

CRACK A HACK

Winning Hand of Cards

(From RookieRank 4)

By :
Vishal M N
USN : 01FE15BCS233

Problem Statement

You are given a number of cards and try to create as many combinations from those cards as possible that result in a *winning hand*. A winning hand is the one where the product of the numbers on the cards modulo a given value, the *modulo divisor* is equal to another given value, the *target value*.

Solution

```
#include <iostream>
using namespace std;
typedef long long ll;

int main() {

    int n,m,x;
    cin>>n>>m>>x;
    ll a[n];
    for(int i=0;i<n;i++) cin>>a[i];

    ll count[m]={0},count1[m];

    for(int i=0;i<n;i++){ //loop 1
        for(int j=0;j<m;j++) count1[j] = count[j]; //loop 2

        for(int j=0;j<m;j++) count[((ll)j*a[i])%m] += count1[j];
        //loop 3

        count[(a[i])%m]+=1;
    }
    cout<<count[x];
    return 0;
}
```

Explanation

Using Brute force approach to generate all the combinations requires time complexity of $O(2^n)$.

Generally problems having exponential growth in terms of time complexity are solved using dynamic programming which has time complexity in terms of a polynomial equation.

The above problem can be solved using dynamic programming with a time complexity of $O(m*n)$.

N is the size of array elements.

M is the modulo divisor.

X is the target value.

The array `count[m]` stores the count of combinations that have remainder equal to 'j' where $0 \leq j < m$.

The array `count1[m]` stores count of combinations that remainder equal to 'j' where $0 \leq j < m$ of the previous state.

Loop1 is to select a value from the array.

Loop2 is to store the previous state of the `count[m]` array into `count1[m]` array.

Loop3 : In this array $0 \leq j < m$ and j represents remainder which used as index in this array.

Considering i'th value from the array, the result of $(j * A[i]) \% m$ is obtained and count array is updated as `count[(j * A[i]) \% m] += count1[j]`.

The operation $j * A[i]$ indicates obtaining combinations of all the values in the array `A[n]` that give j as remainder after modulo operation.

Addition of $\text{count1}[j]$ indicates inclusion of all the previous combinations that resulted into remainder $= j$.

Finally, the result that should be outputted is count of all combinations that resulted in remainder x after modulo operation.

Hence the final output is value in $\text{count}[x]$.