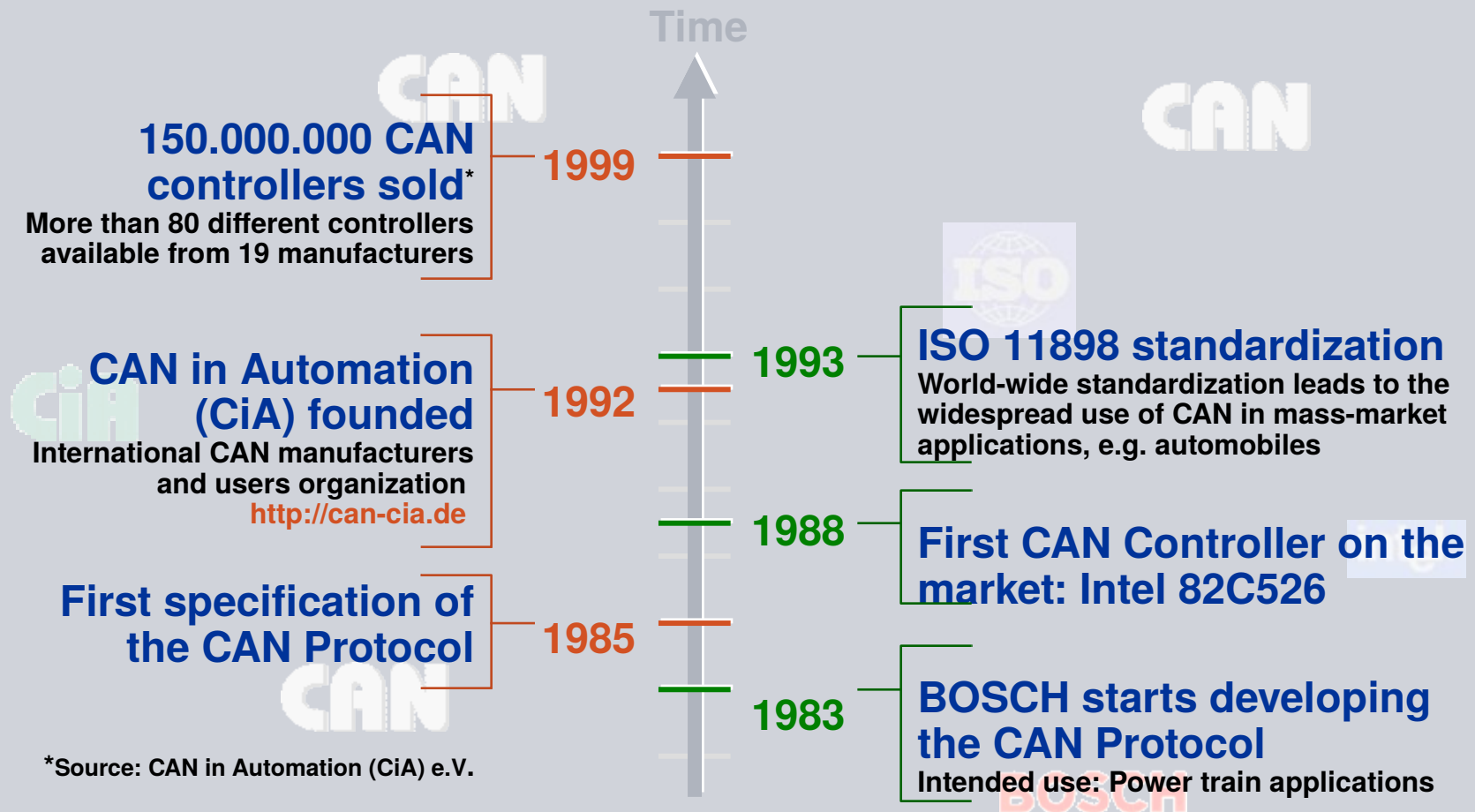


CAN (Control Area Network) Common Training

Marcel Liviu Sodinca – SV I IP TSR PV INT

Nur für
internen
Gebrauch

CAN History Timeline



Applications examples for CAN Bus



Construction Equipment



Domestic Appliances



Ships & Boats



Agricultural Machines

90 %
of all
applications



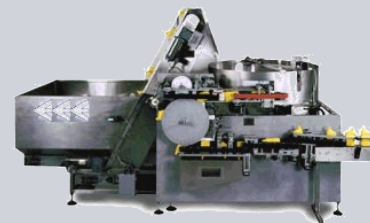
Automotive Sector

many uses
CAN
of the

Building
Automation



Trains

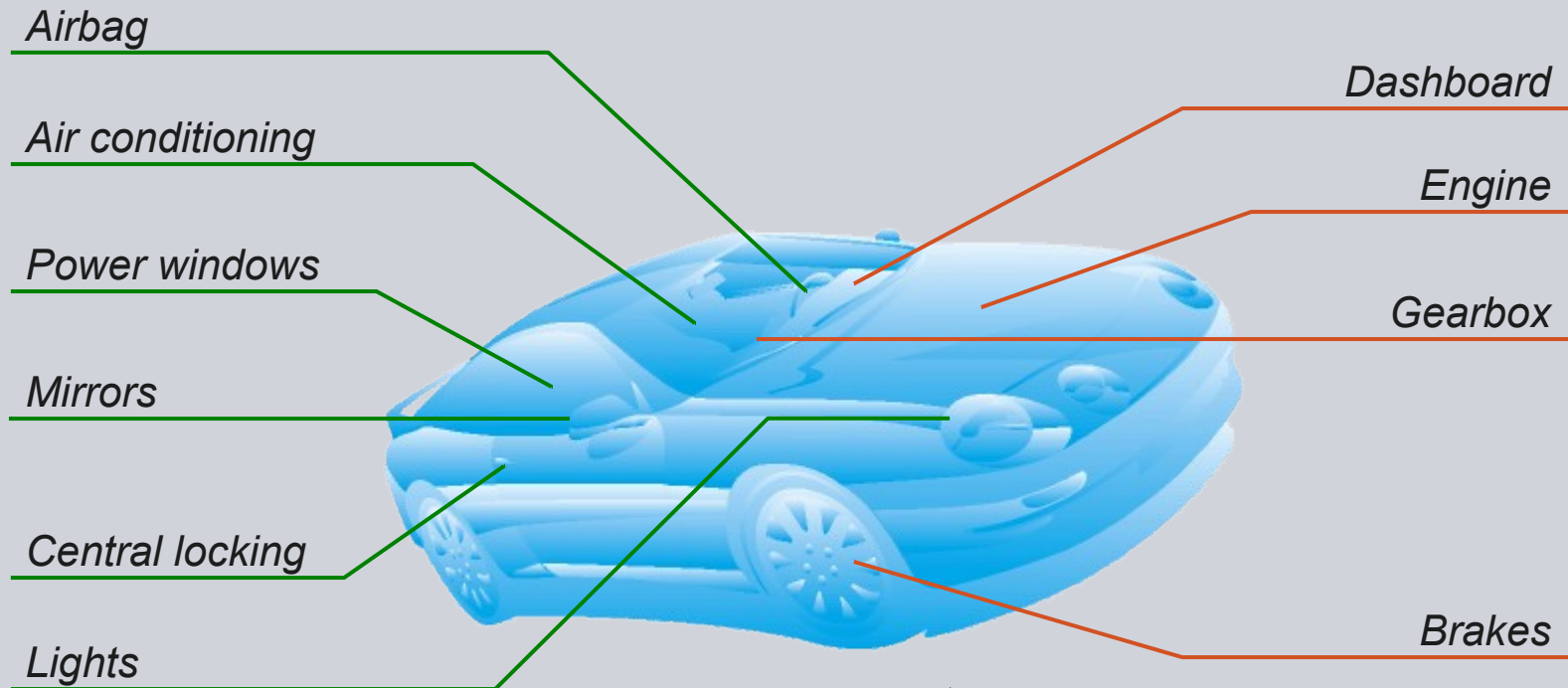


Industry Automation



Medical
Equipment

Typical examples for CAN in automobiles



Car Body components

Typically **low-speed** CAN Bus

Power Train components

Typically **high-speed** CAN Bus

Bus characteristics

➤ **Serial data communications bus**

Inexpensive and simple, but slower than parallel bus.

➤ **“Good” real-time capabilities**

Small latency (“fast enough”)

➔ indispensable for automotive applications.



➤ **Data rate dependent on bus length**

Data rate: 1 Mbit/sec	➔	Bus length: 40 meters,
Data rate: 125 kBit/sec	➔	Bus length: 500 meters,
Data rate: 50 kBit/sec	➔	Bus length: 1000 meters.

Typical definitions:

Low-speed: 25 kBit/sec up to 125 kBit/sec.

High-speed: 500 kBit/sec up to 1 Mbit/sec.

Bus characteristics

➤ Multicast / broadcast philosophy

CAN messages do not include references to sender or receivers, but to information contents.



➤ Bus access principle: CSMA/CA Carrier Sense Multiple Access with Collision Avoidance

- | | |
|------------------------------------|---|
| <i>Carrier Sense:</i> | Every node monitors the bus level, all the time. |
| <i>Multiple Access:</i> | Every node can start a transmission any time when the bus is free. |
| <i>Collision Avoidance:</i> | When several nodes start transmission at the same time, <u>all but one</u> withdraw from sending. |

Hardware characteristics

➤ Hardware message acceptance filtering

Only leaves messages through which are of interest for the node

➔ reduces CPU load.

➤ Sophisticated hardware error management

Combination of different error prevention and detection methods, automatic re-transmission of messages detected as erroneous

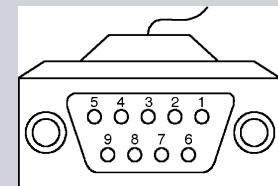
➔ **very high transmission security** among field bus systems.

➤ Various transmission media

Twisted-pair (dual-wire) cable, single-wire cable or optical fiber.

➤ Standardized connector

Recommended: 9-pin D-sub connectors (DIN 41652).



Bus Systems

Types of Serial Data Transmission

Serial data transmissions may be classified in:

- **Synchronous Serial Data Transmission**

Data is synchronous transmitted to a clock signal
(e.g. data shift register, I²C bus, etc.)



- **Asynchronous Serial Data Transmission**

Transmitter and Receiver are synchronized via a start bit (e.g. RS232, CAN, etc.)

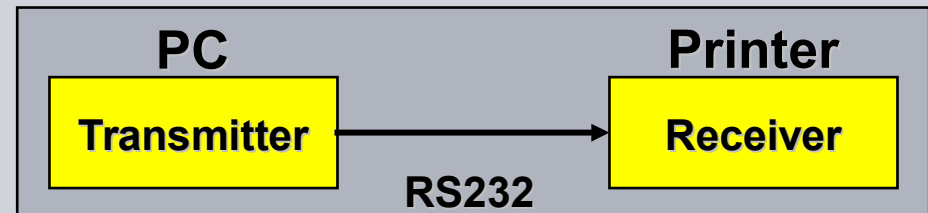


Bus Systems

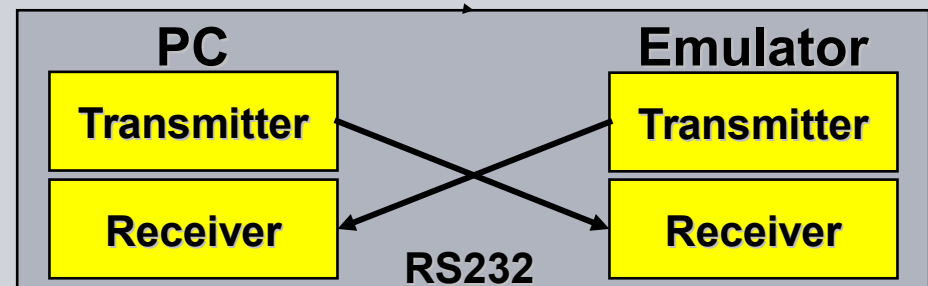
Types of Serial Data Transmission :

Serial data transmissions may also be classified in:

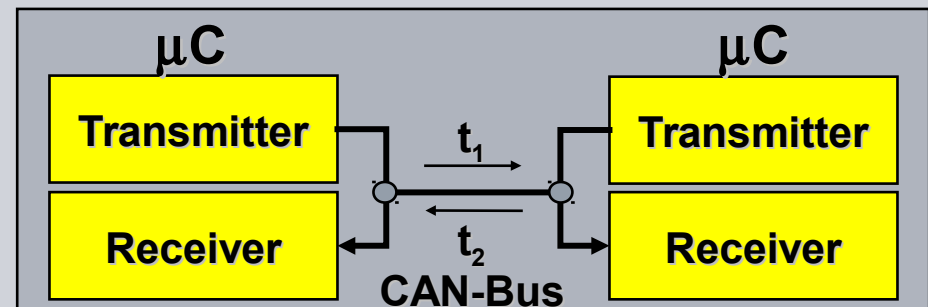
⇒ unidirectional



⇒ bidirectional full duplex



⇒ bidirectional half duplex



Serial Data Transmission with and without Protocol:

- **Data Transmission without Protocol**

=> only data frames:

- ⇒ **RS232** e.g. UART/USART plus Transceiver
- ⇒ **RS485** e.g. UART/USART plus Transceiver

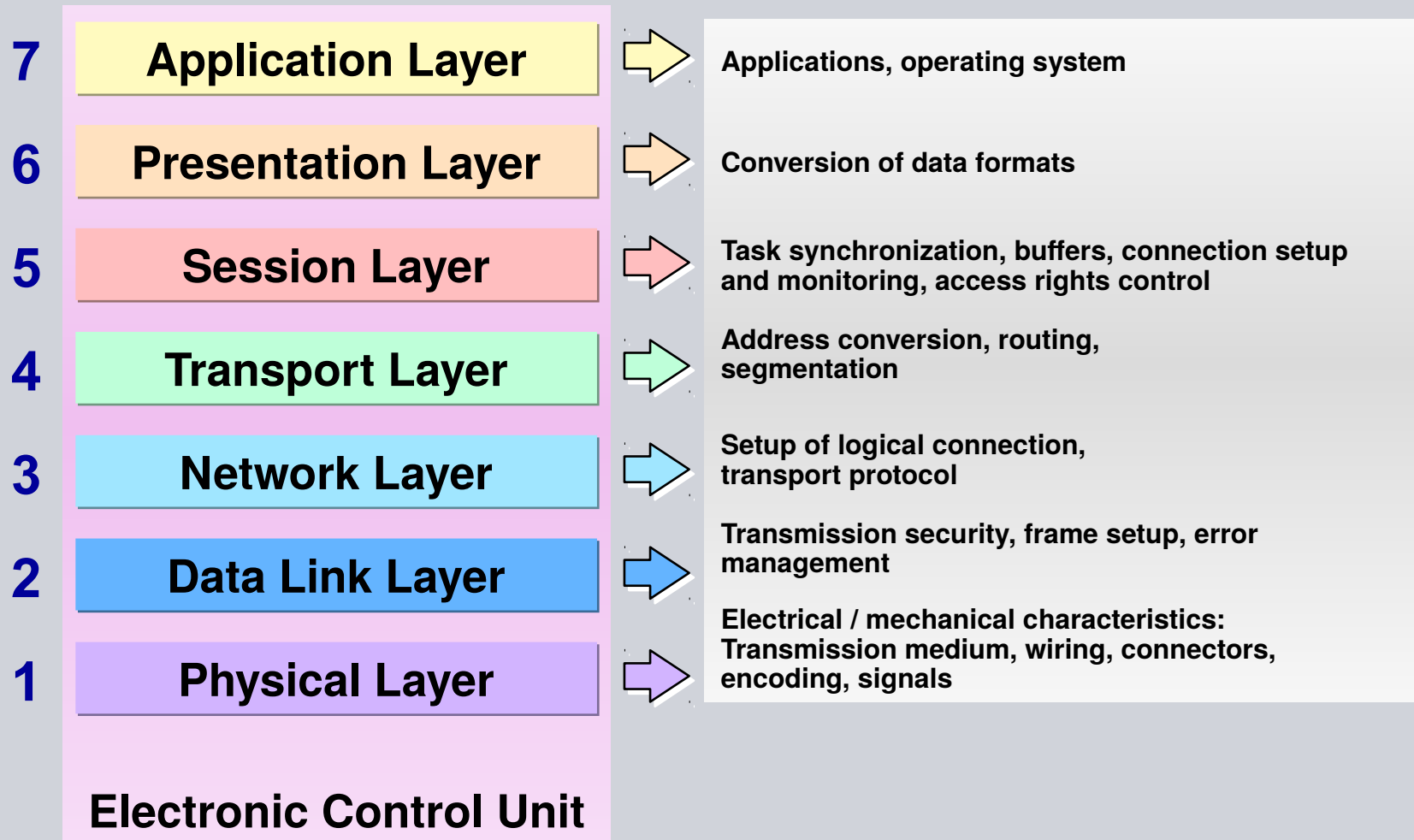
- **Data Transmission with Protocol**

=> **Address + Data + Data-check + Signaling-information:**

- ⇒ **I²C** e.g. I²C-Controller plus Transceiver
- ⇒ **CAN** e.g. CAN-Controller plus Transceiver

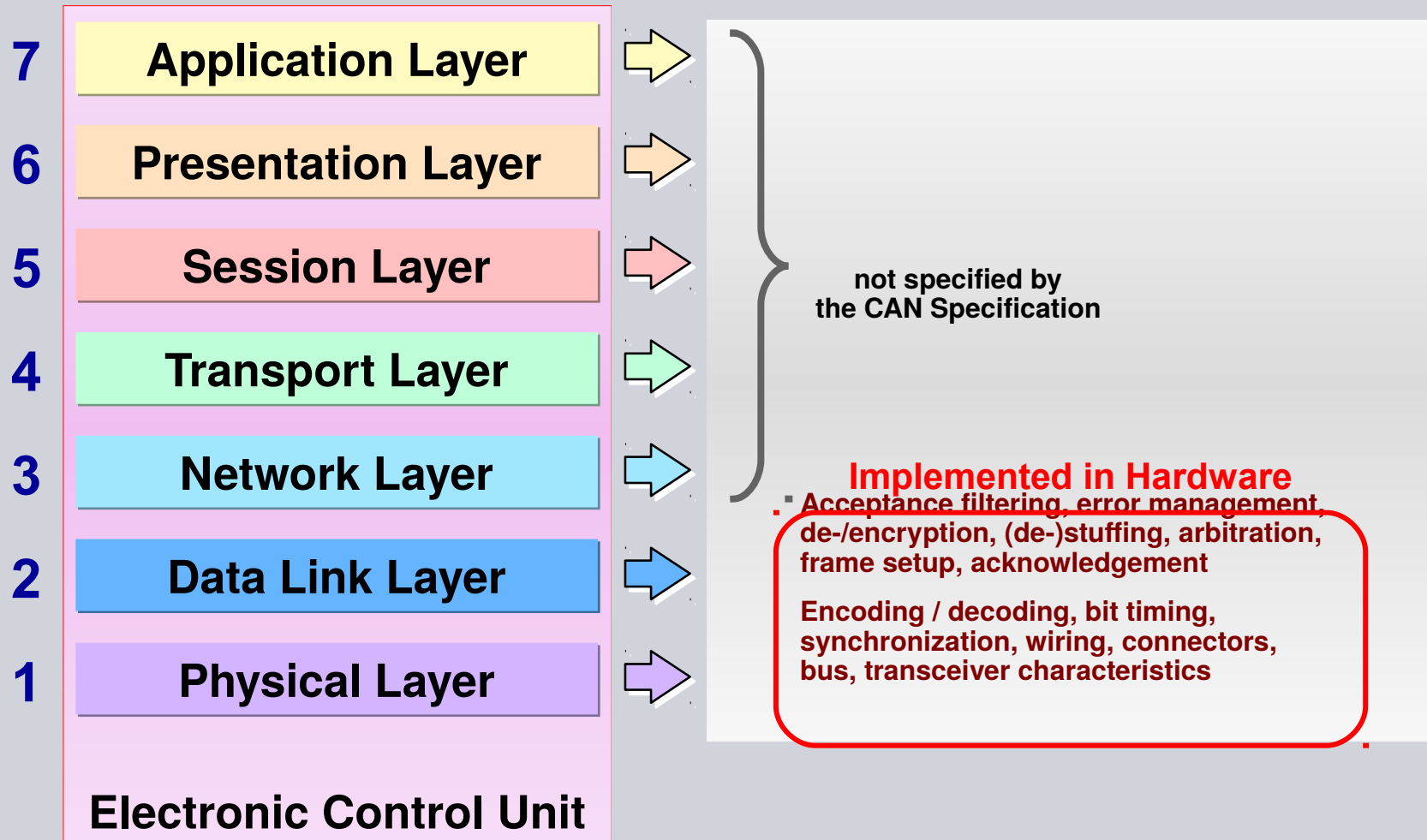
ISO/OSI Layer structure of CAN node

(International Standard Organization's Open System Interconnect)




ISO/OSI Layer structure of CAN node

(International Standard Organization's Open System Interconnect)



ISO/OSI Layer structure of CAN node

(International Standard Organization's Open System Interconnect)

7	Application Layer: Interface between the data communication environment and the application.		
6 5 4 3	Presentation Layer: Session Layer: Transport Layer: Network Layer:	 Not explicitly specified in CAN	
2	Data Link Layer:	Logical Link Control (Object Handling) Medium Access Control (Transfer Handling)	Acceptance Filtering Overload Notification Error Recovery Management Data Encapsulation / Decapsulation Message Framing incl. Stuffing / Destuffing Medium Access Control incl. Arbitration Error Detection / Signalling / Confinement Acknowledgement Handling
1	Physical Layer:	Physical Signalling Physical Medium Attachment Transmission Medium	Bit Encoding / Decoding Bit Timing Frame / Bit Synchronization Driver / Receiver Characteristics Cable Connectors

How to remember the seven layers

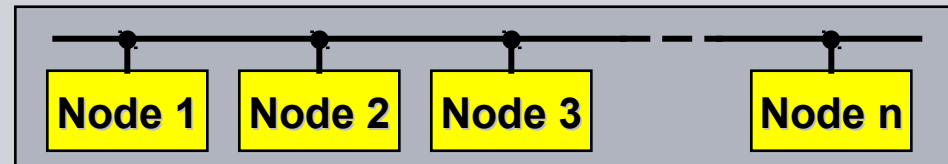
7	Application Layer	All
6	Presentation Layer	People
5	Session Layer	Seem
4	Transport Layer	To
3	Network Layer	Need
2	Data Link Layer	Data
1	Physical Layer	Processing

Bus structure

The most used bus structures:

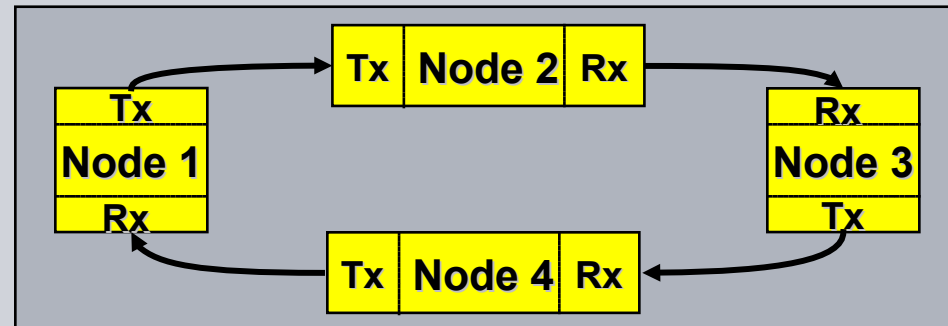
⇒ linear bus

e.g. Ethernet, PROFIBUS



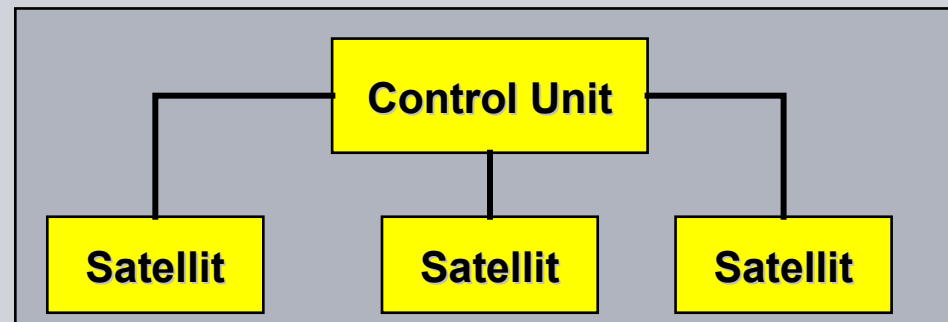
⇒ ring bus

e.g. InterBus-S



⇒ bus star

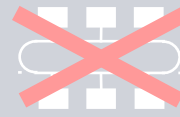
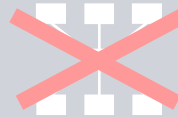
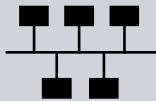
e.g. USB



