

# Logic

## For any P, Q:

$$(P \rightarrow Q) \vee (Q \rightarrow P) \quad (\text{always True})$$

## An arg is valid iff:

If premises are True, then conclusion is True

So a counter-example is when premises are True but conclusion is False

## Ideompotence

$$P \equiv (P \vee P)$$

$$P \equiv (P \wedge P)$$

## Commutativity

$$(P \vee Q) \equiv (Q \vee P)$$

$$(P \wedge Q) \equiv (Q \wedge P)$$

## Associativity

$$(P \vee Q) \vee R \equiv P \vee (Q \vee R)$$

$$(P \wedge Q) \wedge R \equiv P \wedge (Q \wedge R)$$

## DeMorgan

$$\neg(P \vee Q) \equiv (\neg P) \wedge (\neg Q)$$

$$\neg(P \wedge Q) \equiv (\neg P) \vee (\neg Q)$$

## Distributivity

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)$$

$$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)$$

## With True / False

$$(P \vee \text{True}) \equiv \text{True}$$

$$(P \wedge \text{False}) \equiv \text{False}$$

$$(P \vee \text{False}) \equiv P$$

$$(P \wedge \text{True}) \equiv P$$

$$(P \vee \neg P) \equiv \text{True}$$

$$(P \wedge \neg P) \equiv \text{False}$$

## Double Negation

$$P \equiv \neg(\neg P)$$

## Implication

$$(P \rightarrow Q) \equiv (\neg P \vee Q) \equiv \neg(P \wedge \neg Q)$$

*DM*

$$(P \rightarrow Q) \equiv (\neg Q \rightarrow \neg P)$$

*Contrapos.*

$$\neg(P \rightarrow Q) \equiv (P \wedge \neg Q)$$

## Equivalence

$$(P \leftrightarrow Q) \equiv (P \rightarrow Q) \wedge (Q \rightarrow P)$$

$$(P \leftrightarrow Q) \equiv (P \wedge Q) \vee [(\neg P) \wedge (\neg Q)]$$

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$$\neg(P \leftrightarrow Q) \equiv (P \wedge \neg Q) \vee (\neg P \wedge Q)$$

## Exportation

$$(P \wedge Q) \rightarrow R \equiv P \rightarrow (Q \rightarrow R)$$

## Absurdity

$$(P \rightarrow Q) \wedge (P \rightarrow \neg Q) \equiv \neg P$$

# Modal Logic

$$\forall x \phi \equiv \neg [\exists x \neg \phi]$$

$$\exists x \phi \equiv \neg [\forall x \neg \phi]$$

## DeMorgan

$$\neg \forall x \phi \equiv \exists x \neg \phi$$

$$\forall x \neg \phi \equiv \neg \exists x \phi$$

## If x is the set of {a,b,c}, then

$$\forall x P(x) \equiv P(a) \wedge P(b) \wedge P(c)$$

$$\exists x P(x) \equiv P(a) \vee P(b) \vee P(c)$$

## Applying DeMorgan:

$$P(a) \wedge P(b) \wedge P(c) \equiv \neg(\neg P(a) \vee \neg P(b) \vee \neg P(c))$$

$$P(a) \vee P(b) \vee P(c) \equiv \neg(\neg P(a) \wedge \neg P(b) \wedge \neg P(c))$$