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   20
   9 Peppa
   1
     Basic
     createFile
 1.1
// Linux
 for i in {A..Z}; do cp tem.cpp $i.cpp; done
// Windows
'A'..'Z' | % { cp tem.cpp "$_.cpp" }
 1.2 run
g++ -std=c++20 -DPEPPA -Wall -Wextra -Wshadow -02 -fsanitize=
     address,undefined $1.cpp -o $1 && ./$1
 1.3 tem
#include <bits/stdc++.h>
 #pragma GCC optimize("Ofast,unroll-loops,no-stack-protector")
 #pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
 using namespace std;
 using i64 = long long;
 #define int i64
 #define all(a) a.begin(), a.end()
 #define rep(a, b, c) for (int a = b; a < c; a++)
 bool chmin(auto& a, auto b) { return (b < a and (a = b, true));</pre>
 bool chmax(auto& a, auto b) { return (a < b and (a = b, true));</pre>
     }
 void solve() {
  //
 }
 int32_t main() {
  std::ios::sync_with_stdio(false);
   std::cin.tie(nullptr);
   int t = 1;
  std::cin >> t;
  while (t--) {
    solve();
  return 0;
}
 1.4 debug
 #ifdef PEPPA
 template <typename R>
 concept I = ranges::range<R> && !std::same_as<ranges::</pre>
 range_value_t<R>, char>;
template <typename A, typename B>
 std::ostream& operator<<(std::ostream& o, const std::pair<A, B</pre>
     >6 p) {
  return o << "(" << p.first << ", " << p.second << ")";</pre>
 template <I T>
 std::ostream& operator<<(std::ostream& o, const T& v) {</pre>
  o << "{";
   int f = 0;
  for (auto &&i : v) o << (f++ ? " " : "") << i;
  return o << "}";
 void debug__(int c, auto&&... a) {
  std::cerr << "\e[1;" << c << "m";
(..., (std::cerr << a << ""));
  std::cerr << "\e[0m" << std::endl;</pre>
 }
 #define debug_(c, x...) debug__(c, __LINE__, "[" + std::string
```

1.7 TempleHash

```
| cat file.cpp | cpp -dD -P -fpreprocessed | tr -d "[:space:]" |
| md5sum | cut -c-6
```

2 Misc

2.1 FastIO

```
#include <unistd.h>
int OP;
char OB[65536]:
inline char RC() {
 static char buf[65536], *p = buf, *q = buf;
 return p == q && (q = (p = buf) + read(0, buf, 65536)) == buf
      ? -1: *p++;
inline int R() {
 static char c;
 while ((c = RC()) < '0');
int a = c ^ '0';</pre>
  while ((c = RC()) >= '0') a *= 10, a += c ^ '0';
  return a;
inline void W(int n) {
  static char buf[12], p;
  if (n == 0) OB[OP++] = '0';
 p = 0:
 while (n) buf[p++] = '0' + (n % 10), n /= 10;
  for (--p; p >= 0; --p) OB[OP++] = buf[p];
 if (OP > 65520) write(1, OB, OP), OP = 0;
// another FastIO
char buf[1 << 21], *p1 = buf, *p2 = buf;</pre>
inline char getc() {
  return p1 == p2 && (p2 = (p1 = buf) + fread(buf, 1, 1 << 21,
     stdin), p1 == p2) ? 0 : *p1++;
template<typename T> void Cin(T &a) {
 T res = 0; int f = 1;
  char c = getc();
  for (; c < '0' || c > '9'; c = getc()) {
  if (c == '-') f = -1;
  for (; c >= '0' && c <= '9'; c = getc()) {
   res = res * 10 + c - '0';
  a = f * res:
}
template<typename T, typename... Args> void Cin(T &a, Args &...
    args) {
  Cin(a), Cin(args...);
template<typename T> void Cout(T x) { // there's no '\n' in
    output
  if (x < 0) putchar('-'), x = -x;
  if (x > 9) Cout(x / 10);
 putchar(x % 10 + '0');
```

2.2 stress.sh

```
| #!/usr/bin/env bash
| g++ $1.cpp -o $1
| g++ $2.cpp -o $2
| g++ $3.cpp -o $3
| for i in {1..100} ; do
| ./$3 > input.txt
```

```
# st=$(date +%s%N)
  ./$1 < input.txt > output1.txt
  # echo "$((($(date +%s%N) - $st)/1000000))ms"
  ./$2 < input.txt > output2.txt
  if cmp --silent -- "output1.txt" "output2.txt" ; then
    continue
  fi
  echo Input:
  cat input.txt
  echo Your Output:
  cat output1.txt
  echo Correct Output:
  cat output2.txt
  exit 1
done
echo OK!
./stress.sh main good gen
```

2.3 stress.bat

@echo off

```
setlocal EnableExtensions
g++ -std=c++20 -03 "%1.cpp" -o "%1.exe"
g++ -std=c++20 -03 "%2.cpp" -0 "%2.exe"
g++ -std=c++20 -03 "%3.cpp" -o "%3.exe"
for /l %%i in (1,1,100) do (
 "%3.exe" > input.txt
 "%1.exe" < input.txt > output1.txt
 "%2.exe" < input.txt > output2.txt
 fc /b output1.txt output2.txt >nul
 if errorlevel 1 (
  echo Input:
  type input.txt
  echo Your Output:
  type output1.txt
  echo Correct Output:
  type output2.txt
  exit /b 1
)
@REM ./stress main good gen
```

2.4 Timer

```
| struct Timer {
   int t;
   bool enable = false;

   void start() {
      enable = true;
      t = std::clock();
   }
   int msecs() {
      assert(enable);
      return (std::clock() - t) * 1000 / CLOCKS_PER_SEC;
   }
};
```

2.5 MinPlusConvolution

```
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(vector<int> &a, vector<int> &b
    ) {
    int n = ssize(a), m = ssize(b);
    vector<int> c(n + m - 1, INF);
    auto dc = [&auto Y, int l, int r, int jl, int jr) {
        if (l > r) return;
        int mid = (l + r) / 2, from = -1, &best = c[mid];
        for (int j = jl; j <= jr; ++j)
        if (int i = mid - j; i >= 0 && i < n)
            if (best > a[i] + b[j])
            best = a[i] + b[j], from = j;
        Y(Y, l, mid - 1, jl, from), Y(Y, mid + 1, r, from, jr);
    };
    return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
}
```

2.6 PyTrick

```
import sys
input = sys.stdin.readline

from itertools import permutations
op = ['+', '-', '*', '']
a, b, c, d = input().split()
```

```
for (x,y,z,w) in permutations([a, b, c, d]):
 for op1 in op:
    for op2 in op:
      for op3 in op:
        val = eval(f''\{x\}\{op1\}\{y\}\{op2\}\{z\}\{op3\}\{w\}'')
        if (op1 == '' and op2 == '' and op3 == '') or val < 0:
          continue
        ans.add(val)
print(len(ans))
map(int,input().split())
arr2d = [ [ list(map(int,input().split())) ] for i in range(N)
from decimal import *
from fractions import *
s = input()
n = int(input())
f = Fraction(s)
g = Fraction(s).limit_denominator(n)
h = f * 2 - g
if h.numerator <= n and h.denominator <= n and h < g:</pre>
print(g.numerator, g.denominator)
from fractions import Fraction
x = Fraction(1, 2), y = Fraction(1)
print(x.as_integer_ratio()) # print 1/2
print(x.is_integer())
print(x.__round__())
print(float(x))
r = Fraction(input())
N = int(input())
r2 = r - 1 / Fraction(N) ** 2
ans = r.limit_denominator(N)
ans2 = r2.limit_denominator(N)
if ans2 < ans and 0 <= ans2 <= 1 and abs(ans - r) >= abs(ans2 -
     r):
 ans = ans2
print(ans.numerator,ans.denominator)
```

Data Structure

3.1 Fenwick Tree

```
template<class T>
struct Fenwick {
  int n;
   vector<T> a;
  Fenwick(\verb"int _n") : n(\_n), a(\_n) \{ \}
  void add(int p, T x) {
    for (int i = p; i < n; i = i | (i + 1)) {
      a[i] = a[i] + x;
    }
  T qry(int p) { // sum [0, p]
    T s{};
    for (int i = p; i >= 0; i = (i & (i + 1)) - 1) {
      s = s + a[i]:
    return s;
  T qry(int l, int r) { // sum [l, r)
    return qry(r - 1) - qry(l - 1);
  pair<int, T> select(T k) { // [first position >= k, sum [0, p
     )
    T s{}:
    int p = 0;
     for (int i = 1 << __lg(n); i; i >>= 1) {
      if (p + i \le n \text{ and } s + a[p + i - 1] \le k) {
        p += i;
         s = s + a[p - 1];
    return {p, s};
| }:
3.2 Li Chao
```