CAN Bus structure

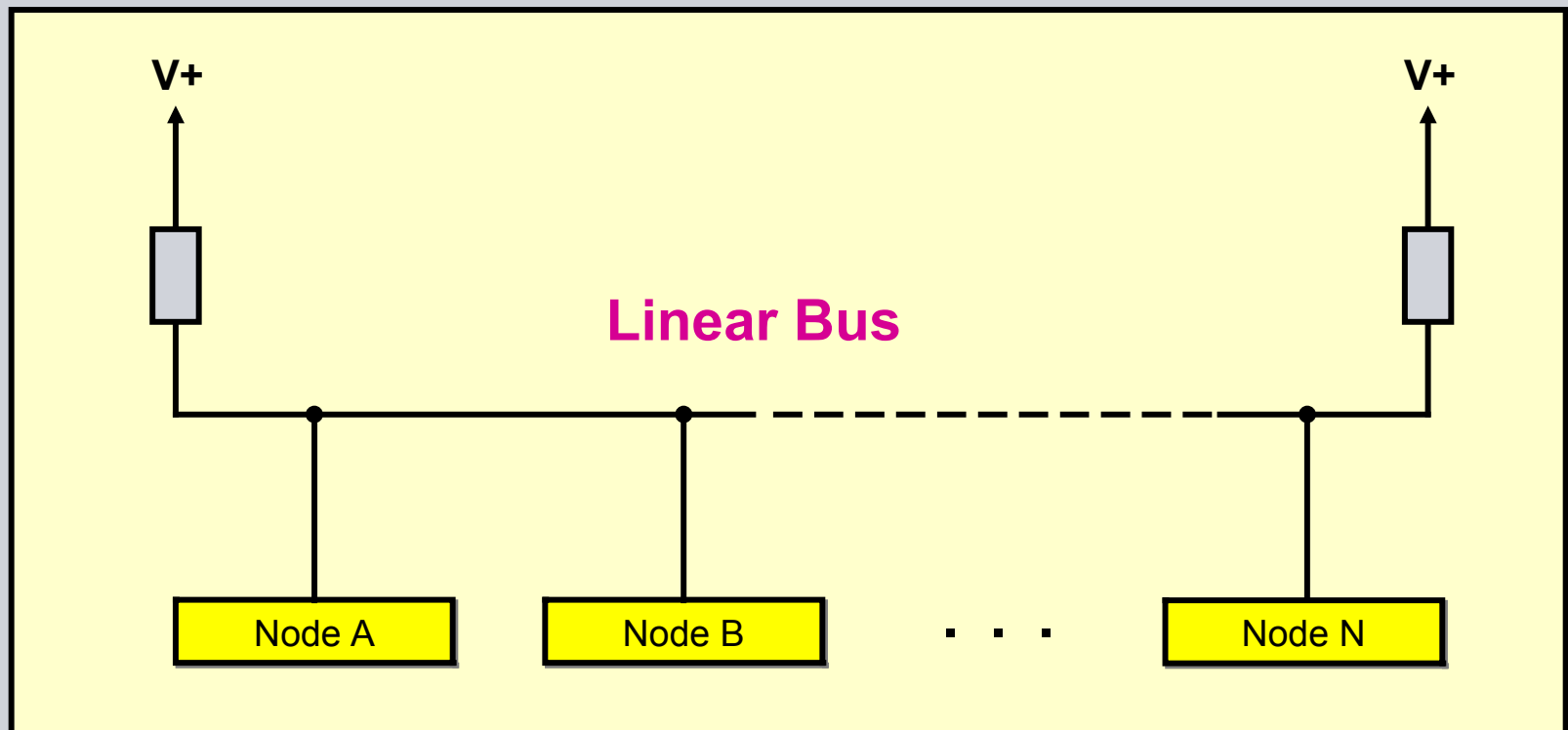
➤ Linear bus structure

No star structure, **no** ring structure, **no** tree structure !



How the CAN networks works

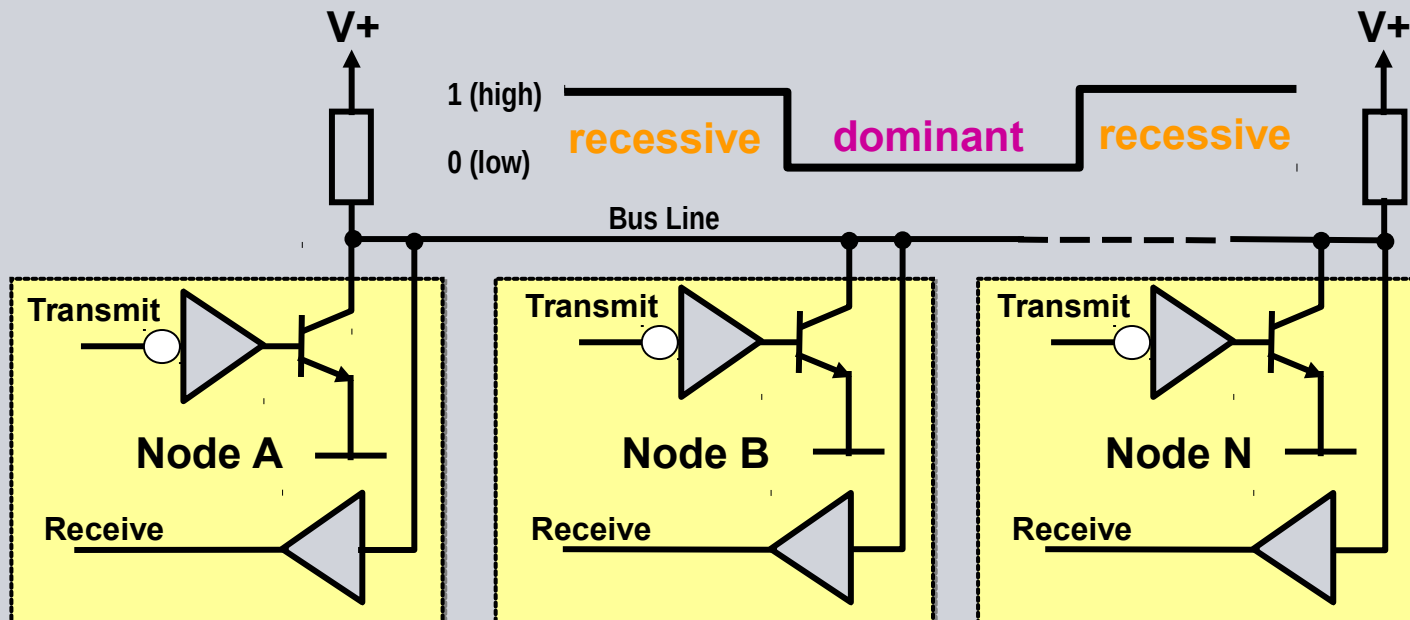
- CAN uses serial half duplex data communication with a linear bus structure



How the CAN networks works

Physical Transmission

- ⇒ Transmission with Carrier Sense Multiple Access with Collision Avoidance **CSMA/CA**
- ⇒ **dominante** Bits (low) win against **recessive** Bits (high) on the bus

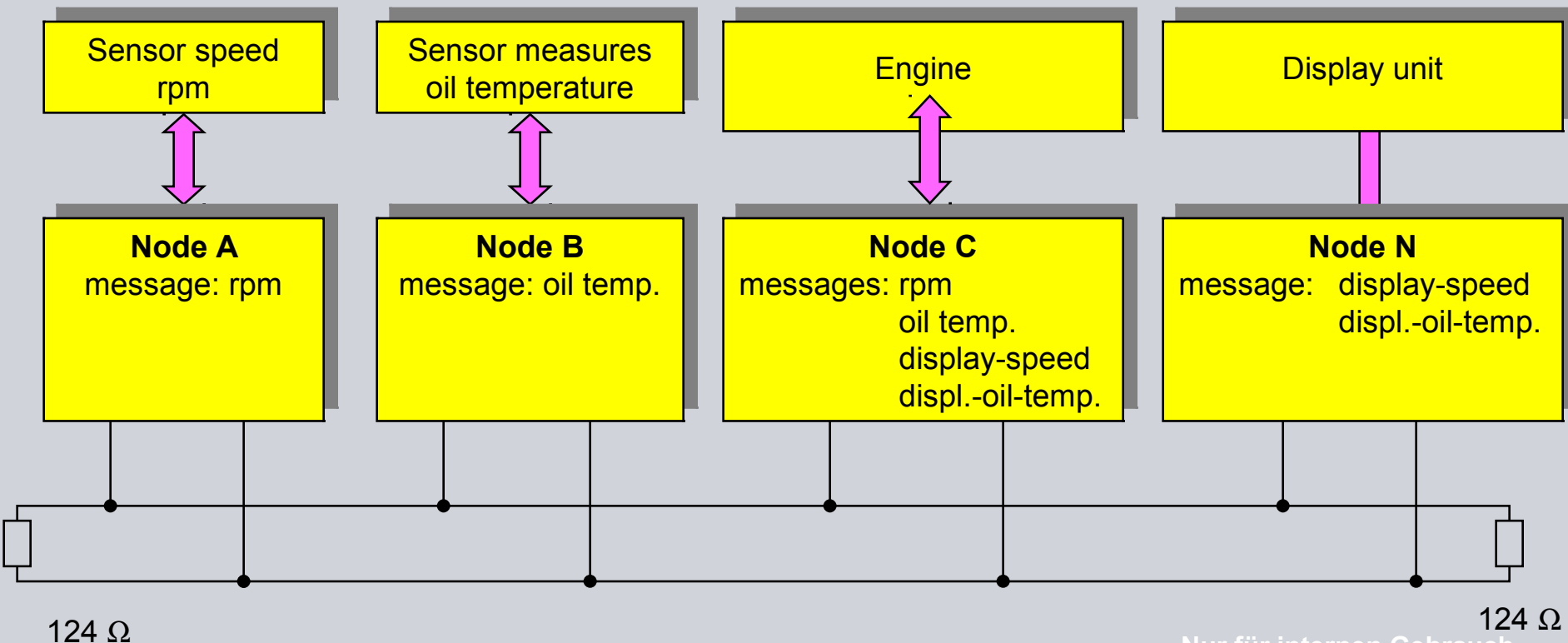


Nur für internen Gebrauch

How the CAN networks works

→ CAN doesn't address stations with physical addresses, but instead of this the message gets an identifier

An identifier is used as in form of symbol on the programming level.



Nur für internen Gebrauch

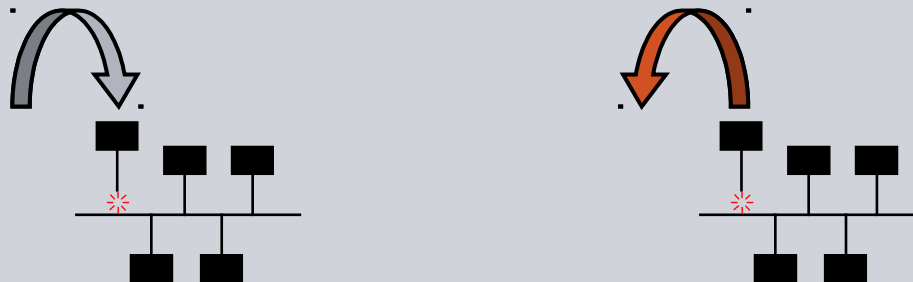
Bus stations (Nodes)

➤ Typically 3 to 40 nodes per bus

No limit for node number defined in CAN specification.
Number of nodes depends on capabilities of CAN transceivers.
Bus load usually gets higher with more nodes.

➤ Hot plug-in / plug-out

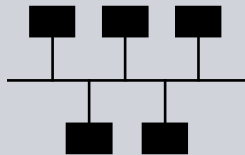
Connect / disconnect nodes while the bus is up and running.



