

ESc201 : Introduction to Electronics

Logic Gates

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Example (recap)

Adding or subtracting numbers with addition operation alone

To get a negative number, 2's complement of positive number is taken

+ 5

+ 2

+ 7

0 1 0 1

+ 0 0 1 0

0 1 1 1

+ 5

- 2

+ 3

0 1 0 1

+ 1 1 1 0

0 0 1 1

2's comp.
of +2

- 5

+ 2

- 3

2's comp.
of +5

→ 1 0 1 1

+ 0 0 1 0

1 1 0 1

↓

2's complement is 0011 = 3

- 5

- 2

- 7

1 0 1 1

+ 1 1 1 0

1 0 0 1

↓

2's comp.
of +5

2's comp.
of +2

2's complement is 0111 = 7

Boolean Algebra (recap)

Algebra on Binary numbers

A variable x can take two values $\{0,1\}$

0

- False
- No
- Low voltage

Basic operations:

$$\text{AND: } y = x_1 \cdot x_2$$

1

- True
- Yes
- High voltage

y is 1 if and only if both x_1 and x_2 are 1, otherwise zero

Truth Table

x_1	x_2	y
0	0	0
0	1	0
1	0	0
1	1	1

Basic operations (recap)

$$\text{OR: } y = x_1 + x_2$$

y is 1 if either x_1 or x_2 is 1. $y = 0$ if and only if both variables are zero

x_1	x_2	y
0	0	0
0	1	1
1	0	1
1	1	1

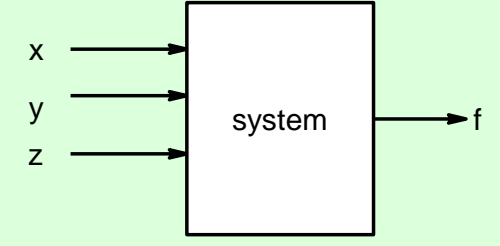
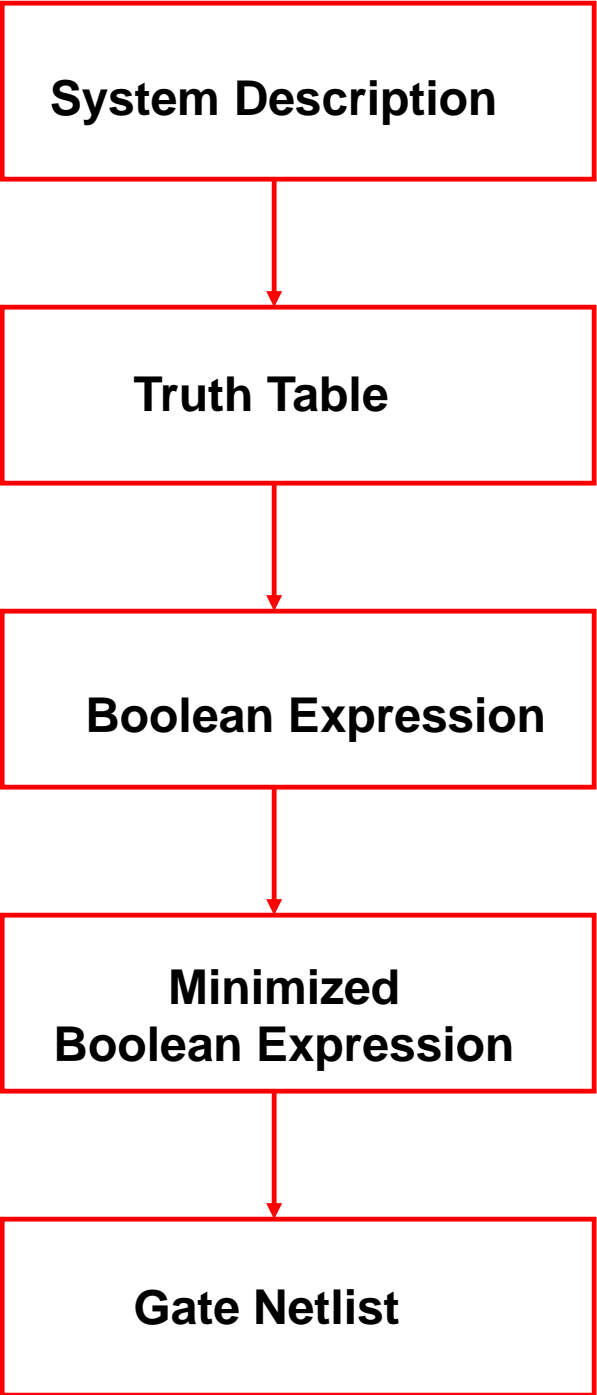
$$\text{NOT: } y = \bar{x}$$

x	y
0	1
1	0



How do we get the chocolate?

