A Linear Time Active Learning Algorithm for Link Classification

positive edge
negative edge

deges with binary labels

Model

Initial labelling; consistent

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We present three efficient active learning algorithms for link classification in signed networks.

Signed networks are graphs whose edges carry a sign representing the **positive or negative nature** of the relation.

Applicative domains such e-commerce, social networks and biology offer examples of this kind of networks.

Most of the heuristics for link classification are summarized by the motto "the enemy of my enemy is my friend" (i.e. see [1,2,3]).

In this paper we assume a 2-cluster structure with a p-stochastic label assignment. Each edges label has P(sign_is_flipped) ≤ p, where the initial labelling is a consistent 2-cluster.

Setup

We used different training set sizes for ASymExp (indicated on the x-axis) and a fixed amount (|V|-1 edges) for TreeCutter. Since classes are very unbalanced, we report the F-Measure averaged on 10 runs for each combination of parameters.

Experiments

Algorithms

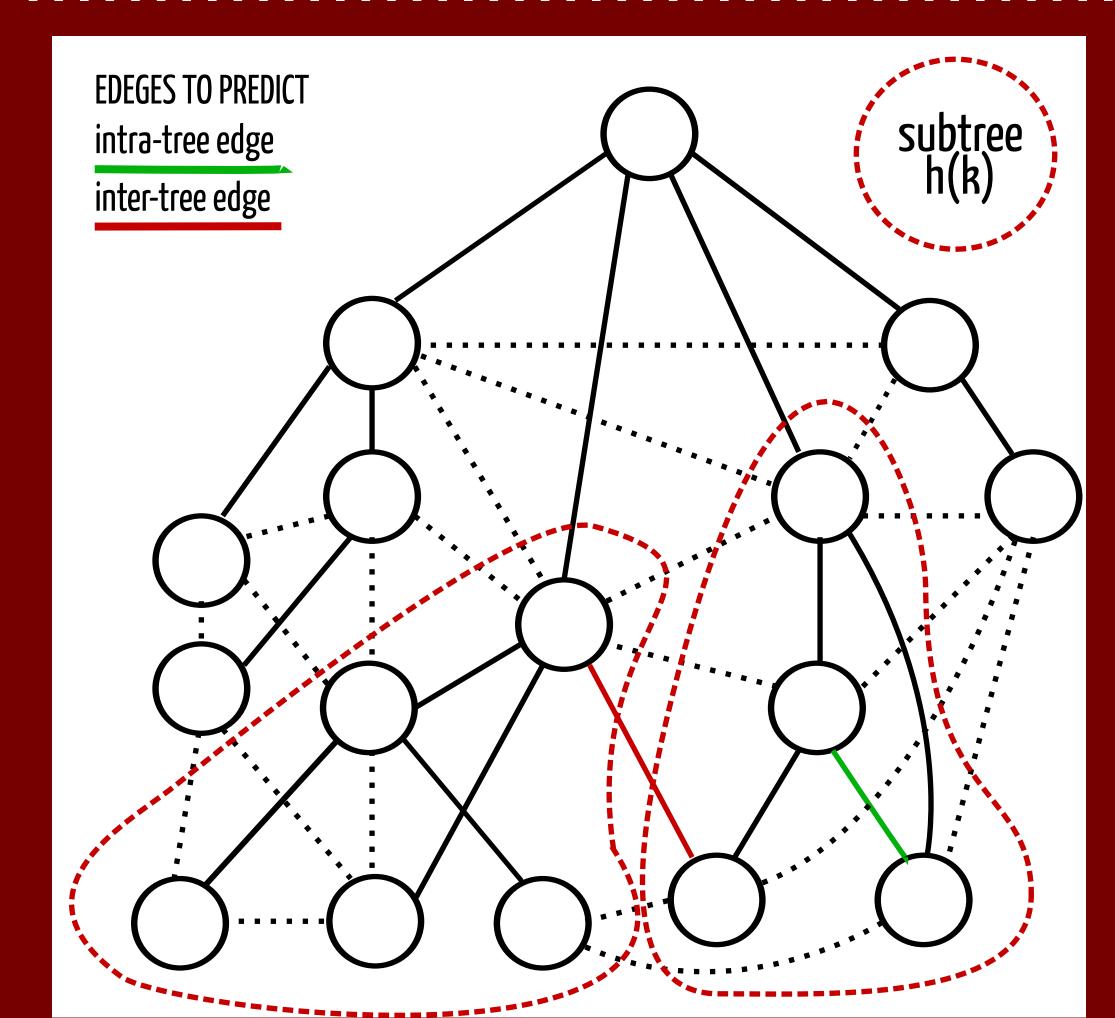
Giovanni Zappella*

We compared our simplest algorithm (**TreeCutter**) to the heuristics presented in [1], but we reported only the best of those algorithms (**ASymExp**) with different settings of the parameter

Datasets

DELTA* are synthetic datasets created in order to control the number of flipped labels. (e.g. DELTA100 means 100 labels have been flipped in the original consistent 2-cluster). **Movielens** is a graph of users created from movies (normalized) co-rating. **SlashDot** and **Epinions** are real-world social networks with natively signed edges.

