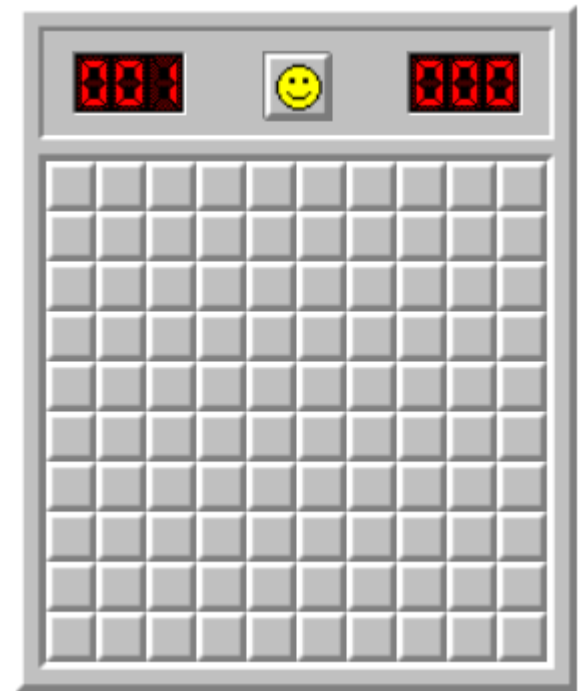


**COSC76/276 Artificial Intelligence**  
**Fall 2022**  
**First Order Logic**

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# Problems with propositional logic

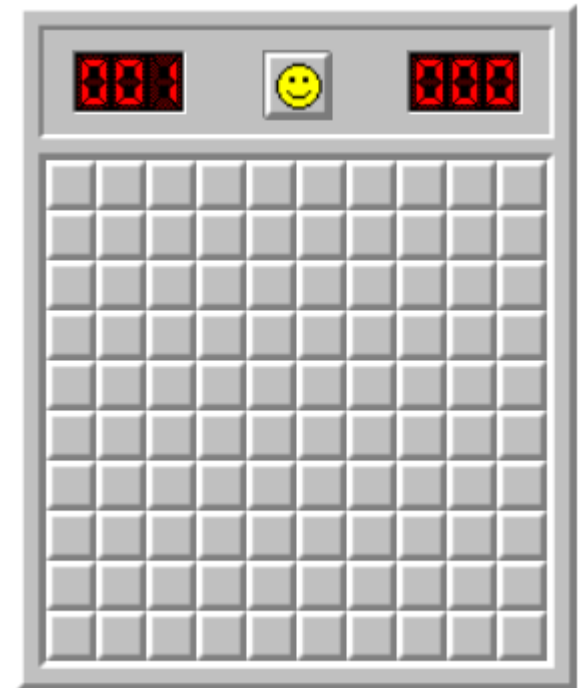
- With the game “minesweeper” on a 10x10 grid with only one landmine, how do we express in propositional logic the knowledge that squares adjacent to the landmine should display the number 1?