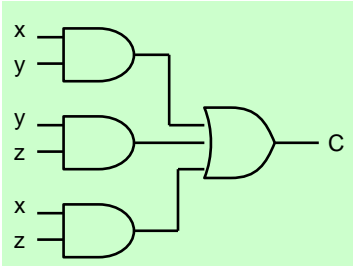
Design Flow



x	y	z	f
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

$$f = \bar{x}.\bar{y}.z + \bar{x}.y.z + x.\bar{y}.z + x.y.z$$

$$\Rightarrow f = \bar{x}.\bar{z} + x.z$$



Proving Theorems

$$\text{P1.a: } x + 0 = x$$

$$\text{P2.a: } x + y = y + x$$

$$\text{P3.a: } x.(y+z) = x.y+x.z$$

$$\text{P4.a: } x + \bar{x} = 1$$

$$\text{P1.b: } x . 1 = x$$

$$\text{P2.b: } x . y = y . x$$

$$\text{P3.b: } x+y.z = (x+y).(x+z)$$

$$\text{P4.b: } x . \bar{x} = 0$$

$$\text{Prove : } x + 1 = 1$$

$$x + 1 = x + (x + \bar{x})$$

$$= (x+x) + \bar{x}$$

$$= x + \bar{x}$$

$$= 1$$

$$x + x . y = x$$

$$= x . 1 + x . y$$

$$= x . (1 + y)$$

$$= x . 1$$

$$= x$$

DeMorgan's Theorem

$$\overline{(x_1 + x_2 + x_3 + \dots)} = \bar{x}_1 . \bar{x}_2 . \bar{x}_3 .$$

$$\overline{(x_1 . x_2 . x_3 \dots)} = (\bar{x}_1 + \bar{x}_2 + \bar{x}_3 + \dots)$$

Simplification of Boolean expressions

$$\overline{(X_1 + X_2 + X_3 + \dots)} = \overline{X_1} \cdot \overline{X_2} \cdot \overline{X_3} \cdot \dots$$

$$\overline{(X_1 \cdot X_2 \cdot X_3 \dots)} = (\overline{X_1} + \overline{X_2} + \overline{X_3} + \dots)$$

$$\overline{\overline{(X_1 \cdot X_2 + X_2 \cdot X_3)}} = ?$$

$$= \overline{\overline{(X_1 \cdot X_2)}} \cdot \overline{\overline{(X_2 \cdot X_3)}}$$

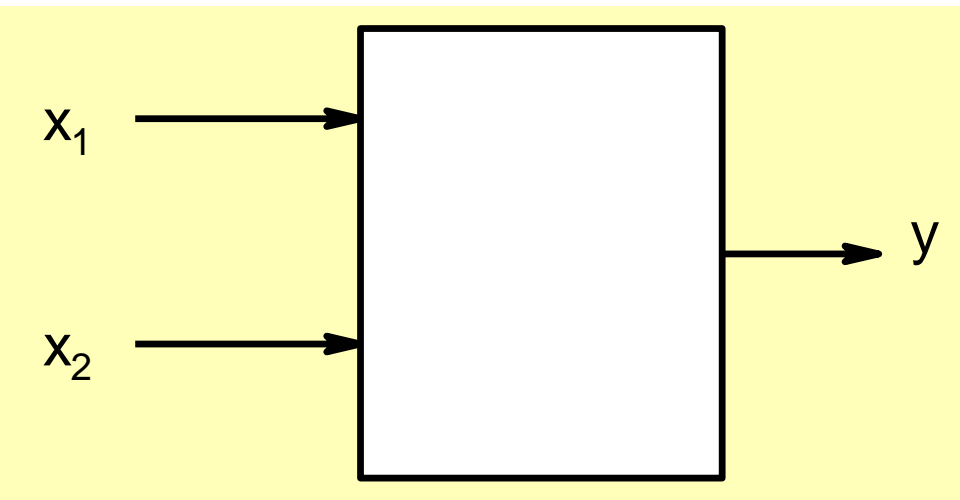
$$= (\overline{\overline{X_1}} + \overline{\overline{X_2}}) \cdot (\overline{\overline{X_2}} + \overline{\overline{X_3}})$$