```
struct Line {
  // y = ax + b
 i64 a{0}, b{-inf<i64>};
```

```
i64 operator()(i64 x) {
     return a * x + b;
};
// max LiChao
struct Seg {
   int l. r
   Seg *ls{}, *rs{};
   Line f{};
   Seg(int l, int r) : l(l), r(r) {}
   void add(Line g) {
     int m = (l + r) / 2;
     if (g(m) > f(m)) {
       swap(g, f);
     if (g.b == -inf<i64> or r - l == 1) {
       return;
     if (g.a < f.a) {
       if (!ls) {
         ls = new Seg(l, m);
       ls->add(g);
     } else {
       if (!rs) {
         rs = new Seg(m, r);
       rs->add(g);
     }
   i64 qry(i64 x) {
     if (f.b == -inf<i64>) {
       return -inf<i64>;
     int m = (l + r) / 2;
     i64 y = f(x);
     if (x < m \text{ and } ls) {
       chmax(y, ls->qry(x));
     } else if (x >= m and rs) {
       chmax(y, rs->qry(x));
     return y;
   }
};
```

3.3 PBDS

map<int, int> odt;

};

odt[x] = it->second;

// do something

// initialize edges odt[1] and odt[n + 1] auto split = [&](const int &x) -> void {

const auto it = prev(odt.upper_bound(x));

for (; itl != itr; itl = odt.erase(itl)) {

auto merge = [8](const int 8l, const int 8r) -> void {

auto itl = odt.lower_bound(l), itr = odt.lower_bound(r + 1);

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename T> using RBT = tree<T, null_type, less<T>,
     rb_tree_tag, tree_order_statistics_node_update>;
.find_by_order(k) 回傳第 k 小的值 (based-0)
.order_of_key(k) 回傳有多少元素比 k 小
struct custom_hash {
  static uint64_t splitmix64(uint64_t x) {
    x += 0x9e3779b97f4a7c15;
    x = (x ^(x >> 30)) * 0xbf58476d1ce4e5b9;
    x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
    return x ^ (x >> 31);
   size_t operator()(uint64_t x) const {
     static const uint64_t FIXED_RANDOM = chrono::steady_clock::
     now().time_since_epoch().count();
     return splitmix64(x + FIXED_RANDOM);
// gp_hash_table<int, int, custom_hash> ss;
 3.4 ODT
```

```
// assign value to odt[l]
                                                                      void splay(Node *x) {
                                                                        while (nroot(x)) {
3.5 Sparse Table
                                                                          Node *y = x->p
                                                                          if (nroot(y)) {
template<class T>
                                                                            rotate(dir(x) == dir(y) ? y : x);
struct SparseTable{
  function<T(T, T)> F;
                                                                          rotate(x);
   vector<vector<T>> sp;
                                                                        }
  SparseTable(vector<T> &a, const auto &f) {
    F = f;
                                                                      Node *nth(Node *x, int k) {
    int n = a.size();
                                                                        assert(size(x) > k);
    sp.resize(n, vector<T>(__lg(n) + 1));
    for (int i = n - 1; i >= 0; i--) {
                                                                        while (true) {
                                                                          push(x);
       sp[i][0] = a[i];
                                                                          int left = size(x->ch[0]);
       for (int j = 1; i + (1 << j) <= n; j++) {
                                                                          if (left > k) {
         sp[i][j] = F(sp[i][j - 1], sp[i + (1 << j - 1)][j - 1])
                                                                            x = x->ch[0];
                                                                          } else if (left < k) {
   k -= left + 1;</pre>
      }
    }
  }
                                                                            x = x->ch[1];
                                                                          } else {
  T query(int l, int r) \{ // [l, r) \}
                                                                            break;
    int k = __lg(r - l);
    return F(sp[l][k], sp[r - (1 << k)][k]);</pre>
                                                                        }
  }
                                                                        splay(x);
};
                                                                        return x;
3.6 Splay
                                                                      Node *split(Node *x) {
struct Node {
                                                                        assert(x);
  Node *ch[2]{}, *p{};
                                                                        push(x);
  Info info{}, sum{};
                                                                        Node *l = x->ch[0];
                                                                        if (l) l->p = x->ch[0] = nullptr;
  Tag tag{};
  int size{};
                                                                        pull(x);
                                                                        return l;
  bool rev{};
} pool[int(1E5 + 10)], *top = pool;
Node *newNode(Info a) {
                                                                      Node *join(Node *x, Node *y) {
  Node *t = top++;
                                                                        if (!x or !y) return x ? x : y;
   t->info = t->sum = a;
                                                                        y = nth(y, 0);
  t->size = 1;
                                                                        push(y);
  return t;
                                                                        y->ch[0] = x;
                                                                        if (x) x->p = y;
int size(const Node *x) { return x ? x->size : 0; }
                                                                        pull(y);
Info get(const Node *x) { return x ? x->sum : Info{}; }
                                                                        return y;
int dir(const Node *x) { return x->p->ch[1] == x; }
bool nroot(const Node *x) { return x->p and x->p->ch[dir(x)] ==
                                                                      Node *find_first(Node *x, auto &&pred) {
                                                                        Info pre{};
void reverse(Node *x) { if (x) x->rev = !x->rev; }
                                                                        while (true) {
void update(Node *x, const Tag &f) {
                                                                          push(x);
  if (!x) return;
                                                                          if (pred(pre + get(x->ch[0]))) {
  f(x->tag);
                                                                            x = x->ch[0];
  f(x->info);
                                                                          } else if (pred(pre + get(x->ch[0]) + x->info) or !x->ch
  f(x->sum);
                                                                           [1]) {
                                                                            break:
void push(Node *x) {
                                                                          } else {
  if (x->rev) {
                                                                            pre = pre + get(x->ch[0]) + x->info;
     swap(x->ch[0], x->ch[1]);
                                                                            x = x->ch[1];
    reverse(x->ch[0]);
                                                                        }
    reverse(x->ch[1]);
    x->rev = false;
                                                                        splay(x);
                                                                        return x;
                                                                     }
  update(x->ch[0], x->tag);
  update(x->ch[1], x->tag);
                                                                      3.7 Treap
   x->tag = Tag{};
                                                                      struct Treap {
                                                                        Treap *l, *r;
void pull(Node *x) {
                                                                        int key, size;
  x \rightarrow size = size(x \rightarrow ch[0]) + 1 + size(x \rightarrow ch[1]);
                                                                        Treap(int k) : l(nullptr), r(nullptr), key(k), size(1) {}
  x->sum = get(x->ch[0]) + x->info + get(x->ch[1]);
                                                                        void pull();
                                                                        void push() {};
void rotate(Node *x) {
  Node *y = x->p, *z = y->p;
                                                                      };
                                                                      inline int SZ(Treap *p) {
  push(y);
                                                                        return p == nullptr ? 0 : p->size;
  int d = dir(x);
                                                                      }
  push(x);
                                                                      void Treap::pull() {
  Node *w = x - > ch[d ^ 1];
                                                                        size = 1 + SZ(1) + SZ(r);
  if (nroot(y)) {
    z \rightarrow ch[dir(y)] = x;
                                                                      Treap *merge(Treap *a, Treap *b) {
  }
                                                                        if (!a || !b) return a ? a : b;
  if (w) {
                                                                        if (rand() % (SZ(a) + SZ(b)) < SZ(a)) {
    w->p = y;
                                                                          return a->push(), a->r = merge(a->r, b), a->pull(), a;
  (x->ch[d ^1] = y)->ch[d] = w;
  (y->p = x)->p = z;
                                                                        return b->push(), b->l = merge(a, b->l), b->pull(), b;
  pull(y);
                                                                     // <= k, > k
  pull(x);
```

```
void split(Treap *p, Treap *&a, Treap *&b, int k) { // by key
  if (!p) return a = b = nullptr, void();
  p->push();
  if (p->key <= k) {
    a = p, split(p->r, a->r, b, k), a->pull();
  } else {
    b = p, split(p->l, a, b->l, k), b->pull();
  }
void split2(Treap *p, Treap *&a, Treap *&b, int k) { // by size
  if (!p) return a = b = nullptr, void();
  p->push();
  if (SZ(p->l) + 1 <= k) {
    a = p, split2(p->r, a->r, b, k - SZ(p->l) - 1);
  } else {
    b = p, split2(p->l, a, b->l, k);
  p->pull();
}
void insert(Treap *&p, int k) {
  Treap *l, *r;
  p->push(), split(p, l, r, k);
  p = merge(merge(l, new Treap(k)), r);
  p->pull();
bool erase(Treap *&p, int k) {
  if (!p) return false;
  if (p->key == k) {
    Treap *t = p;
    p->push(), p = merge(p->l, p->r);
    delete t;
    return true;
  Treap *&t = k < p->key ? p->l : p->r;
  return erase(t, k) ? p->pull(), true : false;
int Rank(Treap *p, int k) { // # of key < k</pre>
  if (!p) return 0;
  if (p->key < k) return SZ(p->l) + 1 + Rank(p->r, k);
  return Rank(p->l, k);
Treap *kth(Treap *p, int k) { // 1-base
  if (k <= SZ(p->l)) return kth(p->l, k);
  if (k == SZ(p\rightarrow l) + 1) return p;
  return kth(p->r, k - SZ(p->l) - 1);
// pref: kth(Rank(x)), succ: kth(Rank(x+1)+1)
tuple<Treap*, Treap*, Treap*> interval(Treap *&o, int l, int r)
      { // 1-based
  Treap *a, *b, *c; // b: [l, r]
  split2(o, a, b, l - 1), split2(b, b, c, r - l + 1);
  return make_tuple(a, b, c);
```

4 Matching and Flow

4.1 Dinic

```
template <typename T>
struct Dinic {
  const T INF = numeric_limits<T>::max() / 2;
  struct edge {
    int v, r; T rc;
  vector<vector<edge>> adi:
  vector<T> dis, it;
  Dinic(int n) : adj(n), dis(n), it(n) {}
  void add_edge(int u, int v, T c) {
    adj[u].pb({v, adj[v].size(), c});
    adj[v].pb({u, adj[u].size() - 1, 0});
  bool bfs(int s, int t) {
    fill(all(dis), INF);
    queue<int> q;
    q.push(s);
    dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (const auto& [v, r, rc] : adj[u]) {
        if (dis[v] < INF || rc == 0) continue;</pre>
        dis[v] = dis[u] + 1;
        q.push(v);
```

```
}
     return dis[t] < INF;</pre>
   T dfs(int u, int t, T cap) {
     if (u == t || cap == 0) return cap;
     for (int &i = it[u]; i < (int)adj[u].size(); ++i) {</pre>
       auto &[v, r, rc] = adj[u][i];
       if (dis[v] != dis[u] + 1) continue;
       T tmp = dfs(v, t, min(cap, rc));
       if (tmp > 0) {
  rc -= tmp;
         adj[v][r].rc += tmp;
         return tmp:
       }
     return 0;
   T flow(int s, int t) {
  T ans = 0, tmp;
     while (bfs(s, t)) {
       fill(all(it), 0);
       while ((tmp = dfs(s, t, INF)) > 0) {
         ans += tmp;
       }
     return ans;
   bool inScut(int u) { return dis[u] < INF; }</pre>
};
         General Matching
struct GeneralMatching { // n <= 500</pre>
   const int BLOCK = 10;
```

```
int n;
  vector<vector<int> > g;
  vector<int> hit, mat;
  priority_queue<pair<int, int>, vector<pair<int, int>>,
     greater<pair<int, int>>> unmat;
  g[a].push_back(b);
     g[b].push_back(a);
  int get_match() {
     for (int i = 0; i < n; i++) if (!g[i].empty()) {</pre>
      unmat.emplace(0, i);
     // If WA, increase this
     // there are some cases that need >=1.3*n^2 steps for BLOCK
     =1
     // no idea what the actual bound needed here is.
     const int MAX_STEPS = 10 + 2 * n + n * n / BLOCK / 2;
     mt19937 rng(random_device{}());
     for (int i = 0; i < MAX_STEPS; ++i) {</pre>
       if (unmat.empty()) break;
      int u = unmat.top().second;
       unmat.pop();
       if (mat[u] != -1) continue;
       for (int j = 0; j < BLOCK; j++) {</pre>
        ++hit[u];
        auto &e = g[u];
        const int v = e[rng() % e.size()];
        mat[u] = v;
        swap(u, mat[v]);
        if (u == -1) break;
      if (u != -1) {
        mat[u] = -1;
        unmat.emplace(hit[u] * 100ULL / (g[u].size() + 1), u);
     int siz = 0;
     for (auto e : mat) siz += (e != -1);
     return siz / 2;
|};
```

4.3 KM

template<class T>

```
T KM(const vector<vector<T>> &w) {
  const T INF = numeric_limits<T>::max() / 2;
  const int n = w.size();
  vector<T> lx(n), ly(n);
  vector<int> mx(n, -1), my(n, -1), pa(n);
  auto augment = [8](int y) {
    for (int x, z; y != -1; y = z) {
      x = pa[y];
       z = mx[x];
      my[y] = x;
      mx[x] = y;
    }
  }:
  auto bfs = [8](int s) {
    vector<T> sy(n, INF);
    vector<bool> vx(n), vy(n);
    queue<int> q;
    q.push(s);
    while (true) {
      while (q.size()) {
         int x = q.front();
         q.pop();
         vx[x] = 1;
         for (int y = 0; y < n; y++) {</pre>
           if (vy[y]) continue;
          T d = lx[x] + ly[y] - w[x][y];
           if (d == 0) {
             pa[y] = x;
             if (my[y] == -1) {
               augment(y);
               return;
             vy[y] = 1;
             q.push(my[y]);
           } else if (chmin(sy[y], d)) {
             pa[y] = x;
        }
       T cut = INF;
       for (int y = 0; y < n; y++)
         if (!vy[y])
          chmin(cut, sy[y]);
       for (int j = 0; j < n; j++) {
         if (vx[j]) lx[j] -= cut;
         if (vy[j]) ly[j] += cut;
         else sy[j] -= cut;
       for (int y = 0; y < n; y++)
         if (!vy[y] \text{ and } sy[y] == 0) {
           if (my[y] == -1) {
             augment(y);
             return;
           vy[y] = 1;
          q.push(my[y]);
    }
  }:
  for (int x = 0; x < n; x++)
    lx[x] = ranges::max(w[x]);
  for (int x = 0; x < n; x++)
    bfs(x);
  T ans = 0;
  for (int x = 0; x < n; x++)
    ans += w[x][mx[x]];
  return ans;
4.4 MCMF
```

```
template<class T>
struct MCMF {
  const T INF = numeric_limits<T>::max() / 2;
  struct edge { int v, r; T f, w; };
  vector<vector<edge>> adj;
  const int n;
  MCMF(int n) : n(n), adj(n) {}
  void addEdge(int u, int v, T f, T c) {
    adj[u].push_back({v, ssize(adj[v]), f, c});
adj[v].push_back({u, ssize(adj[u]) - 1, 0, -c});
  vector<T> dis;
  vector<bool> vis;
```

```
bool spfa(int s, int t) {
     queue<int> que;
     dis.assign(n, INF);
     vis.assign(n, false);
     que.push(s);
     vis[s] = 1;
     dis[s] = 0;
     while (!que.empty()) {
       int u = que.front(); que.pop();
       vis[u] = 0;
       for (auto [v, _, f, w] : adj[u])
  if (f && chmin(dis[v], dis[u] + w))
            if (!vis[v]) {
             que.push(v);
              vis[v] = 1;
     return dis[t] != INF;
   T dfs(int u, T in, int t) {
     if (u == t) return in;
     vis[u] = 1;
     T out = 0;
     for (auto &[v, rev, f, w] : adj[u])
       if (f && !vis[v] && dis[v] == dis[u] + w) {
         T x = dfs(v, min(in, f), t);
         in -= x;
         out += x;
         adj[v][rev].f += x;
         if (!in) break;
     if (in) dis[u] = INF;
     vis[u] = 0;
     return out;
   pair<T, T> flow(int s, int t) { // {flow, cost}
     T a = 0, b = 0;
     while (spfa(s, t)) {
       T x = dfs(s, INF, t);
       a += x;
       b += x * dis[t];
     return {a, b};
};
```

4.5 Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.