Advantages of the CAN bus

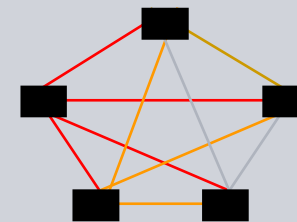
Advantages of





the
CAN Bus



over

conventional
cabling



-  Less wires
-  Less connections
-  Less weight
-  Less EMI problems

-  Less space requirements
-  Easier addition of new nodes
-  Lower assembly costs
-  **Increased transmission reliability**

CAN Physical layer

The Physical Layer is responsible for:

Physical Signalling

Bit Encoding / Decoding

Bit Timing

Frame / Bit Synchronization

Physical Medium Attachment

Driver / Receiver Characteristics

Transmission Medium Cable

Connectors

There specifications concerning both the Physical Signalling and the Physical Medium Attachment:

CAN High Speed Transmission

ISO/DIS 11898

CAN Low Speed Transmission

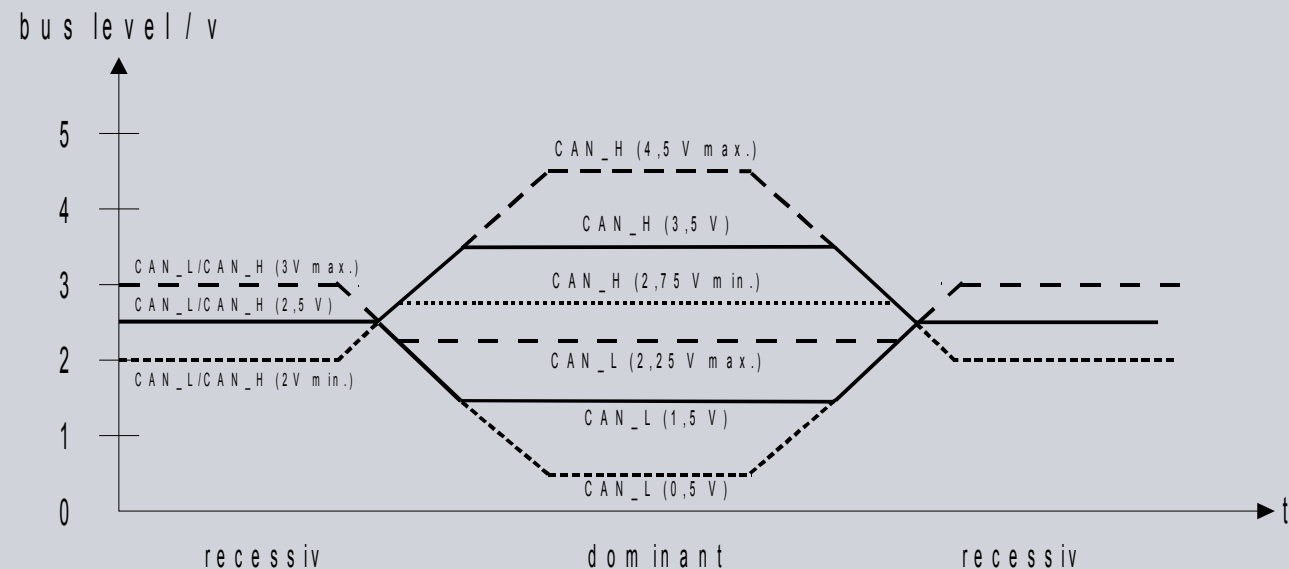
ISO/DIS 11519-2

These Specifications do **not define** the **Transmission Medium** (Cable and Connectors)

CAN Physical layer

CAN - Physical aspects

- **Transmission rate up to 1 MBit/s needs bus driver circuits (transceiver) according to ISO/DIS 11898**
- **CAN High Speed Transmission (ISO/DIS 11898)**
 - ◆ Transmission rate 125 Kbit/s up to 1 MBit/s
 - ◆ max. bus length depend on transmission rate (e.g. 1 MBit/s max. 40 m bus length)
 - ◆ up to 30 nodes

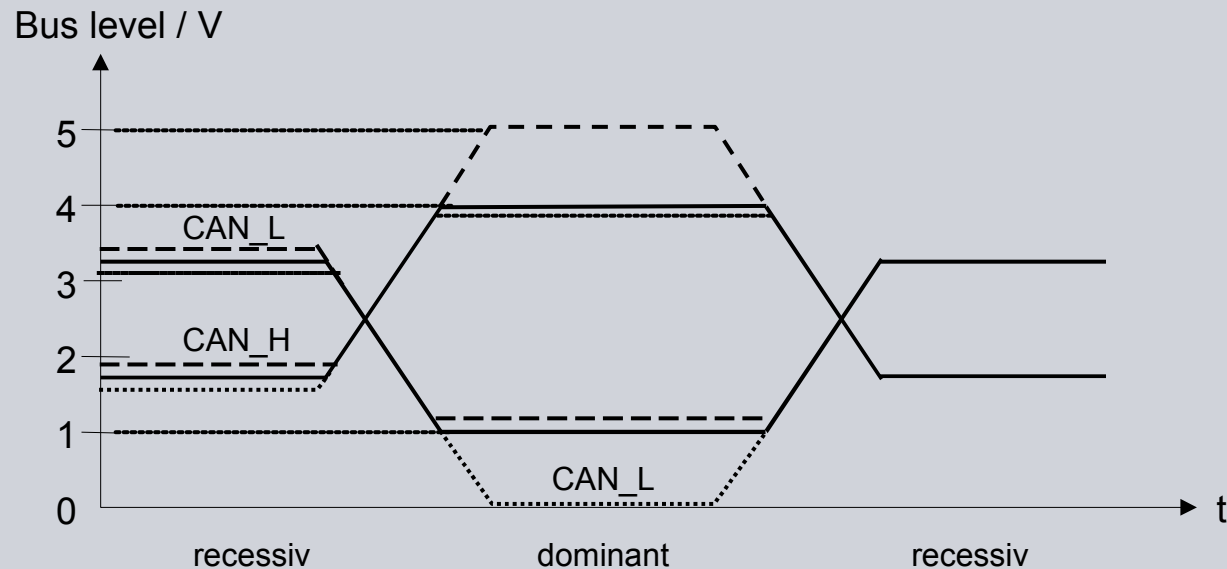


Nur für internen Gebrauch

CAN Physical layer

→ CAN Low Speed Transmission (ISO / DIS 11519-2)

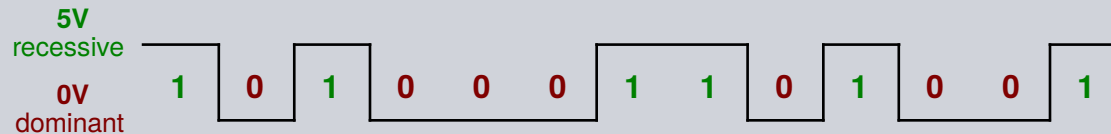
- ◆ Transmission rate 10 KBit/s up to 125 KBit/s
- ◆ max. bus length depend on the distributed capacity of the line
- ◆ up to 20 nodes



Signal coding

➤ NRZ (non-return-to-zero) coding

Example: Voltage levels: **0V (dominant)**, **5V (recessive)**



Characteristics of NRZ coding:

Voltage level stays the same for consecutive bits of same polarity.

Note:

Different voltage levels are defined for different purposes.

CAN Protocol and Frame Types

CAN Protocol and Frame Types: overview

➤ **Frame: “Envelope” for transmission data**

Exact frame format is defined in CAN specification

➤ **Note: CAN Frame \neq CAN Message !!!**

A CAN *message* can be spread out over several CAN *frames*

➤ **Four different frame types:**

Data Frame:	Transmission of regular data
Remote Frame:	Remote request for data transmission
Overload Frame:	Indication of bus overload situations
Error Frame:	Indication of transmission errors

CAN Protocol and Frame Types

CAN Message Format



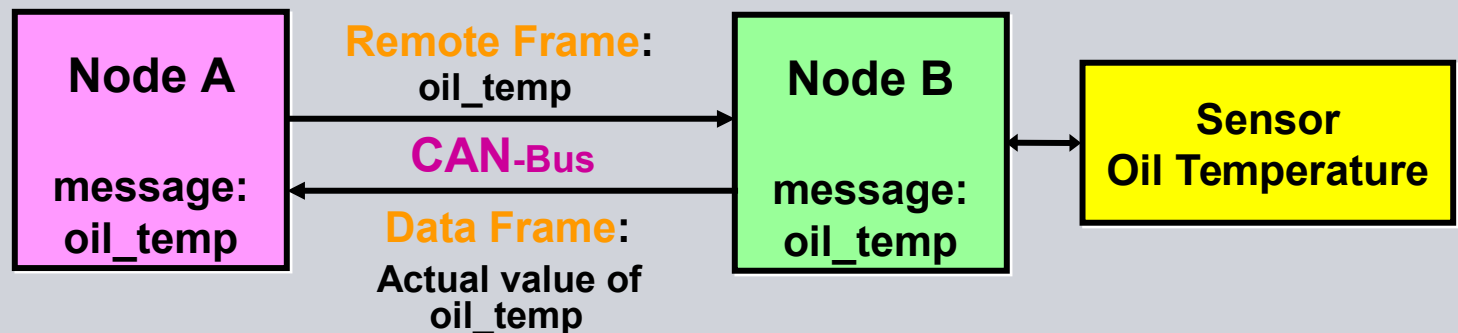
- ⇒ **IFS/IDLE** Idle Level of Data Line (high)
- ⇒ **SOF** Start of Frame (low)
- ⇒ **ID** Standard CAN: 11 Bit Identifier
Extended CAN: 29 Bit Identifier
- ⇒ **Control** Remote, ID-Extension, reserved Bit(s), Data Length Code
- ⇒ **Data** 0 bis 8 Data bytes
- ⇒ **CRC** Cyclic Redundancy Check + CRC Delimiter
- ⇒ **ACK** Acknowledgement - Receiver Response + Delimiter
- ⇒ **EOF** End of Frame (7 bits recessive)
- ⇒ **IFS** Inter Frame Space (3 bits recessive)
- ⇒ **IDLE/SOF** Bus Idle Level or SOF of the next message

CAN Protocol and Frame Types

CAN Transmission Types

- ⇒ **Data Frame** Transmission of up to 8 byte user data with Standard or Extended Frame (Remote-Bit **RTR = 0**)
- ⇒ **Remote Frame** Destination node is able to require data from another node by sending a Remote Frame (available with Standard and Extended Frames) (Remote-Bit **RTR = 1**)
- ⇒ **Error Frame** Each CAN node detecting an transmission error sends an “Error Frame” (Sequence of 6 bit with the same level)

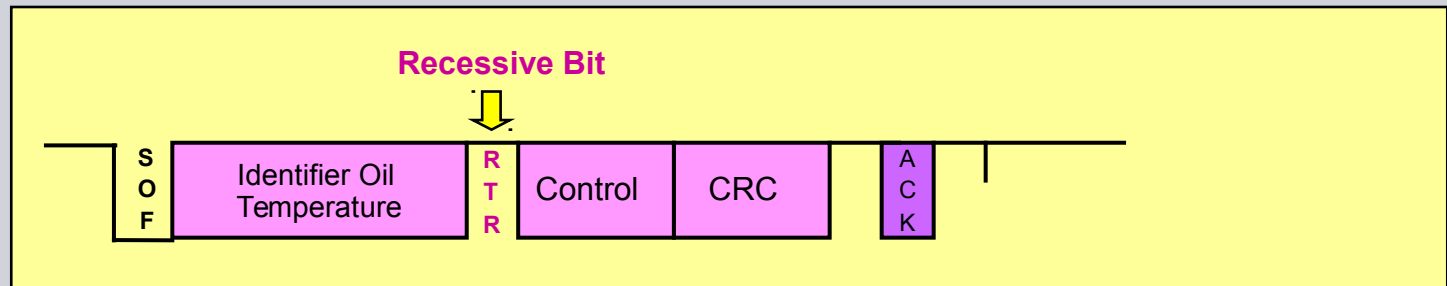
Example:



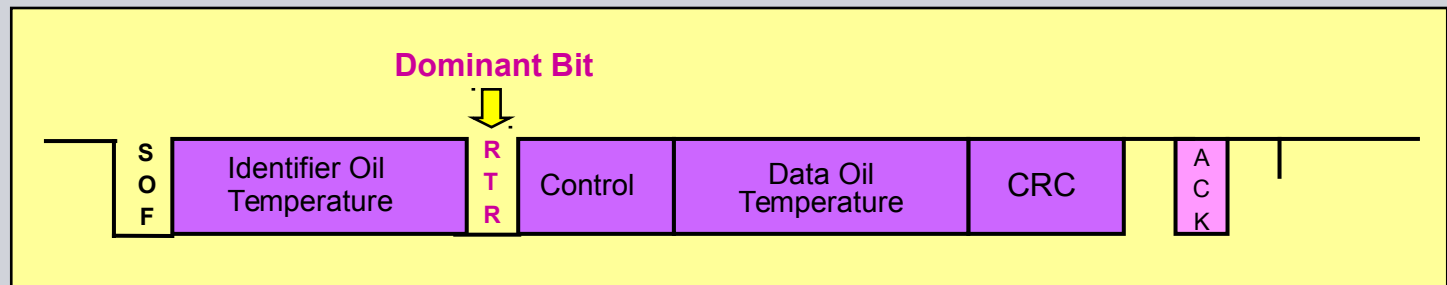
CAN Protocol and Frame Types

CAN Remote and Data Frame

Node A transmits
Remote Frame
(request data)