Discussion

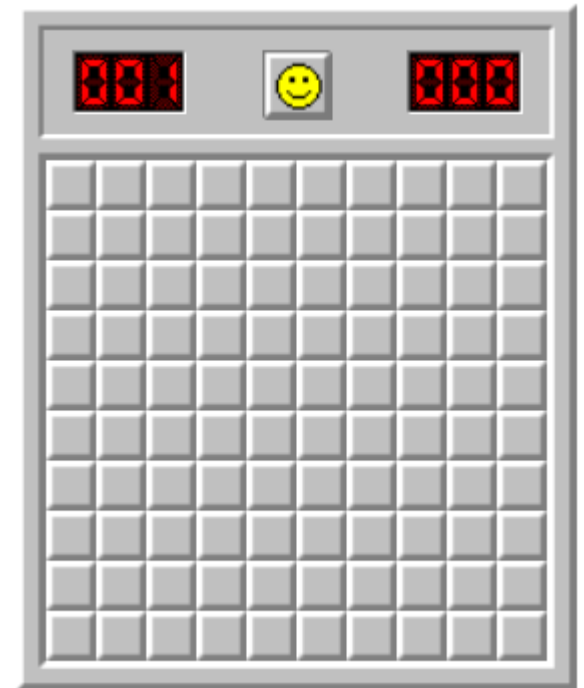
# Problems with propositional logic

- For example, for cell (2,3)
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_1\_2}$
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_1\_3}$
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_1\_4}$
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_2\_2}$
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_2\_4}$
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_3\_2}$
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_3\_3}$
  - $\text{Landmine\_2\_3} \Rightarrow \text{number1\_3\_4}$
- Similarly for other cells, resulting in explosion of symbols



# Today's learning objectives

- We will discover the **first order logic** which allows to write
  - `landmine(x,y)=>number1(neighbors(x,y))`



# Why not natural language?

- Can be imprecise and depend on context, e.g., "Spring":
  - mechanical?
  - a season?
  - flowing water?
- Using natural language as communication assumes large knowledge base, and seems to require some probability-based reasoning

# Why not programming languages?

- Most formal languages are procedural rather than declarative. You can have objects, but don't expect Java to reason about them automatically
- There are exceptions. **Prolog** can reason about statements like "All cars are red." (Constraint satisfaction and some formal logic tools are built-in to Prolog)

# Knowledge representation

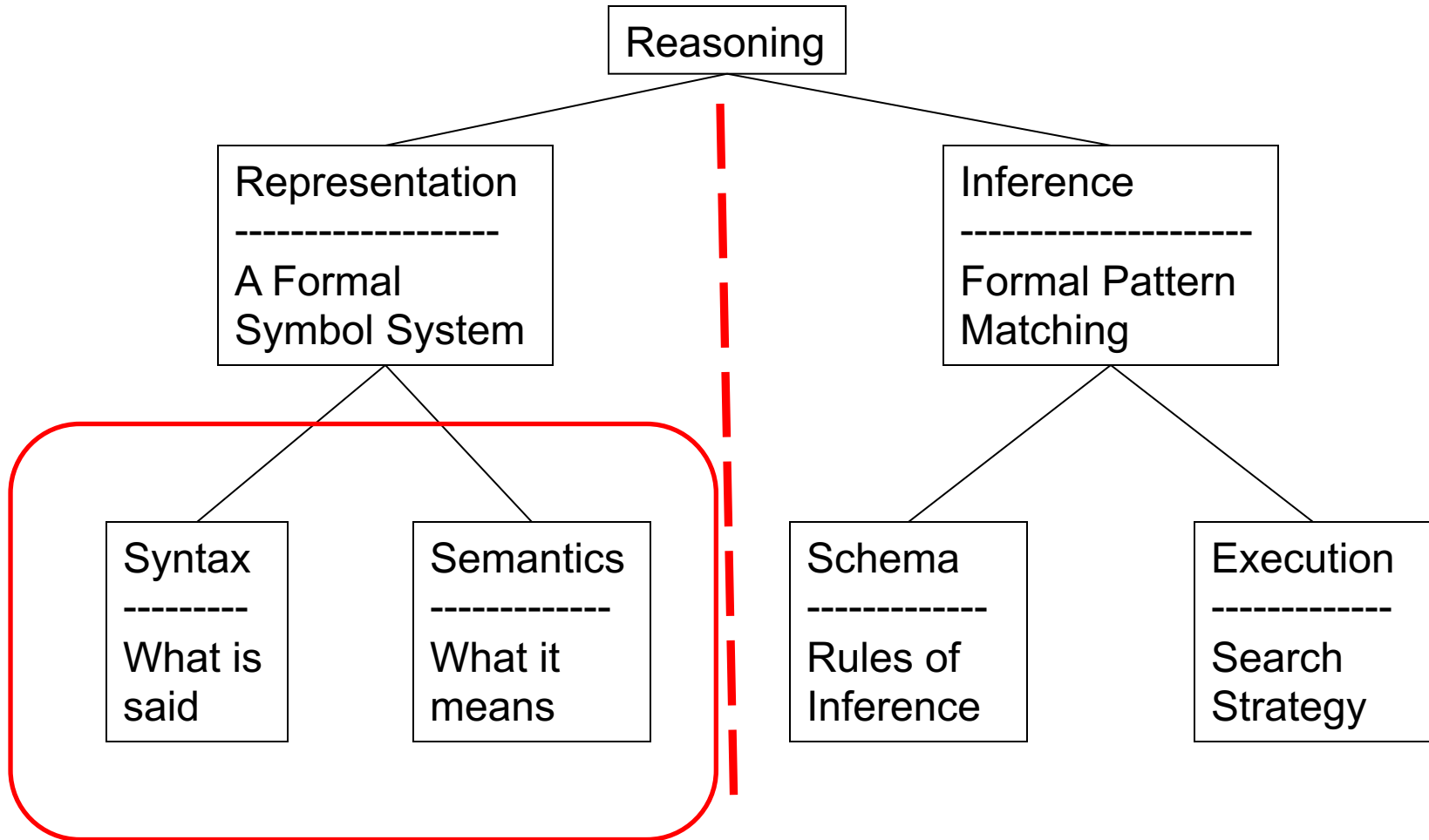
- We want something that:
  - has objects (like NL)
  - is declarative (like NL, propositional logic)
  - is context-independent (unlike NL, no "hidden" KB)
  - has a precise syntax

