

Delimited overloading

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Outline

- 1 Standard Overloadings
- 2 Defining overloadings
- 3 Priority & associativity
- 4 Macros
- 5 Some technical details

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Basic use

```
Float.(1 + x * f 4)  
Float.(4 * u**2 / sqrt(abs alpha))  
Hashtbl.(h.(key) <- x)
```

- ☞ Literals, functions, and “constant constructions” substitution.
- ☞ Better readability.

```
Big_int.(if x > 0 then x else 0)
```

- ☞ Can use the usual comparison operators.

```
Int32.(4 + a.(Int.(1 + x)))
```

- ☞ Embedded overloadings.

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Static checks & simple optimizations

```
Num. ("12345678" + x)
```

☞ Compile time check.

If one writes `Num. ("a12")`, when compiling, the following error is issued

```
Parse error: The string "a12" does not represent a  
valid Num.
```

```
Preprocessing error on file foo.ml
```

```
Float. ((x+1)**2)
```

☞ Simple optimization. The whole expression is substituted by (binding introduced only if needed):

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let tmp = x +. 1.0 in tmp *. tmp
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Complex numbers

```
Complex.(let z = 3 + 2 I in sin(z * z))
```

- ☞ “I” notation.
- ☞ Let binding are allowed.
- ☞ Complex functions like `sin`, `cos`,... are inlined.

For example,

```
Complex.((2 + 3 I) * f x)
```

is turned into

```
let tmp = f x in
{ Complex.re = (tmp.Complex.re *. 2.0) -.
                  (tmp.Complex.im *. 3.0);
  Complex.im = (tmp.Complex.re *. 3.0) +.
                  (tmp.Complex.im *. 2.0); }
```

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Summary

`pa_do.cmo` provides overloadings for

<code>Int</code>	<code>Float</code>	<code>Hashtbl</code>
<code>Int32</code>	<code>Complex</code>	<code>String</code>
<code>Int64</code>		
<code>Nativeint</code>		

`pa_do_nums.cmo` provides overloadings for

`Num`
`Ratio`
`Big_int`

⚠ Requires `nums.cmo` to be loaded by `camlp4` for static checks.

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Concrete syntax v.s. API

Concrete syntax	API
In the source file	In a separate file
Must be repeated in each file	Can be bundled with a library
No possibility of overloading general expressions	One can overload general expressions and perform some optimizations



Demo

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Demo

“Set of overloadings”

Constructions that can be overloaded:

literals

`'x`

`a.(i)`

`a.(i) <- x`

`a.[i]`

`a.[i] <- x`

`a.{i}`

`a.{i} <- x`

`a := x`

`a <- x`

`a.p`

functions & operators overloadings

general substitutions

How to overload one's own module (1/2)

```
module Foo :  
sig  
  type t  
  val of_int : int -> t  
  val compare : t -> t -> int  
  val add : t -> t -> t  
  val mul : t -> t -> t  
end
```

Overloading with the concrete syntax:

```
let int_of_int : int -> Foo.t = Foo.of_int  
let compare : 'a -> 'a -> int = Foo.compare  
let add : 'a -> 'a -> 'a = Foo.add  
let mul : 'a -> 'a -> 'a = Foo.mul
```

How to overload one's own module (1/2)

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```

Overloading with the concrete syntax:

- int literals: `OVERLOAD_INT Foo (of_int)`
- comparison: `OVERLOAD_COMPARISON Foo (compare)`
- functions:
`OVERLOAD Foo ((+) -> add; (*) -> mul)`

How to overload one's own module (1/2)

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```
OVERLOAD Foo ( ( + ) -> add; ( * ) -> mul )
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Overloading with the concrete syntax:

- int literals: `OVERLOAD_INT Foo (of_int)`
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`OVERLOAD Foo ((+) -> add; (*) -> mul)`

How to overload one's own module (2/2)

Remarks

- If Foo implements all standard functions add, sub, mul, div and neg (unary negation):
 - `OVERLOAD_ARITHMETIC Foo.`
- If a new module is implemented:

```
module Special_foo
  use Foo
  include Foo
  overload add: t -> t -> t
  and
```

```
➤ OVERLOAD Special_foo inherit Foo
➤ OVERLOAD Special_foo ( ( - ) -> sub )
```

How to overload one's own module (2/2)

