

$$T(n) \leq T\left(\frac{n}{5}\right) + T\left(\frac{7 \cdot 2n}{10}\right) + cn > \left(\frac{1}{5}\right)^j \left(\frac{7 \cdot 2}{10}\right)^j$$

Claim: $T(n) \leq dn$ for some constant d .
prof: By induction.
choose d large enough $\geq \frac{c}{0.08}$

I.H. Claim is true for $1, \dots, n-1$

$$\begin{aligned} T(n) &\leq T\left(\frac{n}{5}\right) + T\left(\frac{7 \cdot 2n}{10}\right) + cn \\ &\leq d \cdot \frac{n}{5} + d \cdot \frac{7 \cdot 2n}{10} + cn \leq dn \end{aligned}$$

$$0.92dn + cn \leq dn.$$

$$\begin{aligned} &\equiv cn \leq 0.08dn \\ &\equiv d \geq \frac{c}{0.08} \end{aligned}$$

Therefore, $T(n) \leq dn$.

[Induction step holds]

□

$$T(n) = \Theta(n), \quad T(n) = \mathcal{O}(n)$$

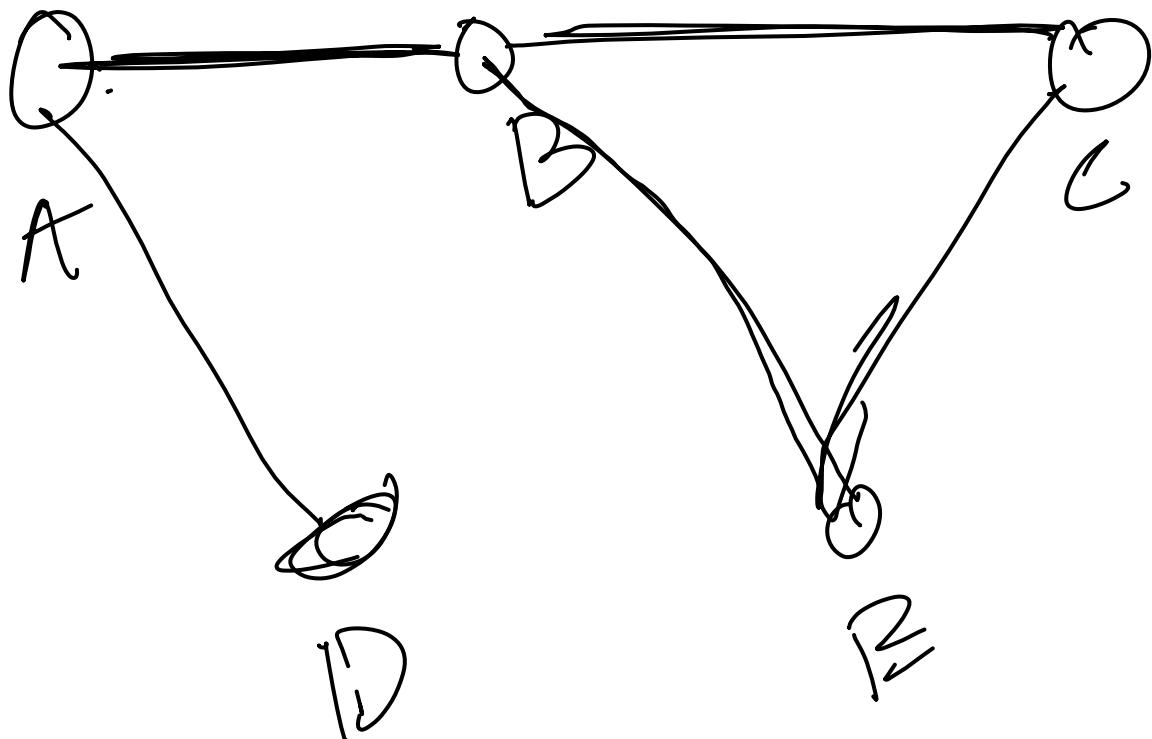
$$T(n) = \Theta(n)$$

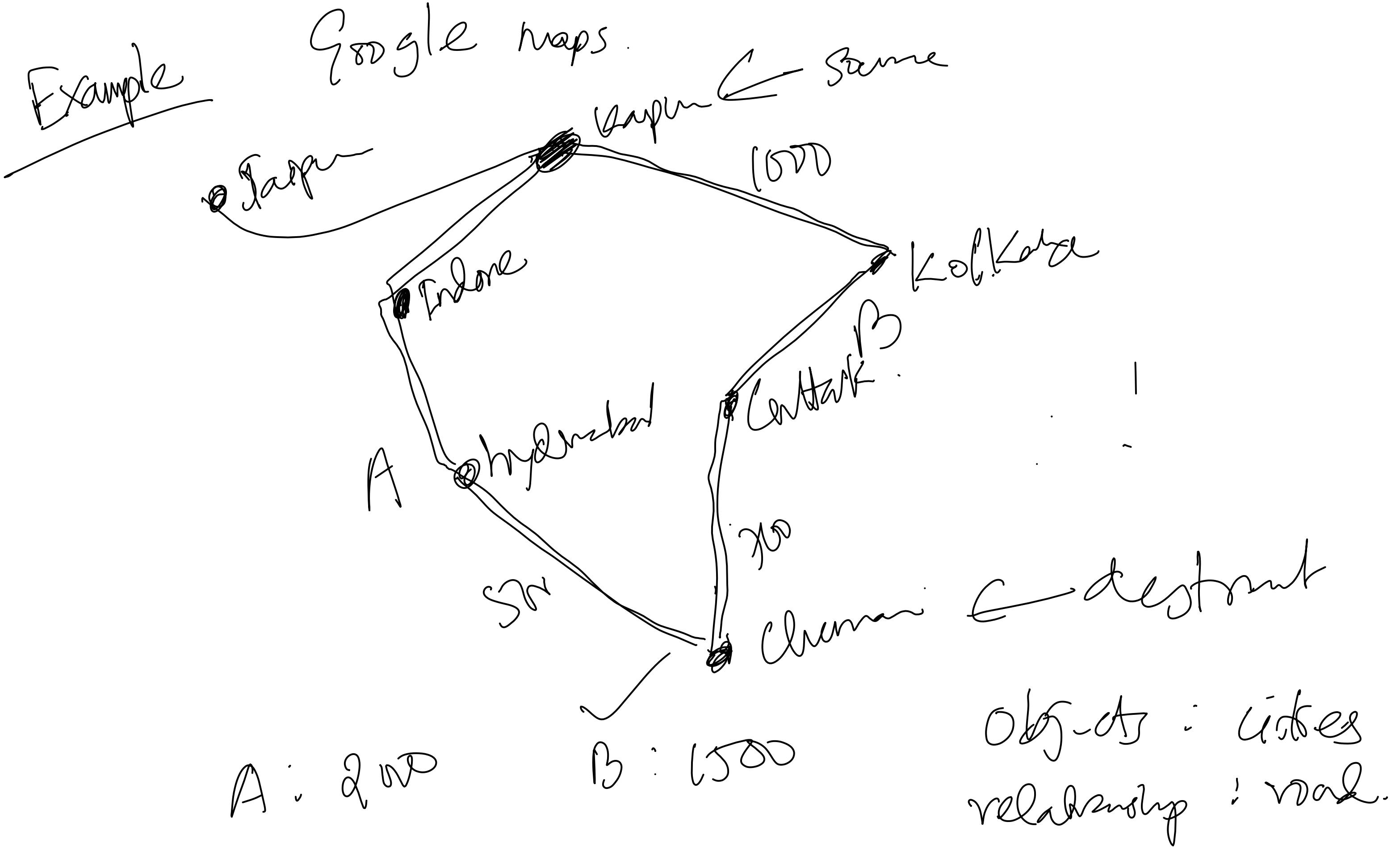
Module 9 : Graph Algorithm

(object, relationship)

Examples: Facebook graph.

objects: users
relationship: friendship.

