

# Introduction to Data Structure



# Data Structures

- **Data** is just collection of facts or set of values that are in particular format.
- **Information** is the processed data.
- If the **data** is not organized effectively, it is very difficult to perform any task on **large** amount of **data**.
- A **data structure** is a particular way of organizing a large amount of data more efficiently in a computer so that any operation on that data becomes easy.
- **Data structures** is about rendering data elements in terms of some relationship, for better **performance**, **organization** and **storage**.



# Types of Data Structures

Data structures are divided into two types:

## 1. Primitive Data Structures (Built-In Data Structures)

- Primitive data structures are char, int, double and float.

## 2. Non-primitive Data Structures (User-defined Data Structures)

- Non-primitive data structures are used to store large and connected data. Some example of non-primitive data structures are: **Linked List, Tree, Graph, Stack and Queue.**

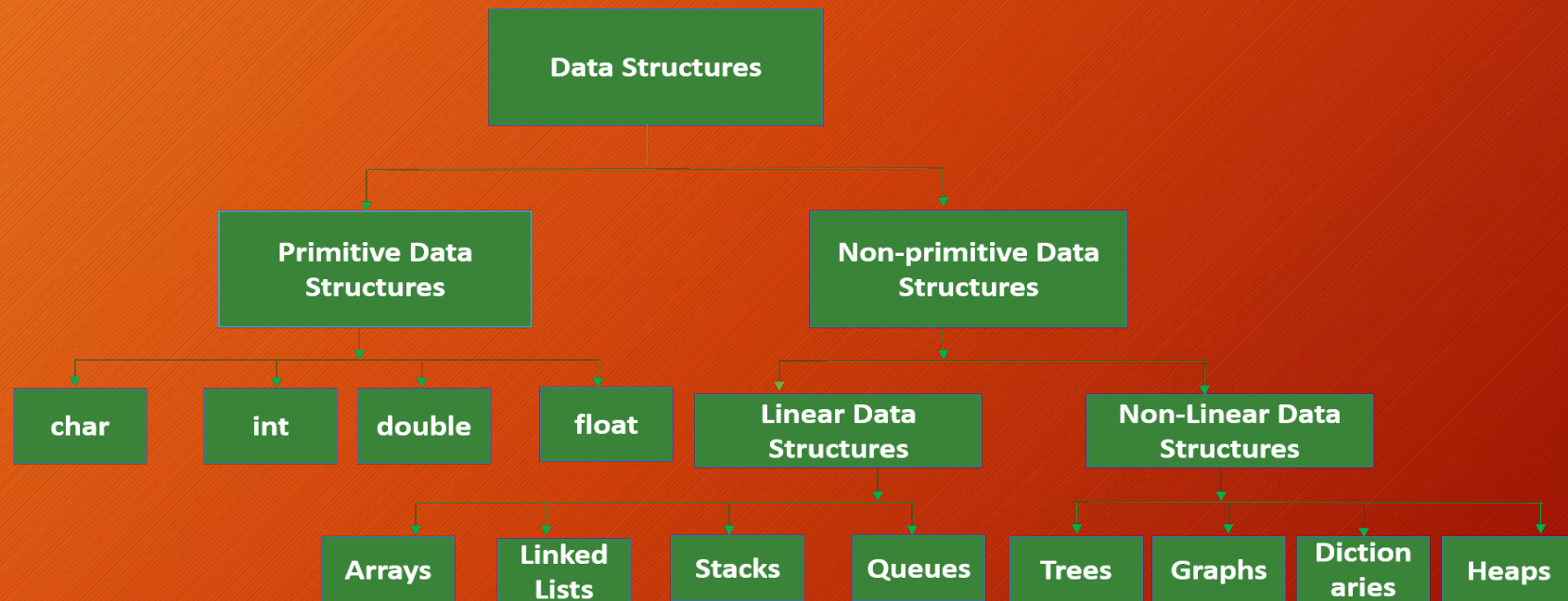


# Types of Data Structures

- The **non-primitive data structures** are subcategorized into two ways: **Linear data structures** and **Non-linear data structures**.
- If a data structure is organizing the data in **sequential order** then that data structure is called a linear data structure.
  - Some of the examples are Arrays, Linked Lists, Stacks and Queues.
- If a data structure is organizing the data in **random order** or **hierarchical order**, not in sequential order, then that data structure is called as non-linear data structure.
  - Some of the examples are Trees, Graphs, Dictionaries and Heaps.



# Types of Data Structures





# Introduction to Algorithm

- An Algorithm is a finite set of instructions or logic, written in order, to accomplish a certain predefined task.
- An **Algorithm** is independent of the programming language. An Algorithm is the core logic to solve a given problem.
- An **Algorithm** is expressed generally as **flow chart** or as an informal high level description called as **pseudocode Algorithm** can be defined as “a sequence of steps to be performed for getting the desired output for a given input.”
- There may be more than one way to solve a problem, hence there may be more than one **algorithms** for a given problem.



# Characteristics of an Algorithm

1. Input: Algorithm should be accepting 0 or more inputs supplied externally.
2. Output: Algorithm should be generating at least one output.
3. Definiteness: Each step of an algorithm must be precisely defined. Meaning the step should perform a clearly defined task without much complication.
4. Finiteness: An algorithm must always terminate after a finite number of steps.
5. Effectiveness: The efficiency of the steps and the accuracy of the output determine the effectiveness of the algorithm.
6. Correctness: Each step of the algorithm must generate a correct output.



