

MIDTERM #3 REVIEW

JULY 14TH, 5:30 – 6:45PM + EXTRA

- LHRR
- Recursive Form → Closed Form
- Closed Form → Recursive Form
- Big – O
- Binary Relations

ABOUT THE EXAM...

Detailed logistics
In Moodle Announcement

- CLOSED-BOOK, CLOSED NOTE, 110/100
- Show up before 5:27pm, Camera ON, you and a clean desk
- Deliver via Moodle at 5:30pm, June 14th
- 90 minutes in total (5:30 – 7:00pm)
 - 5 min for setting up
 - 75 min for exam (by 6:50pm)
 - 10 min for scanning and submission (by 7:00pm)
- Submit via Moodle as assignment

unityID.pdf

RECURRENCE RELATION

A rule that defines a_n as a function of previous terms in the sequence is called a recurrence relation.

LINEAR HOMOGENOUS RECURSIVE RELATION (LHRR)

- Linear
 - There are **no product** of recursive terms
- Homogenous
 - There are **no additional terms** besides previous terms in the sequence

$$f_n = c_1 f_{n-1} + c_2 f_{n-2} + \dots + c_k f_{n-k}$$

Step 2: Substitute $s_n = c * x^n$ to find x

$$\text{Hence, } c * x^n = c * x^{n-1} + 12 * c * x^{n-2}$$

$$x^n = x^{n-1} + 12 x^{n-2} \text{ (divide } c \text{ for both side)}$$

$$x^2 = x + 12 \text{ (divide } x^{n-2} \text{ for both side)}$$

$$x^2 - x - 12 = 0$$

$$(x + 3)(x - 4) = 0$$

Therefore, we have $x = -3$ and $x = 4$.

$$\text{Hence, } s_n = c_1 * (-3)^n + c_2 * 4^n$$

Step 0: Is it a LHRR?

Yes, LHRR, degree = 2

Step 1: Assume $s_n = c * x^n$

Step 3: Use the base case to find c

$$s_0 = c_1 * (-3)^0 + c_2 * 4^0 = -2 \rightarrow c_1 + c_2 = -2$$

$$s_1 = c_1 * (-3)^1 + c_2 * 4^1 = 20 \rightarrow -3c_1 + 4c_2 = 20$$

Therefore, $c_1 = -4$, $c_2 = 2$

That is, $s_n = -4 * (-3)^n + 2 * 4^n$

Q: Find the recursive form given the closed form: $s_n = 5 * 2^n + 2 * 4^n$

Therefore, $x_1 = 2, x_2 = 4$

So we have $(x - 2)(x - 4) = 0$

$$x^2 - 6x + 8 = 0$$

$$x^2 = 6x - 8$$

Then we have $x^n = 6x^{n-1} - 8x^{n-2}$ (multiply both sides by x^{n-2})

$$s_n = 6s_{n-1} - 8s_{n-2} \text{ for } n \geq 2$$

$$s_0 = 5 * 2^0 + 2 * 4^0 = 7$$

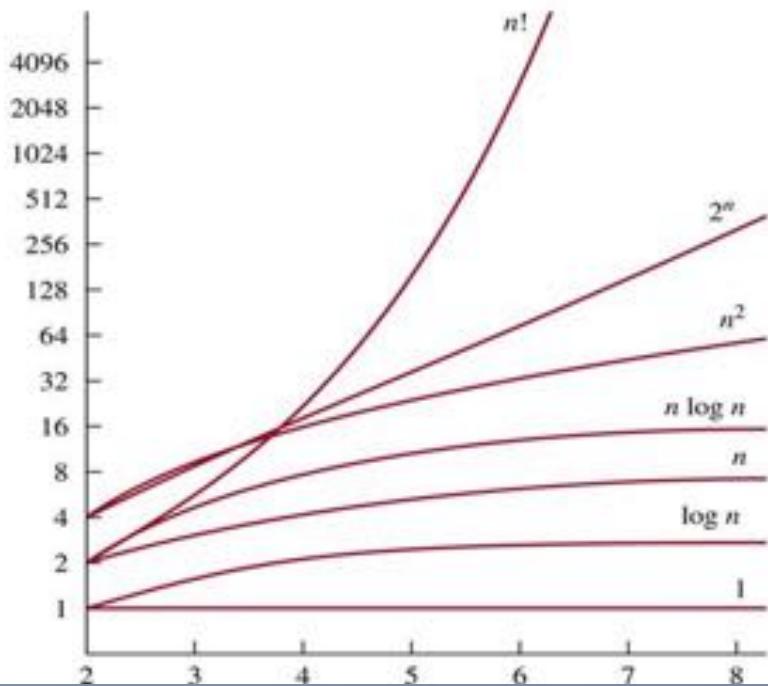
$$s_1 = 5 * 2^1 + 2 * 4^1 = 18$$

STEPS TO FIND BIG – O OF A FUNCTION

- 1) Find the dominant term (as a function of the variable)
- 2) Replace all other functions with the variable with 1)
- 3) Simplify
- 4) Remove constants

The Growth of Combinations of Functions

- 1
- $\log n$
- n
- $n \log n$
- n^2
- 2^n
- $n!$



$$n! > 2^n > n^2 > n \log n > n > \log n > 1$$

SOME BIG – O PROPERTIES

$$n! > 2^n > n^2 > n \log n > n > \log n > 1$$

- When **adding** functions, the Big – O of the **fastest** growing function dominates
 - E.g., $3n \log n + 2000 \log n + 20n + 10000 \in O(n \log n)$
- When **multiplying** functions, the Big – O of the functions are **multiplied**
 - E.g., $n^3(n^2 \log n + n)$
 - $n^3 \in O(n^3)$
 - $n^2 \log n + n \in O(n^2 \log n)$
 - Therefore, $n^3(n^2 \log n + n) \in O(n^3 \cdot n^2 \log n) = O(n^5 \log n)$

RECAP

- A **binary relation** between two sets A and B is a **subset R** of $A \times B$ (aRb).
- Representation
 - Set, Matrix, Directed Graph
- Reflexive: $\forall x \in S, xRx$
- Irreflexive / anti-reflexive $\forall x \in S, x \neq x$
- Transitive $xRy \wedge yRz \rightarrow xRz$
- Symmetric $xRy \Leftrightarrow yRx$
- Anti-symmetric $xRy \wedge yRx \rightarrow x = y$
- Asymmetric $xRy \rightarrow yRx$

OPERATIONS ON BINARY RELATIONS

- Union
- Intersection
- XOR
- Composition
 - The pair $(a, c) \in S \circ R$ if and only if there is a $b \in A$ such that $(a, b) \in R$ and $(b, c) \in S$.