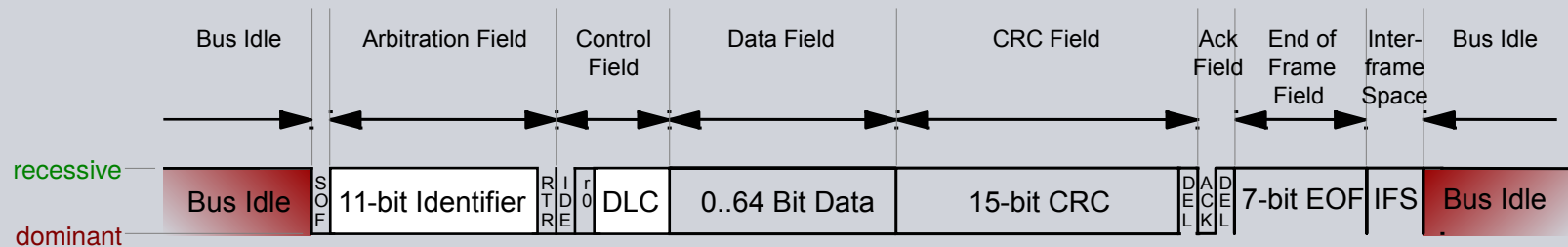
Node B transmits
requested
Data Frame



CAN Protocol and Frame Types

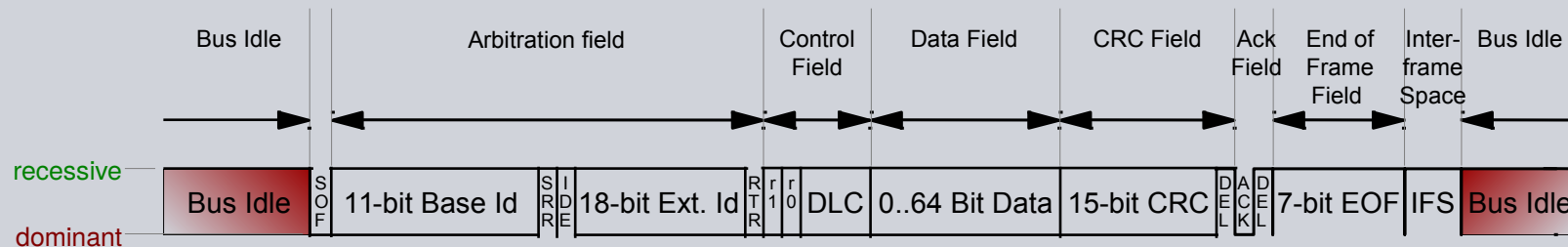
➤ Standard Format (CAN 2.0A): 11-bit Identifier

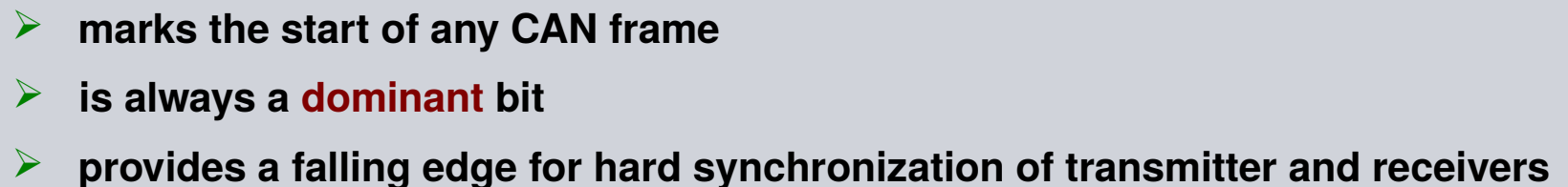
$2^{11} = 2048$ (in reality only 2032) identifiers possible



➤ Extended Format (CAN 2.0B): 29-bit Identifier

$2^{29} = 536.870.912$ identifiers possible

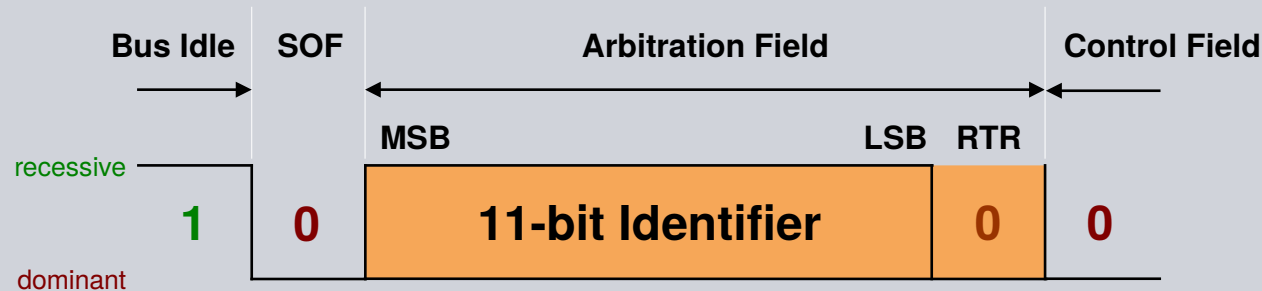




Data Frame Format



Arbitration Field



- contains the Identifier (11 bit for CAN 2.0A) which is used for arbitration
- Identifier determines frame priority: **low** identifier = **high** priority
- the highest seven bits of the identifier must **not** be all **recessive**
- Remote Transmission Request (RTR) bit is always **dominant** in a Data Frame

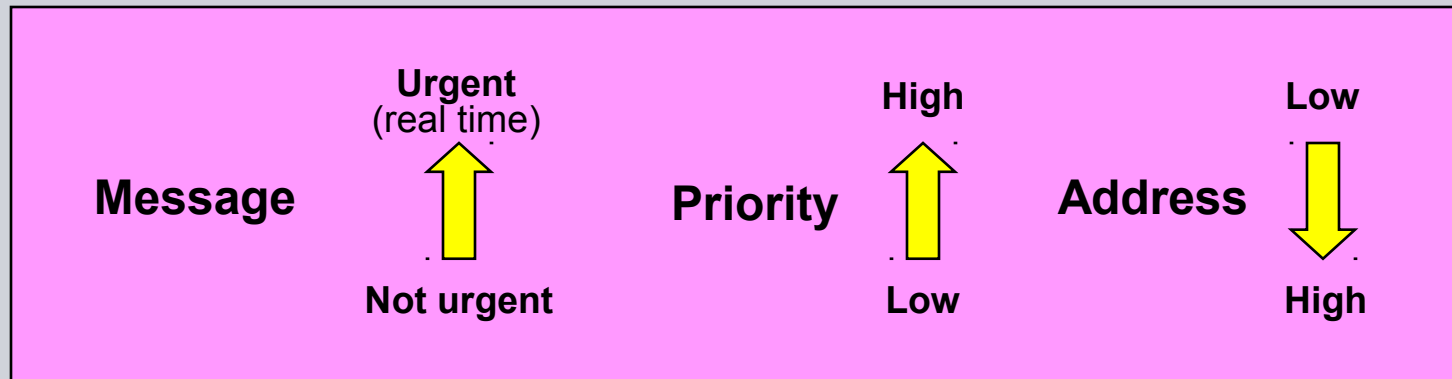
Arbitration on CAN

- **Arbitration = the allocation of bus access rights**
- **Arbitration occurs when several nodes start transmission on the bus at the same time**
- **Arbitration procedure:**
 1. All controllers monitor the bus while transmitting simultaneously
 2. A **dominant** bit (“0”) pulls the bus voltage level to zero
 3. When a controller transmits “1”, but observes “0” on the bus, it has lost arbitration
 4. Controllers who lost arbitration retreat immediately and retry later
 5. Arbitration is won by frame with **lowest** identifier = **highest** priority

Arbitration on CAN

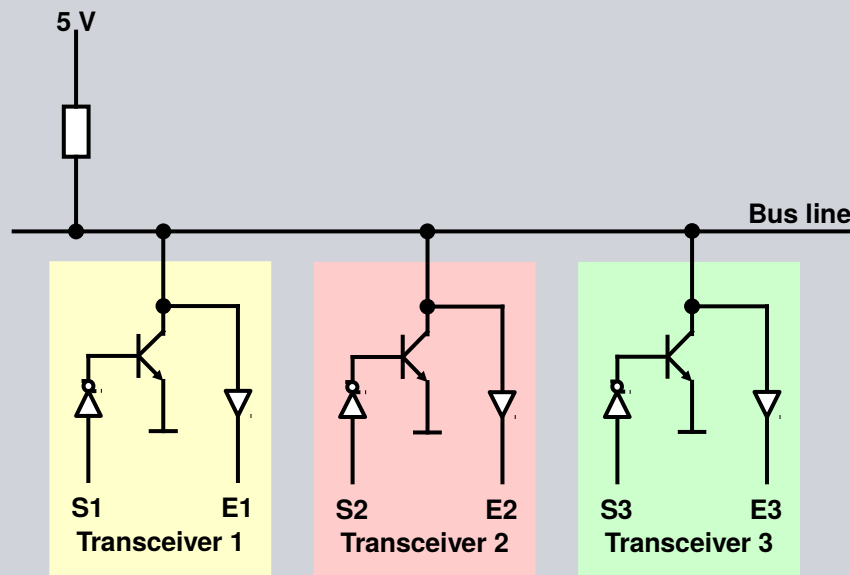
Message Prioritization

Each message includes an identifier. This identifier is used to receive selected messages from the CAN bus data stream and for a simultaneous priority access of different messages.



Recessive and Dominant bus levels

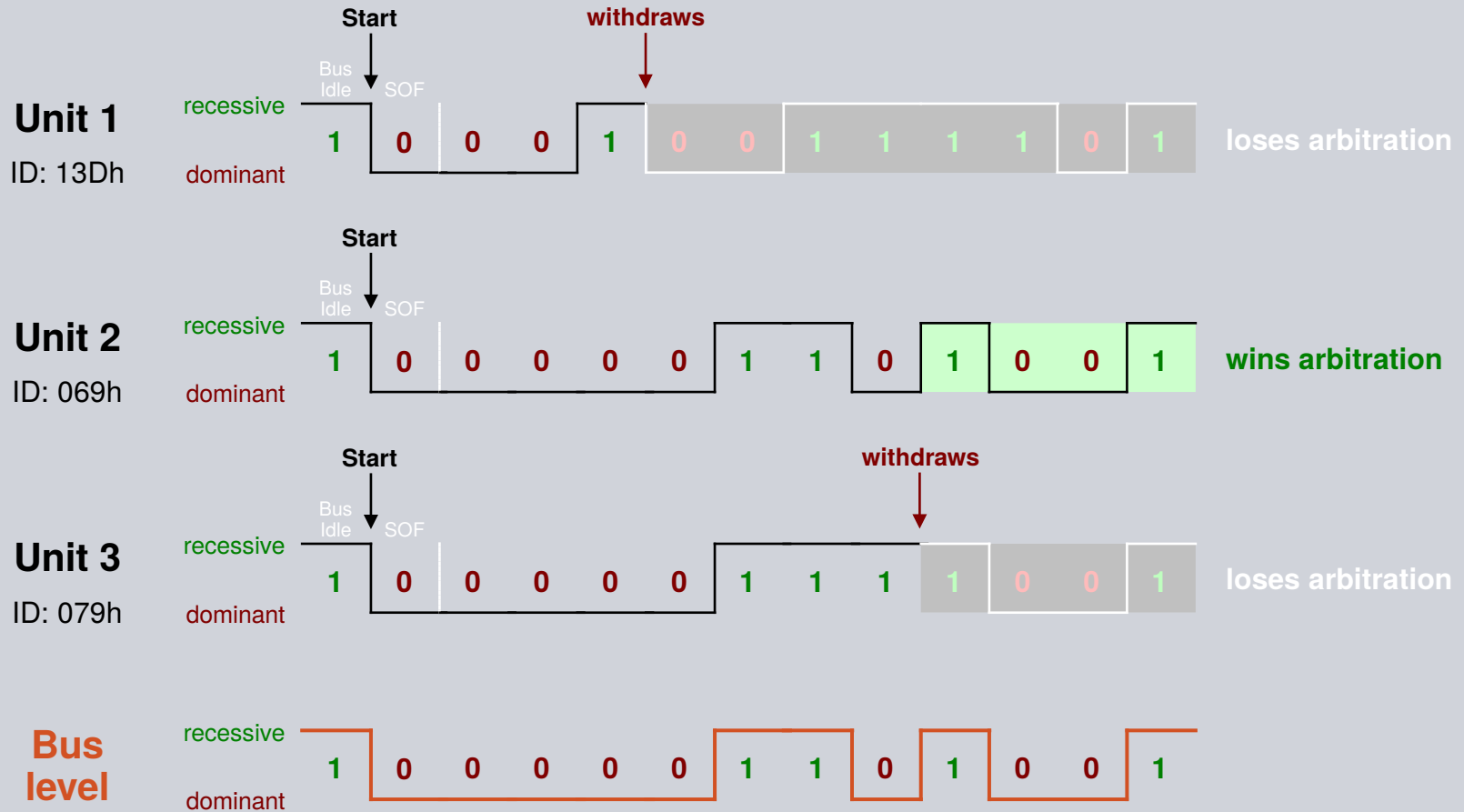
Hardware representation of recessive and dominant bus levels



