

09.09.2024

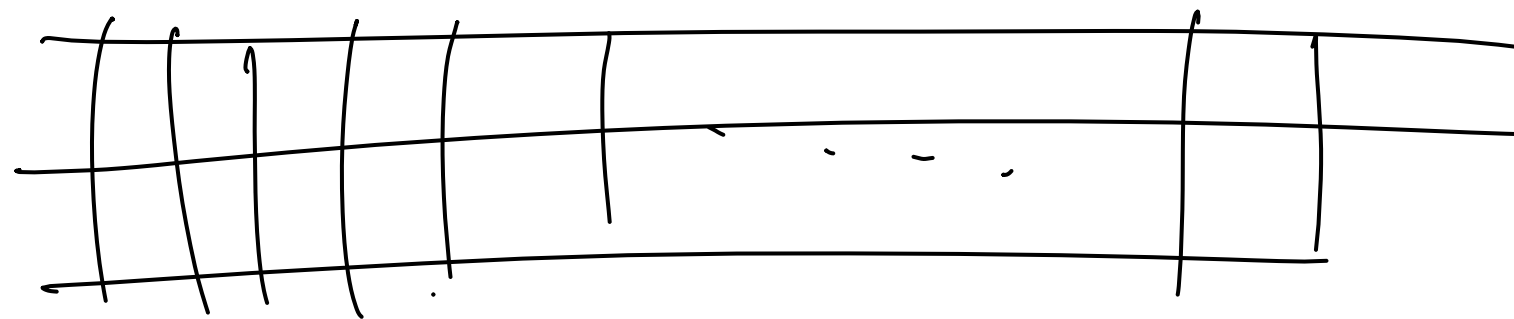
ED(i, j) :

Input:

String A: $[1 \dots n]$

B: $[1 \dots m]$.

Edit distance.



$A[1 \dots i]$

$1 \leq i \leq n$

$B[1 \dots j]$

$1 \leq j \leq m$

$i = 0 / j = 0 \Rightarrow$ empty string.

$$\begin{aligned}
 ED(i, j) &= \frac{ED(i, j-1) + 1}{ED(i-1, j) + 1} \\
 &= \frac{ED(i-1, j-1) + \underbrace{O(i, j)}}{ED(i-1, j-1) + O(i, j)}
 \end{aligned}$$

$$\frac{-}{\beta(j)} \checkmark$$

$$\frac{A(i)}{-} \checkmark$$

$$\frac{A(i)}{\beta(j)} \checkmark$$

0 if agree
1 o/w.

