

The Algebra of Algebraic Data Types

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Abusing the algebra of algebraic data types - why does this work?

86



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The 'algebraic' expression for algebraic data types looks very suggestive to someone with a background in mathematics. Let me try to explain what I mean.

Having defined the basic types

- Product •
- Union +
- Singleton X
- Unit 1

and using the shorthand X^2 for $X \cdot X$ and $2X$ for $X + X$ et cetera, we can then define algebraic expressions for e.g. linked lists

```
data List a = Nil | Cons a (List a) ↔  $L = 1 + X \cdot L$ 
```

and binary trees:

```
data Tree a = Nil | Branch (Tree a) (Tree a) ↔  $T = 1 + X \cdot T^2$ 
```

Now, my first instinct as a mathematician is to go nuts with these expressions, and try to solve for L and T . I could do this through repeated substitution, but it seems much easier to abuse the notation horrifically and pretend I can rearrange it at will. For example, for a linked list:

*“Haskell's algebraic data types are named such since they correspond to an **initial algebra in category theory**, giving us some laws, some operations and some symbols to manipulate.”*

– Don Stewart

*“Haskell's algebraic data types are named such since they correspond to an initial algebra in category theory, giving us some **laws**, some **operations** and some **symbols** to manipulate.”*

– Don Stewart

Algebra

Symbols

Operations

Laws

Algebra

Symbols $0, 1, 2, x, y, z, \dots$

Operations $+, -, \times, \div, \dots$

Laws $0 + x = x, \dots$

Algebra

Symbols	Types <code>()</code> , <code>Int</code> , <code>Bool</code> , ...)
Operations	Type constructors (<code>Maybe</code> , <code>Either</code>)
Laws	?

Algebra

Symbols

Things

Operations

Ways to make new things

Laws

Rules the things follow

Prelude

```
{-# LANGUAGE EmptyDataDecls  
      , TypeOperators #-}
```

One

`data Unit = Unit`

One

`data Unit = Unit`

`data () = ()`

Addition

`data a :+ b = AddL a | AddR b`

Addition

`data a :+ b = AddL a | AddR b`

`data Either a b = Left a | Right b`

Multiplication

```
data a :* b = Mul a b
```

Multiplication

`data a :* b = Mul a b`

`data (a,b) = (a,b)`

Zero

“In general, an algebraic type specifies a sum of one or more alternatives, where each alternative is a product of zero or more fields.

It might have been useful to permit a sum of zero alternatives, which would be a completely empty type, but at the time the value of such a type was not appreciated.”

– Hudak, Hughes, Peyton Jones, Wadler

Zero

data Void

Two

```
type Two = Unit :+: Unit
```

Two

```
type Two = Unit :+: Unit
```

```
data Bool = False | True
```

Notation

Void $\Rightarrow \emptyset$

Unit $\Rightarrow 1$

Bool $\Rightarrow 2$

Addition $\Rightarrow a + b$

Multiplication $\Rightarrow a \cdot b$

Laws

$$0 + x = x$$

Either Void $x \cong x$

Laws

$$0 \cdot x = 0$$

$$(Void, x) \cong Void$$

Laws

$$1 \cdot x = x$$

$$(\text{()}, x) \cong x$$

Laws

$$x + y = y + x$$

$$\text{Either } x \ y \cong \text{Either } y \ x$$

Laws

$$x \cdot y = y \cdot x$$

$$(x, y) \cong (y, x)$$

Functions

data $a \rightarrow b = ?$

Functions

Domain:	True	False
---------	------	-------

Range:	True	False
--------	------	-------

Functions

Domain:

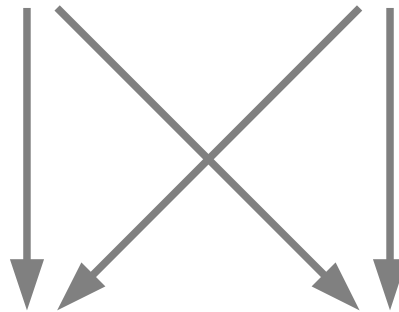
True

False

Range:

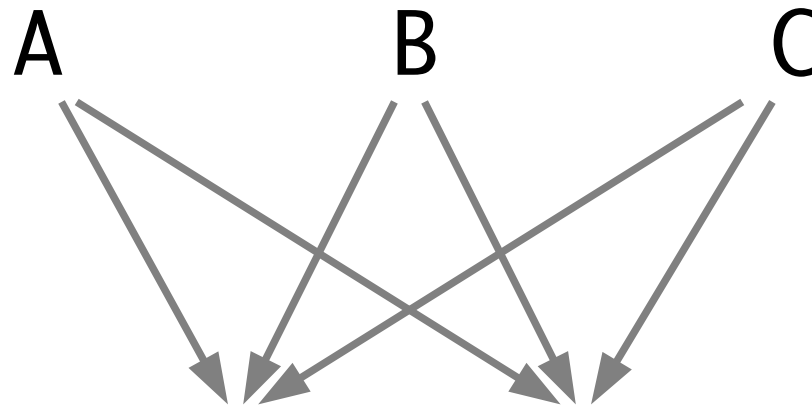
True

False



Functions

Domain:



