## **CRACK A HACK**

# Sherlock's Array Merging Algorithm

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### **Problem statement:-**

Watson gave Sherlock a collection of arrays V. Here each  $V_i$  is an array of variable length. It is guaranteed that if you merge the arrays into one single array, you'll get an array, M, of n distinct integers in the range [1, n].

Watson asks Sherlock to merge V into a sorted array. Sherlock is new to coding, but he accepts the challenge and writes the following algorithm:

- $M \leftarrow [\ ]$  (an empty array).
- $k \leftarrow$  number of arrays in the collection V.
- While there is at least one non-empty array in V:
  - $\bullet \ \ T \leftarrow [\ ] \ \text{(an empty array) and} \ i \leftarrow 1.$
  - While  $i \leq k$ :
    - If  $V_i$  is not empty:
      - ullet Remove the first element of  $V_i$  and push it to T.
    - $i \leftarrow i+1$ .
  - ullet While T is not empty:
    - ullet Remove the minimum element of T and push it to M.
- ullet Return M as the *output*.

#### **Explanation:-**

We need to use Dynamic Programming. First, we define  $\operatorname{cnt}(i)$  as count of rows with  $\operatorname{length} \geq i$ . For example, rows can be like this:

```
So, \operatorname{cnt}(3)=2 and \operatorname{cnt}(2)=3 and \operatorname{cnt}(1)=4. Note, that \operatorname{cnt}(i)\geq\operatorname{cnt}(i+1)\ \forall i.
```

Now, we will instead of building solution column wise, we store current  $\operatorname{cnt}(i)$  and also how many elements in array A left can be used.

f(i,j) is number of distinct matrices possible using elements  $A_1,A_2,\ldots,A_i$  if  $\operatorname{cnt}(\operatorname{last\ column})=j$ . Our answer will be  $\sum_{i=1}^n f(N-1,i)$ .

Let's see how we define recurrences. We want to build a column using j values  $A_{i-j+1}, A_{i-j}, \ldots, A_i$ . If this sequence is increasing(because matrices is column wise sorted) then we can choose these j elements to form a column, else there are 0 number of ways.

Now, if we have used j elements, state i, gets reduced to i-j. Also, for the next colmumn c we are going to build  $\operatorname{cnt}(c) \geq j$ .

```
//returns f(i, j)
def rec(i, j):
    ret=0;

for k = j; k <= i - j + 1; k++
    if(A[i - j + 1, i] is increasing){
        //now k is cnt(c), where c is the previous column
        //now, out of these k positions we can choose any j positions
        //and put any of the value from A[i - j + 1, i]
        //so we multiply number of ways by
        //choose(k, j) * factorial(j)
        int nextways = rec(i-j, k) * choose(k, j) * factorial(j);
        ret += nextways;
    }
}
return ret;</pre>
```

This is current  $O(N^3)$  solution. I couldn't think of anything better than this. But we should note the small constant factor helps us set constraints  $N=10^3$ . If you calculate it would turn out to be  $O(\frac{N^3}{12})$ .

#### CODE:-

```
import java.io.OutputStream;
import java.io.IOException;
import java.io.InputStream;
import java.io.OutputStream;
import java.io.PrintWriter;
import java.util.Arrays;
import java.io.BufferedWriter;
import java.io.IOException;
import java.io.InputStreamReader;
import java.util.StringTokenizer;
import java.io. Writer;
import java.io.OutputStreamWriter;
import java.io.BufferedReader;
import java.io.InputStream;
/**
* Built using CHelper plug-in
* Actual solution is at the top
*/
public class Solution {
  public static void main(String[] args) {
    InputStream inputStream = System.in;
    OutputStream outputStream = System.out;
    InputReader in = new InputReader(inputStream);
    OutputWriter out = new OutputWriter(outputStream);
    SherlocksArrayMergingAlgorithm solver = new SherlocksArrayMergingAlgorithm();
```

```
solver.solve(1, in, out);
  out.close();
}
static\,class\,Sherlocks Array Merging Algorithm\,\{
  public int mod = 1000000007;
  public boolean[][] inc;
  public long[] fact;
  public long[] ifact;
  public int n;
  public long[][] dp;
  public void solve(int testNumber, InputReader in, OutputWriter out) {
     n = in.nextInt();
     long[][] w = Factorials.getFIF(2 * n, mod);
    fact = w[0];
     ifact = w[1];
     int[] arr = in.readIntArray(n);
     inc = new boolean[n][n];
    for (int i = 0; i < n; i++) {
       inc[i][i] = true;
       for (int j = i + 1; j < n; j++) {
         inc[i][j] = inc[i][j - 1] && (arr[j] > arr[j - 1]);
       }
     }
     dp = new long[n][n];
    for (long[] x : dp) Arrays.fill(x, -1);
```

```
long ans = 0;
    for (int i = 0; i < n \&\& inc[0][i]; i++) {
       ans = (ans + dfs(i + 1, i + 1)) \% mod;
    }
    out.println(ans);
  }
  public long dfs(int curidx, int lastblock) {
    if (curidx == n) return 1;
    if (dp[curidx][lastblock] != -1) return dp[curidx][lastblock];
    long ret = 0;
    for (int next = curidx; next < n && next < curidx + lastblock && inc[curidx][next]; next++) {
       int cblock = (next - curidx + 1);
       ret = (ret + fact[lastblock] * ifact[lastblock - cblock] % mod * dfs(next + 1, cblock)) % mod;
    }
    return dp[curidx][lastblock] = ret;
  }
static class InputReader {
  public BufferedReader reader;
  public StringTokenizer tokenizer;
  public InputReader(InputStream stream) {
    reader = new BufferedReader(new InputStreamReader(stream), 32768);
    tokenizer = null;
```

}

```
}
public String next() {
  while (tokenizer == null || !tokenizer.hasMoreTokens()) {
    try {
       tokenizer = new StringTokenizer(reader.readLine());
    } catch (IOException e) {
       throw new RuntimeException(e);
    }
  }
  return tokenizer.nextToken();
}
public int[] readIntArray(int tokens) {
  int[] ret = new int[tokens];
  for (int i = 0; i < tokens; i++) {
     ret[i] = nextInt();
  }
  return ret;
}
public int nextInt() {
  return Integer.parseInt(next());
}
```

}

```
static class OutputWriter {
  private final PrintWriter writer;
  public OutputWriter(OutputStream outputStream) {
    writer = new PrintWriter(new BufferedWriter(new OutputStreamWriter(outputStream)));
  }
  public OutputWriter(Writer writer) {
    this.writer = new PrintWriter(writer);
  }
  public void close() {
    writer.dose();
  }
  public void println(long i) {
    writer.println(i);
  }
}
static class Factorials {
  public static long[][] getFIF(int max, int mod) {
    long[] fact = new long[max];
    long[] ifact = new long[max];
    long[] inv = new long[max];
    inv[1] = 1;
```

```
for (int i = 2; i < max; i++) {
    inv[i] = (mod - mod / i) * inv[mod % i] % mod;
}

fact[0] = 1;
ifact[0] = 1;
for (int i = 1; i < max; i++) {
    fact[i] = fact[i - 1] * i % mod;
    ifact[i] = ifact[i - 1] * inv[i] % mod;
}

return new long[][]{fact, ifact, inv};
}
</pre>
```

#### **REFENCES:**

Question :- <a href="https://www.hackerrank.com/contests/university-codesprint-2/challenges/sherlocks-array-merging-algorithm">https://www.hackerrank.com/contests/university-codesprint-2/challenges/sherlocks-array-merging-algorithm</a>

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