$$= (X_1 + \overline{X_2}) \cdot (X_2 + \overline{X_3})$$

$$= X_1 \cdot X_2 + \overline{X_2} \cdot X_2 + X_1 \cdot \overline{X_3} + \overline{X_2} \cdot \overline{X_3}$$

$$= X_1 \cdot X_2 + X_1 \cdot \overline{X_3} + \overline{X_2} \cdot \overline{X_3}$$

Function of Boolean variables



x_1	x_2	y
0	0	0
0	1	1
1	0	0
1	1	0

$y = 1$ when x_1 is 0 and x_2 is 1

$$y = \overline{x_1} \cdot x_2$$

Boolean expression

Obtaining Boolean expressions from truth Table

x_1	x_2	y
0	0	1
0	1	0
1	0	0
1	1	0

$$y = \overline{x_1} \cdot \overline{x_2}$$

x_1	x_2	y
0	0	0
0	1	0
1	0	1
1	1	0

$$y = x_1 \cdot \overline{x_2}$$

x_1	x_2	y
0	0	1
0	1	0
1	0	0
1	1	1

$$\overline{x_1} \cdot \overline{x_2}$$

$$x_1 \cdot x_2$$

$$y = \overline{x_1} \cdot \overline{x_2} + x_1 \cdot x_2$$

Obtaining Boolean expressions from truth Table

x_1	x_2	y
0	0	0
0	1	1
1	0	1
1	1	0

$$y = \overline{x_1} \cdot x_2 + x_1 \cdot \overline{x_2}$$

Instead of writing expressions as sum of terms that make y equal to 1, we can also write expressions using terms that make y equal to 0

x_1	x_2	y
0	0	1
0	1	1
1	0	1
1	1	0

$$y = \overline{x_1} \cdot \overline{x_2} + \overline{x_1} \cdot x_2 + x_1 \cdot \overline{x_2}$$

$$y = \overline{x_1} + \overline{x_2}$$

x_1	x_2	y
0	0	1
0	1	0
1	0	1
1	1	1

$$y = x_1 + \overline{x_2}$$

x_1	x_2	y
0	0	0
0	1	1
1	0	1
1	1	1

$$y = x_1 + x_2$$

x_1	x_2	y
0	0	0
0	1	1
1	0	1
1	1	0

$$x_1 + x_2$$

$$y = (x_1 + x_2) \cdot (\overline{x_1} + \overline{x_2})$$

$$\overline{x_1} + \overline{x_2}$$

Obtaining Boolean expressions from truth Table

x_1	x_2	x_3	y
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

$$y = \overline{x_1} \cdot \overline{x_2} \cdot x_3 + \overline{x_1} \cdot x_2 \cdot x_3 + x_1 \cdot \overline{x_2} \cdot x_3 + x_1 \cdot x_2 \cdot x_3$$

Sum of Products (SOP) form

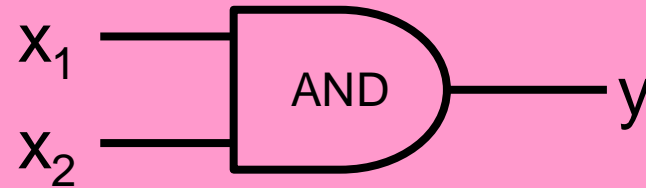
$$y = (x_1 + x_2 + x_3) \cdot (x_1 + \overline{x_2} + x_3) \cdot (\overline{x_1} + x_2 + x_3) \cdot (\overline{x_1} + \overline{x_2} + x_3)$$

