An Experimental Comparison of Concurrent Data Structures

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What is the Problem?

- Concurrent Data Structure
 - Designed for access by multiple threads
 - Potential to be highly scalable
- Plenty of work done on how to implement concurrent data structures
- Not much data on comparing the different types of concurrent data structure

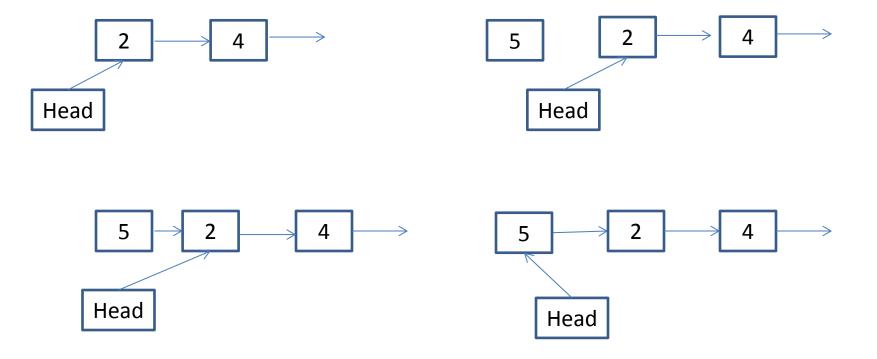
Atomic Instructions

- Either complete fully or not at all
- Are used to implement locked and lockless algorithms
- Example: compare-and-swap instruction

```
If(*lock == 0)
{
    *lock = 1;
}
```

Atomic Instructions

 Can use compare-and-swap to atomically add a node to the head of a linked list:



Atomic Instructions

Compare-and-swap can be used locklessly as follows:

```
if(std::atomic_compare_exchange_strong(Atomic *Obj, *Expected, Desired))
{
     //Value of Obj was equal to Expected, now equal to Desired
}
```

 Similarly we can use compare-and-swap to implement a lock which will allow us to do the same thing:

```
while(true){
          if(std::atomic_compare_exchange_weak(Atomic * Obj, 0, 1))break;//Try and acquire lock
          _mm_pause();
}
//Lock acquired, change value of head to new node
```

Locking Algorithms

Locked

- Uses mutexes and semaphores to acquire a lock
- Blocks threads that do not have the lock

Lock Free

- Uses atomic instructions such as compare-and-swap
- Guarantees system-wide throughput with the chance of starvation

Wait Free

Similar to lock free but is also starvation free

Background Work

- Some of my references:
 - "The Art of Multiprocessor Programming" Herlihy
 & Shavit 2008
 - "Designing Concurrent Data Structures" Moir & Shavit – 2001
 - "Implementing Concurrent Data Objects" –Herlihy 1993
 - Locklessinc.com

What I have Done

- Implemented 3 concurrent data structures
 - Ring Buffer
 - Linked List
 - Hash Table
- Implemented both locked and lockless variations
- Compared them on 3 different systems

Data Structure Implementation

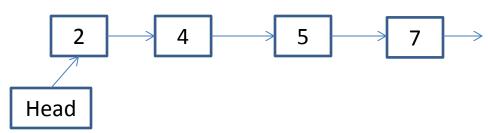
- Locked Variations
 - Mutex
 - Test-and-set
 - Ticket Lock
- Lockless Variations
 - C++11 atomic library operations
 - Atomic Compare-and-swap
 - Atomic Fetch-and-add
- Multi Producer Multi Consumer (MPMC)
- Single Producer Single Consumer (SPSC)

Ring Buffer

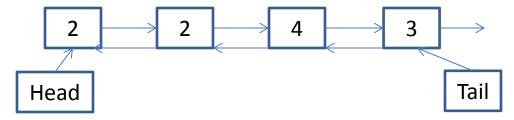
- Used to get to grips with the C++11 library
- Variations
 - MPMC Locked
 - SPSC Lockless

Linked List

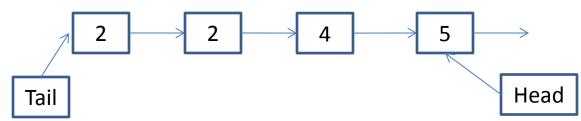
MPMC Singly Linked List – Locked/Lockless



MPMC Doubly Linked Buffer – Locked/Lockless



MPMC Singly Linked Buffer – Locked/Lockless



Hash Table

- Closed Addressing
 - Collisions are added onto a linked list
- Functionality
 - Contains
 - Add
 - Remove
 - Resize
- Variations
 - MPMC Globally Locked Hash Table
 - MPMC Lock Per List
 - MPMC Lockless