 - 2. For each edge (x,y,l,u), connect $x\to y$ with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph(X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X 3. $x \in X$ is chosen iff x is unvisited.

 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y)by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph

 - 1. Binary search on answer, suppose we're checking answer T 2. Construct a max flow model, let K be the sum of all weights 3. Connect source $s \to v, \, v \in G$ with capacity K

```
4. For each edge (u,v,w) in G, connect u \to v and v \to u with capacity
      5. For v \in {\it G}, connect it with sink v \to t with capacity K + 2T -
         \left(\sum_{e \in E(v)} w(e)\right) - 2w(v)
      6. T is a valid answer if the maximum flow f < K|V|
· Minimum weight edge cover
      1. Change the weight of each edge to \mu(u) + \mu(v) - w(u, v), where
         \mu(v) is the cost of the cheapest edge incident to v.
      2. Let the maximum weight matching of the graph be x, the answer will
         be \sum \mu(v) - x.
```

Geometry

5.1 Point

```
using numbers::pi;
template<class T> inline constexpr T eps = numeric_limits<T>::
     epsilon() * 1E6;
using Real = long double;
struct Pt {
  Real x\{\}, y\{\};
  Pt operator+(Pt a) const { return {x + a.x, y + a.y}; }
  Pt operator-(Pt a) const { return {x - a.x, y - a.y}; }
  Pt operator*(Real k) const { return {x * k, y * k}; }
  Pt operator/(Real k) const { return {x / k, y / k}; }
  Real operator*(Pt a) const { return x * a.x + y * a.y; }
  Real operator^(Pt a) const { return x * a.y - y * a.x; }
  auto operator<=>(const Pt&) const = default;
  bool operator==(const Pt&) const = default;
};
int sgn(Real x) { return (x > -eps<Real>) - (x < eps<Real>); }
Real ori(Pt a, Pt b, Pt c) { return (b - a) ^ (c - a); }
bool argcmp(const Pt &a, const Pt &b) { // arg(a) < arg(b)
  int f = (Pt{a.y, -a.x} > Pt{} ? 1 : -1) * (a != Pt{});
  int g = (Pt{b.y, -b.x} > Pt{} ? 1 : -1) * (b != Pt{});
  return f == g ? (a ^ b) > 0 : f < g;
Pt rotate(Pt u) { return {-u.y, u.x}; }
Real abs2(Pt a) { return a * a; }
// floating point only
Pt rotate(Pt u, Real a) {
  Pt v{sinl(a), cosl(a)};
  return {u ^ v, u * v};
Real abs(Pt a) { return sqrtl(a * a); }
Real arg(Pt x) { return atan2l(x.y, x.x); }
Pt unit(Pt x) { return x / abs(x); }
5.2 Line
struct Line {
  Pt a, b;
  Pt dir() const { return b - a; }
};
int PtSide(Pt p, Line L) {
  return sgn(ori(L.a, L.b, p)); // for int
  return sgn(ori(L.a, L.b, p) / abs(L.a - L.b));
bool PtOnSeg(Pt p, Line L) {
  return PtSide(p, L) == 0 and sgn((p - L.a) * (p - L.b)) <= <math>0;
Pt proj(Pt p, Line l) {
  Pt dir = unit(l.b - l.a);
  return l.a + dir * (dir * (p - l.a));
5.3 Circle
struct Cir {
  Pt o;
  double r:
bool disjunct(const Cir &a, const Cir &b) {
  return sgn(abs(a.o - b.o) - a.r - b.r) >= 0;
bool contain(const Cir &a, const Cir &b) {
  return sgn(a.r - b.r - abs(a.o - b.o)) >= 0;
```

5.4 Point to Segment Distance

```
double PtSegDist(Pt p, Line l) {
  double ans = min(abs(p - l.a), abs(p - l.b));
  if (sgn(abs(l.a - l.b)) == 0) return ans;
  if (sgn((l.a - l.b) * (p - l.b)) < 0) return ans;</pre>
 if (sgn((l.b - l.a) * (p - l.a)) < 0) return ans;</pre>
```

```
double SegDist(Line l, Line m) {
  return PtSegDist({0, 0}, {l.a - m.a, l.b - m.b});
}
```

return min(ans, abs(ori(p, l.a, l.b)) / abs(l.a - l.b));

5.5 Point In Polygon

```
int inPoly(Pt p, const vector<Pt> &P) {
   const int n = P.size();
   int cnt = 0;
   for (int i = 0; i < n; i++) {
     Pt a = P[i], b = P[(i + 1) \% n];
     if (PtOnSeg(p, {a, b})) return 1; // on edge
     if ((sgn(a.y - p.y) == 1) ^ (sgn(b.y - p.y) == 1))
       cnt += sgn(ori(a, b, p));
  return cnt == 0 ? 0 : 2; // out, in
```

5.6 Intersection of Line

```
bool isInter(Line l, Line m) {
  if (PtOnSeg(m.a, l) or PtOnSeg(m.b, l) or
    PtOnSeg(l.a, m) or PtOnSeg(l.b, m))
    return true
  return PtSide(m.a, l) * PtSide(m.b, l) < 0 and
      PtSide(l.a, m) * PtSide(l.b, m) < 0;
Pt LineInter(Line l, Line m) {
  double s = ori(m.a, m.b, l.a), t = ori(m.a, m.b, l.b);
  return (l.b * s - l.a * t) / (s - t);
bool strictInter(Line l, Line m) {
  int la = PtSide(m.a, l);
  int lb = PtSide(m.b, l);
  int ma = PtSide(l.a, m);
  int mb = PtSide(l.b, m);
  if (la == 0 and lb == 0) return false;
  return la * lb < 0 and ma * mb < 0;</pre>
```

5.7 Intersection of Circles

```
| vector<Pt> CircleInter(Cir a, Cir b) {
   double d2 = abs2(a.o - b.o), d = sqrt(d2);
   if (d < max(a.r, b.r) - min(a.r, b.r) || d > a.r + b.r)
     return {};
   Pt u = (a.o + b.o) / 2 + (a.o - b.o) * ((b.r * b.r - a.r * a.
     r) / (2 * d2));
   double A = sqrt((a.r + b.r + d) * (a.r - b.r + d) * (a.r + b.
     r - d) * (-a.r + b.r + d));
   Pt v = rotate(b.o - a.o) * A / (2 * d2);
   if (sgn(v.x) == 0 \text{ and } sgn(v.y) == 0) \text{ return } \{u\};
   return {u - v, u + v}; // counter clockwise of a
```

5.8 Intersection of Circle and Line

```
vector<Pt> CircleLineInter(Cir c, Line l) {
   Pt H = proj(c.o, l);
   Pt dir = unit(l.b - l.a);
   double h = abs(H - c.o);
   if (sgn(h - c.r) > 0) return {};
   double d = sqrt(max((double)0., c.r * c.r - h * h));
   if (sgn(d) == 0) return {H};
   return {H - dir *d, H + dir * d};
   // Counterclockwise
}
```

5.9 Area of Circle Polygon

```
| double CirclePoly(Cir C, const vector<Pt> &P) {
   auto arg = [\delta](Pt p, Pt q) \{ return atan2(p ^ q, p * q); \};
   double r2 = C.r * C.r / 2;
   auto tri = [8](Pt p, Pt q) {
     Pt d = q - p;
     auto a = (d * p) / abs2(d), b = (abs2(p) - C.r * C.r)/ abs2
     (d);
     auto det = a * a - b;
     if (det <= 0) return arg(p, q) * r2;</pre>
     auto s = max(0., -a - sqrt(det)), t = min(1., -a + sqrt(det))
     ));
     if (t < 0 or 1 <= s) return arg(p, q) * r2;</pre>
     Pt u = p + d * s, v = p + d * t;
     return arg(p, u) * r2 + (u ^ v) / 2 + arg(v, q) * r2;
  };
```

if (r < l) r += n;

int s = PtSide(A[l % n], L);

```
return *ranges::partition_point(views::iota(l, r),
  double sum = 0.0;
  for (int i = 0; i < P.size(); i++)</pre>
                                                                          [8](int m) {
  sum += tri(P[i] - C.o, P[(i + 1) % P.size()] - C.o);
                                                                            return PtSide(A[m % n], L) == s;
  return sum;
                                                                          }) - 1;
                                                                      };
                                                                      // Line A_x A_x+1 interset with L
5.10 Convex Hull
                                                                      vector<int> intersect(Line L) {
                                                                       int l = tangent(L.a - L.b). r = tangent(L.b - L.a):
vector<Pt> BuildHull(vector<Pt> pt) {
                                                                        if (PtSide(A[l], L) * PtSide(A[r], L) >= 0) return {};
 sort(all(pt)):
                                                                        return {find(l, r, L) % n, find(r, l, L) % n};
  pt.erase(unique(all(pt)), pt.end());
                                                                     }
  if (pt.size() <= 2) return pt;</pre>
  vector<Pt> hull;
                                                                  };
  int sz = 1;
  rep (t, 0, 2) {
                                                                    5.12 Half Plane Intersection
    rep (i, t, ssize(pt)) {
                                                                   bool cover(Line L, Line P, Line Q) {
      while (ssize(hull) > sz && ori(hull.end()[-2], pt[i],
                                                                      // for double, i128 => Real
    hull.back()) >= 0)
                                                                      i128 u = (Q.a - P.a) ^ Q.dir();
i128 v = P.dir() ^ Q.dir();
        hull.pop_back();
      hull.pb(pt[i]);
                                                                      i128 x = P.dir().x * u + (P.a - L.a).x * v;
    }
                                                                      i128 y = P.dir().y * u + (P.a - L.a).y * v;
    sz = ssize(hull);
                                                                      return sgn(x * L.dir().y - y * L.dir().x) * sgn(v) >= 0;
   reverse(all(pt));
                                                                   vector<Line> HPI(vector<Line> P) {
 hull.pop_back();
                                                                      sort(all(P), [&](Line l, Line m) {
  return hull;
                                                                        if (argcmp(l.dir(), m.dir())) return true;
                                                                        if (argcmp(m.dir(), l.dir())) return false;
                                                                        return ori(m.a, m.b, l.a) > 0;
5.11 Convex Trick
                                                                      });
struct Convex {
                                                                      int n = P.size(), l = 0, r = -1;
  int n:
                                                                      for (int i = 0; i < n; i++) {</pre>
  vector<Pt> A, V, L, U;
                                                                        if (i and !argcmp(P[i - 1].dir(), P[i].dir())) continue;
  Convex(const vector<Pt> \delta_A) : A(_A), n(_A.size()) { // n >=
                                                                        while (l < r \text{ and } cover(P[i], P[r - 1], P[r])) r--;
                                                                        while (l < r and cover(P[i], P[l], P[l + 1])) l++;
    auto it = max_element(all(A));
                                                                       P[++r] = P[i];
    L.assign(A.begin(), it + 1);
    U.assign(it, A.end()), U.push_back(A[0]);
                                                                     while (l < r and cover(P[l], P[r - 1], P[r])) r--;
    rep (i, 0, n) {
                                                                      while (l < r and cover(P[r], P[l], P[l + 1])) l++;
      V.push_back(A[(i + 1) % n] - A[i]);
                                                                      if (r - l <= 1 or !argcmp(P[l].dir(), P[r].dir()))</pre>
    }
                                                                        return {}; // empty
                                                                      if (cover(P[l + 1], P[l], P[r]))
  int inside(Pt p, const vector<Pt> &h, auto f) {
                                                                       return {}; // infinity
    auto it = lower_bound(all(h), p, f);
                                                                      return vector(P.begin() + l, P.begin() + r + 1);
    if (it == h.end()) return 0;
    if (it == h.begin()) return p == *it;
    return 1 - sgn(ori(*prev(it), p, *it));
                                                                    5.13 Minimal Enclosing Circle
                                                                   struct Cir {
  // 0: out, 1: on, 2: in
                                                                      Pt o;
  int inside(Pt p) {
                                                                      double r;
   return min(inside(p, L, less{}), inside(p, U, greater{}));
                                                                      bool inside(Pt p) {
                                                                        return sgn(r - abs(p - o)) >= 0;
  static bool cmp(Pt a, Pt b) { return sgn(a ^ b) > 0; }
  // A[i] is a far/closer tangent point
                                                                   };
  int tangent(Pt v, bool close = true) {
    assert(v != Pt{});
                                                                   Pt Center(Pt a, Pt b, Pt c) {
    auto l = V.begin(), r = V.begin() + L.size() - 1;
                                                                     Pt x = (a + b) / 2;
    if (v < Pt{}) l = r, r = V.end();</pre>
                                                                      Pt v = (b + c) / 2;
    if (close) return (lower_bound(l, r, v, cmp) - V.begin()) %
                                                                      return LineInter({x, x + rotate(b - a)}, {y, y + rotate(c - b
                                                                        )});
    return (upper_bound(l, r, v, cmp) - V.begin()) % n;
 }
                                                                   Cir MEC(vector<Pt> P) {
  // closer tangent point
                                                                      mt19937 rng(time(0));
  array<int, 2> tangent2(Pt p) {
                                                                      shuffle(all(P), rng);
    array<int, 2> t{-1, -1};
                                                                      Cir C{};
    if (inside(p) == 2) return t;
                                                                      for (int i = 0; i < P.size(); i++) {</pre>
    if (auto it = lower_bound(all(L), p); it != L.end() and p
                                                                        if (C.inside(P[i])) continue;
    == *it) {
                                                                        C = {P[i], 0};
      int s = it - L.begin();
                                                                        for (int j = 0; j < i; j++) {
      return {(s + 1) % n, (s - 1 + n) % n};
                                                                          if (C.inside(P[j])) continue;
                                                                          C = \{(P[i] + P[j]) / 2, abs(P[i] - P[j]) / 2\};
    if (auto it = lower_bound(all(U), p, greater{}); it != U.
                                                                          for (int k = 0; k < j; k++) {
    end() and p == *it) {
                                                                            if (C.inside(P[k])) continue
      int s = it - U.begin() + L.size() - 1;
                                                                            C.o = Center(P[i], P[j], P[k]);
      return {(s + 1) % n, (s - 1 + n) % n};
                                                                            C.r = abs(C.o - P[i]);
    for (int i = 0; i != t[0]; i = tangent((A[t[0] = i] - p),
                                                                       }
    0));
    for (int i = 0; i != t[1]; i = tangent((p - A[t[1] = i]),
                                                                      return C;
    1)):
    return t;
                                                                    5.14 Minkowski
  int find(int l, int r, Line L) {
```