S1	0	1	0	1	0	1	0	1
S2	0	0	1	1	0	0	1	1
S3	0	0	0	0	1	1	1	1
Bus	0	0	0	0	0	0	0	1

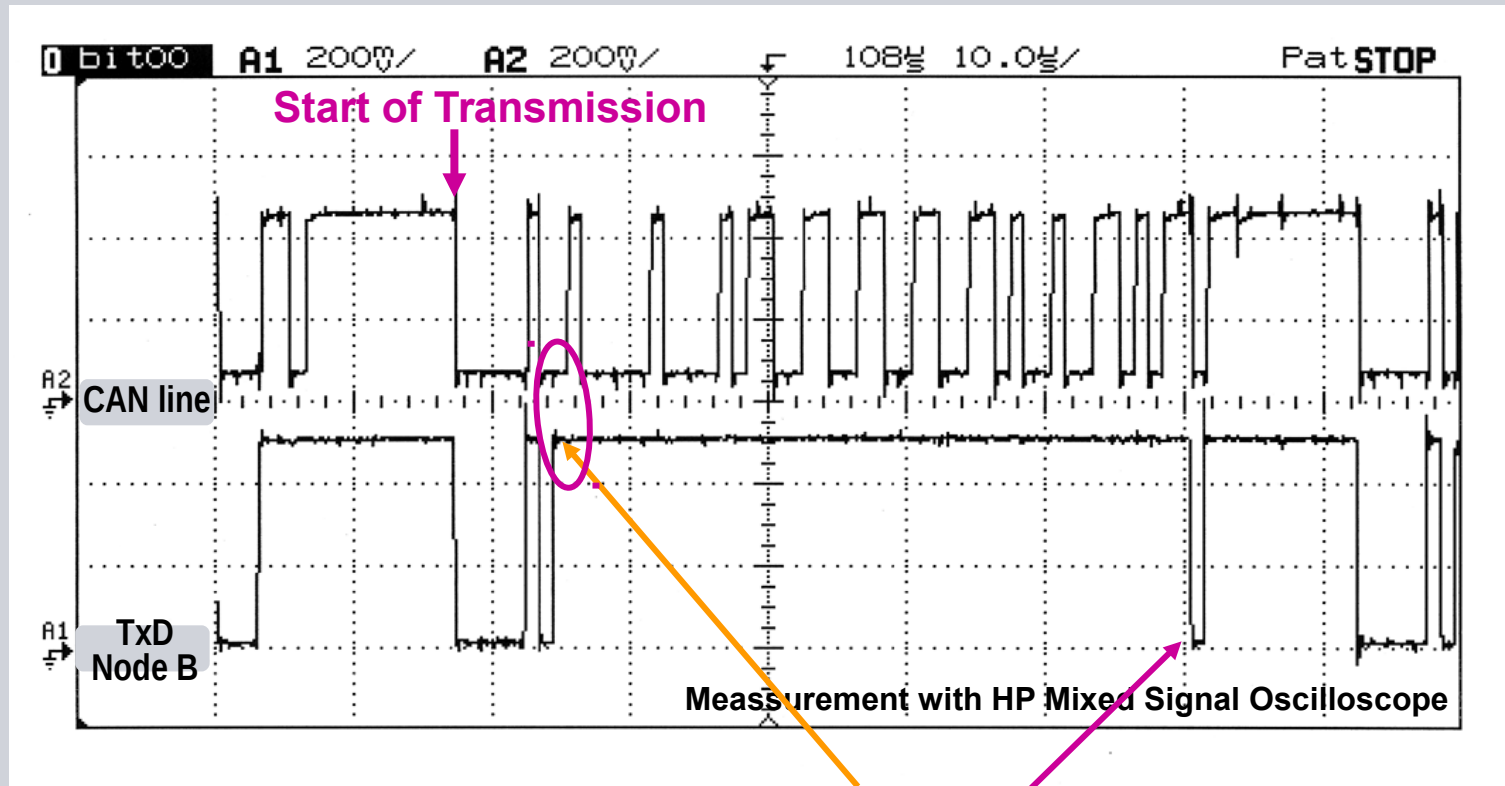
Wired-AND / Open-Collector circuit

Arbitration example.



Arbitration example.

Example: Node A1 and A2 are accessing the CAN bus synchronously.



- Node B: Loses arbitration because of a **higher identifier** than winning node.
- Node B: Switches to receive mode and **acknowledges** the transmitted message.
- Winning node: Transmits with a **lower identifier** **wins** arbitration on the CAN bus.

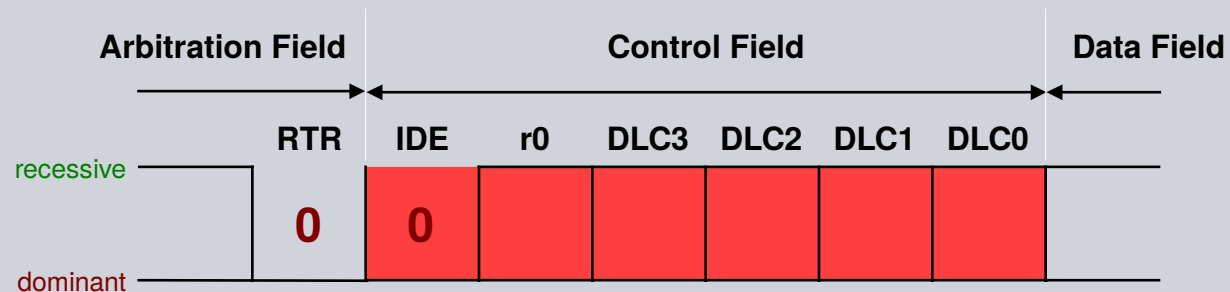


Consequence:
Frames
carrying information
of **high priority**
should have
a **low identifier**

Data Frame Format



Control Field

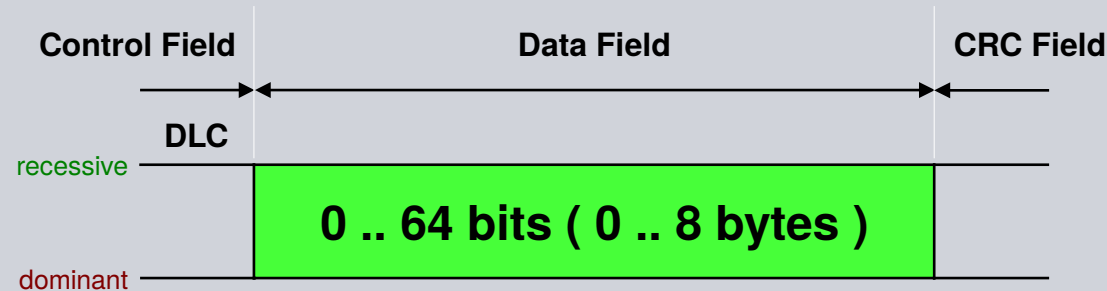


- Identifier Extension (IDE) bit is **dominant** for Standard Frames and **recessive** for Extended Frames
- r0 bit is not used (“reserved for future extensions”)
- Data Length Code (DLC, 4 bits) indicates number of data bytes in Data Field; may take values ranging from 0 to 8, other values are not allowed

Data Frame Format



Data Field

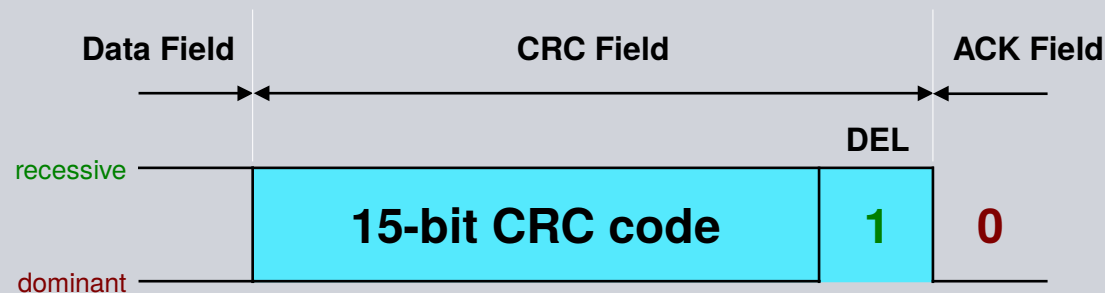


- contains the actual information which is transmitted
- number of data bytes may range from 0 to 8 in units of bytes
- number of data bytes is given in the Data Length Code (DLC)
- transmission starts with the first data byte (byte 0), MSB first

Data Frame Format



CRC Field



- contains the 15-bit Cyclic Redundancy Check (CRC) code
- CRC is a complex, but fast and effective error detection method
- the CRC Field Delimiter (DEL) marks the end of the CRC field
- the CRC Field Delimiter is always **recessive**



The diagram illustrates the Manchester-Bellman protocol timing. It shows three fields: CRC Field, ACK Field, and EOF Field. The signal level is indicated by a horizontal line, with 'recessive' (high) and 'dominant' (low) states. The CRC Field starts with a recessive level (1) and transitions to dominant (0). The ACK Field starts with a dominant level (0) and transitions to recessive (1). The EOF Field starts with a recessive level (1) and transitions to dominant (0).

Field	Signal Level	Value
CRC Field	recessive	1
CRC Field	dominant	0
ACK Field	dominant	0
ACK Field	recessive	1
EOF Field	recessive	1
EOF Field	dominant	0

- contains the Acknowledgement (ACK) bit
- the Acknowledgement bit can be **dominant** or **recessive**
- the ACK Field Delimiter (DEL) marks the end of the ACK field
- the ACK Field Delimiter is always **recessive**

Data Frame Format

⇒ ACK Field

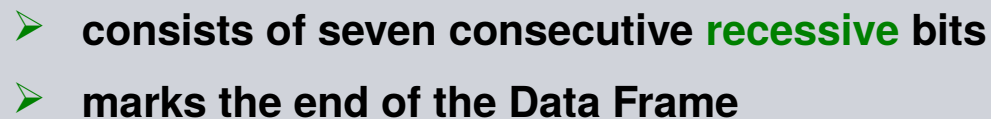
2 bit: **ACK Slot** and **ACK Delimiter**
Transmitter: **recessive** **recessive**
Receiver: **dominant**

⇒ Acknowledgement

Each **Receiver** which received a valid message sends during the ACK slot a dominant bit.
Each **Transmitter** get the information that minimum one node received the message correct. If the transmitter gets no acknowledgement a retransmission is started.

➤ A recessive ACK bit

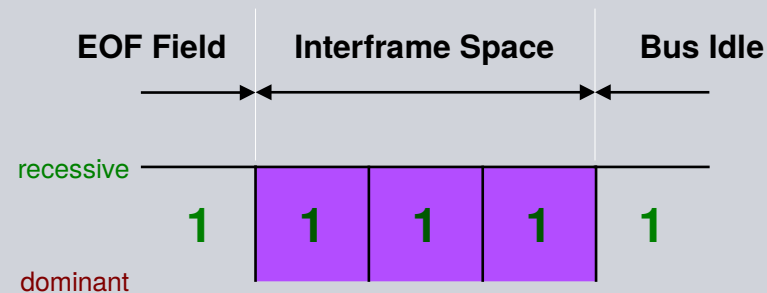
1. could mean that **no** node received the frame without an error or
2. that **no** other node is connected to the bus



Data Frame Format



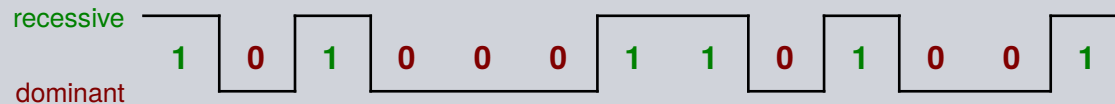
Interframe Space (IFS)



- consists of at least three consecutive **recessive** bits
- no transmission is allowed during the Interframe Space (IFS)
- is needed by controllers to copy received frames from their Rx buffers
- ACK Field Delimiter + EOF + IFS = 11 consecutive **recessive** bits

Bit Stuffing: Motivation

Problems when using NRZ coding



- long sequences of bits of the same polarity
- no changes in voltage level for a longer time
- no falling edges for synchronization
- synchronization between sender and receiver may be lost

Bit Stuffing : Solution

Solution: Bit stuffing

- after **five** consecutive bits of same polarity, insert **one** bit of reverse polarity
- CRC code is calculated before bit stuffing is done
- bit stuffing is done by the sender directly before transmission
- de-stuffing is done by the receiver directly after reception
- CRC code check is done after de-stuffing the frame
- bit stuffing is applied to part of the frame from SOF to CRC field



