

Preliminar notes on maximum-flow problem solved in FCPP.

Notation:

- $G$  will denote the directed weighted graph of capacities,  $s$  and  $t$  will denote respectively source and sink of our graph.
- For a path  $\mathbf{p}$  we will denote with  $|\mathbf{p}|$  its lenght.

When it will be convenient, we will treat a directed weighted graph  $G$  as a function  $G : V \times V \rightarrow \mathbb{R}_+$ , where  $V$  is its set of vertices. With a little abuse of notation we will say that the edge from a node  $\delta$  to a node  $\delta'$  is in  $G$ , or  $(\delta, \delta') \in G$ , in place of  $G(\delta, \delta') \neq 0$ .

Likewise we will say that a path  $\mathbf{p}$  is contained in  $G$ , or  $\mathbf{p} \subset G$ , if every edge in  $\mathbf{p}$  is in  $G$ .

- For a field  $\Phi$ ,  $\mathcal{R}_\Phi$  will denote the directed weighted graph of residual capacities respect to  $\Phi$ :

$$\mathcal{R} = \begin{cases} G + \Phi & \text{where } \Phi < 0 \\ 0 & \text{otherwise} \end{cases} \quad (1)$$

- We define admissible path respect to  $\Phi$  a path  $\gamma$  in  $G$  from source to sink such that

$$\Phi|_\gamma < G|_\gamma$$

- We say that a flow  $\Phi$  is maximal if there are no admissible paths in  $G$  respect to  $\Phi$ .

**Proposizione 0.1.**  $\Phi$  is a maximum flow iff  $\mathcal{R}_\Phi$  does not have paths from source to sink.

*Proof.* Let's show that a path  $\gamma$  in  $G$  is admissible iff  $\gamma$  is a path in  $\mathcal{R}_\Phi$ . On every edge  $e$  in  $G$  we have

$$\begin{aligned} 0 < \mathcal{R}_\Phi(e) &\Leftrightarrow \begin{cases} \Phi(e) \leq 0 < \mathcal{R}_\Phi(e) = G(e) & \text{or} \\ 0 < \mathcal{R}_\Phi(e) = G(e) - \Phi(e) \end{cases} \\ &\Leftrightarrow \Phi(e) < G(e) \quad \text{and} \quad 0 < G(e) \end{aligned}$$

Hence

$$0 < \mathcal{R}_\Phi|_\gamma \Leftrightarrow \Phi|_\gamma < G|_\gamma \quad \text{and} \quad 0 < G|_\gamma$$

This is what we wanted to prove.  $\square$

For a field-value  $f$  in  $\delta$  we define  $|f| := \sum_{\delta \mapsto \delta'} f(\delta')$ .

For a pre-flow  $\Phi$  we define

$$\|\Phi\| := \sum_{\substack{\delta \text{ connected to} \\ t \text{ in } \mathcal{R}_\Phi}} |\mathcal{R}_\Phi(\delta)|$$

**Lemma 0.2.** Let  $\Phi$  be a pre-flow. Then  $\|\mathbf{update}(\Phi)\| \leq \|\Phi\|$  with equality holding iff  $\mathbf{update}(\Phi) = \Phi$ .