Product of Sum (POS) form

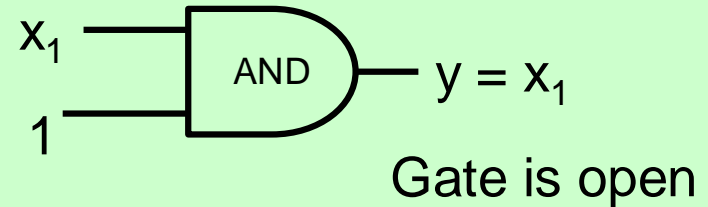
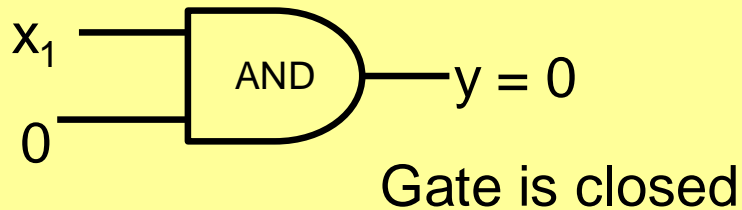
Implementing Boolean expressions

Elementary Gates

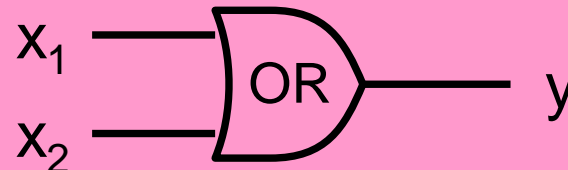
$$\text{AND: } y = x_1 \cdot x_2$$



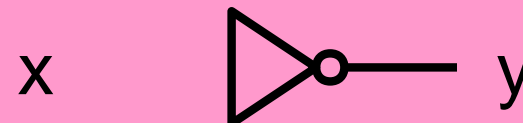
Why call it a gate?



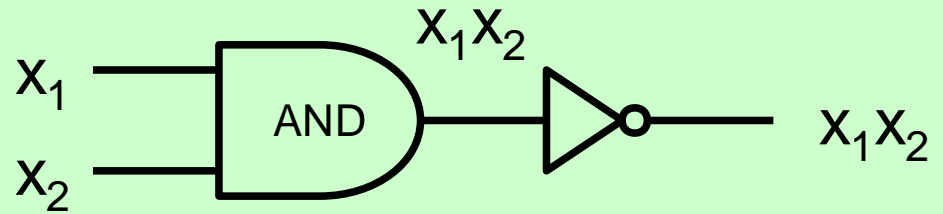
$$\text{OR: } y = x_1 + x_2$$



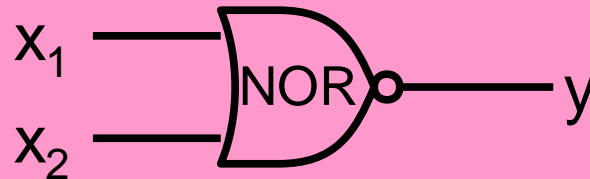
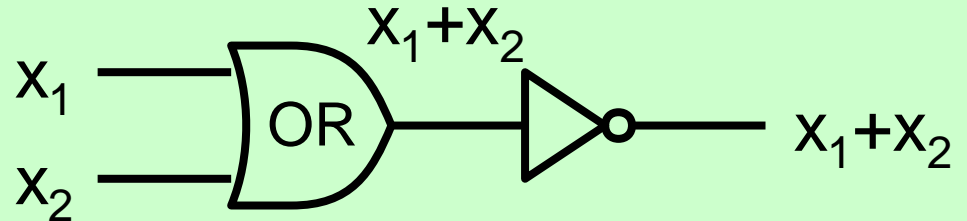
$$\text{NOT: } y = \bar{x}$$