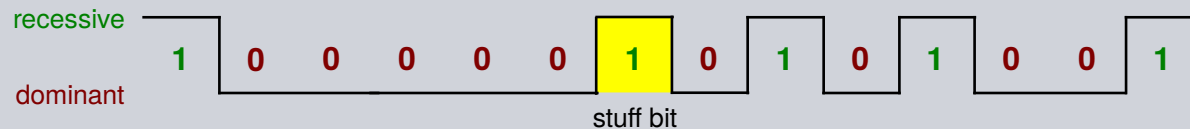
Bit Stuffing: example

Example 1

original
sequence



stuffed
sequence

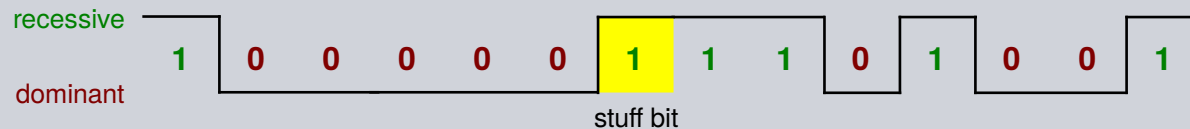


Example 2

original
sequence



stuffed
sequence



Example 3

original
sequence



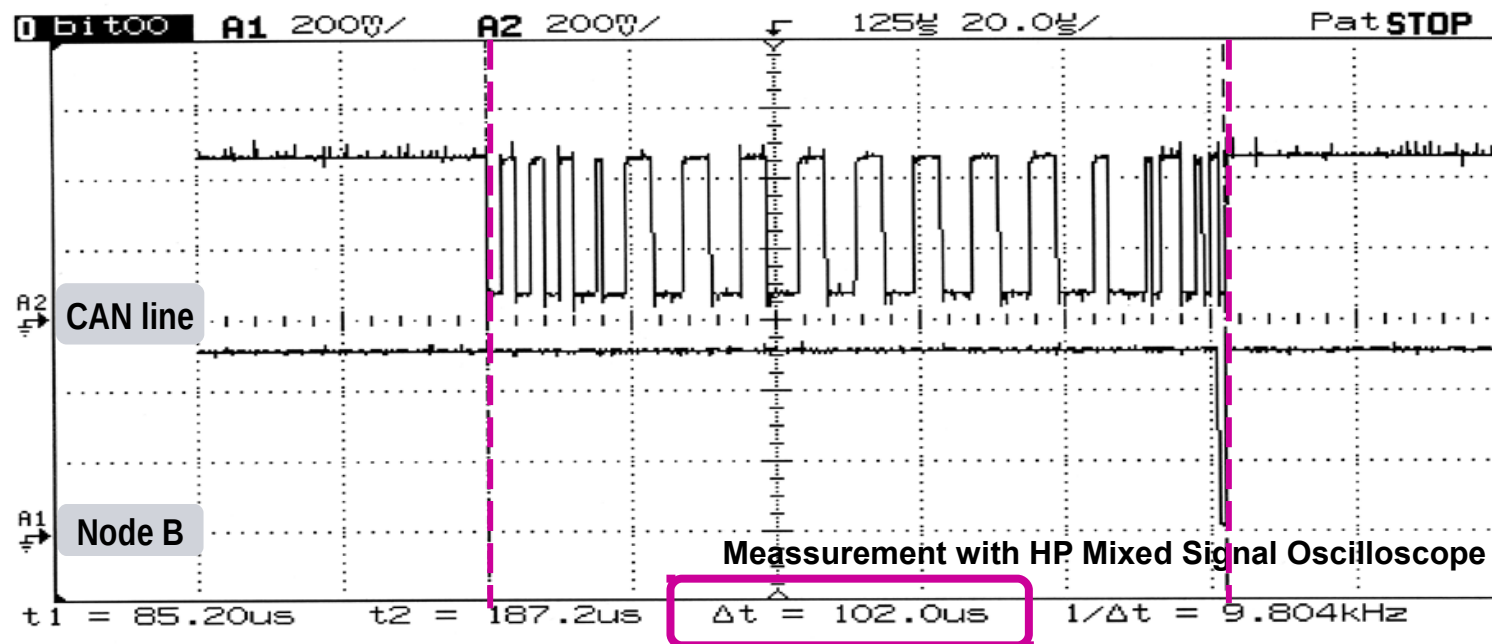
stuffed
sequence



Bit Stuffing: example

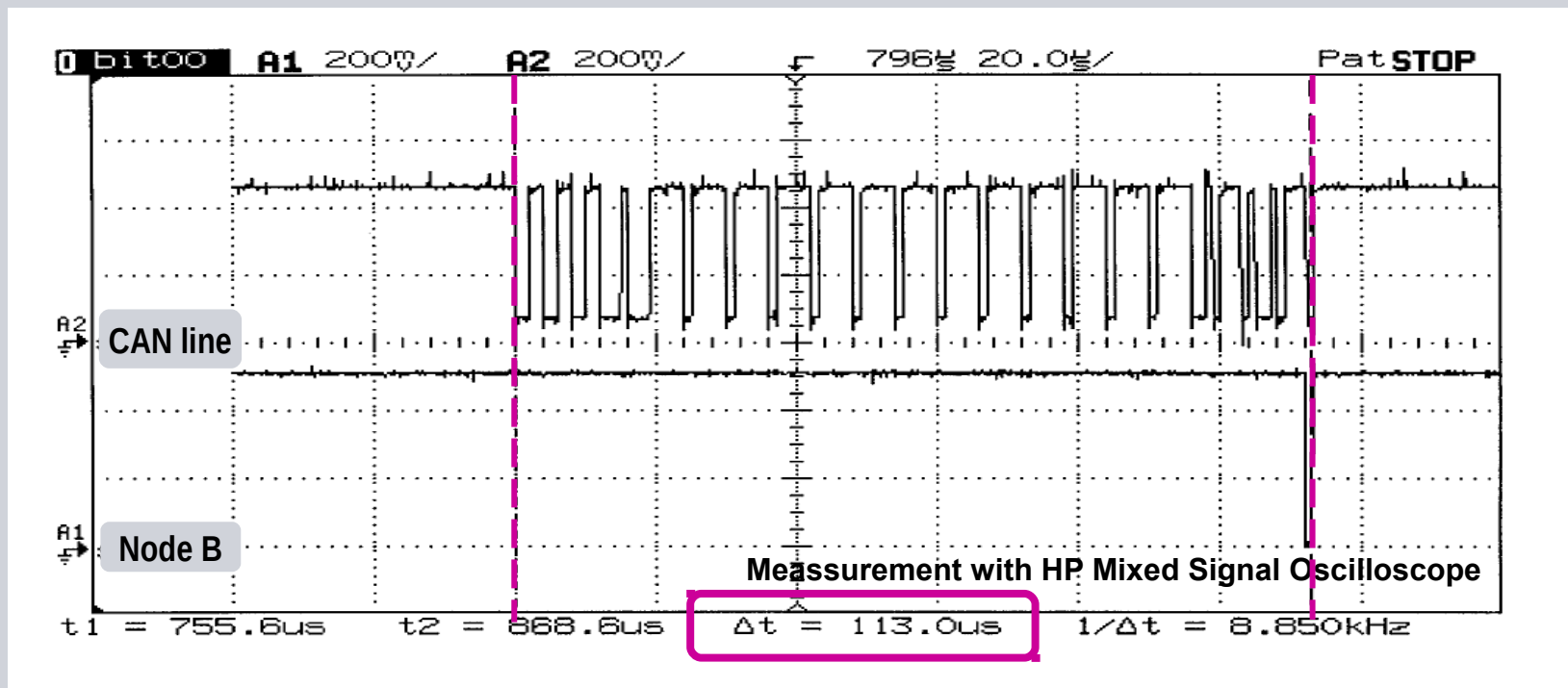
The contents of a frame - address, data, control, CRC bits - defines the absolute bit count.

Transmission of a frame with 8 byte data (F0 F0 F0 F0 F0 F0 F0 F0)
and address 333H. Absolute bit count: **102 Bits**.



Bit Stuffing: example

Transmission of a frame with 8 byte data (FF FF FF FF FF FF FF FF)
and address 333H. Absolute bit count: **113 Bits** (11 additional stuff bits).



Data Frame Format : maximum size

Maximum frame size (for 8 Data Bytes)

	Standard Frame CAN 2.0A (11-bit identifier)	Extended Frame CAN 2.0B (29-bit identifier)
without stuff bits	111 bits	131 bits
with stuff bits	130 bits	154 bits

Error Management



Error Management

Fault Types in a CAN Bus System

There are different types of faults in a CAN bus system:

- ⇒ **Global Faults:**
 - lead to a breakdown of the system wide communication (e.g. short circuits or interrupts of the line, deadlock of the bus due to a permanent transmitting node)
- ⇒ **Local Faults:**
 - are limited to malfunctions of a node or node components (e.g. interrupts in a bus connector, transmit problems of a transmitter, etc.)
- ⇒ **Temporary Faults:**
 - disturb the CAN bus system for a limited time
- ⇒ **Permanent Faults:**
 - disturb the CAN bus system all the time, they are not reversible

Error Management

Following error types are detected:

- ◆ **Bit error** **transmitted and received bit are different**
(except in arbitration, acknowledgement and passive error)
- ◆ **Bit stuffing error** **more than five bits of equal polarity inside of a frame are detected**
- ◆ **CRC error** **receive CRC code doesn't match with the calculated code**
- ◆ **ACK error** **transmitting node receives no dominant acknowledgement bit**
(no receiver accepts the transmission message)
- ◆ **Form error** **fixed-form bit field contains one ore more illegal bits**
e. g. violation of end of frame EOF format, CRC- or ACK-delimiter

Error type: Bit error

Error description

- The bit actually appearing on the bus is different from the one transmitted

Method of detection

- Sending unit constantly monitors the bus while transmitting

Possible cause

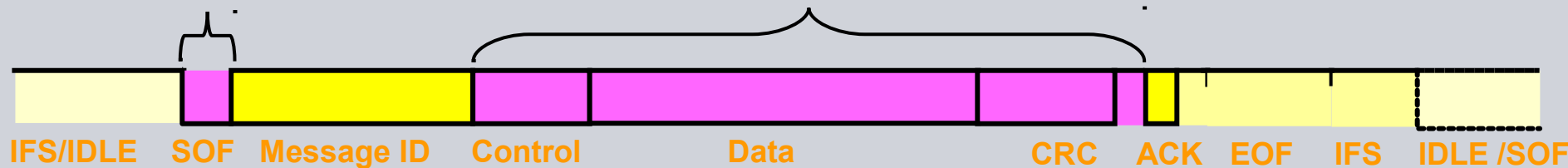
- Sending unit hardware is defective

Exceptions

- Arbitration phase (unit loses arbitration)
- Acknowledgement bit (unit gets positive ACK from at least one receiver)
- Sending of a Passive Error Frame

Error type: Bit error

Bit Error Monitoring Area



Error type: Bit stuffing error

Error description

- Data Frame contains six or more consecutive bits of the same polarity

Method of detection

- Receiver detects error when de-stuffing a received frame

Possible causes

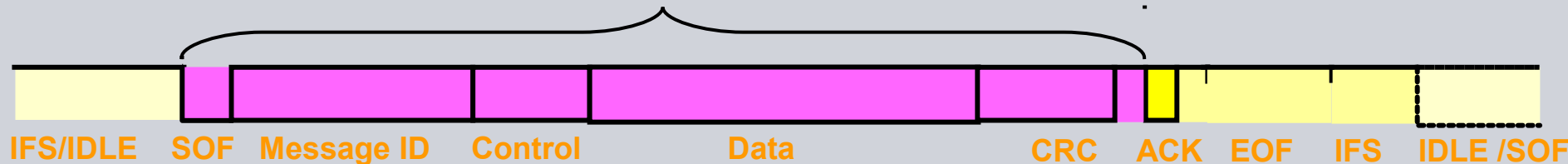
- Error of sending unit during bit stuffing and/or transmission
- Bit changed value during transmission, possibly due to EMI/RFI
- An Active Error Frame was sent

Exceptions

- Transmission of ACK delimiter, EOF field and Interframe Space (IFS):
11 consecutive **recessive** bits, bit stuffing does not apply to this section

Error type: Bit stuffing error

Bit Error Stuffing Monitoring Area



Error type: CRC error

Error description

- CRC code calculated by receiver does not match received CRC code

Method of detection

- Receiver calculates CRC code immediately after reception of the Data Field
- Receiver compares calculated CRC code with the one contained in frame

Efficiency

- Recognizes up to 5 single bit errors per frame → Hamming distance: 6
- Recognizes burst errors with lengths of up to 14 bits
- Recognizes all odd numbers of bit errors
- The more bits the CRC code has, the more efficient it is

Error type: CRC error

CRC Error Monitoring Area



Error type: ACK error

Error description

- Acknowledge (ACK) bit is **recessive**

Method of detection

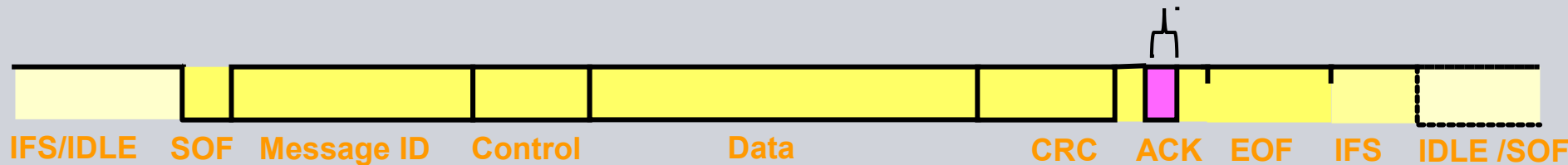
- Sender monitors the bus while transmitting **recessive** ACK bit
- Sender expects to observe **dominant** ACK bit on bus
- Acknowledgement error when ACK bit on bus remains **recessive**

