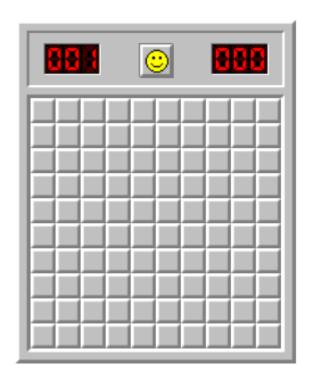
# COSC76/276 Artificial Intelligence Fall 2022 First Order Logic

Soroush Vosoughi
Computer Science
Dartmouth College
Soroush@Dartmouth.edu

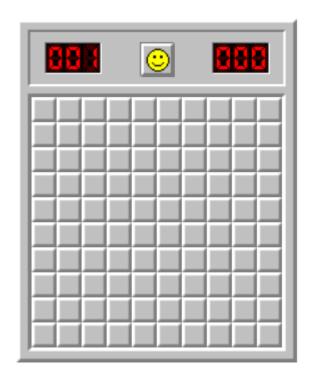
## Problems with propositional logic

- For example, for cell (2,3)
  - Landmine\_2\_3=>number1\_1\_2
  - Landmine\_2\_3=>number1\_1\_3
  - Landmine 2 3=>number1 1 4
  - Landmine\_2\_3=>number1\_2\_2
  - Landmine\_2\_3=>number1\_2\_4
  - Landmine 2 3=>number1 3 2
  - Landmine\_2\_3=>number1\_3\_3
  - Landmine\_2\_3=>number1\_3\_4
- Similarly for other cells, resulting in explosion of symbols



## **First Order Logic**

- We will discover the first order logic which allows to write
  - landmine(x,y)=>number1(
     neighbors(x,y))



## **Syntax of FOL**

- Three types of symbols:
  - Constant symbols (capture objects): KingJohn, 2,
     Dartmouth
  - Predicate symbols (capture relations): Brother, >,...
  - function symbols (capture functions): Sqrt, LeftLegOf
- Connectives:  $\land |\lor| \Rightarrow | \Leftrightarrow | \neg \text{ (standard)}$
- Equality: = Two symbols refer to the same object
- Variables: x, y, z
- Quantifiers: ∀,∃; ways to refer to groups of objects.

## Syntax of FOL: Connectives & Complex Sentences

- Complex Sentences are formed in the same way, using the same logical connectives, as in propositional logic
- The Logical Connectives:
  - − ⇔ biconditional
  - $\Rightarrow$  implication
  - $\wedge and$
  - $\vee or$
  - − ¬ negation
- Semantics for these logical connectives are the same as we already know from propositional logic.

## **Examples**

- Brother(Richard, John) ∧ Brother(John, Richard)
- King(Richard) ∨ King(John)
- King(John) => ¬ King(Richard)

(Semantics of complex sentences are the same as in propositional logic)

## **Syntax of FOL**

- Three types of symbols:
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     Dartmouth
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- Connectives:  $\land |\lor| \Rightarrow |\Leftrightarrow| \neg \text{ (standard)}$
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## **Syntax of FOL: Variables**

- Variables range over objects in the world.
- (A variable not bound by a quantifier is called free.)
  - All variables we will use are bound by a quantifier.

## **Universal quantification**

- Universal (∀)
  - Sentence is true for all values of x in the domain of variable x.
  - Conjunction of all sentences obtained by substitution of an object for the quantified variable
- $\forall x \text{ human}(x) \Rightarrow \text{mammal}(x)$ 
  - What it really means (universal instantiation):

```
human(John)⇒mammal(John)
```

 $(\land)$  human(Alice) $\Rightarrow$ mammal(Alice)

 $(\land)$  human(laptop) $\Rightarrow$ mammal(laptop)

• • •

## Is this a correct sentence?

•  $\forall x \text{ human}(x) \land \text{mammal}(x)$ 



## Common mistake for universal quantification

- Common mistake is to use AND as main connective
  - $\forall$  x human(x)  $\land$  mammal(x)
    - This means everything is human and a mammal!
    - (human(Jerry) ∧ mammal(Jerry ∧ (human(laptop)∧ mammal(laptop)) ∧ ...
- Note that => is the natural connective to use with ∀.