NAND: $y = \overline{x_1 \cdot x_2}$



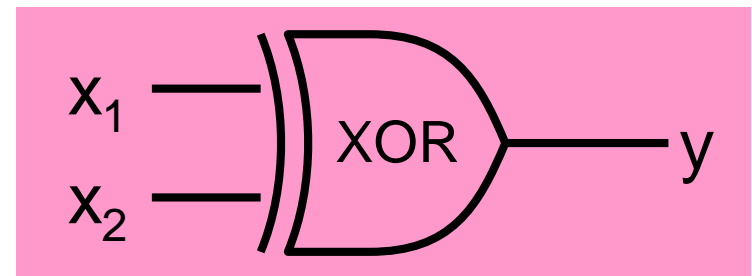
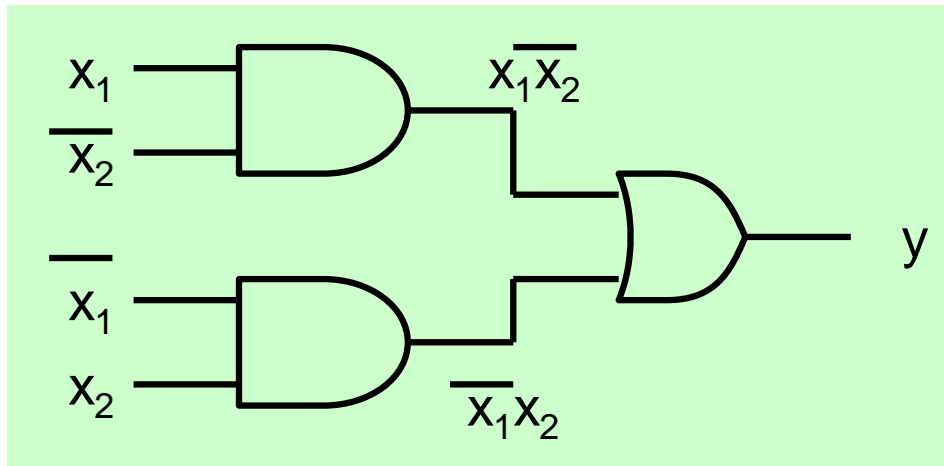
NOR: $y = \overline{x_1 + x_2}$