|// P, Q, R(return) are counterclockwise order convex polygon

vector<Pt> Minkowski(vector<Pt> P, vector<Pt> Q) {

```
assert(P.size() >= 2 && Q.size() >= 2);
auto cmp = [&](Pt a, Pt b) {
    return Pt{a.y, a.x} < Pt{b.y, b.x};
};
auto reorder = [&](auto &R) {
    rotate(R.begin(), min_element(all(R), cmp), R.end());
    R.push_back(R[0]), R.push_back(R[1]);
};
const int n = P.size(), m = Q.size();
reorder(P), reorder(Q);
vector<Pt> R;
for (int i = 0, j = 0, s; i < n || j < m; ) {
    R.push_back(P[i] + Q[j]);
    s = sgn((P[i + 1] - P[i]) ^ (Q[j + 1] - Q[j]));
    if (s >= 0) i++;
    if (s <= 0) j++;
}
return R; // May not be a strict convexhull
}</pre>
```

5.15 Point In Circumcircle

```
|// p[0], p[1], p[2] should be counterclockwise order
|int inCC(const array<Pt, 3> &p, Pt a) {
| i128 det = 0;
| for (int i = 0; i < 3; i++)
| det += i128(abs2(p[i]) - abs2(a)) * ori(a, p[(i + 1) % 3],
| p[(i + 2) % 3]);
| return (det > 0) - (det < 0); // in:1, on:0, out:-1
|}</pre>
```

5.16 Tangent Lines of Circle and Point

```
vector<Line> CircleTangent(Cir c, Pt p) {
  vector<Line> z;
  double d = abs(p - c.o);
  if (sgn(d - c.r) == 0) {
    Pt i = rotate(p - c.o);
    z.push_back({p, p + i});
  } else if (d > c.r) {
    double o = acos(c.r / d);
    Pt i = unit(p - c.o);
    Pt j = rotate(i, o) * c.r;
    Pt k = rotate(i, -o) * c.r;
    z.push_back({c.o + j, p});
    z.push_back({c.o + k, p});
  }
  return z;
}
```

5.17 Tangent Lines of Circles

```
vector<Line> CircleTangent(Cir c1, Cir c2, int sign1) {
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
 double d sq = abs2(c1.o - c2.o);
 if (sgn(d_sq) == 0) return ret;
 double d = sqrt(d_sq);
 Pt v = (c2.0 - c1.0) / d;
 double c = (c1.r - sign1 * c2.r) / d;
 if (c * c > 1) return ret;
 double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
   Pt n = Pt(v.x * c - sign2 * h * v.y, v.y * c + sign2 * h *
   v.x);
Pt p1 = c1.o + n * c1.r;
   Pt p2 = c2.o + n * (c2.r * sign1);
   if (sgn(p1.x - p2.x) == 0 \& sgn(p1.y - p2.y) == 0)
     p2 = p1 + rotate(c2.o - c1.o);
   ret.push_back({p1, p2});
return ret;
```

5.18 Triangle Center

```
Pt TriangleCircumCenter(Pt a, Pt b, Pt c) {
  Pt res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y) / 2;
  double bx = (c.x + b.x) / 2;
  double by = (c.y + b.y) / 2;
  double r1 = (sin(a2) * (ax - bx) + cos(a2) * (by - ay)) / (sin (a1) * cos(a2) - sin(a2) * cos(a1));
```

5.19 Union of Circles

```
|// Area[i] : area covered by at least i circle
vector<double> CircleUnion(const vector<Cir> &C) {
   const int n = C.size();
   vector<double> Area(n + 1);
   auto check = [&](int i, int j) {
     if (!contain(C[i], C[j]))
       return false
     return sgn(C[i].r - C[j].r) > 0 or (sgn(C[i].r - C[j].r) ==
       0 and i < j);</pre>
   };
   struct Teve {
     double ang; int add; Pt p;
     bool operator<(const Teve &b) { return ang < b.ang; }</pre>
   auto ang = [8](Pt p) { return atan2(p.y, p.x); };
   for (int i = 0; i < n; i++) {</pre>
     int cov = 1;
     vector<Teve> event;
     for (int j = 0; j < n; j++) if (i != j) {
       if (check(j, i)) cov++;
       else if (!check(i, j) and !disjunct(C[i], C[j])) {
         auto I = CircleInter(C[i], C[j]);
         assert(I.size() == 2);
         double a1 = ang(I[0] - C[i].o), a2 = ang(I[1] - C[i].o)
         event.push_back({a1, 1, I[0]});
         event.push_back(\{a2, -1, I[1]\});
         if (a1 > a2) cov++;
     if (event.empty()) {
       Area[cov] += pi * C[i].r * C[i].r;
       continue:
     sort(all(event));
     event.push_back(event[0]);
     for (int j = 0; j + 1 < event.size(); j++) {</pre>
       cov += event[j].add;
       Area[cov] += (event[j].p ^ event[j + 1].p) / 2.;
       double theta = event[j + 1].ang - event[j].ang;
       if (theta < 0) theta += 2 * pi;</pre>
       Area[cov] += (theta - sin(theta)) * C[i].r * C[i].r / 2.;
     }
   return Area;
}
```

6 Graph

6.1 Block Cut Tree

```
struct BlockCutTree {
  int n;
  vector<vector<int>> adj;
  BlockCutTree(int _n) : n(_n), adj(_n) {}
  void addEdge(int u, int v) {
    adj[u].push_back(v);
    adj[v].push_back(u);
}
  pair<int, vector<pair<int, int>>> work() {
    vector<int> dfn(n, -1), low(n), stk;
    vector<pair<int, int>> edg;
  int cnt = 0, cur = 0;
  function<void(int)> dfs = [&](int x) {
```

```
stk.push_back(x);
       dfn[x] = low[x] = cur++;
       for (auto y : adj[x]) {
         if (dfn[y] == -1) {
           dfs(v);
           low[x] = min(low[x], low[y]);
           if (low[y] == dfn[x]) {
             int v;
             do {
               v = stk.back();
               stk.pop_back();
               edg.emplace_back(n + cnt, v);
             } while (v != y);
             edg.emplace_back(x, n + cnt);
             cnt++;
         } else {
          low[x] = min(low[x], dfn[y]);
         }
     };
     for (int i = 0; i < n; i++) {
       if (dfn[i] == -1) {
         stk.clear();
         dfs(i);
     return {cnt, edg};
  }
|};
```

6.2 Count Cycles

```
// ord = sort by deg decreasing, rk[ord[i]] = i
// D: undirected to directed edge from rk small to rk big
vector<int> vis(n, 0);
int c3 = 0, c4 = 0;
for (int x : ord) { // c3
  for (int y : D[x]) vis[y] = 1;
  for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
  for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
  for (int y : D[x]) for (int z : adj[y])
    if (rk[z] > rk[x]) c4 += vis[z]++
  for (int y : D[x]) for (int z : adj[y])
    if (rk[z] > rk[x]) --vis[z];
```

Dominator Tree 6.3

```
vector<int> BuildDomTree(vector<vector<int>> adj, int rt) {
 int n = adj.size();
  // buckets: list of vertices y with sdom(y) = x
 vector<vector<int>> buckets(n), radj(n);
 // rev[dfn[x]] = x
 vector<int> dfn(n, -1), rev(n, -1), pa(n, -1);
 vector<int> sdom(n, -1), dom(n, -1);
 vector<int> fa(n, -1), val(n, -1);
 int stamp = 0;
 \ensuremath{\text{//}} re-number in DFS order
 auto dfs = [8](auto self, int u) -> void {
    rev[dfn[u] = stamp] = u;
    fa[stamp] = sdom[stamp] = val[stamp] = stamp;
   stamp++;
    for (int v : adj[u]) {
      if (dfn[v] == -1) {
       self(self, v);
        pa[dfn[v]] = dfn[u];
      radj[dfn[v]].pb(dfn[u]);
 };
  function<int(int, bool)> Eval = [&](int x, bool fir) {
   if (x == fa[x]) return fir ? x := -1;
    int p = Eval(fa[x], false);
   // x is one step away from the root
   if (p == -1) return x;
    if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[x]];
    fa[x] = p;
```

```
return fir ? val[x] : p;
auto Link = [\delta](int x, int y) \rightarrow void \{ fa[x] = y; \};
dfs(dfs, rt);
// compute sdom in reversed DFS order
for (int x = stamp - 1; x >= 0; --x) {
  for (int y : radj[x]) {
    // sdom[x] = min({y | (y, x) in E(G), y < x}, {sdom[z] | }
   (y, x) in E(G), z > x && z is y's ancestor<math>)
    chmin(sdom[x], sdom[Eval(y, true)]);
  if (x > 0) buckets[sdom[x]].pb(x);
  for (int u : buckets[x]) {
    int p = Eval(u, true);
    if (sdom[p] == x) dom[u] = x;
    else dom[u] = p;
  if (x > 0) Link(x, pa[x]);
}
// idom[x] = -1 if x is unreachable from rt
vector<int> idom(n, -1);
idom[rt] = rt;
rep (x, 1, stamp) {
  if (sdom[x] != dom[x]) dom[x] = dom[dom[x]];
rep (i, 1, stamp) idom[rev[i]] = rev[dom[i]];
return idom;
```