## **Existential quantifiers**

- Existential (∃)
  - Sentence is true for some value of x in the domain of variable x
  - Is equivalent to disjunction of all sentences obtained by substitution of an object for the quantified variable.

- "some humans are male"
  - $-\exists x human(x) \land male(x)$
  - Means there is an x who is a human and is a male
  - What it really means (existential instantiation):
     (human(Jerry) ∧ male(Jerry)) ∨

```
(human(laptop) \( \text{male(laptop)} \( \text{v} \)...
```

"Some pig can fly" ∃ x pig(x) => fly(x)
 (correct?)

## Common mistake for existential quantifiers

- Common mistake is to use => as main connective
- "Some pig can fly"  $\exists x pig(x) => fly(x)$  (wrong)
  - This is true if there is something not a pig! (pig(Jerry) => fly(Jerry)) V (pig(laptop) => fly(laptop)) V ...
- Note that ∧ is the natural connective to use with ∃.

### **Combining Quantifiers – Order (Scope)**

The order of "like" quantifiers does not matter.

$$\forall x \ \forall y \ P(x, y) \equiv \forall y \ \forall x \ P(x, y)$$
  
 $\exists x \ \exists y \ P(x, y) \equiv \exists y \ \exists x \ P(x, y)$ 

Like nested ANDs and ANDs in a logical sentence

#### **Combining Quantifiers – Order (Scope)**

The order of "unlike" quantifiers is important. Like nested ANDs and ORs in a logical sentence.

```
\forall x \exists y Loves(x,y)
```

- For everyone ("all x") there is someone ("exists y") whom they love.
- There might be a different y for each x (y is inside the scope of x)

```
\exists y \forall x Loves(x,y)
```

- There is someone ("exists y") whom everyone loves ("all x").
- Every x loves the same y (x is inside the scope of y)

Parentheses can clarify:  $\exists y ( \forall x \text{ Loves}(x,y))$ 

## **Properties of quantifiers**

- $\forall x P(x)$  when negated becomes ?
- $\exists x P(x)$  when negated becomes?

## **Properties of quantifiers**

•  $\forall x P(x)$  when negated becomes  $\exists x \neg P(x)$ 

•  $\exists x P(x)$  when negated becomes  $\forall x \neg P(x)$ 

- Example
  - $\forall x \text{ sleep(x)}$ 
    - It means everybody sleeps
  - If negated, it becomes  $\exists x \neg sleep(x)$ 
    - There is somebody who doesn't sleep

## Properties of quantifiers

•  $\forall x P(x)$  is logically equivalent to  $\equiv \neg \exists x \neg P(x)$ 

•  $\exists x P(x)$  is logically equivalent to  $\equiv \neg \forall x \neg P(x)$ 

- Example
  - $\forall x \text{ sleep(x)}$ 
    - It means everybody sleeps
  - $-\neg\exists x \neg sleep(x)$ 
    - There is nobody who doesn't sleep

## **Connections between Quantifiers**

#### In effect:

- $\forall$  is a conjunction over the universe of objects
- ∃ is a disjunction over the universe of objects
  Thus, DeMorgan's rules can be applied

## De Morgan's Law for Quantifiers

#### De Morgan's Rule

#### Generalized De Morgan's Rule

$$P \wedge Q \equiv \neg (\neg P \vee \neg Q) \qquad \forall x P(x) \equiv \neg \exists x \neg P(x)$$

$$P \vee Q \equiv \neg (\neg P \wedge \neg Q) \qquad \exists x P(x) \equiv \neg \forall x \neg P(x)$$

$$\neg (P \wedge Q) \equiv (\neg P \vee \neg Q) \qquad \neg \forall x P(x) \equiv \exists x \neg P(x)$$

$$\neg (P \vee Q) \equiv (\neg P \wedge \neg Q) \qquad \neg \exists x P(x) \equiv \forall x \neg P(x)$$

**AND/OR Rule is simple:** if you bring a negation inside a disjunction or a conjunction, always switch between them ( $\neg$  OR  $\rightarrow$  AND  $\neg$ ;  $\neg$  AND  $\rightarrow$  OR  $\neg$ ).

**QUANTIFIER Rule is similar:** if you bring a negation inside a universal or existential, always switch between them  $(\neg \exists \rightarrow \forall \neg; \neg \forall \rightarrow \exists \neg)$ .

"All persons are mortal."