A

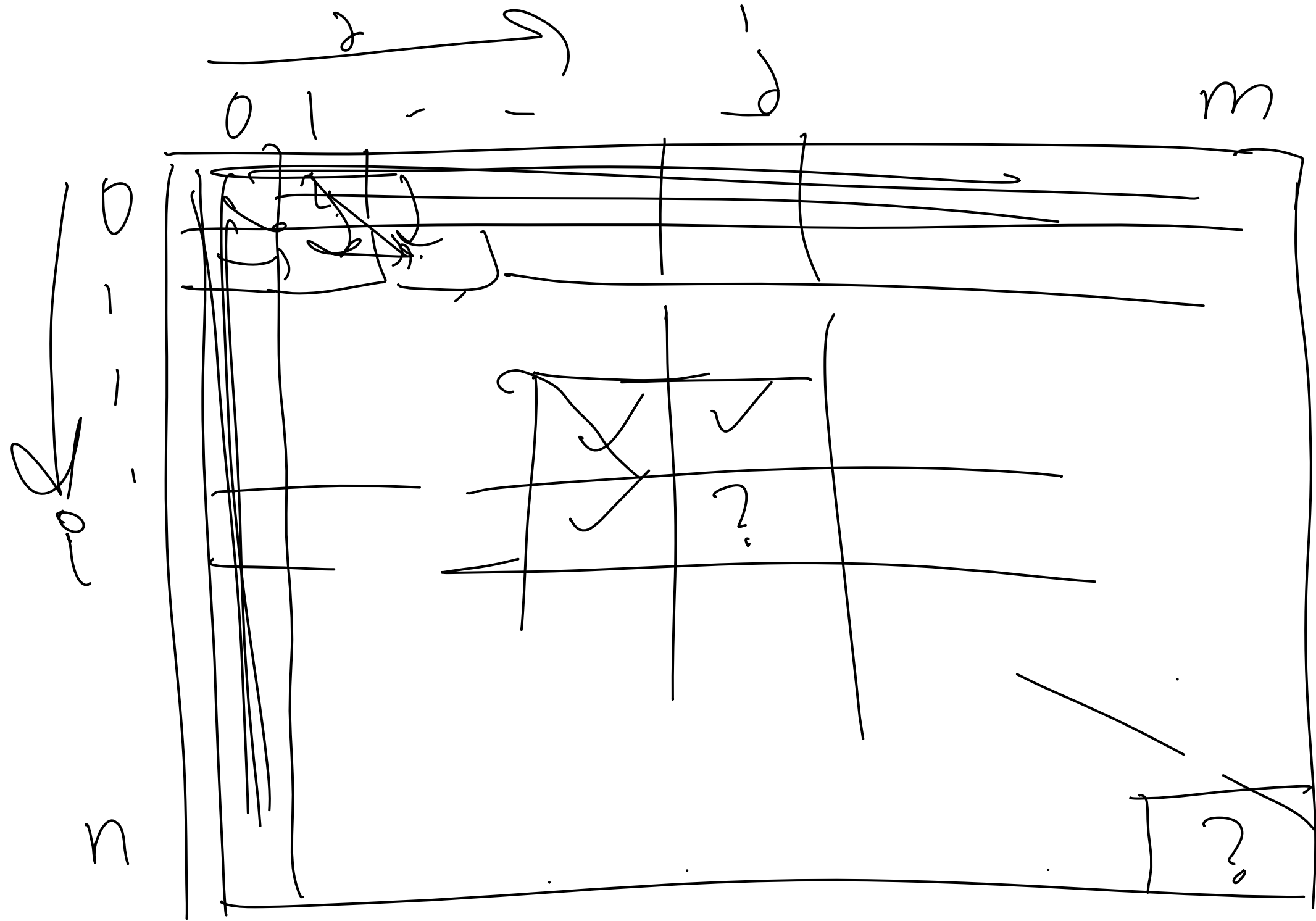
B

base Case: $i = 0, j = 0, 1, \dots, m.$

$$ED(0, j) = j$$

$$ED(i, 0) = i$$

6



$$ED. \left[\begin{matrix} n \\ \cancel{m} \end{matrix}, \begin{matrix} m \\ \cancel{n} \end{matrix} \right]$$

	0	F	A	V	O	R	1	T	E	
→	0	1	2	3	4	5	6	7	8	
→	F	1	2	2	3	4	5	6	7	
L	2	1	2	3	4	5	6	7		
A	3	2	1	2	3	4	5	6	7	
V	4	3	2	1	2	3	4	5	6	
O	5	4	3	2	1	2	3	4	5	
R	6	5	4	3	2	1	2	3	4	
Y	7	6	5	4	3	2	3	3	4	
	F	F	A	V	O	R	T	T	E	
	F	L	A	V	O	R	Y	L	E	

Procedure ^{EditDist.} ~~ED~~ (A, B, n, m)

~~ED~~ ~~← empty~~ // $n = \text{length}(A)$, $m = \text{length}(B)$.

ED \leftarrow empty table of size $(n+1) \times (m+1)$.

for $j = 0$ to $(m+1)$

ED(0, j) $\leftarrow j$

for $i = 0$ to $(n+1)$

ED(i, 0) $\leftarrow i$

for $i = 1$ to $(n+1)$

for $j = 1$ to $(m+1)$

$$\underline{\underline{ED(i, j)}} = \min \left\{ \begin{array}{l} \underline{ED(i-1, j)} + 1 \\ \underline{ED(i, j-1)} + 1 \\ \underline{ED(i-1, j-1)} + 1 \quad \text{if } A[i] \neq B[j] \\ \underline{ED(i-1, j-1)} \quad \text{if } A[i] = B[j] \end{array} \right.$$

Exercise 6 Modify the pseudocode so that it also returns an alignment?