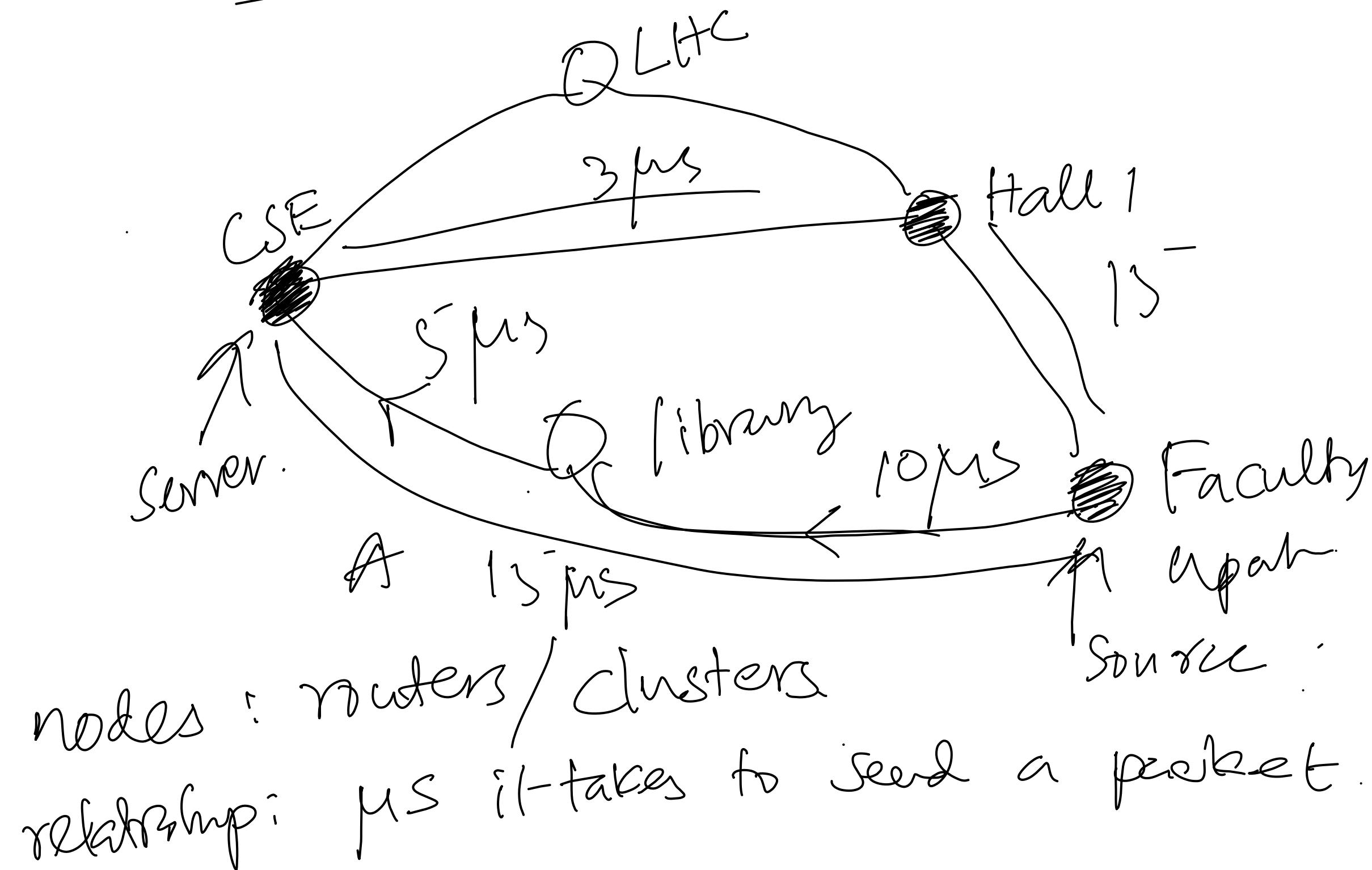




obj-obj : cities
relationship : road.