## FOL (or FOPC) Ontology:

What kind of things exist in the world?

What do we need to describe and reason about?

Objects --- with their relations, functions, predicates, properties, and general rules.





# Ontological commitment (what exists in the world)

- Looking at what propositional logic and first order logic assumes about the nature of reality

Logic	Primitives	Available knowledge
Propositional	Facts	True/false/unknown
First Order	Facts, objects, relations	True/false/unknown

# First-order logic (FOL)

- First-order logic includes:
  - Objects: people, houses, numbers, ...
    - Generally correspond to English nouns
  - Relations, properties, or maps take a tuple and return true or false:
    - Generally correspond to English verbs
    - First argument is generally the subject, the second the object, i.e., `Verb(Noun1, Noun2)` usually means “Noun1 verb noun2.”
  - Functions take in any number of objects and return one object – not true/false.

# Star Wars Examples

- Objects: Leia, Luke, The Empreor, Darth Vader
- Relation (binary): siblings(Leia,Luke) -> true
- Property (a unary relation): evil(emperor) -> true
- Function: father(Luke) is DarthVader
  - or equivalently, expressed as relation,
    - father(Obiwon,Luke)=false
    - father(Leia,Luke)=false
    - father(DarthVader,Luke)=true

# Syntax of FOL

- Three types of symbols:
  - Constant symbols (capture objects): KingJohn, 2, Dartmouth
  - Predicate symbols (capture relations): Brother, >,...
  - function symbols (capture functions): Sqrt, LeftLegOf
- Connectives:  $\wedge$  |  $\vee$  |  $\Rightarrow$  |  $\Leftrightarrow$  |  $\neg$  (standard)
- Equality: = Two symbols refer to the same object
- Variables: x, y, z
- Quantifiers:  $\forall$ ,  $\exists$ ; ways to refer to groups of objects.

# Syntax of FOL: Terms

- **Term** = logical expression that **refers to an object**
- **There are two kinds of terms:**
  - **Constant Symbols** stand for (or name) objects:
    - E.g., KingJohn, 2, UCI, Wumpus, ...
  - **Function Symbols** map tuples of objects to an object:
    - E.g., LeftLeg(KingJohn), Mother(Mary), Sqrt(x)

# Syntax of FOL: Atomic Sentences

- **Atomic Sentences** state facts (logical truth values).
  - An **atomic sentence** is a Predicate symbol, followed by a parenthesized list of any argument terms
    - E.g., *Married( Father(Richard), Mother(John) )*
  - An **atomic sentence** asserts that some relationship (some predicate) holds among the objects that are its arguments.
- An **Atomic Sentence is true** if the relation referred to by the predicate symbol holds among the objects (terms) referred to by the arguments.

