i_2 i_3} \dots (\mathbf{X})_{i_p i_1}] = trace(\mathbf{A}^p) = \| \mathbf{B} \|_{2p}^{2p}$$

And now define the estimator as below, where C(k, p) is the binomial coefficient with k and p:

$$V_p = C(k, p)^{-1} \sum_{1 \le i_1 \le i_p \le k} (\mathbf{X})_{i_1 i_2} (\mathbf{X})_{i_2 i_3} \dots (\mathbf{X})_{i_p i_1}$$

Finally, we can reformulate  $V_p$ . Let  $T: \mathbb{H}_k \to \mathbb{R}^{kxk}$  be the linear map that reports the strict upper triangle of a symmetric matrix. Then:

$$V_p = C(k, p)^{-1} \operatorname{trace}(T(\mathbf{X})^{p-1}\mathbf{X})$$
(3)

The algorithm to compute  $V_p$  follows.

Algorithm 2: Schatten 2p-norm estimation by random sampling

**Input:** Matrix  $\mathbf{B} \in \mathbb{R}^{m \times n}$ , the order p of the norm to estimate, and the number k of samples to take **Output:** Schatten 2p-norm estimate  $V_p$ 

- 1. Draw the test matrix  $\Omega \in \mathbb{R}^{nxk}$  with iid isotropic columns
- 2. Compute the sample matrix  $Y = B\Omega$
- 3. Form the Gram matrix  $\mathbf{X} = \mathbf{Y'Y} \in \mathbb{R}^{kxk}$
- 4. Extract the strict upper triangle T = T(X)
- 5. Compute  $T^{p-1}$  by repeated squaring
- 6. Return  $V_p = trace(\mathbf{T}^{p-1}\mathbf{X})$

The runtime of this algorithm is dominated by the  $O(k^2n)$  arithmetic required to form **X** given **Y**.