[Use: Person(x), Mortal (x)]

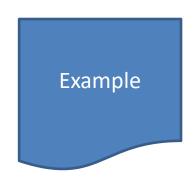
 $\forall x \ Person(x) \Rightarrow Mortal(x)$ 

• Equivalent Forms:

 $\forall x \neg Person(x) \lor Mortal(x)$ 

Common Mistakes:

 $\forall x \ Person(x) \land Mortal(x)$ 



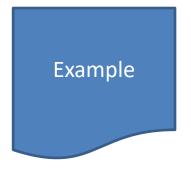
"Sissy has a sister who is a cat."

[Use: Sister(Sissy, x), Cat(x)]

 $\exists x \ Sister(Sissy, x) \land Cat(x)$ 

Common Mistakes:

 $\exists x \; Sister(Sissy, x) \Rightarrow Cat(x)$ 



"For every food, there is a person who eats that food."

[Use: Food(x), Person(y), Eats(y, x)]  

$$\forall x \exists y \text{ Food}(x) \Rightarrow [\text{ Person}(y) \land \text{ Eats}(y, x)]$$

Equivalent Forms:

```
\forall x \ \mathsf{Food}(x) \Rightarrow \exists y \ [ \ \mathsf{Person}(y) \land \mathsf{Eats}(y, x) \ ]

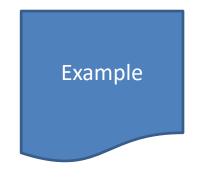
\forall x \ \exists y \ \neg \mathsf{Food}(x) \ \lor \ [ \ \mathsf{Person}(y) \land \mathsf{Eats}(y, x) \ ]

\forall x \ \exists y \ [ \ \neg \mathsf{Food}(x) \ \lor \ \mathsf{Person}(y) \ ] \land [ \ \neg \mathsf{Food}(x) \Rightarrow \mathsf{Eats}(y, x) \ ]

\forall x \ \exists y \ [ \ \mathsf{Food}(x) \Rightarrow \mathsf{Person}(y) \ ] \land [ \ \mathsf{Food}(x) \Rightarrow \mathsf{Eats}(y, x) \ ]
```

Common Mistakes:

```
\forall x \exists y [ Food(x) \land Person(y) ] \Rightarrow Eats(y, x) 
\forall x \exists y Food(x) \land Person(y) \land Eats(y, x)
```



"Every person eats some food."

```
[Use: Person (x), Food (y), Eats(x, y)]
```

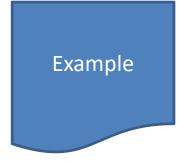
```
\forall x \exists y \ \mathsf{Person}(x) \Rightarrow [\ \mathsf{Food}(y) \land \mathsf{Eats}(x, y) \ ]
```

Equivalent Forms:

```
\forall x \ \mathsf{Person}(x) \Rightarrow \exists y \ [ \ \mathsf{Food}(y) \land \mathsf{Eats}(x, y) \ ]
\forall x \ \exists y \ \neg \mathsf{Person}(x) \ \lor \ [ \ \mathsf{Food}(y) \land \mathsf{Eats}(x, y) \ ]
\forall x \ \exists y \ [ \ \neg \mathsf{Person}(x) \ \lor \ \mathsf{Food}(y) \ ] \land [ \ \neg \mathsf{Person}(x) \ \lor \ \mathsf{Eats}(x, y) \ ]
```

Common Mistakes:

```
\forall x \exists y [ Person(x) \land Food(y) ] \Rightarrow Eats(x, y)
\forall x \exists y Person(x) \land Food(y) \land Eats(x, y)
```



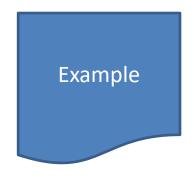
"Some person eats some food."