6.4 Enumerate Planar Face

```
// 0-based
struct PlanarGraph{
   int n, m, id;
   vector<Pt<int>> v;
   vector<vector<pair<int, int>>> adj;
   vector<int> conv, nxt, vis;
   PlanarGraph(int n, int m, vector<Pt<int>> _v):
   n(n), m(m), id(0),
   v(_v), adj(n),
   conv(m << 1), nxt(m << 1), vis(m << 1) {}</pre>
   void add_edge(int x, int y) {
     adj[x].push_back({y, id << 1});</pre>
     adj[y].push_back({x, id << 1 | 1});
     conv[id << 1] = x;
     conv[id << 1 | 1] = y;
     id++;
   vector<int> enumerate_face() {
     for (int i = 0; i < n; i++) {</pre>
       sort(all(adj[i]), [&](const auto &a, const auto & b) {
         return (v[a.first] - v[i]) < (v[b.first] - v[i]);</pre>
       }):
       int sz = adj[i].size(), pre = sz - 1;
       for (int j = 0; j < sz; j++) {</pre>
         nxt[adj[i][pre].second] = adj[i][j].second ^ 1;
     }
     vector<int> ret;
     for (int i = 0; i < m * 2; i++) {
       if (!vis[i]) {
         int area = 0, now = i;
         vector<int> pt;
         while (!vis[now]) {
           vis[now] = true:
           pt.push_back(conv[now]);
           now = nxt[now];
         pt.push_back(pt.front());
         for (int i = 0; i + 1 < ssize(pt); i++) {
  area -= (v[pt[i]] ^ v[pt[i + 1]]);</pre>
         // pt = face boundary
         if (area > 0) {
           ret.push_back(area);
         } else {
           // pt is outer face
```

```
National Yang Ming Chiao Tung University – MutedByPEPPA
      }
     return ret;
  }
| };
                                                                                // }
        Manhattan MST
 6.5
                                                                           }
|// {w, u, v}
 vector<tuple<int, int, int>> ManhattanMST(vector<Pt> P) {
  vector<int> id(P.size());
  iota(all(id), 0);
  vector<tuple<int, int, int>> edg;
for (int k = 0; k < 4; k++) {</pre>
                                                                         return true;
     sort(all(id), [8](int i, int j) {
         return (P[i] - P[j]).ff < (P[j] - P[i]).ss;</pre>
                                                                       M1(), M2();
       });
     map<int, int> sweep;
                                                                    1 }
     for (int i : id) {
       auto it = sweep.lower_bound(-P[i].ss);
       while (it != sweep.end()) {
         int j = it->ss;
         Pt d = P[i] - P[j];
         if (d.ss > d.ff) {
           break;
         edg.emplace_back(d.ff + d.ss, i, j);
         it = sweep.erase(it);
      sweep[-P[i].ss] = i;
     for (Pt &p : P) {
       if (k % 2) {
        p.ff = -p.ff;
       } else {
         swap(p.ff, p.ss);
                                                                         if (i < 4) {
    }
   return edg;
                                                                         }
 6.6 Matroid Intersection
 M1 = xx matroid, M2 = xx matroid
 y<-s if I+y satisfies M1
 y->t if I+y satisfies M2
x<-y if I-x+y satisfies M2
 x->y if I-x+y satisfies M1
 交換圖點權
 -w[e] if e \in I
 w[e] otherwise
vector<int> I(, 0);
 while (true) {
   vector<vector<int>> adj();
  int s = , t = s + 1;
  auto M1 = [&]() \rightarrow void { // xx matroid}
     { // y<-s
       // x->y
     {
    }
  auto M2 = [&]() \rightarrow void { // xx matroid}
    { // y->t
     {
       // x<-y
                                                                            --q;
     }
                                                                         }
                                                                       }
  auto augment = [&]() -> bool { // 註解掉的是帶權版
                                                                       int solve() {
     vector<int> vis( + 2, 0), dis( + 2, IINF), from( + 2, -1);
     queue<int> q;
     vis[s] = 1;
                                                                         ans = q = 0;
     dis[s] = 0;
     q.push(s);
```

while (!q.empty()) {

// vis[u] = 0; for (int v : adj[u]) {

int u = q.front(); q.pop();

int w = ; // no weight -> 1

```
if (chmin(dis[v], dis[u] + w)) {
           from[v] = u:
           // if (!vis[v]) {
             // vis[v] = 1;
             q.push(v);
     if (from[t] == -1) return false;
     for (int cur = from[t];; cur = from[cur]) {
       if (cur == -1 || cur == s) break;
       I[cur] ^= 1;
   if (!augment()) break;
        Maximum Clique
constexpr size_t kN = 150;
 using bits = bitset<kN>;
struct MaxClique {
  bits G[kN], cs[kN];
   int ans, sol[kN], q, cur[kN], d[kN], n;
   void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i) G[i].reset();</pre>
   void addEdge(int u, int v) {
     G[u][v] = G[v][u] = 1;
   void preDfs(vector<int> &v, int i, bits mask) {
       for (int x : v) d[x] = (G[x] \& mask).count();
       sort(all(v), [&](int x, int y) {
         return d[x] > d[y];
     vector<int> c(v.size());
     cs[1].reset(), cs[2].reset();
     int l = max(ans - q + 1, 1), r = 2, tp = 0, k;
     for (int p : v) {
       for (k = 1;
         (cs[k] & G[p]).any(); ++k);
       if (k >= r) cs[++r].reset();
       cs[k][p] = 1;
       if (k < l) v[tp++] = p;</pre>
     for (k = 1; k < r; ++k)
       for (auto p = cs[k]._Find_first(); p < kN; p = cs[k].</pre>
      _Find_next(p))
         v[tp] = p, c[tp] = k, ++tp;
     dfs(v, c, i + 1, mask);
   void dfs(vector<int> &v, vector<int> &c, int i, bits mask) {
     while (!v.empty()) {
       int p = v.back();
       v.pop_back();
       mask[p] = 0;
       if (q + c.back() <= ans) return;</pre>
       cur[q++] = p;
       vector<int> nr;
       for (int x : v)
         if (G[p][x]) nr.push_back(x);
       if (!nr.empty()) preDfs(nr, i, mask & G[p]);
       else if (q > ans) ans = q, copy_n(cur, q, sol);
       c.pop_back();
     vector<int> v(n);
     iota(all(v), 0);
     preDfs(v, 0, bits(string(n, '1')));
     return ans;
  }
|} cliq;
```

6.8 Tree Hash

dfn.assign(2 * n, -1);

```
map<vector<int>, int> id;
                                                                         low.assign(2 * n, -1);
vector<vector<int>> sub;
                                                                         for (int i = 0; i < n * 2; i++)
vector<int> siz;
                                                                           if (dfn[i] == -1) {
int getid(const vector<int> &T) {
                                                                             dfs(i);
  if (id.count(T)) return id[T];
  int s = 1;
                                                                         for (int i = 0; i < n; ++i) {</pre>
  for (int x : T) {
                                                                           if (id[2 * i] == id[2 * i + 1]) {
                                                                             return false;
   s += siz[x];
                                                                           ans[i] = id[2 * i] > id[2 * i + 1];
  sub.push_back(T);
  siz.push_back(s);
                                                                         return true;
 return id[T] = id.size();
                                                                       }
                                                                    };
int dfs(int u, int f) {
  vector<int> S;
                                                                             Virtual Tree
                                                                    6.10
  for (int v : G[u]) if (v != f) {
    S.push_back(dfs(v, u));
                                                                    // need LCA
 }
                                                                    vector<vector<int>> vir(n);
  sort(all(S));
                                                                    auto clear = [&](auto self, int u) -> void {
 return getid(S);
                                                                       for (int v : vir[u]) self(self, v);
                                                                       vir[u].clear();
                                                                    };
6.9
       Two-SAT
                                                                    auto build = [\delta](\text{vector} < \text{int} > \delta v) \rightarrow \text{void} \{ // \text{ be careful of the} \}
                                                                           changes to the array
struct TwoSat {
                                                                       // maybe dont need to sort when do it while dfs
  int n;
                                                                       sort(all(v), [&](int a, int b) {
  vector<vector<int>> G;
                                                                         return dfn[a] < dfn[b];</pre>
  vector<bool> ans;
                                                                       });
  vector<int> id, dfn, low, stk;
  TwoSat(int n) : n(n), G(2 * n) {}
                                                                       clear(clear, 0);
                                                                       if (v[0] != 0) v.insert(v.begin(), 0);
  void addClause(int u, bool f, int v, bool g) { // (u = f) or
    (v = g)
                                                                       int k = v.size();
                                                                       vector<int> st;
    G[2 * u + !f].push_back(2 * v + g);
                                                                       rep (i, 0, k) {
    G[2 * v + !g].push_back(2 * u + f);
                                                                         if (st.empty()) {
                                                                           st.push_back(v[i]);
  void addImply(int u, bool f, int v, bool g) { // (u = f) -> (
                                                                           continue;
     v = g
    G[2 * u + f].push_back(2 * v + g);
                                                                         int p = lca(v[i], st.back());
    G[2 * v + !g].push_back(2 * u + !f);
                                                                         if (p == st.back()) {
                                                                           st.push_back(v[i]);
  int addVar() {
                                                                           continue;
    G.emplace_back();
    G.emplace_back();
                                                                         while (st.size() >= 2 && dep[st.end()[-2]] >= dep[p]) {
    return n++;
                                                                           vir[st.end()[-2]].push_back(st.back());
  void addAtMostOne(const vector<pair<int, bool>> &li) {
                                                                           st.pop_back();
    if (ssize(li) <= 1) return;</pre>
    int pu; bool pf; tie(pu, pf) = li[0];
                                                                         if (st.back() != p) {
    for (int i = 2; i < ssize(li); i++) {</pre>
                                                                           vir[p].push_back(st.back());
      const auto &[u, f] = li[i];
                                                                           st.pop_back();
                                                                           st.push_back(p);
      int nxt = addVar();
                                                                         ļ
      addClause(pu, !pf, u, !f);
      addClause(pu, !pf, nxt, true);
                                                                         st.push_back(v[i]);
      addClause(u, !f, nxt, true);
                                                                       while (st.size() >= 2) {
      tie(pu, pf) = make_pair(nxt, true);
                                                                         vir[st.end()[-2]].push_back(st.back());
                                                                         st.pop_back();
    addClause(pu, !pf, li[1].first, !li[1].second);
                                                                       }
  int cur = 0, scc = 0;
                                                                    ⊥};
  void dfs(int u) {
    stk.push_back(u);
                                                                          Math
                                                                     7
    dfn[u] = low[u] = cur++;
    for (int v : G[u]) {
                                                                          Combinatoric
      if (dfn[v] == -1) {
                                                                     vector<mint> fac, inv;
        dfs(v);
        chmin(low[u], low[v]);
                                                                    inline void init (int n) {
      } else if (id[v] == -1) {
                                                                       fac.resize(n + 1);
        chmin(low[u], dfn[v]);
                                                                       inv.resize(n + 1);
      }
                                                                       fac[0] = inv[0] = 1;
                                                                       rep (i, 1, n + 1) fac[i] = fac[i - 1] * i;
    if (dfn[u] == low[u]) {
                                                                       inv[n] = fac[n].inv();
      int x;
                                                                       for (int i = n; i > 0; --i) inv[i - 1] = inv[i] * i;
        x = stk.back();
        stk.pop_back();
                                                                    inline mint Comb(int n, int k) {
        id[x] = scc;
                                                                       if (k > n || k < 0) return 0;
      } while (x != u);
                                                                       return fac[n] * inv[k] * inv[n - k];
      scc++;
   }
                                                                    inline mint H(int n, int m) {
  bool satisfiable() {
                                                                       return Comb(n + m - 1, m);
    ans.assign(n, 0);
    id.assign(2 * n, -1);
```