Range:

True

False

Functions

$a \rightarrow b$

\Leftrightarrow

b^a

Laws

$$1^a = 1$$

$$a \rightarrow \bigcirc \cong \bigcirc$$

Laws

$$a^1 = a$$

$$() \rightarrow a \cong a$$

Laws

$$(b \cdot c)^a = b^a \cdot c^a$$

$$a \rightarrow (b, c) \cong (a \rightarrow b, a \rightarrow c)$$

Laws

$$c^{ba} = (c^b)^a$$

$$(a, b) \rightarrow c \cong a \rightarrow b \rightarrow c$$

Lists

```
data List x = Nil | Cons x (List x)
```

Lists

```
data List x = Nil | Cons x (List x)
```

$$L(x) = 1 + x \cdot L(x)$$

Lists

```
data List x = Nil | Cons x (List x)
```

```
L = 1 + x L
```

Lists

```
data List x = Nil | Cons x (List x)
```

$$L = 1 + x L$$
$$L = 1 + x (1 + x L)$$

Lists

```
data List x = Nil | Cons x (List x)
```

$$L = 1 + x L$$

$$L = 1 + x (1 + x L)$$

$$L = 1 + x + x^2 (1 + x L)$$

Lists

```
data List x = Nil | Cons x (List x)
```

$$L = 1 + x L$$

$$L = 1 + x (1 + x L)$$

$$L = 1 + x + x^2 (1 + x L)$$

$$L = 1 + x + x^2 + x^3 + x^4 + \dots$$

Lists

```
data List x = Nil | Cons x (List x)
```

```
L = 1 + x L
```

Lists

```
data List x = Nil | Cons x (List x)
```

$$\begin{aligned} L &= 1 + x L \\ L (1 - x) &= 1 \end{aligned}$$

Lists

```
data List x = Nil | Cons x (List x)
```

$$\begin{aligned}L &= 1 + x L \\L (1 - x) &= 1 \\L &= 1 / (1 - x)\end{aligned}$$

Lists

```
data List x = Nil | Cons x (List x)
```

$$L = 1 + x L$$

$$L (1 - x) = 1$$

$$L = 1 / (1 - x)$$

$$L = 1 + x + x^2 + x^3 + x^4 + \dots$$

Trees

```
data Tree x = Tip | Node (Tree x) x (Tree x)
```

Trees

```
data Tree x = Tip | Node (Tree x) x (Tree x)
```

$$T = 1 + x T^2$$

Trees

`data Tree x = Tip | Node (Tree x) x (Tree x)`

$$T = 1 + x T^2$$

$$x T^2 - T + 1 = 0$$

Quadratic Formula (Interlude)

$$ax^2 + bx + c = 0$$

Quadratic Formula (Interlude)

$$ax^2 + bx + c = 0$$

$$x = (-b \pm \sqrt{b^2 - 4ac}) / 2a$$

Trees

```
data Tree x = Tip | Node (Tree x) x (Tree x)
```

$$T = 1 + x T^2$$

$$x T^2 - T + 1 = 0$$

Trees

```
data Tree x = Tip | Node (Tree x) x (Tree x)
```

$$T = 1 + x T^2$$

$$x T^2 - T + 1 = 0$$

$$T = (1 - \sqrt{1 - 4x}) / 2x$$

Trees

```
data Tree x = Tip | Node (Tree x) x (Tree x)
```

$$T = 1 + x T^2$$

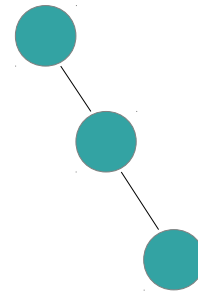
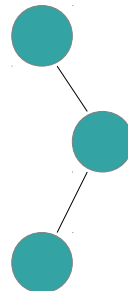
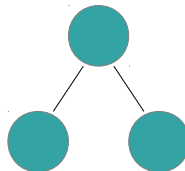
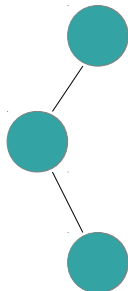
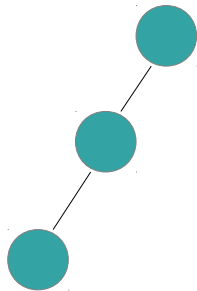
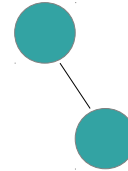
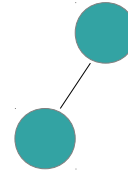
$$x T^2 - T + 1 = 0$$

$$T = (1 - \sqrt{1 - 4x}) / 2x$$

$$T = 1 + x + 2x^2 + 5x^3 + 14x^4 + \dots$$

Trees

$$T = 1 + x + 2x^2 + 5x^3 + 14x^4 + \dots$$



Zippers

Problem: Navigate and modify a data structure (e.g. a list) efficiently

```
> let x = [1, 2, 3, 4, 5, 6]
```