[Use: Person (x), Food (y), Eats(x, y)]

 $\exists x \exists y \ Person(x) \land Food(y) \land Eats(x, y)$ 

Common Mistakes:

 $\exists x \exists y [ Person(x) \land Food(y) ] \Rightarrow Eats(x, y)$ 



"Everyone has a favorite food."

```
[Use: Person(x), Food(y), Favorite(y, x)]
```

#### Equivalent Forms:

```
• \forall x \exists y \, \text{Person}(x) \Rightarrow [\, \text{Food}(y) \land \text{Favorite}(y, x) \,]
```

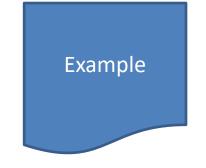
- $\forall x \ \mathsf{Person}(x) \Rightarrow \exists y \ [ \ \mathsf{Food}(y) \land \mathsf{Favorite}(y, x) \ ]$
- ∀x ∃y ¬Person(x) ∨ [ Food(y) ∧ Favorite(y, x) ]
- ∀x ∃y [¬Person(x) ∨ Food(y)] ∧ [¬Person(x)

```
Favorite(y, x) ]
```

•  $\forall x \exists y [Person(x) \Rightarrow Food(y)] \land [Person(x) \Rightarrow Favorite(y, x)]$ 

#### Common Mistakes:

- $\forall x \exists y [ Person(x) \land Food(y) ] \Rightarrow Favorite(y, x)$
- $\forall x \exists y \, Person(x) \land Food(y) \land Favorite(y, x)$



## **Equality**

- term<sub>1</sub> = term<sub>2</sub> is true
   if and only if term<sub>1</sub> and term<sub>2</sub> refer to the
   same object
- E.g., definition of Sibling in terms of Parent, using = is:

```
\forall x,y \ Sibling(x,y) \Leftrightarrow
[\neg(x = y) \land \\ \exists m,f \ \neg (m = f) \land Parent(m,x) \land Parent(f,x) \\ \land Parent(m,y) \land Parent(f,y)]
```

## **Semantics**

- sentences + (model, interpretation) → true/false
- interpretation specifies exactly which objects, relations, and functions are referred to by the constant, predicate, and function symbols.

'=' sign is used

## **Models**

 A set of true/false values for every relation among objects. (Think of a set of directed edges, with different colors for each relation, of graph.)

## **How many models?**

- For each binary relation, there are n^2
  possible object pairs (2-tuples), n^3 possible
  ternary relations, n^k possible k-ary relations.
- That's just the number of tuples. Each can be true or false. So for each relation, we get a factor of 2<sup>(n^k)</sup> models.
- n might be infinite. (Maybe the objects are natural numbers, which can be described in FOL.)

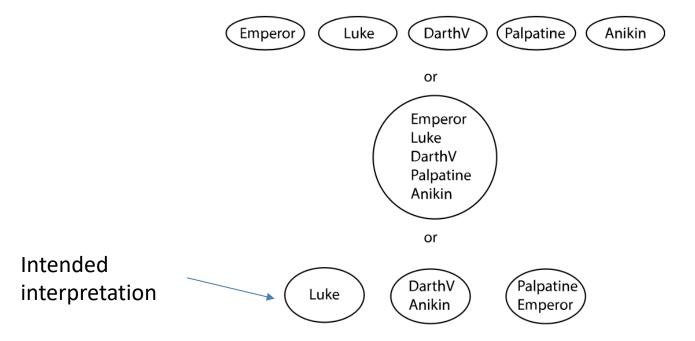
Discussion

## **Interpretation**

Computational complexity gets even worse

## **Example**

Symbols: Luke, DarthVader, Emperor,
 Palpatine, Anikin. Five symbols, but how many objects?



## **Syntactic Ambiguity**

- FOL provides many ways to represent the same thing.
- E.g., "Ball-5 is red."
  - HasColor(Ball-5, Red)
    - Ball-5 and Red are objects related by HasColor.
  - Red(Ball-5)
    - Red is a unary predicate applied to the Ball-5 object.
  - HasProperty(Ball-5, Color, Red)
    - Ball-5, Color, and Red are objects related by HasProperty.
  - ColorOf(Ball-5) = Red
    - Ball-5 and Red are objects, and ColorOf() is a function.
  - HasColor(Ball-5(), Red())
    - Ball-5() and Red() are functions of zero arguments that both return an object, which objects are related by HasColor.
  - **–** ...
- This can GREATLY confuse a pattern-matching reasoner.
  - Especially if multiple people collaborate to build the KB, and they all have different representational conventions.

## **Syntactic Ambiguity – Partial solution**

- FOL can be TOO expressive, can offer TOO MANY choices
- Likely confusion, especially for teams of Knowledge Engineers
- Different team members can make different representation choices
  - E.g., represent "Ball43 is Red." as:
    - a property (= adjective)? E.g., "Red(Ball43)"?
    - an object (= noun)? E.g., "Red = Color(Ball43))"?
    - a predicate (= verb)? E.g., "HasProperty(Ball43, Red)"?