align - A, align - B $\leftarrow (n+m)$ empty array

Time Complexity

$\Theta(nm)$

~~FAVORITE~~

FAVORITE

FLAVORY

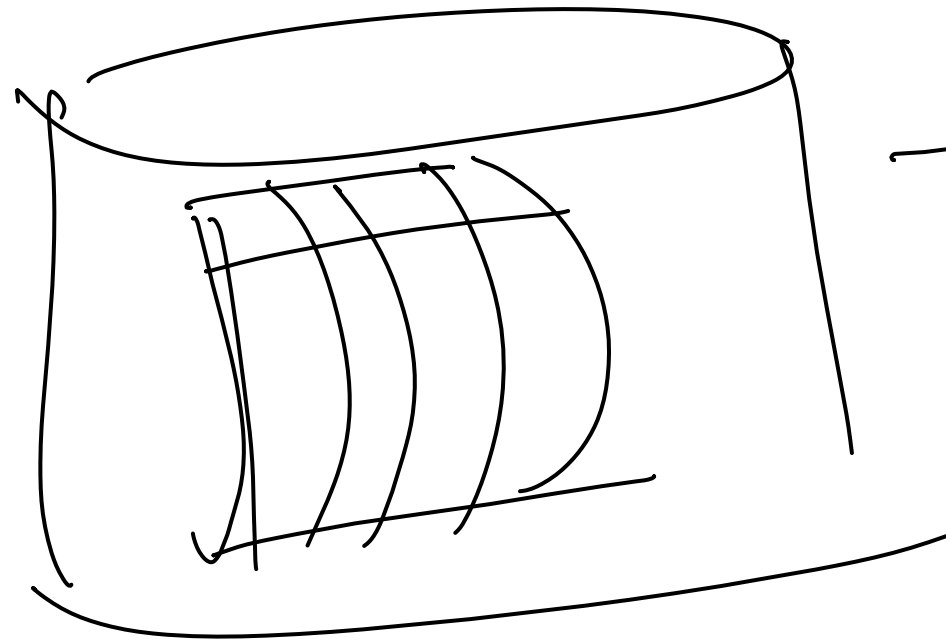
— DP is always associated with DP Table.

(i, j)



— ~~Defining~~ Defining subproblems is extremely important.

Binary Search Tree



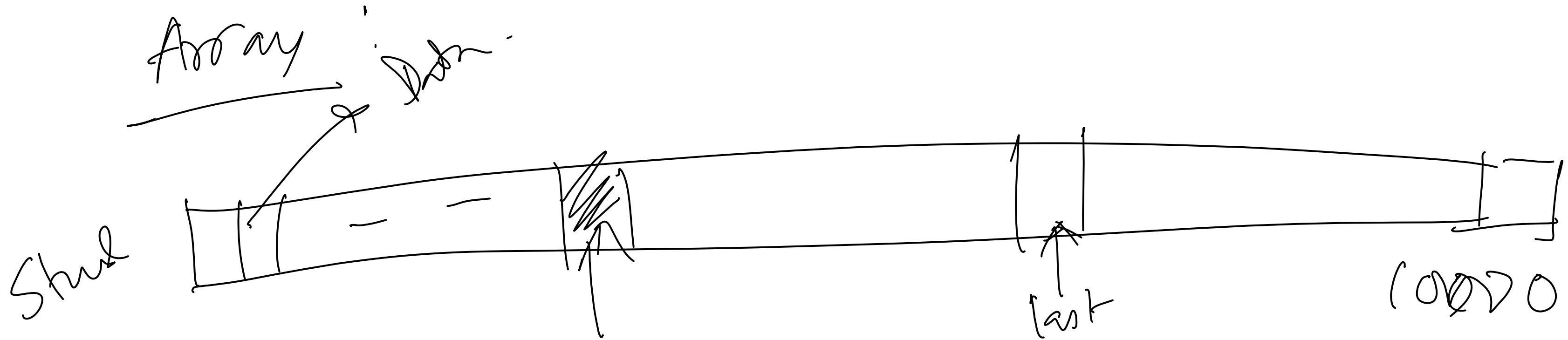
→ { roll no, Contact, CPI, transcript.

✓
Struct
Struct

→ insert.

→ delete.

→ Search.
(roll no, data)



insert:

Index [last] \leftarrow new.

last ++.

~~$O(n)$~~

~~delete~~ ? search.

delete ?

✓ $O(n)$.