Remarks

- If Foo implements all standard functions add, sub, mul, div and neg (unary negation):
 - ☞ `OVERLOAD_ARITHMETIC Foo.`
- If a new module is implemented:

```
module Special_foo :  
sig  
  include Foo  
  val sub : t -> t -> t  
end
```

- ☞ `OVERLOAD Special_foo inherit Foo`
- ☞ `OVERLOAD Special_foo ((-) -> sub)`

Some more examples

Basin of attraction of Newton's method for $z^3 = 1$.

```
Complex.(
  let z = ref z0 in
  for i = Int.(1) to niter do
    z := (2 * !z + 1 / !z**2) / 3
  done;
  if abs(!z - root0) <= r then Some color0
  ... )
```

Let $D = 1.7$. If $p = [p_0; \dots; p_n]$ represents the polynomial $\sum_{i=0}^n p_i z^i$, its norm is (here) defined by $\|p\| := \sum_{i=0}^n |p_i| D^i$

```
let domain = Interval.(17 / 10)
let norm p = Interval.(List.fold_right (fun c n ->
  abs c + domain * n) p 0)
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Priority and associativity of operators

pa_infix.cmo

Concrete syntax:

```
INFIX ( %+ ) RIGHTA HIGHER (+)
INFIX ( ^* ) LEVEL (+)
PREFIX ( /+ / )
POSTFIX ( /// ) LEVEL ( ! )
```

API: treat $a = b \mid > c$ as $a = (b \mid > c)$ and replace $x \mid > f$ by $f\ x$:

```
open Pa_infix
module L = Level
let l = L.binary (L.Higher L.comparison) ~assoc:L.LeftA in
let expr x y _loc = <:expr< $$ $x$ >> in
infix ">" ~expr l
```

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Collaboration with Delimited Overloading

```
DEFINE NEWTON(M,x) = M.( (2 * x + 1 / x**2) / 3 )
```

Use it as

```
NEWTON(Float, r)  
NEWTON(Complex, z)
```

- ☞ Poor man defunctorizer;
- ☞ Contrarily to functors, requalifies constants.

Better error reporting

```
DEFINE A(x) =  
  let s = ref 0.0 in  
  for i = 1 to x do  
    s := !s + float i  
  done;  
  s  
  
let () = print_float(A(100))
```

With the standard macros:

File "...", line 8, characters 21-27:

This expression has type float but is here used with type
int

Better error reporting

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  s
```

```
let () = print_float(A(100))
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With Delimited Overloading macros:

File "...", line 8, characters 21-27:

Expanding of the macro "A" at the previous location yields the error:

File "...", line 4, characters 11-13:

This expression has type float but is here used with type int

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 - Optimization for complex operators
 - Embedded overloading
 - General substitution of expressions

Optimization for complex operators