inline mint catalan(int n){

```
return fac[2 * n] * inv[n + 1] * inv[n];
7.2 Discrete Log
| int power(int a, int b, int p, int res = 1) {
  for (; b; b /= 2, a = 1LL * a * a % p) {
    if (b & 1) {
      res = 1LL * res * a % p;
    }
  return res;
int exbsgs(int a, int b, int p) {
  b %= p;
  if (b == 1 || p == 1) {
    return 0;
  if (a == 0) {
    return b == 0 ? 1 : -1;
  i64 g, k = 0, t = 1; // t : a ^ k / sum{d}
  while ((g = std::gcd(a, p)) > 1) {
    if (b % g) {
      return -1;
    b /= g;
    p /= g;
    k++;
    t = t * (a / g) % p;
    if (t == b) {
      return k;
    }
  }
  const int n = std::sqrt(p) + 1;
  std::unordered_map<int, int> mp;
  mp[b] = 0:
  int x = b, y = t;
  int mi = power(a, n, p);
  for (int i = 1; i < n; i++) {</pre>
    x = 1LL * x * a % p;
    mp[x] = i;
  for (int i = 1; i <= n; i++) {
    t = 1LL * t * mi % p;
    if (mp.contains(t)) {
      return 1LL * i * n - mp[t] + k;
  }
  return -1; // no solution
7.3 Div Floor Ceil
|// b > 0!!!!
int CEIL(int a, int b) {
 return (a >= 0 ? (a + b - 1) / b : a / b);
int FLOOR(int a, int b) {
 return (a >= 0 ? a / b : (a - b + 1) / b);
7.4 exCRT
i64 exgcd(i64 a, i64 b, i64 &x, i64 &y) {
  if (b == 0) {
    x = 1;
y = 0;
    return a;
  i64 g = exgcd(b, a \% b, y, x);
  y = a / b * x;
  return g;
// return {x, T}
// a: moduli, b: remainders
// x: first non-negative solution, T: minimum period
std::pair<i64, i64> exCRT(auto &a, auto &b) {
  auto [m1, r1] = std::tie(a[0], b[0]);
  for (int i = 1; i < std::ssize(a); i++) {</pre>
```

```
auto [m2, r2] = std::tie(a[i], b[i]);
    i64 x, y;
     i64 g = exgcd(m1, m2, x, y);
     if ((r2 - r1) % g) { // no solution
      return {-1, -1};
    }
    x = (i128(x) * (r2 - r1) / g) % (m2 / g);
    if(x < 0) {
      x += (m2 / g);
    r1 = m1 * x + r1;
    m1 = std::lcm(m1, m2);
  r1 %= m1;
  if (r1 < 0) {
    r1 += m1;
  return {r1, m1};
7.5 Factorization
|ull modmul(ull a, ull b, ull M) {
  i64 ret = a * b - M * ull(1.L / M * a * b);
  return ret + M * (ret < 0) - M * (ret >= (i64)M);
ull modpow(ull b, ull e, ull mod) {
  ull ans = 1;
   for (; e; b = modmul(b, b, mod), e /= 2)
     if (e & 1) ans = modmul(ans, b, mod);
}
bool isPrime(ull n) {
  if (n < 2 || n % 6 % 4 != 1) return (n | 1) == 3;
  ull A[] = {2, 325, 9375, 28178, 450775, 9780504, 1795265022},
          _builtin_ctzll(n - 1), d = n >> s;
  for (ull a : A) {
    ull p = modpow(a \% n, d, n), i = s;
    while (p != 1 && p != n - 1 && a % n && i--)
      p = modmul(p, p, n);
     if (p != n - 1 && i != s) return 0;
  return 1:
}
ull pollard(ull n) {
  uniform_int_distribution<ull> unif(0, n - 1);
  auto f = [n, &c](ull x) \{ return modmul(x, x, n) + c % n; \};
  ull x = 0, y = 0, t = 30, prd = 2, i = 1, q;
  while (t++ % 40 || __gcd(prd, n) == 1) {
    if (x == y) c = unif(rng), x = ++i, y = f(x);
     if ((q = modmul(prd, max(x, y) - min(x, y), n))) prd = q;
    x = f(x), y = f(f(y));
  return __gcd(prd, n);
}
vector<ull> factor(ull n) {
  if (n == 1) return {};
  if (isPrime(n)) return {n};
  ull x = pollard(n);
  auto l = factor(x), r = factor(n / x);
  l.insert(l.end(), r.begin(), r.end());
  return 1:
}
7.6 Floor Sum
\frac{1}{\sqrt{sum_0^n}} floor((a * x + b) / c)) in log(n + m + a + b)
int floor_sum(int a, int b, int c, int n) { // add mod if
     needed
   int m = (a * n + b) / c;
  if (a >= c || b >= c)
```

return (a / c) * (n * (n + 1) / 2) + (b / c) * (n + 1) +

return n * m - floor_sum(c, c - b - 1, a, m - 1);

floor_sum(a % c, b % c, c, n);

7.7 FWT

}

return 0;

if (n < 0 || a == 0)

// if (r < m) { // infinite solution

// return std::vector<T>();

```
void fwt(vector<ll> &f, bool inv = false) { // xor-convolution
  const int N = 31 - __builtin_clz(ssize(f)),
        inv2 = (MOD + \overline{1}) / 2;
                                                                       std::vector<T> res(m);
                                                                       for (int i = 0; i < m; i++) {
  rep (i, 0, N) rep (j, 0, 1 << N) {
    if (j >> i & 1 ^ 1) {
                                                                         if (where[i] != -1) {
                                                                           res[i] = b[where[i]];
       ll a = f[j], b = f[j | (1 << i)];
       if (inv) {
         f[j] = (a + b) * inv2 % MOD;
                                                                       return res;
         f[j \mid (1 << i)] = (a - b + MOD) * inv2 % MOD;
                                                                    |}:
       } else {
         f[j] = (a + b) \% MOD;
                                                                     7.9 Lagrange Interpolation
         f[j \mid (1 << i)] = (a - b + MOD) \% MOD;
                                                                    | struct Lagrange {
                                                                       int deg{};
    }
                                                                       vector<int> C;
  }
                                                                       Lagrange(const vector<int> &P) {
| }
                                                                         deg = P.size() - 1;
       Gauss Elimination
                                                                         C.assign(deg + 1, 0);
                                                                         for (int i = 0; i <= deg; i++) {
using Z = ModInt<998244353>;
                                                                            int q = inv[i] * inv[i - deg] % mod;
// using F = long double;
                                                                           if ((deg - i) % 2 == 1) {
  q = mod - q;
using Matrix = std::vector<std::vector<Z>>;
// using Matrix = std::vector<std::vector<F>>; (double)
                                                                           C[i] = P[i] * q % mod;
// using Matrix = std::vector<std::bitset<5000>>; (mod 2)
                                                                         }
template <typename T>
auto gauss(Matrix &A, std::vector<T> &b, int n, int m) {
                                                                       int operator()(int x) \{ // 0 \le x \le mod \}
  assert(std::ssize(b) == n);
                                                                         if (0 \le x \text{ and } x \le \text{deg}) {
  int r = 0;
                                                                           int ans = fac[x] * fac[deg - x] % mod;
  std::vector<int> where(m, -1);
                                                                            if ((deg - x) % 2 == 1) {
                                                                             ans = (mod - ans);
  for (int i = 0; i < m && r < n; i++) {
    int p = r; // pivot
                                                                           return ans * C[x] % mod;
    while (p < n \& A[p][i] == T(0)) {
      p++;
                                                                         vector<int> pre(deg + 1), suf(deg + 1);
    }
                                                                         for (int i = 0; i <= deg; i++) {</pre>
    if (p == n) {
                                                                           pre[i] = (x - i);
       continue;
                                                                           if (i) {
                                                                             pre[i] = pre[i] * pre[i - 1] % mod;
                                                                           }
    std::swap(A[r], A[p]);
                                                                         }
    std::swap(b[r], b[p]);
                                                                         for (int i = deg; i >= 0; i--) {
    where[i] = r;
                                                                           suf[i] = (x - i);
                                                                           if (i < deg) {
    // coef: mod 2 don't need this
                                                                             suf[i] = suf[i] * suf[i + 1] % mod;
    T inv = T(1) / A[r][i];
                                                                           }
     for (int j = i; j < m; j++) {</pre>
                                                                         }
      A[r][j] *= inv;
                                                                         int ans = 0;
                                                                         for (int i = 0; i <= deg; i++) {</pre>
    b[r] *= inv;
                                                                           ans += (i == 0 ? 1 : pre[i - 1]) * (i == deg ? 1 : suf[i])
     for (int j = 0; j < n; j++) { // deduct: mod 2 don't need</pre>
                                                                          + 1]) % mod * C[i];
                                                                           ans %= mod;
      if (j != r) {
                                                                         if (ans < 0) ans += mod;
        T x = A[j][i];
                                                                         return ans;
         for (int k = i; k < m; k++) {
                                                                       }
           A[j][k] -= x * A[r][k];
                                                                    };
        b[j] = x * b[r];
                                                                     7.10 Linear Sieve
      }
                                                                     const int C = 1e6 + 5;
                                                                     int mo[C], lp[C], phi[C], isp[C];
    // for (int j = 0; j < n; ++j) { // (mod 2) -> coef and
                                                                     vector<int> prime;
    // if (j != r && A[j][i]) {
           A[j] ^= A[r];
                                                                     void sieve() {
                                                                       mo[1] = phi[1] = 1;
           b[j] ^= b[r];
    //
                                                                       rep (i, 1, C) lp[i] = 1;
    //
         }
                                                                       rep (i, 2, C) {
    // }
                                                                         if (lp[i] == 1) {
                                                                           lp[i] = i;
                                                                           prime.pb(i);
  }
                                                                           isp[i] = 1;
  for (int i = r; i < n; i++) {
                                                                           mo[i] = -1;
    if (ranges::all_of(A[i] | views::take(m), [](auto x) {
                                                                           phi[i] = i - 1;
     return x == 0; }) && b[i] != T(0)) {
      return std::vector<T>(); // no solution
                                                                         for (int p : prime) {
                                                                           if (i * p >= C) break;
    // if (A[i].none() && b[i]) { // (mod 2)
                                                                            lp[i * p] = p;
    //
        return std::vector<T>();
                                                                           if (i % p == 0) {
                                                                             phi[p * i] = phi[i] * p;
    // }
                                                                             break:
```

phi[i * p] = phi[i] * (p - 1);

mo[i * p] = mo[i] * mo[p];

constexpr ModInt &operator/=(const ModInt &r) { return *this

= *this / r; }

```
constexpr bool operator==(const ModInt &r) const { return v
  }
                                                                         == r.v: }
į }
                                                                      constexpr bool operator!=(const ModInt &r) const { return v
                                                                         != r.v; }
                                                                      explicit constexpr operator bool() const { return v != 0; }
 7.11
       Lucas
                                                                      friend std::ostream &operator<<(std::ostream &os, const</pre>
| // comb(n, m) % M, M = p^k
                                                                         ModInt &r) {
 // O(M)-O(log(n))
                                                                        return os << r.v;
 struct Lucas {
  const int p, M;
                                                                    };
  vector<int> f;
                                                                    using mint = ModInt<998244353>;
  Lucas(int p, int M) : p(p), M(M), f(M + 1) {
                                                                    template <> const mint mint::G = mint(3);
    f[0] = 1;
                                                                    7.13 Primitive Root
    for (int i = 1; i <= M; i++) {
      f[i] = f[i - 1] * (i % p == 0 ? 1 : i) % M;
                                                                   |ull primitiveRoot(ull p) {
                                                                      auto fac = factor(p - 1);
  }
                                                                      sort(all(fac)):
  int CountFact(int n) {
                                                                      fac.erase(unique(all(fac)), fac.end());
    int c = 0;
                                                                      auto test = [p, fac](ull x) {
    while (n) c += (n /= p);
                                                                        for(ull d : fac)
    return c:
                                                                        if (modpow(x, (p - 1) / d, p) == 1)
  }
                                                                          return false;
   // (n! without factor p) % p^k
                                                                        return true;
  int ModFact(int n) {
    int r = 1;
                                                                      uniform_int_distribution<ull> unif(1, p - 1);
    while (n) {
                                                                      ull root;
      r = r * power(f[M], n / M % 2, M) % M * f[n % M] % M;
                                                                      while(!test(root = unif(rng)));
      n /= p;
                                                                      return root:
                                                                   }
    return r;
                                                                    7.14 Simplex
  int ModComb(int n, int m) {
                                                                   | // \max\{cx\}  subject to \{Ax <= b, x >= 0\}
    if (m < 0 or n < m) return 0;
                                                                   // n: constraints. m: vars !!!
     int c = CountFact(n) - CountFact(m) - CountFact(n - m);
                                                                    // x[] is the optimal solution vector
    int r = ModFact(n) * power(ModFact(m), M / p * (p - 1) - 1,
                                                                    // usage :
      M) % M
                                                                    // x = simplex(A, b, c); (A <= 100 x 100)
               * power(ModFact(n - m), M / p * (p - 1) - 1, M) %
                                                                    vector<double> simplex(
                                                                        const vector<vector<double>> &a,
                                                                        const vector<double> &b.
    return r * power(p, c, M) % M;
  }
                                                                        const vector<double> &c) {
};
                                                                      int n = (int)a.size(), m = (int)a[0].size() + 1;
                                                                      vector val(n + 2, vector<double>(m + 1));
 7.12 Mod Int
                                                                      vector<int> idx(n + m);
using u32 = unsigned int;
                                                                       iota(all(idx), 0);
 using u64 = unsigned long long;
                                                                      int r = n, s = m - 1;
 template <class T>
                                                                      for (int i = 0; i < n; ++i) {
 constexpr T power(T a, u64 b, T res = 1) {
                                                                        for (int j = 0; j < m - 1; ++j)
  for (; b != 0; b /= 2, a *= a) {
                                                                          val[i][j] = -a[i][j];
    if (b & 1) {
                                                                        val[i][m - 1] = 1;
      res *= a;
                                                                        val[i][m] = b[i];
    }
                                                                        if (val[r][m] > val[i][m])
                                                                          r = i;
  return res;
                                                                      copy(all(c), val[n].begin());
                                                                      val[n + 1][m - 1] = -1;
 template <u32 P>
                                                                      for (double num; ; ) {
 struct ModInt {
                                                                        if (r < n) {
  u32 v:
                                                                          swap(idx[s], idx[r + m]);
  const static ModInt G;
                                                                          val[r][s] = 1 / val[r][s];
  constexpr ModInt &norm(u32 x) {
                                                                          for (int j = 0; j <= m; ++j) if (j != s)
    v = x < P ? x : x - P;
                                                                            val[r][j] *= -val[r][s];
    return *this;
                                                                          for (int i = 0; i <= n + 1; ++i) if (i != r) {
                                                                            for (int j = 0; j <= m; ++j) if (j != s)</pre>
  constexpr ModInt(i64 x = 0) { norm(x \% P + P); }
                                                                              val[i][j] += val[r][j] * val[i][s];
   constexpr ModInt inv() const { return power(*this, P - 2); }
                                                                            val[i][s] *= val[r][s];
  constexpr ModInt operator-() const { return ModInt() - *this;
                                                                          }
      }
                                                                        }
  constexpr ModInt operator+(const ModInt &r) const { return
                                                                        r = s = -1;
     ModInt().norm(v + r.v); }
                                                                        for (int j = 0; j < m; ++j)
   constexpr ModInt operator-(const ModInt &r) const { return
                                                                          if (s < 0 || idx[s] > idx[j])
     ModInt().norm(v + P - r.v); }
                                                                            if (val[n + 1][j] > eps || val[n + 1][j] > -eps && val[
  constexpr ModInt operator*(const ModInt &r) const { return
                                                                         n][j] > eps)
     ModInt().norm(u64(v) * r.v % P); }
                                                                              s = j;
  constexpr ModInt operator/(const ModInt &r) const { return *
                                                                        if (s < 0) break;
     this * r.inv(); }
                                                                        for (int i = 0; i < n; ++i) if (val[i][s] < -eps) {</pre>
   constexpr ModInt &operator+=(const ModInt &r) { return *this
     = *this + r; }
                                                                            || (num = val[r][m] / val[r][s] - val[i][m] / val[i][s
  constexpr ModInt &operator-=(const ModInt &r) { return *this
                                                                         ]) < -eps
     = *this - r; }
                                                                            \parallel num < eps && idx[r + m] > idx[i + m])
  constexpr ModInt &operator*=(const ModInt &r) { return *this
                                                                            r = i;
     = *this * r; }
```

