

```
fun append (xs,ys) =  
  if xs=[]  
  then ys  
  else (hd xs)::append(tl xs,ys)  
  
fun map (f,xs) =  
  case xs of  
    [] => []  
  | x::xs' => (f x)::(map(f,xs'))  
  
val a = map (increment, [4,8,12,16])  
val b = map (hd, [[8,6],[7,5],[3,0,9]])
```

Programming Languages

Dan Grossman

Definition of Lexical Scope

Very important concept

- We know function bodies can use any bindings in scope
- But now that functions can be passed around: In scope where?

*Where the function was defined
(not where it was called)*

- This semantics is called *lexical scope*
- There are lots of good reasons for this semantics (why)
 - Discussed after explaining what the semantics is (what)
 - Later in course: implementing it (how)
- Must “get this” for homework, exams, and competent programming

Example

Demonstrates lexical scope even without higher-order functions:

```
(* 1 *) val x = 1
(* 2 *) fun f y = x + y
(* 3 *) val x = 2
(* 4 *) val y = 3
(* 5 *) val z = f (x + y)
```

- Line 2 defines a function that, when called, evaluates body $x+y$ in environment where x maps to 1 and y maps to the argument
- Call on line 5:
 - Looks up f to get the function defined on line 2
 - Evaluates $x+y$ in **current environment**, producing 5
 - Calls the function with 5, which evaluates the body in the **old environment**, producing 6

Closures

How can functions be evaluated in old environments that aren't around anymore?

- The language implementation keeps them around as necessary

Can define the semantics of functions as follows:

- A function value has **two parts**
 - The **code** (obviously)
 - The **environment** that was current when the function was defined
- This is a “pair” but unlike ML pairs, you cannot access the pieces
- All you can do is call this “pair”
- This pair is called a ***function closure***
- A call evaluates the code part in the environment part (extended with the function argument)

Example

```
(* 1 *) val x = 1
(* 2 *) fun f y = x + y
(* 3 *) val x = 2
(* 4 *) val y = 3
(* 5 *) val z = f (x + y)
```

- Line 2 creates a closure and binds **f** to it:
 - Code: “take **y** and have body **x+y**”
 - Environment: “**x** maps to 1”
 - (Plus whatever else is in scope, including **f** for recursion)
- Line 5 calls the closure defined in line 2 with 5
 - So body evaluated in environment “**x** maps to 1” extended with “**y** maps to 5”

Coming up:

Now you know the rule: *lexical scope*.

Next steps (rest of section):

- (Silly) examples to demonstrate how the rule works with higher-order functions
- Why the other natural rule, *dynamic scope*, is a bad idea
- Powerful *idioms* with higher-order functions that use this rule
 - Passing functions to iterators like **filter**
 - Several more idioms