XOR: $y = x_1 \oplus x_2 = x_1 \cdot \overline{x_2} + \overline{x_1} \cdot x_2$

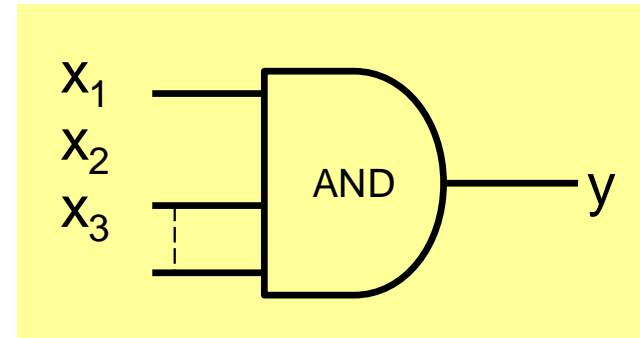
x_1	x_2	y
0	0	0
0	1	1
1	0	1
1	1	0

y is 1 if only one variable is 1 and the other is zero

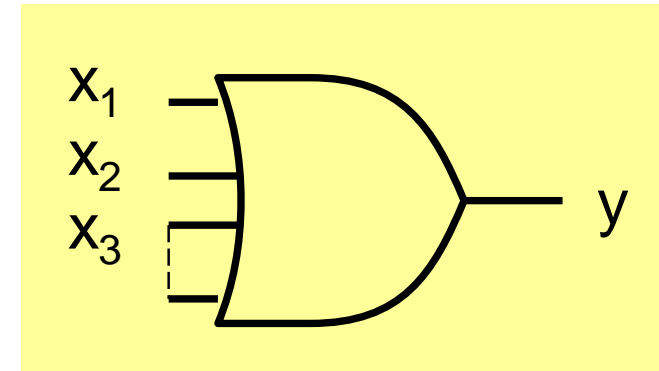


Gates with more than 2 inputs

AND: $y = x_1 \cdot x_2 \cdot x_3 \dots$



OR: $y = x_1 + x_2 + x_3 + \dots$



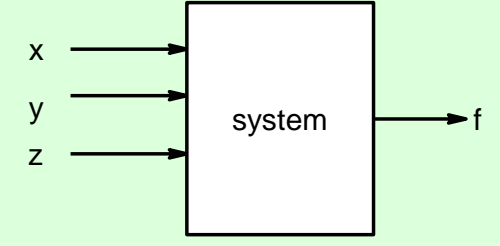
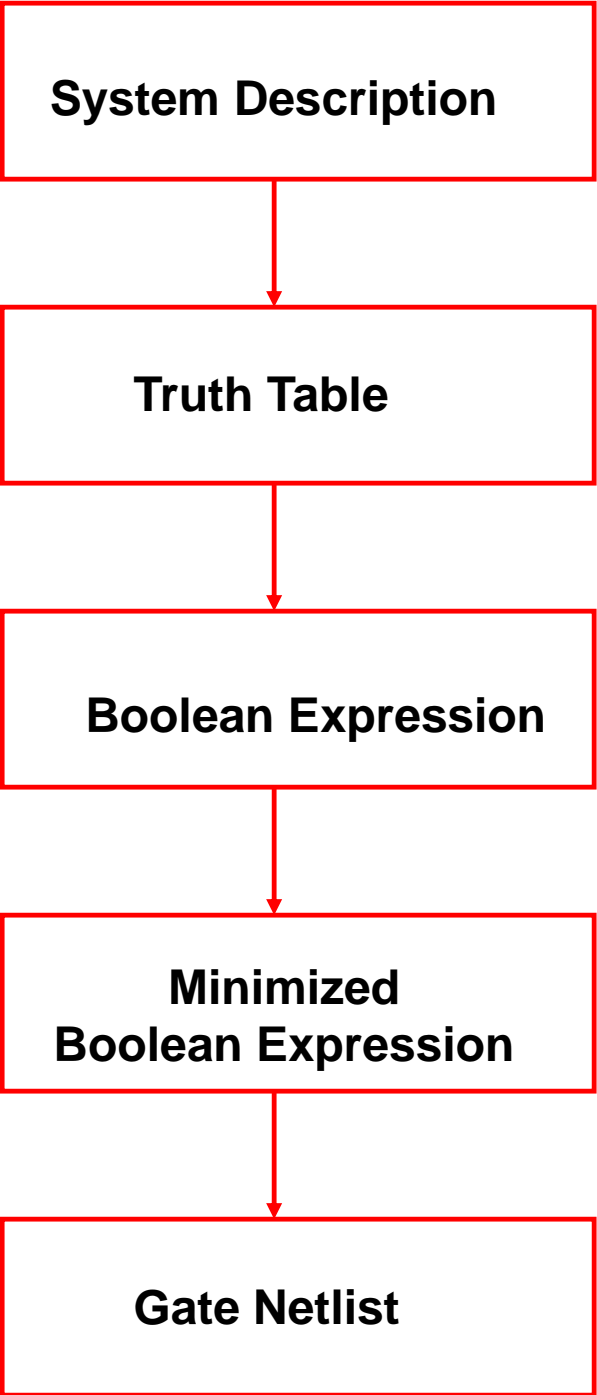
XOR: $y = x_1 \oplus x_2 \oplus x_3 = \overline{\overline{x_1} \cdot \overline{x_2} \cdot x_3} + \overline{x_1} \cdot \overline{x_2} \cdot \overline{x_3} + \overline{x_1} \cdot x_2 \cdot \overline{x_3} + x_1 \cdot x_2 \cdot x_3$

$y = 1$ only if odd number of inputs is 1



How do we get the chocolate?

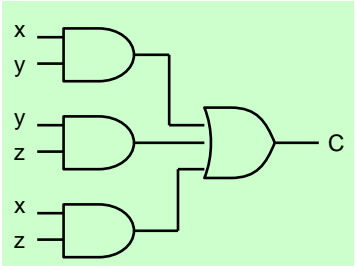
Design Flow



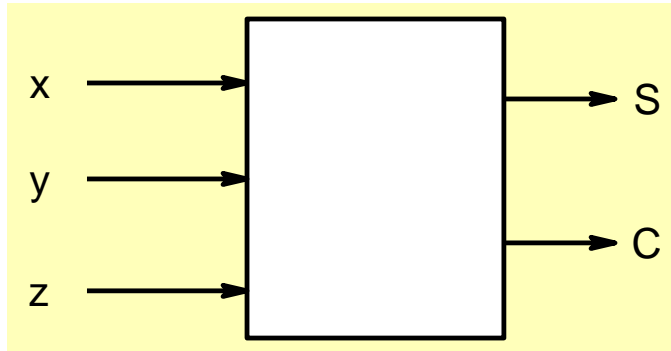
x	y	z	f
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

$$f = \bar{x}.\bar{y}.z + \bar{x}.y.z + x.\bar{y}.z + x.y.z$$

$$\Rightarrow f = \bar{x}.\bar{z} + x.z$$



Implementing Boolean expressions using gates



$$S = \bar{x}.\bar{y}.z + \bar{x}.y.\bar{z} + x.\bar{y}.\bar{z} + x.y.z$$

$$C = x.y + x.z + y.z$$