# Syntax of FOL: Atomic Sentences

- Atomic sentences in logic state facts that are true or false.

LargerThan(2, 3) is false.

BrotherOf(Mary, Pete) is false.

Married(Father(Richard), Mother(John)) could be true or false.

- Note: Functions refer to objects, do not state facts:
  - Brother(Pete) refers to John (his brother) and is neither true nor false.
  - Plus(2, 3) refers to the number 5 and is neither true nor false.
- BrotherOf(Pete, Brother(Pete) ) is True.



Binary relation  
is a truth value.



Function refers to John, an object in the  
world, i.e., John is Pete's brother.  
(Works well iff John is Pete's only brother.)

# Syntax of FOL

- Three types of symbols:
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# Syntax of FOL:

## Connectives & Complex Sentences

- **Complex Sentences** are formed in the same way, using the same logical connectives, as in propositional logic
- The **Logical Connectives**:
  - $\Leftrightarrow$  biconditional
  - $\Rightarrow$  implication
  - $\wedge$  and
  - $\vee$  or
  - $\neg$  negation
- **Semantics** for these logical connectives are the same as we already know from propositional logic.

# Examples

- $\text{Brother}(\text{Richard}, \text{John}) \wedge \text{Brother}(\text{John}, \text{Richard})$
- $\text{King}(\text{Richard}) \vee \text{King}(\text{John})$
- $\text{King}(\text{John}) \Rightarrow \neg \text{King}(\text{Richard})$

(Semantics of complex sentences are the same as in propositional logic)

# Syntax of FOL

- Three types of symbols:
  - Constant symbols (capture objects): KingJohn, 2, Dartmouth
  - Predicate symbols (capture relations): Brother, >,...
  - function symbols (capture functions): Sqrt, LeftLegOf
- Connectives:  $\wedge$  |  $\vee$  |  $\Rightarrow$  |  $\Leftrightarrow$  |  $\neg$  (standard)
- Equality: = Two symbols refer to the same object
- Variables: x, y, z
- Quantifiers:  $\forall$ ,  $\exists$ ; ways to refer to groups of objects.

# Syntax of FOL: Variables

- **Variables** range over objects in the world.
- A **variable** is like a **term** because it represents an object.
- A **variable** may be used wherever a **term** may be used.
  - **Variables** may be arguments to functions and predicates.
- (A **term with NO variables** is called a **ground term**.)
- (A **variable not bound by a quantifier** is called **free**.)
  - All variables we will use are bound by a quantifier.

# Universal quantification

- Universal ( $\forall$ )
  - Sentence is true for all values of  $x$  in the domain of variable  $x$ .
  - Conjunction of all sentences obtained by substitution of an object for the quantified variable
- $\forall x \text{ human}(x) \Rightarrow \text{mammal}(x)$ 
  - What it really means (universal instantiation):
    - $\text{human}(\text{John}) \Rightarrow \text{mammal}(\text{John})$
    - $(\wedge) \text{ human}(\text{Alice}) \Rightarrow \text{mammal}(\text{Alice})$
    - $(\wedge) \text{ human}(\text{laptop}) \Rightarrow \text{mammal}(\text{laptop})$
    - ...

# Is this a correct sentence?

- $\forall x \text{ human}(x) \wedge \text{mammal}(x)$

Discussion

# Common mistake for universal quantification

- Common mistake is to use AND as main connective

$\forall x \text{ human}(x) \wedge \text{mammal}(x)$

- This means everything is human and a mammal!
- $(\text{human}(\text{Jerry}) \wedge \text{mammal}(\text{Jerry}) \wedge (\text{human}(\text{laptop}) \wedge \text{mammal}(\text{laptop}))) \wedge \dots$

- **Note that  $\Rightarrow$  is the natural connective to use with  $\forall$  .**

# Existential quantifiers

- Existential ( $\exists$ )
  - Sentence is true for some value of  $x$  in the domain of variable  $x$
  - Is equivalent to disjunction of all sentences obtained by substitution of an object for the quantified variable.
- “some humans are male”
  - $\exists x \text{ human}(x) \wedge \text{male}(x)$
  - Means there is an  $x$  who is a human and is a male
  - What it really means (existential instantiation):
    - $(\text{human}(\text{Jerry}) \wedge \text{male}(\text{Jerry})) \vee$
    - $(\text{human}(\text{laptop}) \wedge \text{male}(\text{laptop})) \vee \dots$



- “Some pig can fly”  $\exists x \text{ pig}(x) \Rightarrow \text{fly}(x)$   
(correct?)