Example

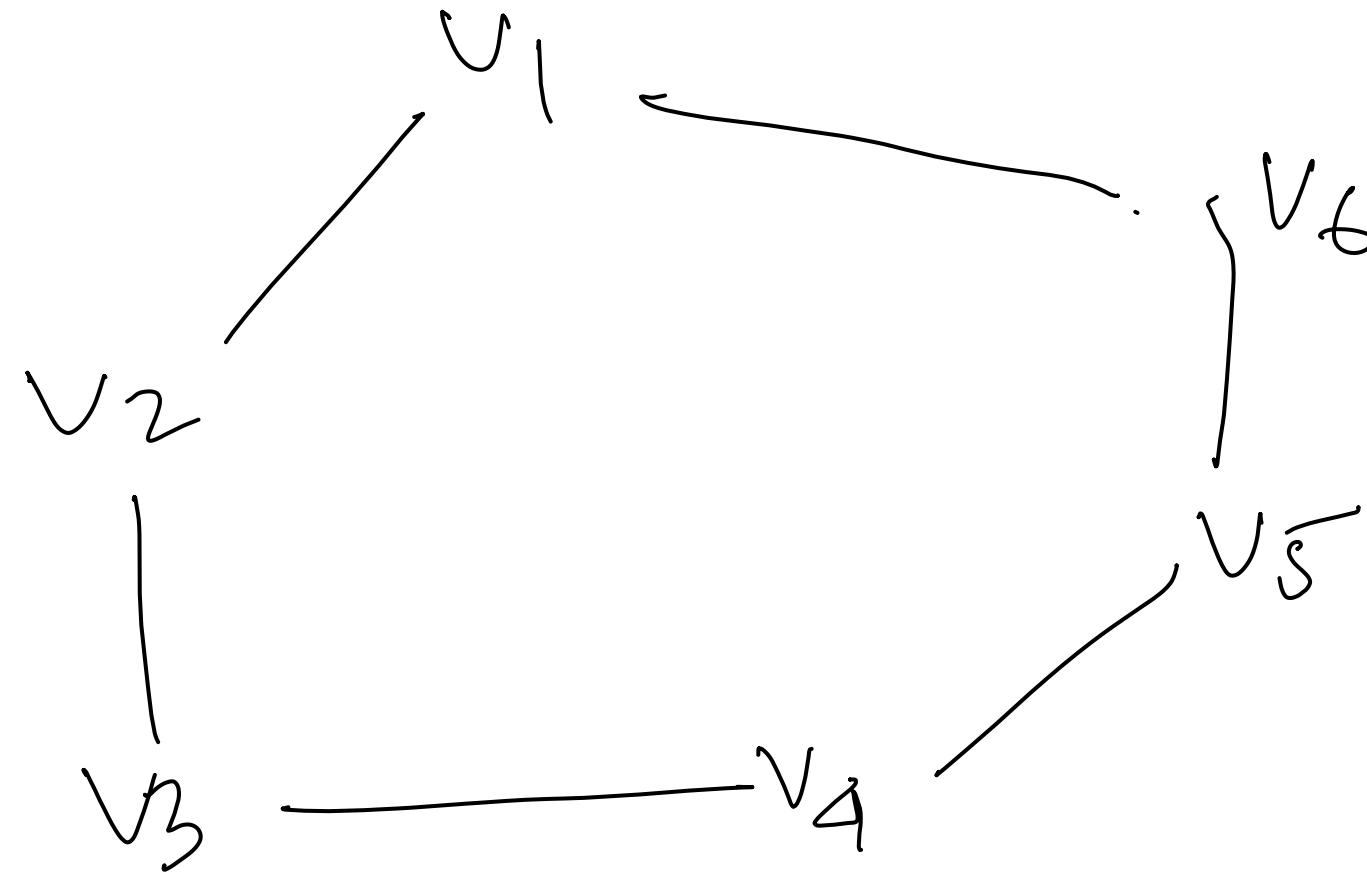
Routing



$G = (V, E)$
 graph
 networks -
 nodes
 vertices
 edges / links

Example 1 -