Zippers

Problem: Navigate and modify a data structure (e.g. a list) efficiently

```
> let x = [1, 2, 3, 4, 5, 6]
```

Solution: The zipper

```
> data Zip a = Zip [a] a [a]
```

Zippers

Problem: Navigate and modify a data structure (e.g. a list) efficiently

```
> let x = [1, 2, 3, 4, 5, 6]
```

Solution: The zipper

```
> let z = Zip [2, 1] 3 [4, 5, 6]  
> right z -- Zip [3, 2, 1] 4 [5, 6]  
> right z -- Zip [4, 3, 2, 1] 5 [6]
```


One-Hole Contexts

One-hole context: Data structure with a hole

$[1, 2, 3] * [5, 6]$

One-Hole Contexts

One-hole context: Data structure with a hole

$$[1, 2, 3] * [5, 6]$$

Zipper: One-hole context paired with data

$$(4, [1, 2, 3] * [5, 6])$$

One-Hole Contexts

x

x

One-Hole Contexts

x

x

*

1

One-Hole Contexts

(x, x)

x^2

One-Hole Contexts

(x, x)

x^2

$(*, x) + (x, *)$

$2x$

One-Hole Contexts

(x, x, x)

x^3

One-Hole Contexts

(x, x, x)

x^3

$(*, x, x) + (x, *, x)$
 $+ (x, x, *)$

$3x^2$

One-Hole Contexts

$$x \Rightarrow 1$$

$$x^2 \Rightarrow 2x$$

$$x^3 \Rightarrow 3x^2$$

One-Hole Contexts

*“The Derivative of a Regular Type is
its Type of One-Hole Contexts”*

– Conor McBride

Type Calculus

∂ = “Take the derivative with respect to x ”