#### PARTIAL SOLUTION:

- An upon-agreed ontology that settles these questions
- Ontology = what exists in the world & how it is represented
- The Knowledge Engineering teams agrees upon an ontology BEFORE they begin encoding knowledge

## **Summary**

- First order logic to represent also objects and relations
  - Syntax includes sentences, predicate symbols, function symbols, constant symbols, variables, quantifiers
- Nested quantifiers
  - Order of unlike quantifiers matters (the outer scopes the inner)
    - Like nested ANDs and ORs
  - Order of like quantifiers does not matter
    - like nested ANDS and ANDs
- Semantics needs also interpretation

## **Next**

How do we make inference with FOL?

## **Logical agent with FOL**

- Sentences are added to a knowledge base using TELL
  - TELL(KB, King(John)).
  - TELL(KB, Person(Richard))
  - TELL(KB,  $\forall$  x King(x)  $\Rightarrow$  Person(x)).
- We can ask questions of the knowledge base using ASK. E.g.,
  - ASK(KB, King(John)) returns true.
- Query that is logically entailed by the knowledge base should be answered affirmatively.
  - E.g., given the two preceding assertions, the query ASK(KB, Person(John)) should also return true.
- We can ask quantified queries, such as
  - ASK(KB,  $\exists$  x Person(x)).
  - True answer, but not very helpful. It is like answering "Can you tell me the time?" with "Yes."
- ASKVARS returns what value of x makes the sentence true
  - ASKVARS(KB, Person(x))
  - E.g., there will be two answers: {x/John} and {x/Richard} -- answer called a substitution or binding list.

## FOL Version of Wumpus World

- Suppose a wumpus-world agent is using an FOL KB and perceives a smell and a breeze (but no glitter) at t = 5:
- Typical percept sentence:

  TELL(KB, Percept([Stench, Breeze, not Glitter, None, None], t=5))
- Actions:
  Turn(Right), Turn(Left), Forward, Shoot, Grab, Release, Climb
- To determine best action, construct query:

```
ASK(∃ a BestAction(a,5))
ASKVARS(BestAction(a,5))
```

Inference to return {a/Grab}

## **Knowledge Base for Wumpus World**

#### Perception

- $\forall$ s,g,x,y,t Percept([s,Breeze,g,x,y],t)  $\Rightarrow$  Breeze(t)
- $\forall$ s,b,x,y,t Percept([s,b,Glitter,x,y],t)  $\Rightarrow$  Glitter(t)

#### Reflex action

-  $\forall$ t Glitter(t)  $\Rightarrow$  BestAction(Grab,t)

#### Reflex action with internal state

-  $\forall$ t Glitter(t)  $\land$ ¬Holding(Gold,t)  $\Rightarrow$  BestAction(Grab,t)

Holding(Gold,t) can not be observed: keep track of change.

## **Deducing hidden properties**

#### **Environment definition:**

```
\forallx,y,a,b Adjacent([x,y],[a,b]) \Leftrightarrow
[a,b] \in {[x+1,y], [x-1,y],[x,y+1],[x,y-1]}
```

#### Properties of locations:

```
\foralls,t At(Agent,s,t) \land Breeze(t) \Rightarrow Breezy(s)
```

#### Squares are breezy near a pit:

- Diagnostic rule---infer cause from effect
   ∀s Breezy(s) ⇔ ∃ r Adjacent(r,s) ∧ Pit(r)
- Causal rule---infer effect from cause (model based reasoning)
   ∀r Pit(r) ⇒ [∀s Adjacent(r,s) ⇒ Breezy(s)]

## Knowledge engineering in FOL

- 1. Identify the task
- 2. Assemble the relevant knowledge
- 3. Decide on a vocabulary of predicates, functions, and constants
- 4. Encode general knowledge about the domain
- 5. Encode a description of the specific problem instance
- 6. Pose queries to the inference procedure and get answers
- 7. Debug the knowledge base

An interpretation maps all symbols in KB onto matching symbols in a possible world. All possible interpretations gives a combinatorial explosion of mappings. Your job, as a Knowledge Engineer, is to write the axioms in KB so <u>they are satisfied only under the</u> intended interpretation in your own real world.