# Efficiency of an Algorithm

- An algorithm is said to be efficient and fast, if it takes less time to execute and consumes less memory space.
- The performance of an algorithm is measured on the basis of following properties:
  - Time complexity
  - Space complexity



# Time Complexity

- Time complexity is a way to represent the amount of time required by the program to run till its completion.
- The amount of time taken by the program depends on lot of things like hardware, operating system, number of processors, processor architecture etc.
- We do not consider all these factors when analysing the algorithms, we only consider the number of operations to be executed for the completion of the algorithm.
- The reason is very simple because we want the algorithm analysis to be system independent.
- **Time complexity** is commonly estimated by counting the number of elementary operations performed by the algorithm, supposing that each elementary operation takes a fixed amount of time to perform.



# Time Complexity

- Let us consider the problem statement of the algorithm is "Find the sum of first n natural numbers".
- Algorithm - 1 :
- // Calculating the sum of first n natural numbers

```
int sum(int n)
{
    int i, total = 0;
    for (i = 1; i <= n; i++)
    {
        total = total + i;
    }
    return total;
}
```



# Time Complexity

Code	No. of Elementary Operations	No of Times Executed	Total Operations
int i, total = 0;	1 (One Assignment Operation)	1	$1 * 1 = 1$
for (i = 1; i <= n; i++) { ... }	1 (One Assignment Operation)	1	$1 * 1 = 1$
for (i = 1; i <= n; i++) { ... }	1 (One Comparison Operation)	n + 1 (n times it evaluates to true, 1 time it evaluates to false)	$1 * (n + 1) = n + 1$
for (i = 1; i <= n; i++) { ... }	1 (One Increment Operation)	n times (i is incremented n times)	$1 * n = n$
for (i = 1; i <= n; i++) { total = total + i; }	2 (One Addition operation, One Assignment operation)	n (n times as loop executes n times)	$2 * n = 2 * n$
return total;	1 (One return operation)	1	$1 * 1 = 1$
<b>Total Number of Operations</b>			...
<b>Time complexity of the algorithm <math>T(n) = 4 * n + 4</math> [ Assuming each operation completes in unit time ]</b>			



# Time Complexity

- The total number of elementary operation executed in the above algorithm is  $4 * n + 4$ .
- As the value of  $n$  increases the time taken will also increase linearly. Thus this algorithm has a **linear time complexity**.
- Algorithm - 2 :  
// Calculating the sum of first  $n$  natural numbers using formula  

```
int sum(int n)
{
    return n * (n + 1) / 2;
}
```



# Time Complexity

Code	No. of Elementary Operations	No of Times Executed	Total Operations
return $n * (n + 1) / 2$	4 (One multiplication, One Addition, One division, One retrun operation)	1	4
Total Number of Operations			4
Time complexity of the algorithm $T(n) = 4$ [Assuming each operation completes in unit time]			

For the above code, the time complexity is **constant**, because it will never be dependent on the value of  $n$ , it will always give the result in 4 elementary operations.

Thus the time complexity of the **Algorithm - 2** is **constant**. In other words, the running time of the algorithm does not change with the input  $n$ .

In the above two simple algorithms, you saw how a single problem can have many solutions



# Space Complexity

- Space complexity is the amount of memory space that is required by the algorithm for its execution.
- An algorithm generally requires space for following:
  - Instruction space: Space required to store the executable version (Also known as Binary Code) of the program. This space is fixed and depends on the number of instructions in the algorithm.
  - Data space: Space required to store all the program variables (Constants, Variables, Temporary Variables etc.).
  - Environment space: Space or Memory that is to be dynamically allocated. Memory that is required for storing data between functions
  - **One heap and one stack is maintained during the program execution.**
    - Heap is the segment where dynamic memory allocation usually takes place.
    - Stack is a segment where automatic variables and function call stack is stored.



# Space Complexity

- When we are measuring the space complexity analysis of an algorithm, we consider only the **data space** of the algorithm and ignore the **instruction space** and **environmental space**
- For calculating the **space complexity**, we need to know the amount of memory used by variables of different data types, which generally varies for different operating systems.
- The method for calculating the **space complexity** remains the same and independent of the operating system.



# Space Complexity

Data Type	Size
bool, char, unsigned char, signed char, __int8	1 byte
__int16, short, unsigned short, wchar_t, __wchar_t	2 bytes
float, __int32, int, unsigned int, long, unsigned long	4 bytes
double, __int64, long double, long long	8 bytes



# Calculating Constant Space Complexity

- Let us see the procedure for calculating the space complexity of a code with an example:

```
int sum(int x, int y)
{
    int z=x+y;
    return z;
}
```
- The above code takes two inputs x and y of type int as formal parameters.
- In the code, another local variable z of type int is used for storing the sum of x and y.
- The int data type takes 2 bytes of memory, so the total space complexity is 3 (number of variables) \* 2 (Size of each variable) = 6 bytes.
- The space requirement of this algorithm is fixed for any input given to the algorithm, hence it is called as **constant space complexity**



# Calculating Linear Space Complexity

- //Function to calculate a sum of n elements in the array //n is the number of elements in the array.

```
int sum(int a[ ], int n)
{
    int x = 0, i = 0;
    for (i = 0; i < n; i++)
    {
        x = x + a[i];
    }
    return x;
}
```



# Calculating Linear Space Complexity

- In the above code,  $2 * n$  bytes of memory (size of int data type is 2) is required by the array `a[ ]` and 2 bytes of memory for each variable of `x`, `n` and `i`.
- Hence the total space requirement for the above code would be  $(2 * n + 6)$ .
- The space complexity of the program is increasing linearly with the size of the array (input) `n` then it is called as **Linear Space Complexity**.
- Similarly, when the memory requirement of the algorithm increases quadratic to the given input then it is called as a "**Quadratic Space Complexity**".
- Similarly, when the memory requirement of the algorithm increases cubic to the given input then it is called as a "**Cubic Space Complexity**" and so on.