 $if (r < 0) {$

// Solution is unbounded.

int i = roughs[k];

P: 985661441, Rt: 3

```
return vector<double>{};
                                                                                 if (skip[i]) continue;
     }
                                                                                i64 d = 1LL * i * p;
                                                                                larges[ns] = larges[k] - (d <= v ? larges[smalls[d] -</pre>
                                                                             pc] : smalls[n / d]) + pc;
   if (val[n + 1][m] < -eps) {</pre>
     // No solution.
                                                                                roughs[ns++] = i;
     return vector<double>{};
                                                                              s = ns;
                                                                              for (int j = v / p; j >= p; --j) {
   vector<double> x(m - 1):
                                                                                int c = smalls[j] - pc, e = min(j * p + p, v + 1);
for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
   for (int i = m; i < n + m; ++i)
     if (idx[i] < m - 1)</pre>
       x[idx[i]] = val[i - m][m];
   return x;
                                                                            }
| }
                                                                          }
                                                                          for (int k = 1; k < s; ++k) {
 7.15 Sgrt Mod
                                                                            const i64 m = n / roughs[k];
                                                                            i64 t = larges[k] - (pc + k - 1);
// the Jacobi symbol is a generalization of the Legendre symbol
                                                                            for (int l = 1; l < k; ++l) {</pre>
                                                                              int p = roughs[l];
 // such that the bottom doesn't need to be prime.
                                                                              if (1LL * p * p > m) break;
 // (n|p) -> same as legendre
                                                                              t = smalls[m / p] - (pc + l - 1);
// (n|ab) = (n|a)(n|b)
// work with long long
                                                                            larges[0] -= t;
 int Jacobi(int a, int m) {
                                                                          }
  int s = 1;
  for (; m > 1; ) {
  a %= m;
                                                                          return larges[0];
                                                                      }
     if (a == 0) return 0;
                                                                        7.17 ModMin
     const int r = __builtin_ctz(a);
     if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
                                                                      | // min\{k \mid l \le ((ak) mod m) \le r\}, no solution \rightarrow -1
     a >>= r:
                                                                       int mod_min(int a, int m, int l, int r) {
    if (a & m & 2) s = -s;
                                                                        if (a == 0) return l ? -1 : 0;
     swap(a, m);
                                                                         if (int k = (l + a - 1) / a; k * a <= r)
   ļ
                                                                         return k;
   return s;
                                                                         int b = m / a, c = m % a;
}
                                                                         if (int y = mod_min(c, a, a - r % a, a - l % a))
                                                                         return (l + y * c + a - 1) / a + y * b;
 // 0: a == 0
 // -1: a isn't a quad res of p
                                                                      }
 // else: return X with X^2 % p == a
 // doesn't work with long long
                                                                        7.18 FFT
 int QuadraticResidue(int a, int p) {
                                                                       | template<typename C = complex<double>>
  if (p == 2) return a & 1;
                                                                       void FFT(vector<C> &P, C w, bool inv = 0) {
  if (int jc = Jacobi(a, p); jc <= 0) return jc;</pre>
                                                                          int n = P.size(), lg = __builtin_ctz(n);
   int b, d;
                                                                          assert(__builtin_popcount(n) == 1);
   for (; ; ) {
     b = rand() \% p;
                                                                          for (int j = 1, i = 0; j < n - 1; ++j) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1); // !!!
     d = (1LL * b * b + p - a) % p;
     if (Jacobi(d, p) == -1) break;
                                                                            if (j < i) swap(P[i], P[j]);</pre>
   int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
   for (int e = (1LL + p) >> 1; e; e >>= 1) {
                                                                          vector<C> ws = {inv ? C{1} / w : w};
     if (e & 1) {
       tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p
                                                                          rep (i, 1, lg) ws.pb(ws[i - 1] * ws[i - 1]);
                                                                          reverse(all(ws));
       g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
                                                                          rep (i, 0, lg) {
                                                                            for (int k = 0; k < n; k += 2 << i) {
     tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
                                                                              C base = C\{1\};
     f1 = (2LL * f0 * f1) % p;
                                                                              rep (j, k, k + (1 << i)) {
     f0 = tmp;
                                                                                auto t = base * P[j + (1 << i)];</pre>
                                                                                auto u = P[j];
   return g0;
                                                                                P[j] = u + t;
                                                                                P[j + (1 << i)] = u - t;
                                                                                base = base * ws[i];
 7.16 PiCount
                                                                              }
 i64 PrimeCount(i64 n) { // n ~ 10^13 => < 2s
                                                                            }
   if (n <= 1) return 0;</pre>
                                                                          }
   int v = sqrt(n), s = (v + 1) / 2, pc = 0;
   vector<int> smalls(v + 1), skip(v + 1), roughs(s);
                                                                          if (inv) rep (i, 0, n) P[i] = P[i] / C(n);
   vector<i64> larges(s);
                                                                       }
   for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
   for (int i = 0; i < s; ++i) {
                                                                        const int N = 1 << 21;</pre>
     roughs[i] = 2 * i + 1;
                                                                       const double PI = acos(-1);
                                                                       const auto w = exp(-complex<double>(0, 2.0 * PI / N));
     larges[i] = (n / (2 * i + 1) + 1) / 2;
                                                                       7.19 NTT prime
   for (int p = 3; p <= v; ++p) {
     if (smalls[p] > smalls[p - 1]) {
                                                                           • P: 7681, Rt: 17
                                                                                                                               P: 12289, Rt: 11
       int q = p * p;
                                                                           • P: 40961, Rt: 3
                                                                                                                               P: 65537, Rt: 3
       ++pc;
                                                                           • P: 786433, Rt: 10
                                                                                                                             P: 5767169, Rt: 3
       if (1LL * q * q > n) break;
       skip[p] = 1;

    P: 7340033, Rt: 3

                                                                                                                            P: 23068673, Rt: 3
       for (int i = q; i <= v; i += 2 * p) skip[i] = 1;
                                                                           • P: 469762049, Rt: 3
                                                                                                                       P: 2061584302081, Rt: 7
       int ns = 0;
                                                                           • P: 2748779069441, Rt: 3
                                                                                                                            P: 167772161, Rt: 3
       for (int k = 0; k < s; ++k) {
```