# Common mistake for existential quantifiers

- Common mistake is to use  $\Rightarrow$  as main connective
- “Some pig can fly”  $\exists x \text{ pig}(x) \Rightarrow \text{fly}(x)$  (wrong)
  - This is true if there is something not a pig!  
 $(\text{pig}(\text{Jerry}) \Rightarrow \text{fly}(\text{Jerry})) \vee$   
 $(\text{pig}(\text{laptop}) \Rightarrow \text{fly}(\text{laptop})) \vee \dots$
- **Note that  $\wedge$  is the natural connective to use with  $\exists$ .**

## Combining Quantifiers – Order (Scope)

The order of “like” quantifiers does not matter.

$$\forall x \forall y P(x, y) \equiv \forall y \forall x P(x, y)$$

$$\exists x \exists y P(x, y) \equiv \exists y \exists x P(x, y)$$

**Like nested ANDs and ANDs in a logical sentence**

## Combining Quantifiers – Order (Scope)

The order of “unlike” quantifiers is important.  
**Like nested ANDs and ORs in a logical sentence.**

$\forall x \exists y \text{ Loves}(x,y)$

- For everyone (“all x”) there is someone (“exists y”) whom they love.
- There might be a different y for each x (y is inside the scope of x)

$\exists y \forall x \text{ Loves}(x,y)$

- There is someone (“exists y”) whom everyone loves (“all x”).
- Every x loves the same y (x is inside the scope of y)

Parentheses can clarify:  $\exists y ( \forall x \text{ Loves}(x,y) )$

# Properties of quantifiers

- $\forall x P(x)$  when negated becomes ?
- $\exists x P(x)$  when negated becomes ?

# Properties of quantifiers

- $\forall x P(x)$  when negated becomes  $\exists x \neg P(x)$
- $\exists x P(x)$  when negated becomes  $\forall x \neg P(x)$
- Example
  - $\forall x \text{ sleep}(x)$ 
    - It means everybody sleeps
  - If negated, it becomes  $\exists x \neg \text{sleep}(x)$ 
    - There is somebody who doesn't sleep

# Properties of quantifiers

- $\forall x P(x)$  is logically equivalent to  $\equiv \neg \exists x \neg P(x)$
- $\exists x P(x)$  is logically equivalent to  $\equiv \neg \forall x \neg P(x)$
- Example
  - $\forall x \text{ sleep}(x)$ 
    - It means everybody sleeps
  - $\neg \exists x \neg \text{sleep}(x)$ 
    - There is nobody who doesn't sleep

# Connections between Quantifiers

In effect:

- $\forall$  is a conjunction over the universe of objects
- $\exists$  is a disjunction over the universe of objects

Thus, DeMorgan's rules can be applied



# De Morgan's Law for Quantifiers

## De Morgan's Rule

$$P \wedge Q \equiv \neg (\neg P \vee \neg Q)$$

$$P \vee Q \equiv \neg (\neg P \wedge \neg Q)$$

$$\neg (P \wedge Q) \equiv (\neg P \vee \neg Q)$$

$$\neg (P \vee Q) \equiv (\neg P \wedge \neg Q)$$

## Generalized De Morgan's Rule

$$\forall x P(x) \equiv \neg \exists x \neg P(x)$$

$$\exists x P(x) \equiv \neg \forall x \neg P(x)$$

$$\neg \forall x P(x) \equiv \exists x \neg P(x)$$

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

**AND/OR Rule is simple:** if you bring a negation inside a disjunction or a conjunction, always switch between them ( $\neg$  OR  $\rightarrow$  AND  $\neg$  ;  $\neg$  AND  $\rightarrow$  OR  $\neg$ ).