Possible cause

- No other nodes are connected to the bus
- Not one single receiver acknowledges that the received frame was error-free
→ Cause of error is very likely to be found in sender

Error type: **ACK error**

ACK Error Monitoring Area



Error type: **Message format error**

Error description

- **Frame integrity is not preserved**

Method of detection

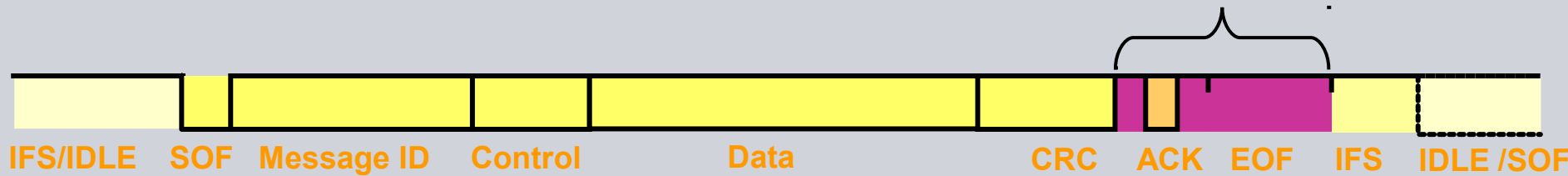
- **Receiver checks received frames for bits or bit fields having a fixed value (e.g. SOF bit, CRC delimiter, ACK delimiter, EOF field, Interframe Space)**
- **Violation of frame integrity when wrong value in one of these fields**

Possible cause

- **Transmission error, error in sender and/or receiver**
- **Transmission of an Active Error Frame during EOF field**
- **Transmission of an Overload Frame during Interframe Space (IFS)**

Error type: **Message format error**

Form Error Monitoring Area

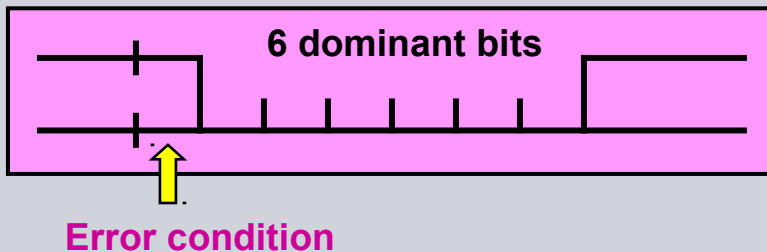


Error Handling

Steps of CAN Error Handling

- Each CAN node which detects an error - **local** or **global error** - sends an error flag to inform all other stations about this error (**globalitization**)
- An error flag is a queue of six consecutive bits

Error active node :



Error passive node :



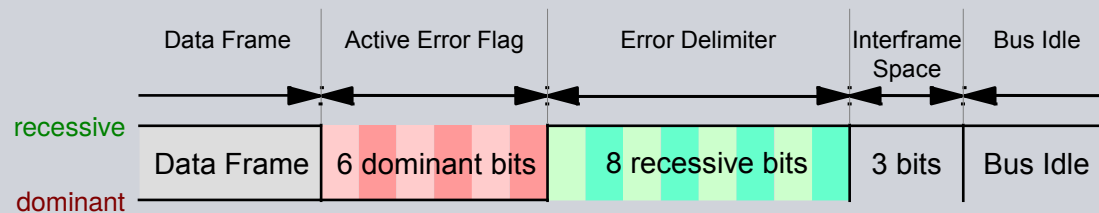
- In the case of one node detects a local error it will send an error flag. Other nodes
- may detect an error because of this error flag. They send also an error flag. After receiving/sending an error flag all nodes cancel the message.
- Each node which detects an error increments his error counter(s).
- The transmitter of a cancelled message **retransmits** automatically this message.

Error management (1)

Procedure observed after detection of an error (1)

- Immediate transmission of an **Error Frame**

Format of an **Active** Error Frame:



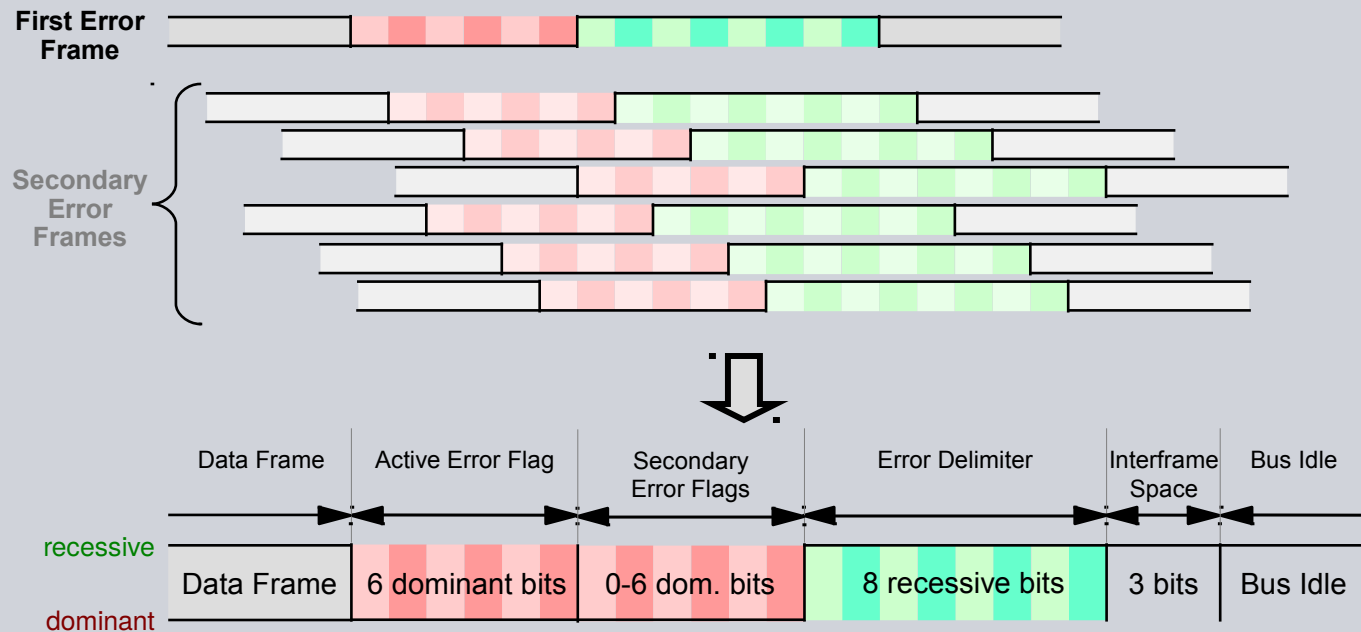
No bit stuffing is applied to Error Frames!

- Error Frame violates the bit stuffing rules → Other receivers are instantly informed that an error has occurred (unless they already found out)

Error management (2)

Procedure observed after detection of an error (2)

- As a result, other units also send Error Frames → Error frames may overlap, resulting in Secondary Error Flags originating from other nodes:



Error management (3)

Procedure observed after detection of an error (3)

- **Sender and receivers reject erroneous frame completely and do not process it any further**
- **Sender retries transmission later**
- **All units on the bus increase their error counters**
- **Recovering from errors can take up to 23 bit cycles
(max. 12 bits Error Flag + 8 bits Error Delimiter + 3 bits Interframe Space)**

Error counters and node state

Two error counters for each unit

- **Transmit Error Counter (TEC)**
- **Receive Error Counter (REC)**

Characteristics of error counters

- TEC and REC start counting at 0 (zero)
- Distinction between sporadic (temporary) and permanent errors possible
- Error-prone units are deactivated automatically after a certain time
- Depending on the values of their error counters, units can assume one of three possible node states: **Error active**, **Error passive**, **Bus off**.

Node state: error active

A unit is in error active state when

- its Transmit Error Counter (TEC) is less than 128: $TEC < 128$ **AND**
- its Receive Error Counter (REC) is less than 128: $REC < 128$

In error active state a unit

- is fully operational
- sends an Active Error Frame when it has detected an error

Node state: error passive

A unit is in error passive state when

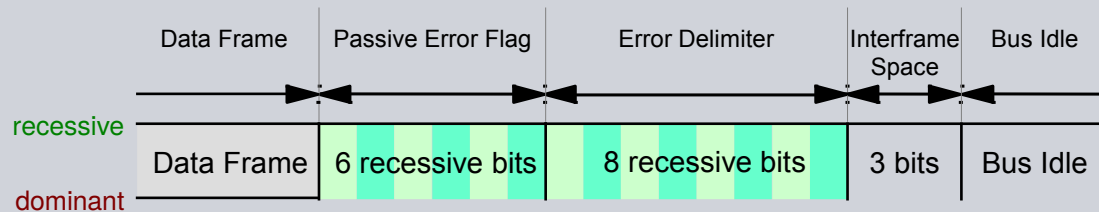
- its Transmit Error Counter (TEC) is greater than 127: $TEC > 127$ **AND / OR**
- its Receive Error Counter (REC) is greater than 127: $REC > 127$

In error passive state a unit

- sends a Passive Error Frame when it has detected an error
- can still receive frames like a unit in **error active** state can
- has to wait after transmission of a Data Frame for 8 additional consecutive **recessive** bit cycles on the bus (Suspend Transmission Field) until it is permitted to transmit another Data Frame
- can go back to **error active** state for $TEC \leq 127$ **AND** $REC \leq 127$

Node state: error passive

Passive Error Frame



- a node in **error passive** state sends a **Passive** Error Frame in case of an error
- a **Passive** Error Frame cannot destroy an ongoing transmission on the bus
- the **Passive** Error Frame might overlap with **Active** Error Frames

Node state: bus off

A unit is in **bus off** state when

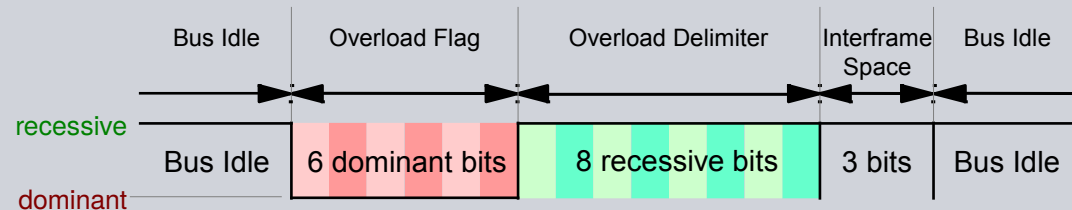
- its Transmit Error Counter (TEC) is greater than 255: $TEC > 255$
- Note: The value of the Receive Error Counter (REC) is of no importance

In **bus off** state a unit

- is practically disconnected from the bus
- cannot receive and transmit anything any more
- can only leave **bus off** mode via a hardware reset **OR** a software reset and subsequent initialization carried through by the host (CAN specification: TEC and REC are set to zero and the unit must receive 128 times a field of 11 consecutive **recessive** bits)

Overload Frame (for delays)

Overload Frame



- Unit sends Overload Frame when at present it cannot receive frames any more due to high workload
- Transmission of an Overload Frame is started during the first two bits of the Interframe Space (IFS) of the preceding frame
- Other units react immediately by also transmitting Overload Frames
⇒ Overload Flags overlap, resulting in up to 12 consecutive **dominant** bits
- Implemented in very few (mostly older) controllers, though controllers must still be able to interpret correctly Overload Frames they receive
- Overload Frames do not influence the error counters (TEC and REC)

Efficiency of the error management

The probability for **not** discovering an error is

$$4.7 * 10^{-11}$$

Example 1*

A CAN bus is used

365 days / year

8 hours / day

with a transmission speed of
and errors arise every

500 kBit / sec

0.7 seconds

⇒ in **1.000** years, only one error remains undiscovered

*Source: CiA

Example 2**

A CAN bus in a car is run at
with an average bus load of
an average data frame size of
for an average operating time of

500 kBit / sec

15 %

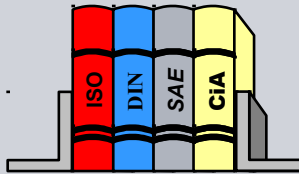
110 bits

4000 hours

⇒ only one error in **100.000** automobiles remains undetected

**Source: Kaiser, Schröder:
"Maßnahmen zur Sicherung
der Daten beim CAN-Bus"

International standards



“The great thing about **standards** is that there are so many to choose from.”

➤ International CAN standards

ISO 11898 “Road vehicles: CAN for **high-speed** communication”

ISO 11519 “Road vehicles: **Low-speed** serial data communication - Low-Speed CAN / VAN”

ISO 11992 “Road vehicles: Electrical connections between towing and towed vehicles”

EIA RS-485 “Electrical Characteristics of Generators and Receivers for Use in Balanced Digital Multipoint Systems” (formerly used for CAN Physical Layer)



Thank you!