```
Complex.((2 + 3 I) * f x)
```

- 1 Classify subexpressions according to

```
type t =  
  | Zero  
  | Re of Ast.expr  
  | Im of Ast.expr  
  | Cplx of Ast.expr * Ast.expr  
  | Unknown of Ast.expr
```

- 2 Specialize complex functions, introducing bindings as needed.

Embedded overloading (1/3)

$$X.(e) \mapsto \sigma_X(e)$$

where $\sigma_X : \text{expression} \rightarrow \text{expression}$

Problem encountered:

`Int32.(a.(Int.(0)) <- 7 + x)`

↓ apply σ_{Int} ; here $\sigma_{\text{Int}}(0) = 0$

`Int32.(a.(0) <- 7 + x)`

↓ apply σ_{Int32}

`a.(01) <- Int32.add 71 x`

Problem!

☞ Protection of already overloaded expressions

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Embedded overloading (2/3)

Requirements:

- $\sigma_{\text{Int}}(0)$ must be a valid expression;
- it must not change the meaning of the program nor its performance;
- locations must not be affected (for correct error reporting).

Solution:

$$X.(e) \mapsto p(\sigma_X(e))$$

where p is an undeclared function name!

⇨ p is removed by the surrounding overloading \Rightarrow global flag to know whether to insert π .