**QUANTIFIER Rule is similar:** if you bring a negation inside a universal or existential, always switch between them ( $\neg \exists \rightarrow \forall \neg$  ;  $\neg \forall \rightarrow \exists \neg$ ).

# Example of sentences with quantifiers

- **“All persons are mortal.”**

[Use: Person(x), Mortal (x) ]

$$\forall x \text{ Person}(x) \Rightarrow \text{Mortal}(x)$$

- **Equivalent Forms:**

$$\forall x \neg \text{Person}(x) \vee \text{Mortal}(x)$$

- **Common Mistakes:**

$$\forall x \text{ Person}(x) \wedge \text{Mortal}(x)$$



Example

# Example of sentences with quantifiers

- **“Sissy has a sister who is a cat.”**

[Use: Sister(Sissy, x), Cat(x) ]

$$\exists x \text{ Sister}(\text{Sissy}, x) \wedge \text{Cat}(x)$$

- **Common Mistakes:**

$$\exists x \text{ Sister}(\text{Sissy}, x) \Rightarrow \text{Cat}(x)$$



Example

# Example of sentences with quantifiers

- “For every food, there is a person who eats that food.”

[Use: Food(x), Person(y), Eats(y, x) ]

$$\forall x \exists y \text{ Food}(x) \Rightarrow [ \text{Person}(y) \wedge \text{Eats}(y, x) ]$$

- **Equivalent Forms:**

$$\forall x \text{ Food}(x) \Rightarrow \exists y [ \text{Person}(y) \wedge \text{Eats}(y, x) ]$$

$$\forall x \exists y \neg \text{Food}(x) \vee [ \text{Person}(y) \wedge \text{Eats}(y, x) ]$$

$$\forall x \exists y [ \neg \text{Food}(x) \vee \text{Person}(y) ] \wedge [ \neg \text{Food}(x) \vee \text{Eats}(y, x) ]$$

$$\forall x \exists y [ \text{Food}(x) \Rightarrow \text{Person}(y) ] \wedge [ \text{Food}(x) \Rightarrow \text{Eats}(y, x) ]$$

- **Common Mistakes:**

$$\forall x \exists y [ \text{Food}(x) \wedge \text{Person}(y) ] \Rightarrow \text{Eats}(y, x)$$

$$\forall x \exists y \text{ Food}(x) \wedge \text{Person}(y) \wedge \text{Eats}(y, x)$$

Example

# Example of sentences with quantifiers

- “Every person eats some food.”

[Use: Person (x), Food (y), Eats(x, y) ]

$$\forall x \exists y \text{ Person}(x) \Rightarrow [ \text{Food}(y) \wedge \text{Eats}(x, y) ]$$

- **Equivalent Forms:**

$$\forall x \text{ Person}(x) \Rightarrow \exists y [ \text{Food}(y) \wedge \text{Eats}(x, y) ]$$

$$\forall x \exists y \neg \text{Person}(x) \vee [ \text{Food}(y) \wedge \text{Eats}(x, y) ]$$

$$\forall x \exists y [ \neg \text{Person}(x) \vee \text{Food}(y) ] \wedge [ \neg \text{Person}(x) \vee \text{Eats}(x, y) ]$$

- **Common Mistakes:**

$$\forall x \exists y [ \text{Person}(x) \wedge \text{Food}(y) ] \Rightarrow \text{Eats}(x, y)$$

$$\forall x \exists y \text{ Person}(x) \wedge \text{Food}(y) \wedge \text{Eats}(x, y)$$



Example

# Example of sentences with quantifiers

- **“Some person eats some food.”**

[Use: Person (x), Food (y), Eats(x, y) ]

$$\exists x \exists y \text{ Person}(x) \wedge \text{Food}(y) \wedge \text{Eats}(x, y)$$

- **Common Mistakes:**

$$\exists x \exists y [ \text{Person}(x) \wedge \text{Food}(y) ] \Rightarrow \text{Eats}(x, y)$$



Example

# Example of sentences with quantifiers

- “Everyone has a favorite food.”