G

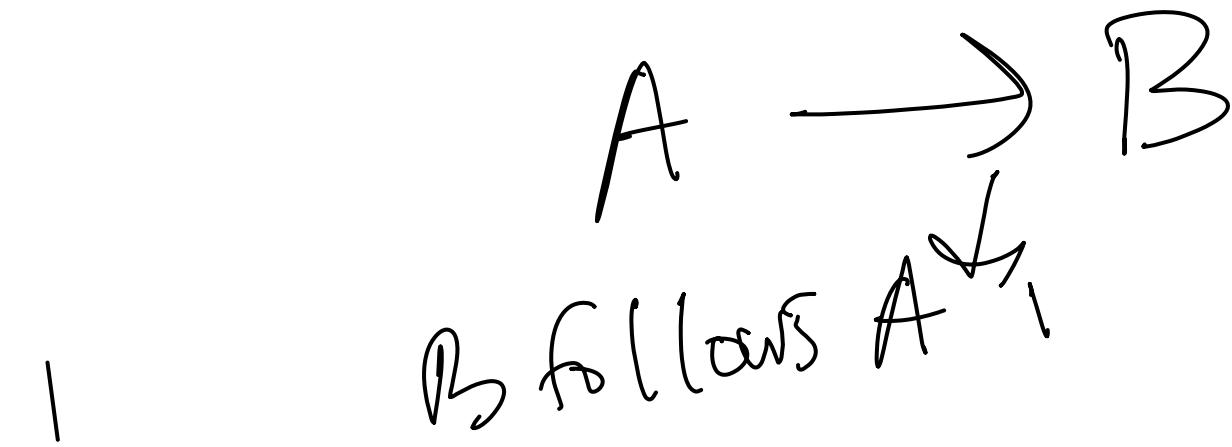


$$V = \{v_1, \dots, v_6\} \equiv (v_2, v_1)$$

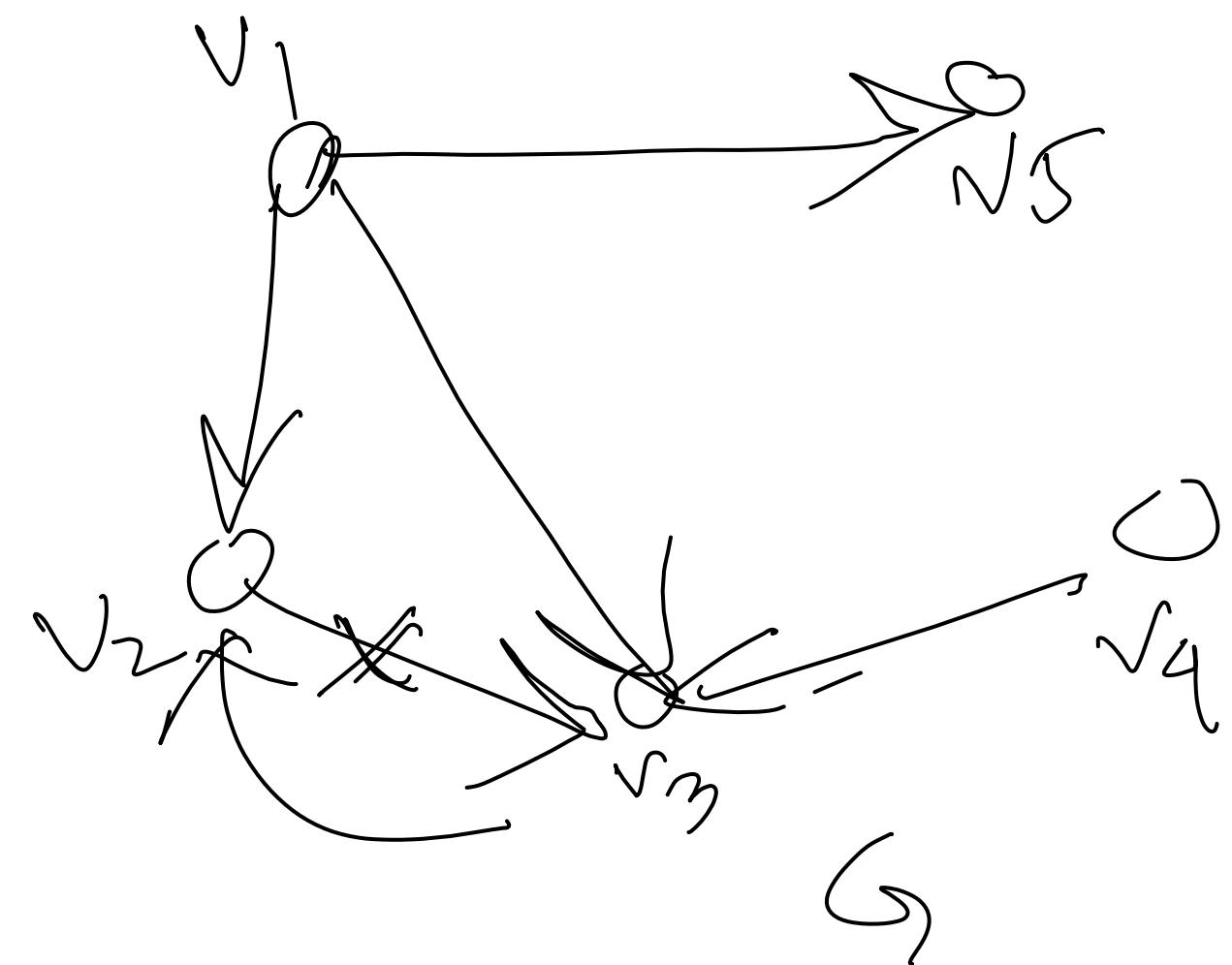
$$E \subseteq V \times V = \{(v_1, v_2), (v_2, v_3), (v_3, v_4), (v_4, v_5), (v_5, v_6), (v_6, v_1)\}$$

Graphs \equiv undirected graphs.