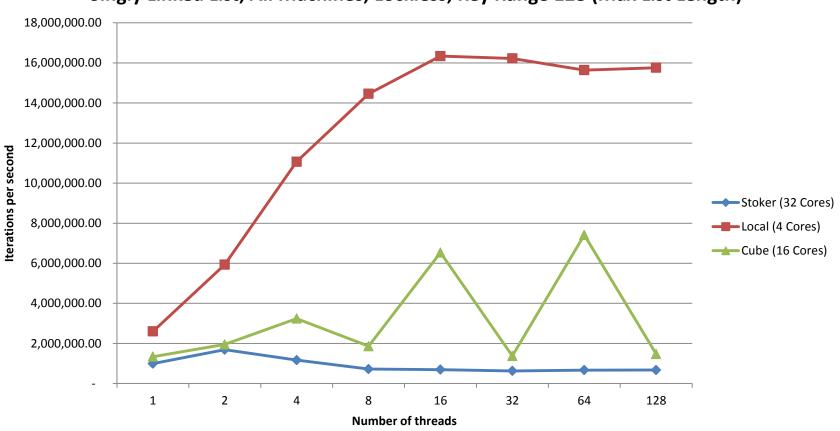
System Details

- My Local Machine (Sandy Bridge 32nm, 4 Cores @ 3.30GHz)
- Stoker (Ivy Bridge EX 22nm, 32 Cores @ 2.00GHz)
- Cube (Gainestown 45nm, 16 Cores @ 2.27 GHz)

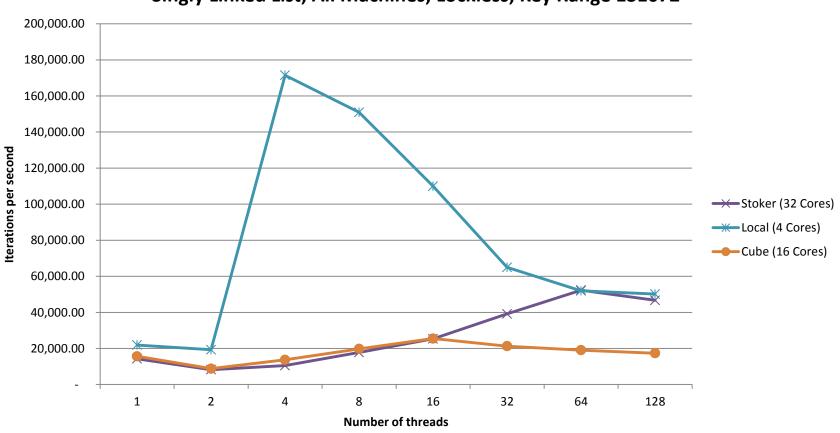
Evaluation

- Compared Data Structure Variations
 - Thread Count 1-128
 - Varied list, buffer and table sizes
 - Locked vs Lockless
- Used Hardware Performance Counters
 - Special registers
 - Performance analysis
 - Record cache misses, cpu cycles etc...

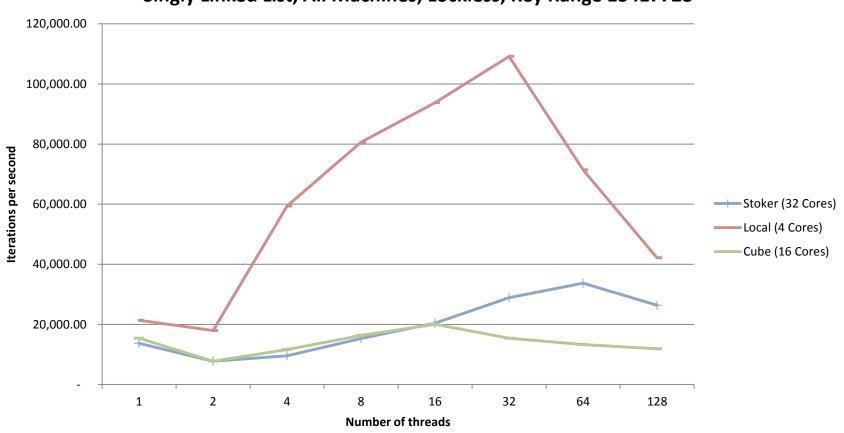
Singly Linked List; All Machines; Lockless; Key Range 128 (Max List Length)