[Use: Person(x), Food(y), Favorite(y, x) ]

- **Equivalent Forms:**

- $\forall x \exists y \text{ Person}(x) \Rightarrow [ \text{Food}(y) \wedge \text{Favorite}(y, x) ]$
- $\forall x \text{ Person}(x) \Rightarrow \exists y [ \text{Food}(y) \wedge \text{Favorite}(y, x) ]$
- $\forall x \exists y \neg \text{Person}(x) \vee [ \text{Food}(y) \wedge \text{Favorite}(y, x) ]$
- $\forall x \exists y [ \neg \text{Person}(x) \vee \text{Food}(y) ] \wedge [ \neg \text{Person}(x) \vee \text{Favorite}(y, x) ]$
- $\forall x \exists y [ \text{Person}(x) \Rightarrow \text{Food}(y) ] \wedge [ \text{Person}(x) \Rightarrow \text{Favorite}(y, x) ]$

- **Common Mistakes:**

- $\forall x \exists y [ \text{Person}(x) \wedge \text{Food}(y) ] \Rightarrow \text{Favorite}(y, x)$
- $\forall x \exists y \text{ Person}(x) \wedge \text{Food}(y) \wedge \text{Favorite}(y, x)$

Example

# Equality

- $term_1 = term_2$  is true under a given **interpretation**

if and only if  $term_1$  and  $term_2$  refer to the same object

- E.g., definition of *Sibling* in terms of *Parent*, using = is:

$$\begin{aligned} \forall x, y \text{ Sibling}(x, y) \Leftrightarrow \\ [\neg(x = y) \wedge \\ \exists m, f \neg(m = f) \wedge \text{Parent}(m, x) \wedge \text{Parent}(f, x) \\ \wedge \text{Parent}(m, y) \wedge \text{Parent}(f, y)] \end{aligned}$$

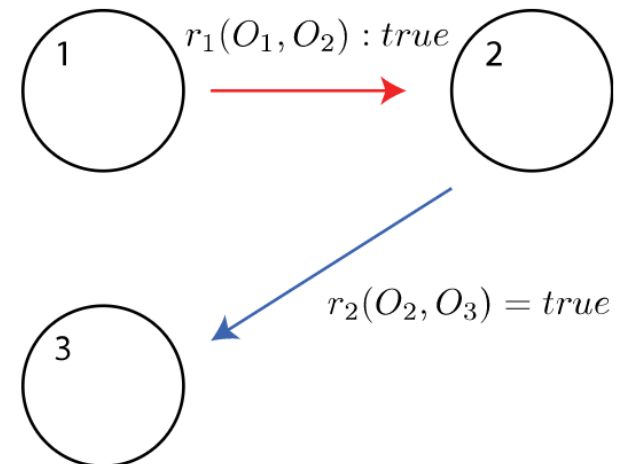


# Semantics

- sentences + (model, interpretation)  $\mapsto$  true/false
- interpretation specifies exactly which objects, relations, and functions are referred to by the constant, predicate, and function symbols.

‘=’ sign is used

- Models, objects, relations



# Models

- A set of true/false values for every relation among objects. (Think of a set of directed edges, with different colors for each relation, of graph.)
- $r_1(O_1, O_2) = tr_1(O_1, O_2) = t$ ,  $r_1(O_1, O_3) = fr_1(O_1, O_3) = f$ ,  $r_1(O_2, O_1) = fr_1(O_2, O_1) = f$ , ...

# How many models?

- For each binary relation (possible edge), there are  $n^2$  possible object pairs (2-tuples),  $n^3$  possible ternary relations,  $n^k$  possible  $k$ -ary relations.
- That's just the number of tuples. Each can be true or false. So for each relation, we get a factor of  $2^{(n^k)}$  models.
- $n$  might be infinite. (Maybe the objects are natural numbers, which can be described in FOL.)