Directed graphs -



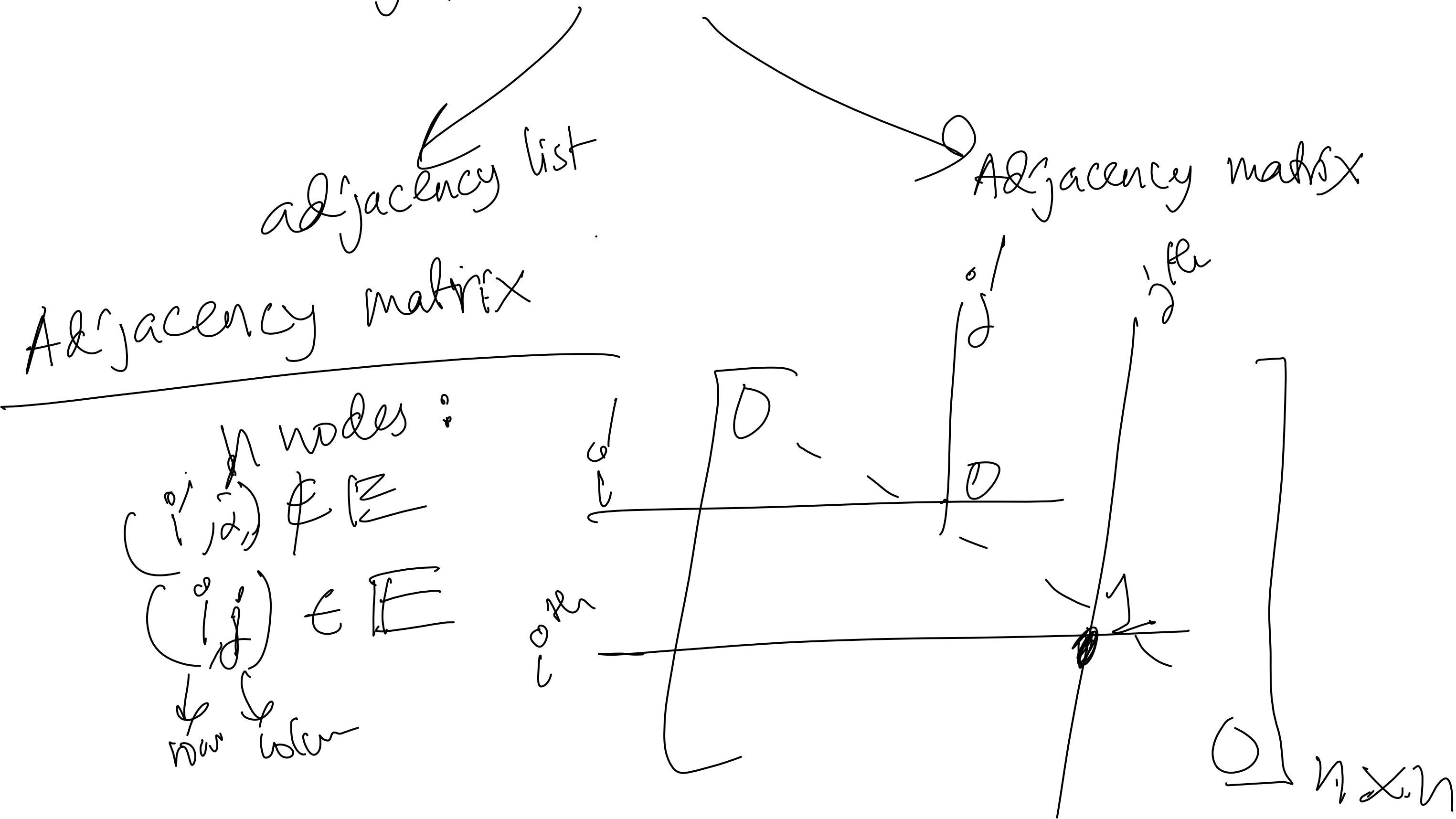
$$\underline{(A,B)} = (B,A)$$



$$V = \{v_1, v_2, \dots, v_5\}$$

$$E \subseteq V \times V$$
$$= \{(v_1, v_2), (v_2, v_3), (v_3, v_2), (v_1, v_3), (v_4, v_3), (v_4, v_3), (v_1, v_5)\}$$

How are graphs stored inside a computer?



	1	2	3	4	5	6
1	0	1	0	0	0	1
2	1	0	1	0	0	0
3	0	1	0	1	0	0
4	0	0	1	0	1	0
5	0	0	0	1	0	1
6	0	0	0	0	1	0