⇨ external $p : \alpha \rightarrow \alpha = \%identity$ forbid some optimizations to take place!

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Embedded overloading (3/3)

`Int32.(a.(Int.(0)) <- 7 + x)`

$\downarrow \text{Int.}(0) = p(\sigma_{\text{Int}}(0)) = p(0)$

`Int32.(a.(p(0)) <- 7 + x)`

$\downarrow \text{apply } \sigma_{\text{Int32}}$

`p(a.(0) <- Int32.add 71 x)`

OK!

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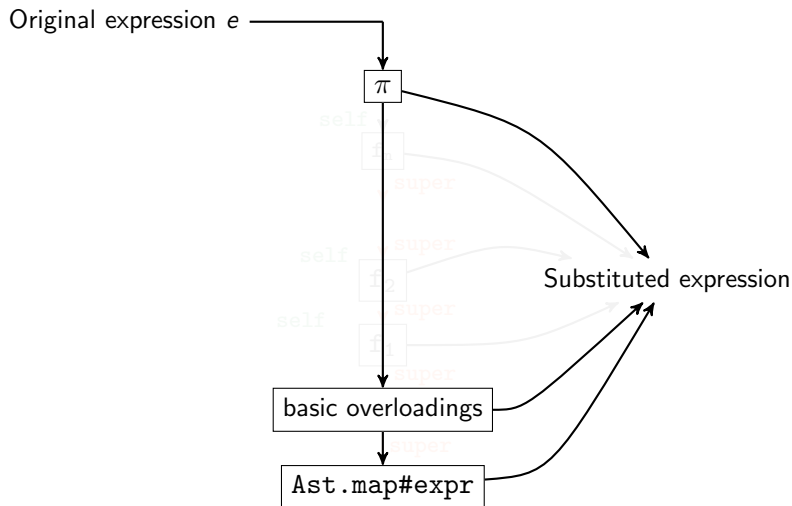
`Int32.(a.(p(0))) <- 7 + x`

\downarrow `apply σ_{Int32}`

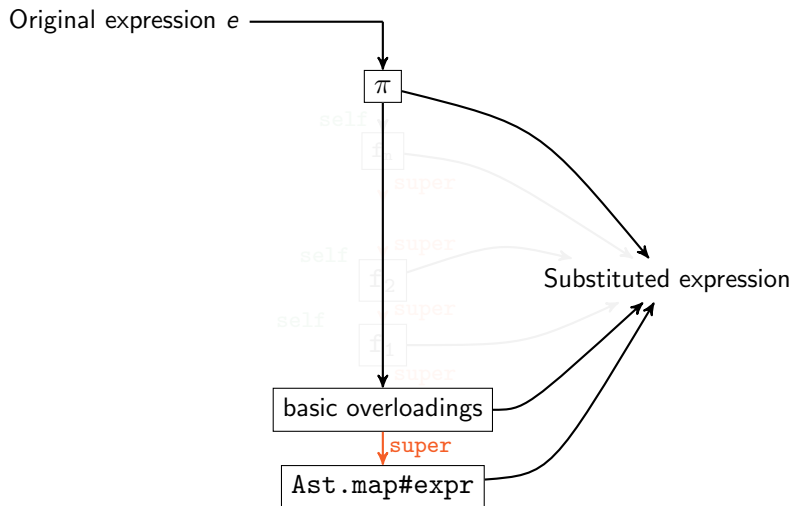
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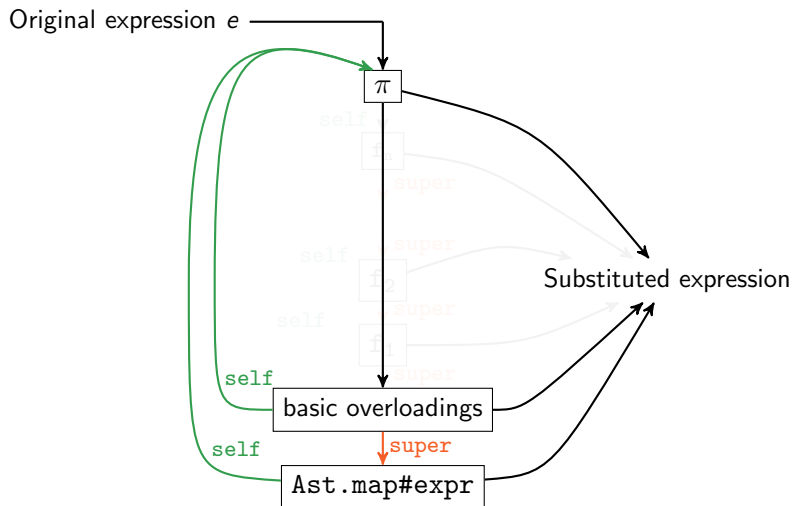
General substitution of expressions X. (e)



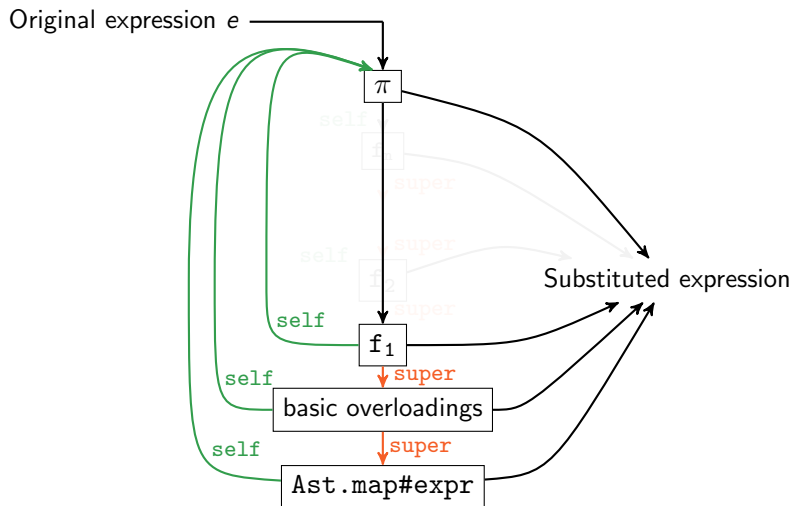
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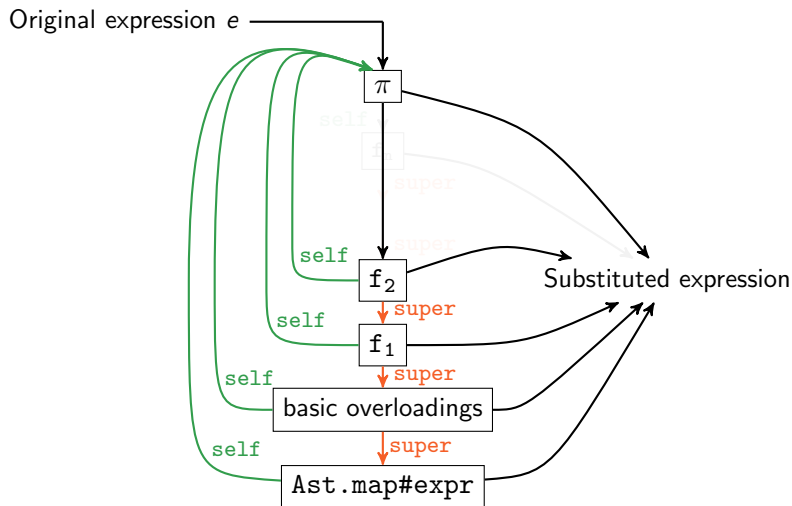
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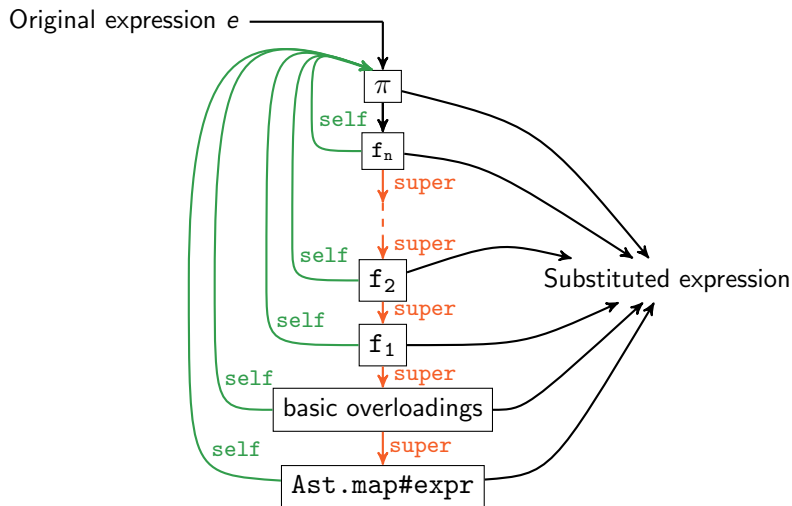
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Thank you for your attention.