Type Calculus

$$\partial 1 = \emptyset$$

Type Calculus

$$\partial 1 = 0$$

$$\partial x = 1$$

Type Calculus

$$\partial 1 = 0$$

$$\partial x = 1$$

$$\partial(f + g) = \partial f + \partial g$$

Type Calculus

$$\partial 1 = 0$$

$$\partial x = 1$$

$$\partial(f + g) = \partial f + \partial g$$

$$\partial(f \cdot g) = \partial f \cdot g + f \cdot \partial g$$

Type Calculus

$$\partial 1 = 0$$

$$\partial x = 1$$

$$\partial(f + g) = \partial f + \partial g$$

$$\partial(f \cdot g) = \partial f \cdot g + f \cdot \partial g$$

$$\partial(f(g)) = \partial f(g) \cdot \partial g$$

Type Calculus

$$L(x) = 1 / (1 - x)$$

Type Calculus

$$L(x) = 1 / (1 - x)$$

$$\partial L(x) = 1 / (1 - x)^2$$

Type Calculus

$$L(x) = 1 / (1 - x)$$

$$\partial L(x) = 1 / (1 - x)^2$$

$$= L(x)^2$$

Type Calculus

$$T = 1 + x \ T^2$$

Type Calculus

$$T = 1 + x T^2$$

$$\partial T = T^2 + 2xT \partial T$$

Type Calculus

$$T = 1 + x T^2$$

$$\partial T = T^2 + 2xT \partial T$$

$$\partial T = T^2 / (1 - 2xT)$$

Type Calculus

$$T = 1 + x T^2$$

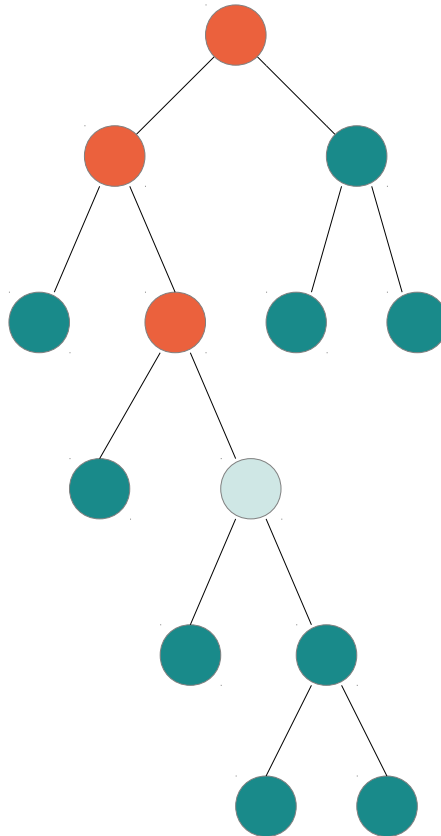
$$\partial T = T^2 + 2xT \partial T$$

$$\partial T = T^2 / (1 - 2xT)$$

$$\partial T = T^2 \cdot L(2xT)$$

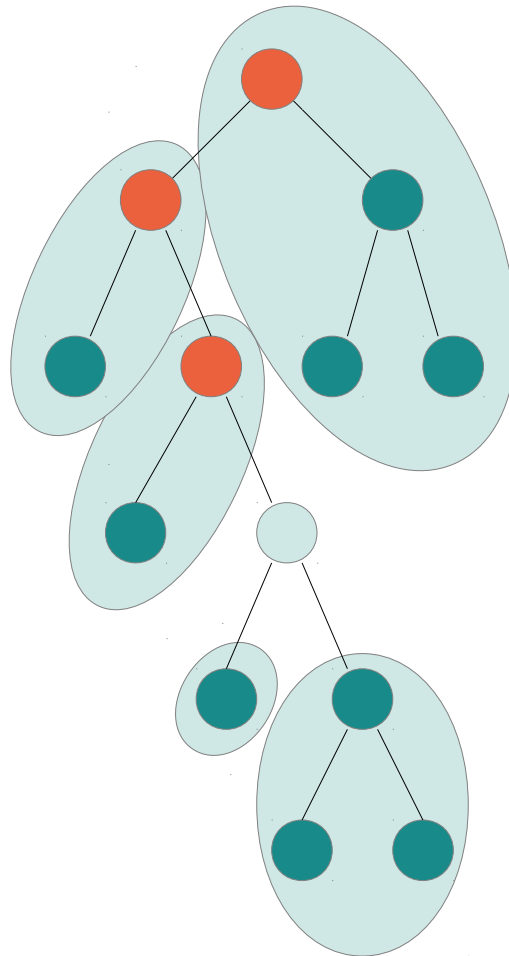
Type Calculus

$$\partial T = T^2 \cdot L(2 \times T)$$



Type Calculus

$$\partial T = T^2 \cdot L(2 \times T)$$



Non-Regular Types

Bags: No ordering

ULists: Unique elements

Sets: Unique elements, no ordering

Cyclic lists, Deques, etc.

Non-Regular Types

$\text{Set}_n = \text{“Sets of size } n\text{”}$

$$\text{Set}_0(x) = 1$$

$$\text{Set}_1(x) = x$$

$$\text{Set}_2(x) = x(x-1) / 2$$

Non-Regular Types

$$\begin{aligned}\text{Set}_n(x) &= x (x-1) \dots (x-n+1) / n! \\ &= x^{\underline{n}} / n!\end{aligned}$$

Non-Regular Types

$$\text{Set}(x) = 1 + x + x^2/2! + x^3/3! + \dots$$

Non-Regular Types

$$\text{Set}(x) = 1 + x + x^2/2! + x^3/3! + \dots$$

$$\Delta f(x) = f(x+1) - f(x)$$

Non-Regular Types

$$\text{Set}(x) = 1 + x + x^2/2! + x^3/3! + \dots$$

$$\Delta f(x) = f(x+1) - f(x)$$

$$\Delta \text{Set}(x) = \text{Set}(x)$$

Non-Regular Types

$$\Delta\text{Set}(x) = \text{Set}(x)$$

Non-Regular Types

$$\Delta\text{Set}(x) = \text{Set}(x)$$

$$\Rightarrow \text{Set}(x+1) - \text{Set}(x) = \text{Set}(x)$$

Non-Regular Types

$$\Delta \text{Set}(x) = \text{Set}(x)$$

$$\Rightarrow \text{Set}(x+1) - \text{Set}(x) = \text{Set}(x)$$

$$\Rightarrow \text{Set}(x+1) = 2 \text{ Set}(x)$$

Non-Regular Types

$$\Delta \text{Set}(x) = \text{Set}(x)$$

$$\Rightarrow \text{Set}(x+1) - \text{Set}(x) = \text{Set}(x)$$

$$\Rightarrow \text{Set}(x+1) = 2 \text{ Set}(x)$$

$$\Rightarrow \text{Set}(x) = 2^x$$

Non-Regular Types

$$\text{Set}(x) = 2^x$$

$$\text{Set } x \cong x \rightarrow \text{Bool}$$

More Type Algebra

github.com/chris-taylor/LondonHUG

More Type Algebra

Combinatorial Species

Andre Joyal, Brent Yorgey

Calculus of Types

Conor McBride, Dan Piponi (sigfpe)