

```
fun append (xs,ys) =  
  if xs=[]  
  then ys  
  else (hd xs)::append(tl xs,ys)  
  
fun map (f,xs) =  
  case xs of  
    [] => []  
  | x::xs' => (f x)::(map(f,xs'))  
  
val a = map (increment, [4,8,12,16])  
val b = map (hd, [[8,6],[7,5],[3,0,9]])
```

# Programming Languages

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What is Type Inference?

# Type-checking

- (Static) **type-checking** can reject a program before it runs to prevent the possibility of some errors
  - A feature of **statically typed languages**
- **Dynamically typed languages** do little (none?) such checking
  - So might try to treat a number as a function at run-time
- Will study relative advantages after some Racket
  - Racket, Ruby (and Python, Javascript, ...) dynamically typed
- ML (and Java, C#, Scala, C, C++) is statically typed
  - Every binding has one type, determined “at compile-time”

# *Implicitly typed*

- ML is statically typed
- ML is **implicitly typed**: rarely need to write down types

```
fun f x = (* infer val f : int -> int *)
  if x > 3
  then 42
  else x * 2

fun g x = (* report type error *)
  if x > 3
  then true
  else x * 2
```

- Statically typed: Much more like Java than Javascript!

# Type inference

- **Type inference** problem: Give every binding/expression a type such that type-checking succeeds
  - Fail if and only if no solution exists
- In principle, could be a pass before the type-checker
  - But often implemented together
- Type inference can be easy, difficult, or *impossible*
  - Easy: Accept all programs
  - Easy: Reject all programs
  - Subtle, elegant, and *not magic*: ML