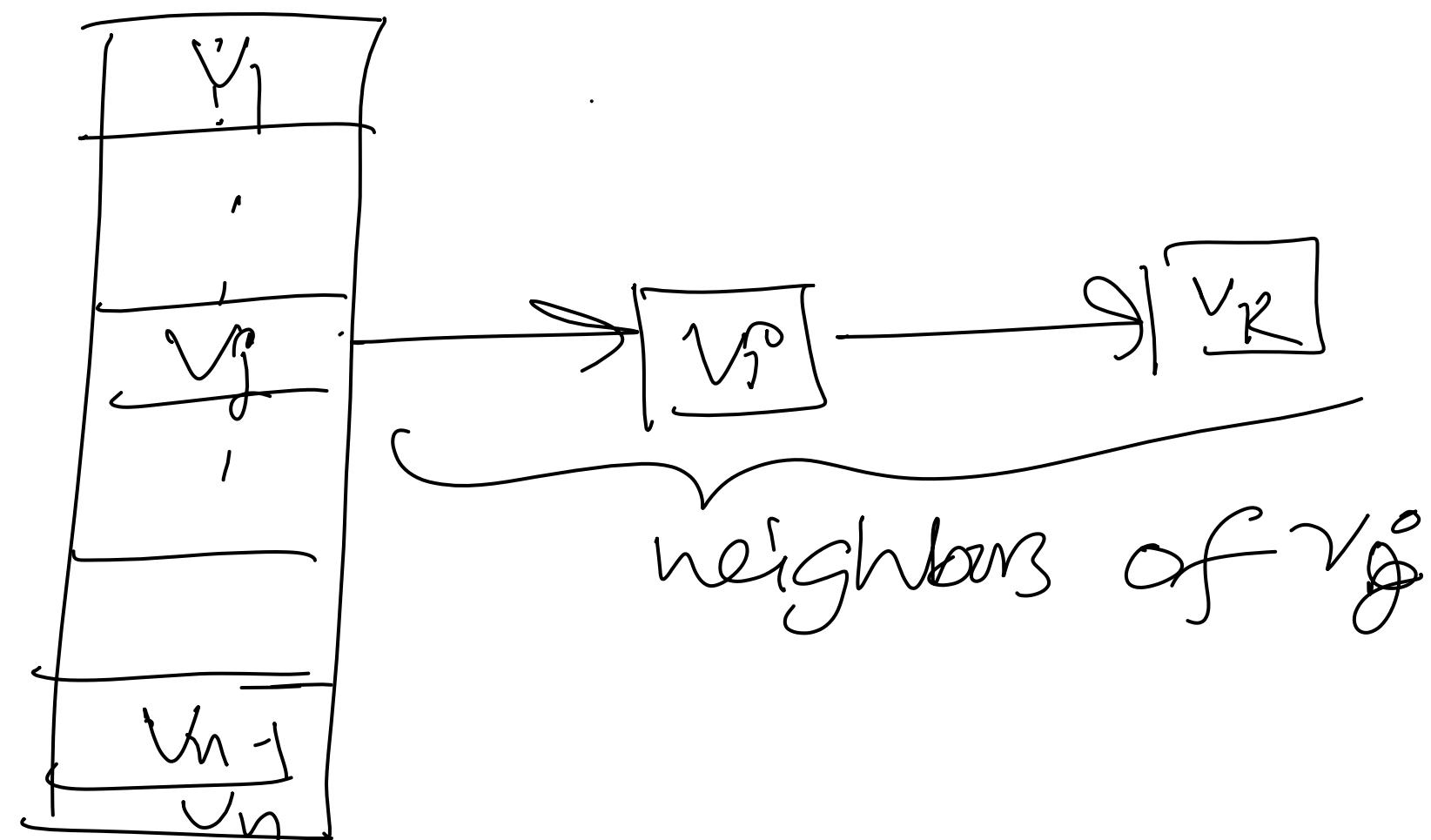
- Symmetric
- stored like a directed graph
where each edge is bidirected.

	1	2	3	4	5
1	0	1	1	0	1
2	0	0	0	0	0
3	0	1	0	0	0
4	0	0	1	0	0
5	0	0	0	0	0

no longer Symmetric.

Remark: We'll focus on designing algorithms for directed graphs. More often, they will extend to undirected graphs.

Adjacency List



$$T(m) \leq T(m/5) + T\left(\frac{7m}{2} + 3\right) + cn$$



$$g = m/5 \quad ? \quad \frac{7m}{2} = \frac{7m}{10}$$
$$m \leq n$$

$$\boxed{T(n) \leq T(n/5) + T\left(\frac{7n}{10} + 3\right) + cn.}$$

$$\leq \frac{7 \cdot 2^n}{10}$$

for n large enough.