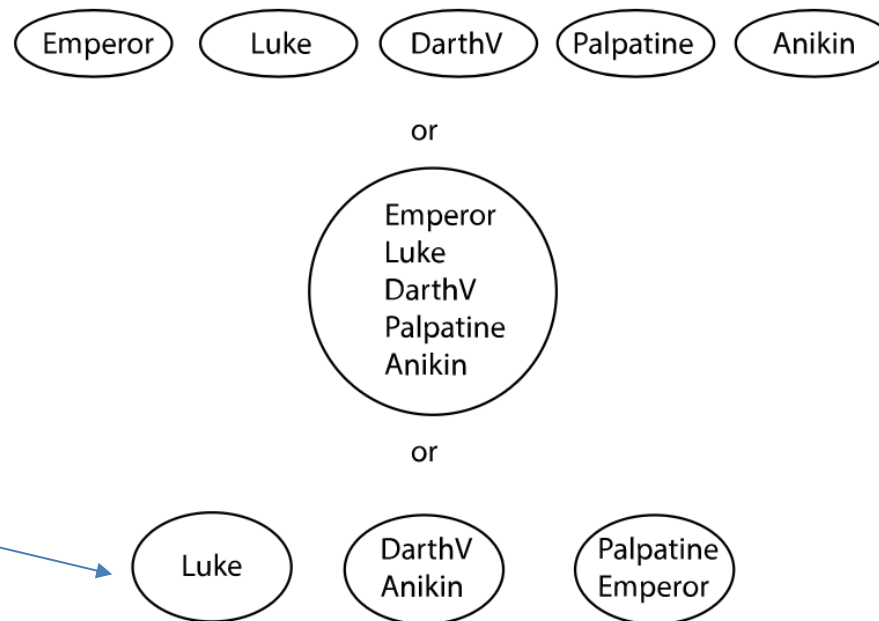
Discussion

# Interpretation

- Computational complexity gets even worse
- The syntax doesn't bind symbols to particular objects

# Example

- Symbols: Luke, DarthVader, Emperor, Palpatine, Anikin. Five symbols, but how many objects?



# Syntactic Ambiguity

- FOL provides many ways to represent the same thing.
- E.g., “Ball-5 is red.”
  - HasColor(Ball-5, Red)
    - Ball-5 and Red are objects related by HasColor.
  - Red(Ball-5)
    - Red is a unary predicate applied to the Ball-5 object.
  - HasProperty(Ball-5, Color, Red)
    - Ball-5, Color, and Red are objects related by HasProperty.
  - ColorOf(Ball-5) = Red
    - Ball-5 and Red are objects, and ColorOf() is a function.
  - HasColor(Ball-5(), Red())
    - Ball-5() and Red() are functions of zero arguments that both return an object, which objects are related by HasColor.
  - ...
- This can GREATLY confuse a pattern-matching reasoner.
  - Especially if multiple people collaborate to build the KB, and they all have different representational conventions.

# Syntactic Ambiguity – Partial solution

- FOL can be TOO expressive, can offer TOO MANY choices
- Likely confusion, especially for **teams** of Knowledge Engineers
- Different team members can make different representation choices
  - E.g., represent “Ball43 is Red.” as:
    - a property (= adjective)? E.g., “Red(Ball43)” ?
    - an object (= noun)? E.g., “Red = Color(Ball43)” ?
    - a predicate (= verb)? E.g., “HasProperty(Ball43, Red)” ?
- PARTIAL SOLUTION:
  - An upon-agreed **ontology** that settles these questions
  - Ontology = what exists in the world & how it is represented
  - The Knowledge Engineering teams agrees upon an ontology BEFORE they begin encoding knowledge

# Summary

- First order logic to represent also objects and relations
  - Syntax includes sentences, predicate symbols, function symbols, constant symbols, variables, quantifiers
- Nested quantifiers
  - Order of unlike quantifiers matters (the outer scopes the inner)
    - Like nested ANDs and ORs
  - Order of like quantifiers does not matter
    - like nested ANDs and ANDs
- Semantics needs also interpretation



# Next

- How do we make inference with FOL?