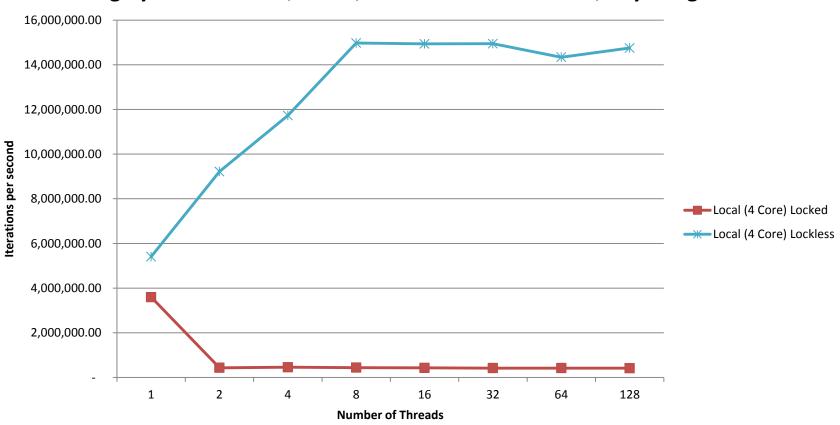
Singly Linked List; All Machines; Lockless; Key Range 131072



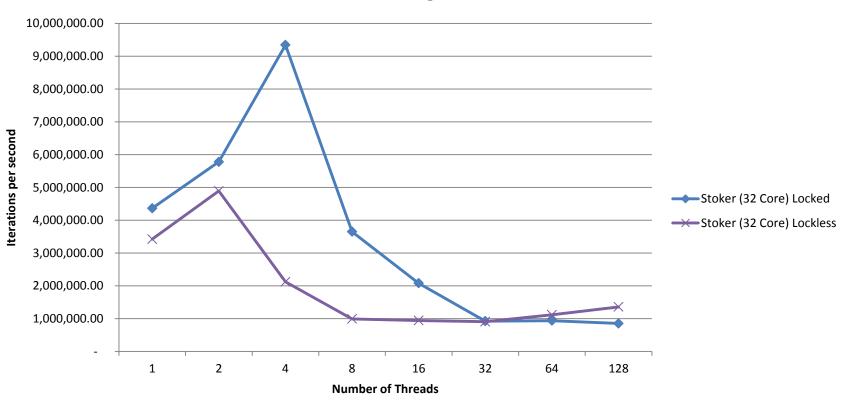
Singly Linked List; All Machines; Lockless; Key Range 13417728



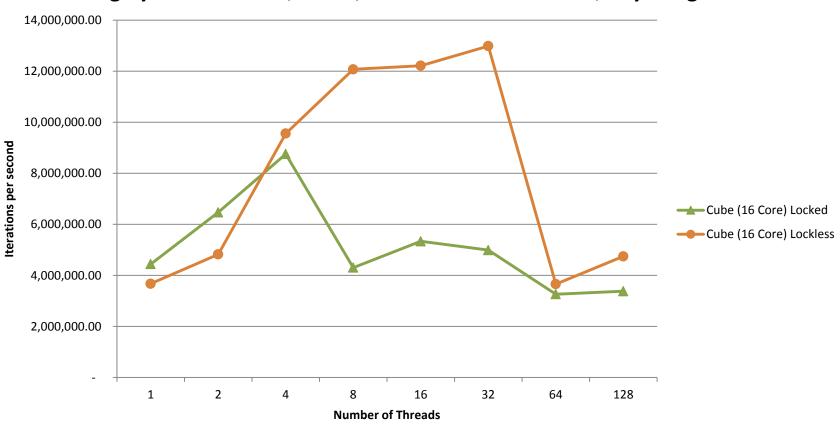
Singely Linked Buffer; Local; Mutex Lock vs Lockless; Key Range 128



Singely Linked Buffer; Stoker; Mutex Lock vs Lockless; Key Range 128



Singely Linked Buffer; Cube; Mutex Lock vs Lockless; Key Range 128



Sample of Hardware Performance Counter
 Data from previous slide (Cube):

	Cube 128	Cube Locked
Stalled Front End Cycles	8,603,339,371	28,157,586,014
Stalled Back End Cycles	4,290,118,281	19,235,301,044
CPU Cycles	16,644,060,219	37,007,949,115

- Lockless algorithms
 - More Difficult to Design & Implement
 - Generally provide a performance boost when compared to locked algorithms
 - There are always exceptions

Ring Buffer

- Locked variations proved to be mostly slower than lockless varieties
- Exceptions came in the form of the TAS and TTAS locks running on the Local Machine
- Stoker proved to be the fastest of the three machines with the lockless algorithm

Linked List

- Local Machine performed well locklessly, outperforming the other two machines
- Local Machine outperformed every lock on all 3 variations with the lockless variation
- Gained the largest performance boost by using a lockless algorithm out of the three data structures

Hash Table

- Globally Locked and Lock per Bucket performed relatively equally
- TestAndSet lock seems to be well suited to my implementation with consistently good performance across variations
- Surprisingly, both locked variations outperformed the lockless variant by a sizeable margin