2-cluster of the graph Each label is flipped with P(flip label) \le p TreeCutter TreeCutter



Given a training set Eo

Lower

Bound

 $M \ge p \mid E \setminus Eo \mid$

for any undirected graph

Breadth-first Spanning tree

Using a simple BF-spanning tree it is possible to bound the expected number of mistakes as

E[M] ≤ 2 Diam p | E \ Eo |

StarMaker

For any undirected graph with $|E| = \Omega(|V|^{3/2})$ the expected number of mistakes made by **StarMaker** is

$E[M] \leq 5p |E \setminus Eo|$

while the query set size is upper bounded by $O(|V|)^{3/2}$.

It is possible to reduce the query set size by a factor of k^{3/2} using **TreeLetStar**(k) with a mistake bound

 $E[M] = O(min\{k,Diam\}) p | E \setminus Eo|$

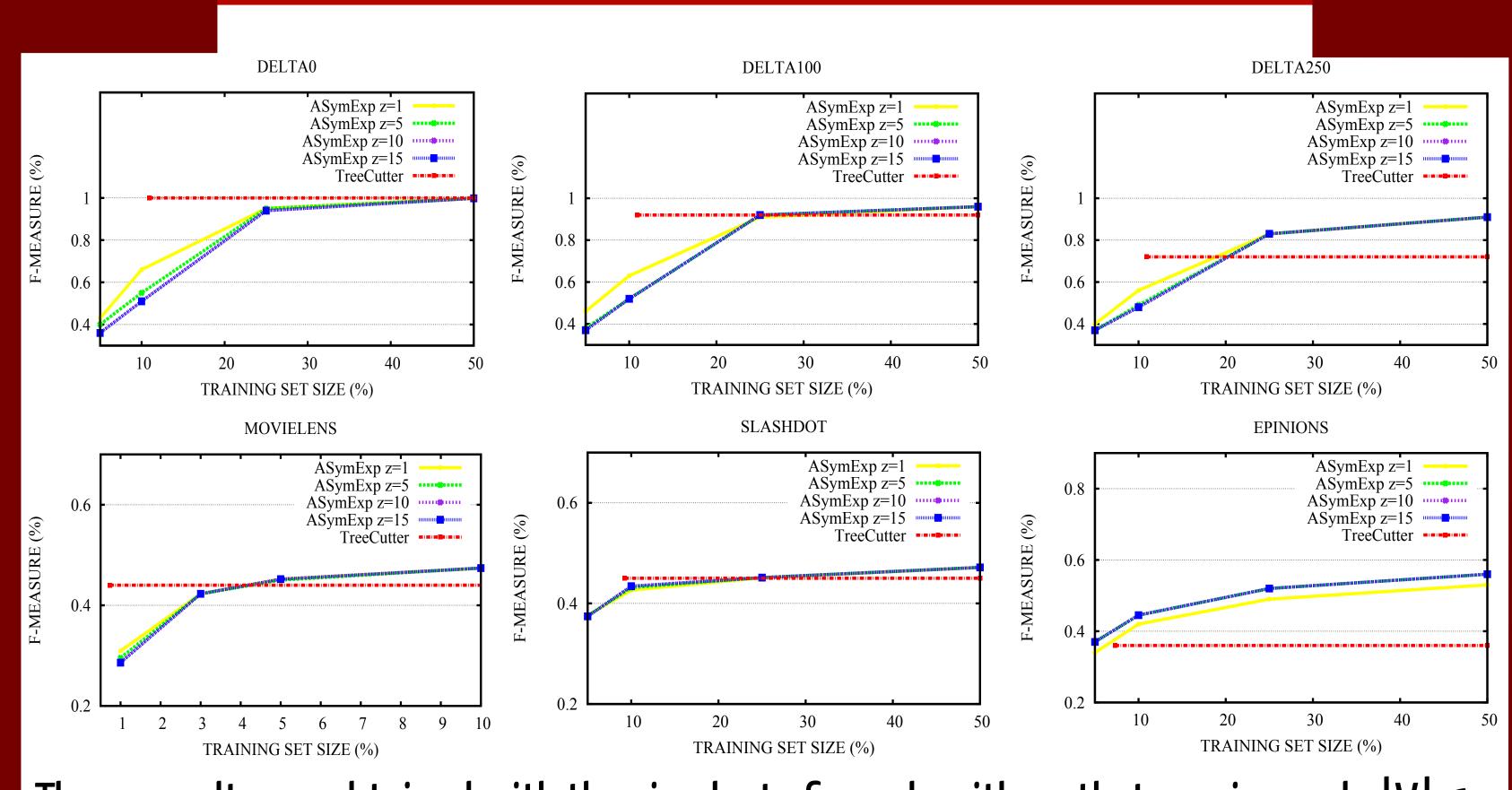
We present 3 algorithms with different tradeoffs between mistakes

and query set size

Running times

When the query set size is not larger than the test set size, the running times for our algorithms are:

O(|E|) for TreeCutter(k) $O(|E| + |V| \log |V|)$ for StarMaker $O(|E| + |V|/k \log |V|/k)$ for TreeLetStar(k)



These results are obtained with the simplest of our algorithms that requires only |V|-1 edges (a spanning tree of the graph) as query set.

Our algorithm is **faster** than its competitors, is **extremely simple to implement** and in most of the cases outperforms its competitors using just a fraction of their training set.

Short Bibliography

- [1] Kunegis et al. "The Slashdot Zoo: Mining a social network with negative edges" -- WWW 2009
- [2] Cesa-Bianchi et al. "A correlation clustering approach to link classification in signed networks" -- COLT 2012
- [3] Leskovec et al. "Predicting positive and negative links in online social networks" -- WWW 2010