Sorted Array

Search: $O(\log n)$

Insert

$O(n)$

Delete



19, 431, 572, 72, 1042, 2, 3

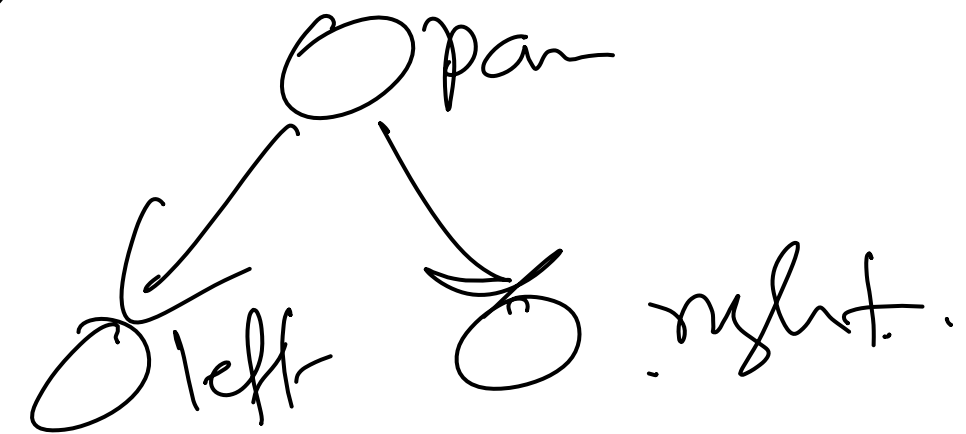
Binary

Tree :: [Directed]

struct Node

Node *left
Node *right
Node *parent
int key

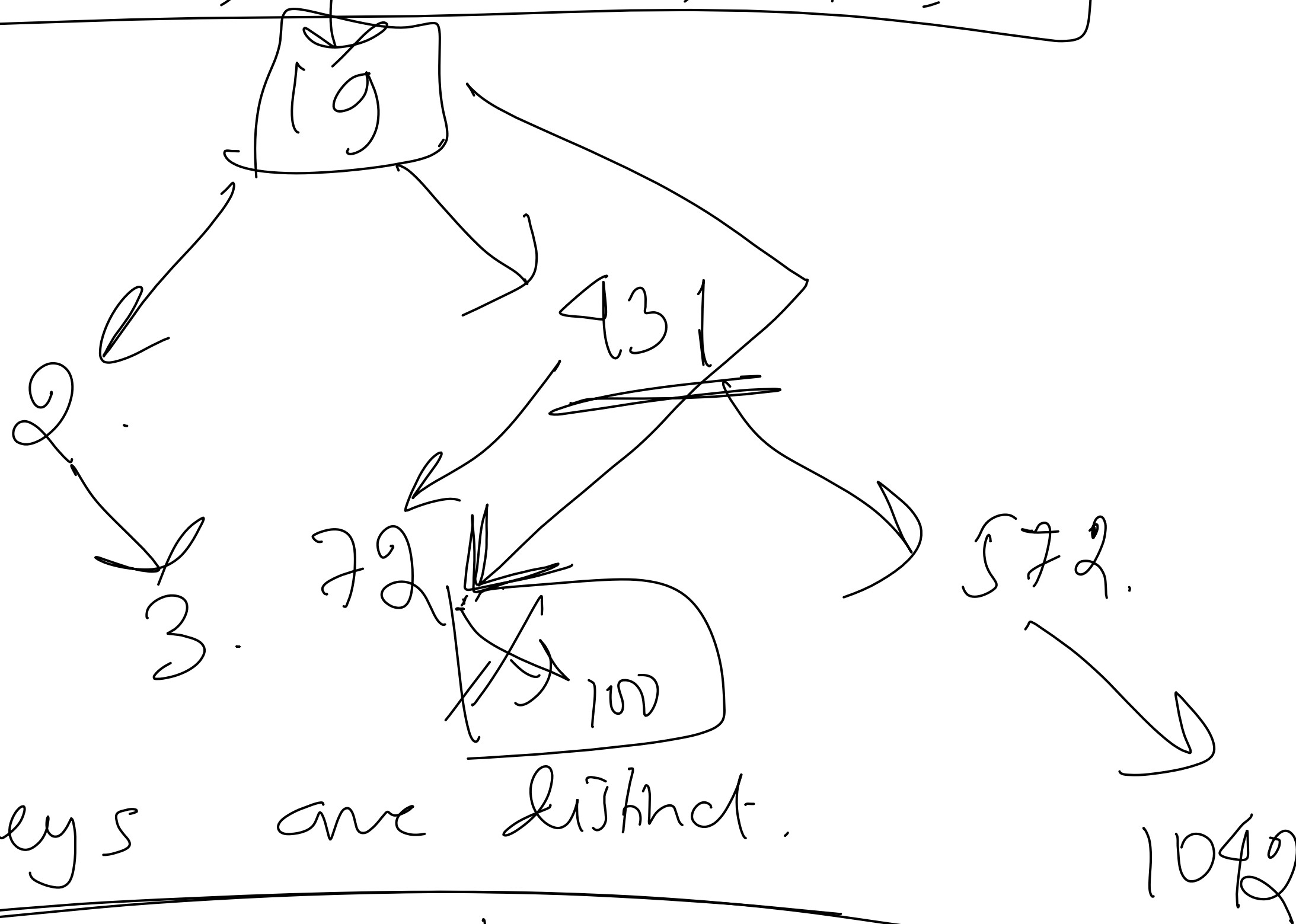
};



Node . left
Node . right
Node . parent
Node . key

19, 431, 572, 72, 1042, 2, 3

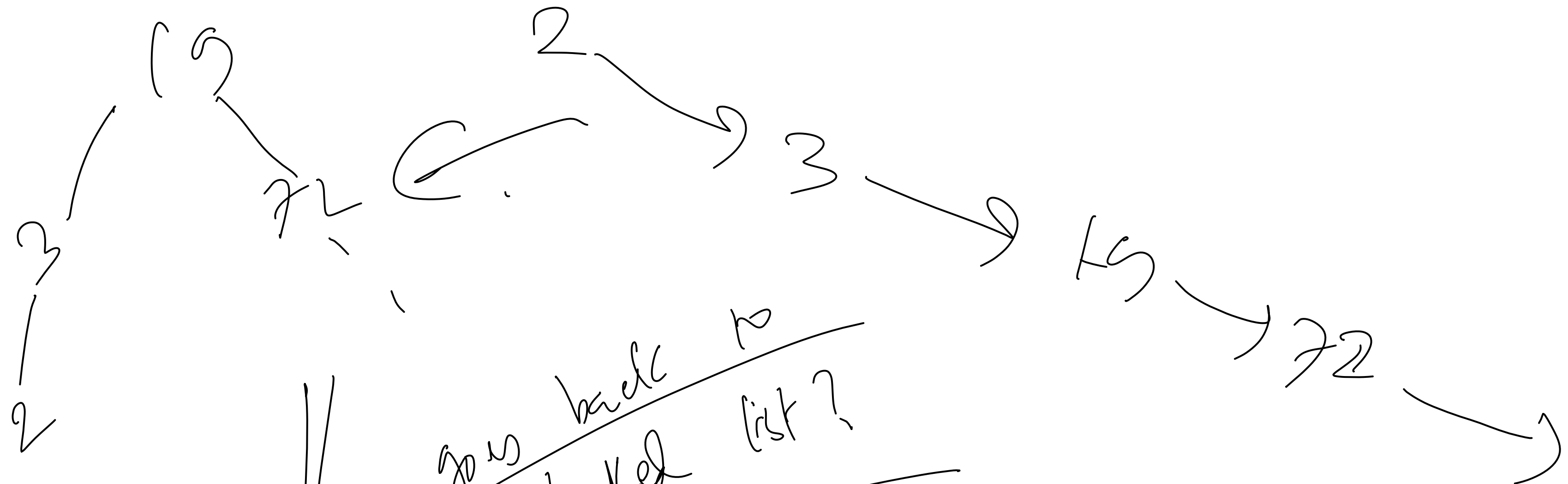
572 2 yls
182 X



Assumptions: - Keys are distinct.

- total ordering between keys

2, 3, 19, 72, 431, 572, 1042.



~~go as back to linked list?~~

Balancing

$O(n)$

- Insert

- ~~Del~~ Search

- Delete