• P: 104857601, Rt: 3

```
• P: 998244353, Rt: 3
                                                  P: 1107296257, Rt: 10

    P: 2013265921, Rt: 31

                                                  P: 2810183681, Rt: 11
                                                  P: 605028353, Rt: 3
  • P: 2885681153, Rt: 3
  • P: 1945555039024054273. Rt: 5
                                        P: 9223372036737335297, Rt: 3
                                                                           std::vector<D> res(n);
7.20 Polynomial
std::mt19937_64 rng(std::chrono::steady_clock::now().
                                                                          return res:
     time_since_epoch().count());
template <class mint>
void nft(bool type, std::vector<mint> &a) {
                                                                           std::vector<D> res(n);
 int n = int(a.size()), s = 0;
  while ((1 << s) < n) {
   s++;
 }
                                                                           return res;
  assert(1 << s == n);
  static std::vector<mint> ep, iep;
  while (int(ep.size()) <= s) {</pre>
    ep.push_back(power(mint::G, mint(-1).v / (1 << int(ep.size
     ()))));
                                                                           int n = size():
    iep.push_back(ep.back().inv());
                                                                           std::vector<D> res(n);
 ļ
  std::vector<mint> b(n);
                                                                            res[i] = v[i] * r;
  for (int i = 1; i <= s; i++) {
    int w = 1 << (s - i);</pre>
                                                                           return res;
    mint base = type ? iep[i] : ep[i], now = 1;
    for (int y = 0; y < n / 2; y += w) {
      for (int x = 0; x < w; x++) {
        auto l = a[y << 1 | x];</pre>
        auto r = now * a[y << 1 | x | w];
        b[y | x] = l + r;
        b[y | x | n >> 1] = l - r;
      now *= base;
    }
                                                                           r * r; }
    std::swap(a, b);
}
                                                                            res[i + s] = v[i];
template <class mint>
std::vector<mint> multiply(const std::vector<mint> &a, const
                                                                           return res;
    std::vector<mint> &b) {
  int n = int(a.size()), m = int(b.size());
  if (!n || !m) return {};
                                                                           if (size() <= s) {</pre>
 if (std::min(n, m) <= 8) {</pre>
                                                                             return Poly();
    std::vector<mint> ans(n + m - 1);
    for (int i = 0; i < n; i++) {</pre>
      for (int j = 0; j < m; j++) {
        ans[i + j] += a[i] * b[j];
                                                                             res[i] = v[i + s];
                                                                          return res:
    return ans;
 }
  int lg = 0;
  while ((1 << lg) < n + m - 1) {
    lg++;
  int z = 1 << lg;</pre>
  auto a2 = a, b2 = b;
  a2.resize(z);
  b2.resize(z):
  nft(false, a2);
                                                                           n; }
  nft(false, b2);
  for (int i = 0; i < z; i++) {</pre>
                                                                           n; }
   a2[i] *= b2[i];
                                                                         Poly pre(int le) const {
 nft(true, a2);
  a2.resize(n + m - 1);
  mint iz = mint(z).inv();
 for (int i = 0; i < n + m - 1; i++) {
                                                                           if (n != -1) {
   a2[i] *= iz;
                                                                             res.resize(n);
  return a2;
                                                                          return res;
template <class D>
                                                                        Polv diff() const {
struct Polv {
  std::vector<D> v;
  Poly(const std::vector<D> \delta v_{-} = \{\}) : v(v_{-}) \{ shrink(); \}
  void shrink() {
    while (v.size() > 1 && !v.back()) {
                                                                          return res:
      v.pop_back();
                                                                        Poly inte() const {
 }
```

```
int size() const { return int(v.size()); }
D freq(int p) const { return (p < size()) ? v[p] : D(0); }</pre>
Poly operator+(const Poly &r) const {
  auto n = std::max(size(), r.size());
  for (int i = 0; i < n; i++)
   res[i] = freq(i) + r.freq(i);
Poly operator-(const Poly &r) const {
  int n = std::max(size(), r.size());
  for (int i = 0; i < n; i++) {</pre>
    res[i] = freq(i) - r.freq(i);
Poly operator*(const Poly &r) const { return {multiply(v, r.v
Poly operator*(const D &r) const {
  for (int i = 0; i < n; i++) {
Poly operator/(const D &r) const { return *this * r.inv(); }
Poly operator/(const Poly &r) const {
  if (size() < r.size()) return {{}};</pre>
  int n = size() - r.size() + 1;
  return (rev().pre(n) * r.rev().inv(n)).pre(n).rev();
Poly operator%(const Poly &r) const { return *this - *this /
Poly operator<<(int s) const {</pre>
  std::vector<D> res(size() + s);
  for (int i = 0; i < size(); i++) {</pre>
Poly operator>>(int s) const {
  std::vector<D> res(size() - s);
  for (int i = 0; i < size() - s; i++) {</pre>
Poly & operator += (const Poly &r) { return *this = *this + r; }
Poly & operator == (const Poly &r) { return *this = *this - r; }
Poly & operator*=(const Poly &r) { return *this = *this * r; }
Poly & Soperator*=(const D &r) { return *this = *this * r; }
Poly &operator/=(const Poly &r) { return *this = *this / r; }
Poly &operator/=(const D &r) { return *this = *this / r; }
Poly & operator %= (const Poly &r) { return *this = *this % r; }
Poly &operator<<=(const size_t &n) { return *this = *this <<
Poly &operator>>=(const size_t &n) { return *this = *this >>
 return {{v.begin(), v.begin() + std::min(size(), le)}};
Poly rev(int n = -1) const {
  std::vector<D> res = v;
 std::reverse(res.begin(), res.end());
  std::vector<D> res(std::max(0, size() - 1));
  for (int i = 1; i < size(); i++) {</pre>
    res[i - 1] = freq(i) * i;
```

```
std::vector<D> res(size() + 1);
                                                                       for (auto i = 0; i < p.size(); i++) {</pre>
  for (int i = 0; i < size(); i++) {</pre>
                                                                         if (p.v[i]) {
                                                                           os << p.v[i] << "x^" << i;
    res[i + 1] = freq(i) / (i + 1);
                                                                           if (i != p.size() - 1) {
                                                                             os << "+";
}
                                                                           }
// f * f.inv() = 1 + g(x)x^m
                                                                         }
Poly inv(int m) const {
  Poly res = Poly(\{D(1) / freq(0)\});
                                                                       return os;
                                                                    }
  for (int i = 1; i < m; i *= 2) {
    res = (res * D(2) - res * res * pre(2 * i)).pre(2 * i);
                                                                  };
                                                                  template <class mint>
                                                                  struct MultiEval {
  return res.pre(m);
                                                                     using NP = MultiEval *;
                                                                     NP l, r;
Poly exp(int n) const {
                                                                     int sz;
  assert(freq(0) == 0);
                                                                     Poly<mint> mul;
  Poly f(\{1\}), g(\{1\});
                                                                     std::vector<mint> que;
  for (int i = 1; i < n; i *= 2) {
                                                                     MultiEval(const std::vector<mint> &que_, int off, int sz_) :
    g = (g * 2 - f * g * g).pre(i);
                                                                        sz(sz_) {
    Poly q = diff().pre(i - 1);
                                                                       if (sz <= 100) {
    Poly w = (q + g * (f.diff() - f * q)).pre(2 * i - 1);
                                                                         que = {que_.begin() + off, que_.begin() + off + sz};
    f = (f + f * (*this - w.inte()).pre(2 * i)).pre(2 * i);
                                                                         mul = {{1}};
                                                                         for (auto x : que) {
  return f.pre(n);
                                                                          mul *= {{-x, 1}};
Poly log(int n) const {
                                                                         return;
  assert(freq(0) == 1);
  auto f = pre(n);
                                                                      l = new MultiEval(que_, off, sz / 2);
  return (f.diff() * f.inv(n - 1)).pre(n - 1).inte();
                                                                       r = new MultiEval(que_, off + sz / 2, sz - sz / 2);
                                                                      mul = l->mul * r->mul;
Poly pow(int n, i64 k) const {
  int m = 0;
                                                                     MultiEval(const std::vector<mint> &que ) : MultiEval(que . 0.
  while (m < n && freq(m) == 0) m++;</pre>
                                                                         int(que_.size())) {}
  Poly f(std::vector<D>(n, 0));
                                                                     void query(const Poly<mint> &pol_, std::vector<mint> &res)
  if (k && m && (k >= n || k * m >= n)) return f;
                                                                       const {
  f.v.resize(n);
                                                                       if (sz <= 100) {
  if (m == n) return f.v[0] = 1, f;
                                                                         for (auto x : que) {
  int le = m * k;
                                                                           mint sm = 0, base = 1;
  Poly g({v.begin() + m, v.end()});
                                                                           for (int i = 0; i < pol_.size(); i++) {</pre>
  D base = power<D>(g.freq(0), k), inv = g.freq(0).inv();
                                                                             sm += base * pol_.freq(i);
  g = ((g * inv).log(n - m) * D(k)).exp(n - m);
                                                                             base *= x;
  for (int i = le; i < n; i++) f.v[i] = g.freq(i - le) * base</pre>
                                                                           }
                                                                          res.push_back(sm);
  return f;
                                                                         return;
Poly Getsgrt(int n) const {
  if (size() == 0) return {{0}};
                                                                       auto pol = pol_ % mul;
  int z = QuadraticResidue(freq(0).v, 998244353);
                                                                      l->query(pol, res);
  if (z == -1) return Poly{};
                                                                       r->query(pol, res);
  Poly f = pre(n + 1);
  Poly g({z});
                                                                     std::vector<mint> query(const Poly<mint> &pol) const {
  for (int i = 1; i < n; i *= 2) {
                                                                       std::vector<mint> res;
    g = (g + f.pre(2 * i) * g.inv(2 * i)) / 2;
                                                                       query(pol, res);
                                                                       return res;
  return g.pre(n + 1);
                                                                  };
                                                                  template <class mint>
Polv sart(int n) const {
                                                                  Poly<mint> berlekampMassey(const std::vector<mint> &s) {
  int m = 0:
                                                                     int n = int(s.size());
  while (m < n && freq(m) == 0) m++;</pre>
                                                                     std::vector<mint> b = {mint(-1)}, c = {mint(-1)};
  if (m == n) return {{0}};
                                                                     mint y = mint(1);
  if (m & 1) return Poly{};
  Poly s = Poly(std::vector<D>(v.begin() + m, v.end())).
                                                                     for (int ed = 1; ed <= n; ed++) {</pre>
                                                                       int l = int(c.size()), m = int(b.size());
  Getsart(n):
                                                                       mint x = 0;
  if (s.size() == 0) return Poly{};
                                                                       for (int i = 0; i < l; i++) {
  std::vector<D> res(n);
                                                                        x += c[i] * s[ed - l + i];
  for (int i = 0; i + m / 2 < n; i++) res[i + m / 2] = s.freq
  (i);
                                                                       b.push_back(0);
  return Poly(res);
                                                                      m++:
                                                                       if (!x) {
Poly modpower(u64 n, const Poly &mod) {
                                                                         continue;
  Poly x = *this, res = {\{1\}};
  for (; n; n \neq 2, x = x * x % mod) {
                                                                       mint freq = x / y;
    if (n & 1) {
                                                                       if (l < m) {
      res = res * x % mod;
                                                                         // use b
    }
                                                                         auto tmp = c;
                                                                         c.insert(begin(c), m - l, mint(0));
  return res:
                                                                         for (int i = 0; i < m; i++) {</pre>
                                                                           c[m - 1 - i] -= freq * b[m - 1 - i];
friend std::ostream &operator<<(std::ostream &os, const Poly</pre>
  8p) {
                                                                         b = tmp;
  if (p.size() == 0) {
    return os << "0";
                                                                       } else {
```

```
// use c
       for (int i = 0; i < m; i++) {</pre>
         c[l - 1 - i] -= freq * b[m - 1 - i];
    }
  }
   return c;
 template <class E, class mint = decltype(E().f)>
 mint sparseDet(const std::vector<std::vector<E>>> &g) {
   int n = int(g.size());
   if (n == 0) {
    return 1;
   }
   auto randV = [&]() {
     std::vector<mint> res(n);
     for (int i = 0; i < n; i++) {
       res[i] = mint(std::uniform_int_distribution<i64>(1, mint
      (-1).v)(rng)); // need rng
     }
     return res;
   }:
   std::vector<mint> c = randV(), l = randV(), r = randV();
   // l * mat * r
   std::vector<mint> buf(2 * n);
   for (int fe = 0; fe < 2 * n; fe++) {
     for (int i = 0; i < n; i++) {</pre>
       buf[fe] += l[i] * r[i];
     for (int i = 0; i < n; i++) {</pre>
       r[i] *= c[i];
     std::vector<mint> tmp(n);
     for (int i = 0; i < n; i++) {
       for (auto e : g[i]) {
         tmp[i] += r[e.to] * e.f;
       = tmp;
  }
   auto u = berlekampMassey(buf);
   if (u.size() != n + 1) {
     return sparseDet(g);
   auto acdet = u.freq(0) * mint(-1);
   if (n % 2) {
     acdet *= mint(-1);
   if (!acdet) {
     return 0;
   mint cdet = 1;
   for (int i = 0; i < n; i++) {
     cdet *= c[i];
   return acdet / cdet;
| }
```

7.21 Theorem

• Pick's Theorem

 $A=i+\frac{b}{2}-1$ A: Area ` i: grid number in the inner ` b: grid number on the side

• Matrix–Tree theorem undirected graph $D_{ii}(G) = \deg(i), D_{ij} = 0, i \neq j \\ A_{ij}(G) = A_{ji}(G) = \#e(i,j), i \neq j \\ L(G) = D(G) - A(G) \\ t(G) = \det L(G)\binom{1,2,\cdots,i-1,i+1,\cdots,n}{1,2,\cdots,i-1,i+1,\cdots,n}$ leaf to root $D_{ii}^{out}(G) = \deg^{out}(i), D_{ij}^{out} = 0, i \neq j \\ A_{ij}(G) = \#e(i,j), i \neq j \\ L^{out}(G) = D^{out}(G) - A(G) \\ t^{root}(G,k) = \det L^{out}(G)\binom{1,2,\cdots,k-1,k+1,\cdots,n}{1,2,\cdots,k-1,k+1,\cdots,n}$ root to leaf $L^{in}(G) = D^{in}(G) - A(G) \\ t^{leaf}(G,k) = \det L^{in}(G)\binom{1,2,\cdots,k-1,k+1,\cdots,n}{1,2,\cdots,k-1,k+1,\cdots,n}$

• Derangement $D_n = (n-1)(D_{n-1} + D_{n-2}) = nD(n-1) + (-1)^n$

• Möbius Inversion $f(n) = \sum_{d \mid n} g(d) \Leftrightarrow g(n) = \sum_{d \mid n} \mu(\tfrac{n}{d}) f(d)$

• Euler Inversion $\sum_{i|n} \varphi(i) = n$

· Binomial Inversion

$$f(n) = \sum_{i=0}^{n} \binom{n}{i} g(i) \Leftrightarrow g(n) = \sum_{i=0}^{n} (-1)^{n-i} \binom{n}{i} f(i)$$

· Subset Inversion

$$f(S) = \sum_{T \subseteq S} g(T) \Leftrightarrow g(S) = \sum_{T \subseteq S} (-1)^{|S| - |T|} f(T)$$

Min-Max Inversion

$$\max_{i \in S} x_i = \sum_{T \subseteq S} (-1)^{|T|-1} \min_{j \in T} x_j$$

• Ex Min-Max Inversion

kthmax
$$x_i = \sum_{T \subseteq S} (-1)^{|T|-k} {|T|-1 \choose k-1} \min_{j \in T} x_j$$

• Lcm-Gcd Inversion

$$\lim_{i \in S} x_i = \prod_{T \subseteq S} \left(\gcd_{j \in T} x_j \right)^{(-1)|T|-1}$$

· Sum of powers

$$\begin{array}{l} \sum_{k=1}^{n}k^{m}=\frac{1}{m+1}\sum_{k=0}^{m}\binom{m+1}{k}\,B_{k}^{+}\,n^{m+1-k}\\ \sum_{j=0}^{m}\binom{m+1}{j}B_{j}^{-}=0\\ \text{note: }B_{1}^{+}=-B_{1}^{-},B_{i}^{+}=B_{i}^{-} \end{array}$$

Cayley's formula

number of trees on n labeled vertices: n^{n-2} Let $T_{n,k}$ be the number of labelled forests on n vertices with k connected components, such that vertices 1, 2, ..., k all belong to different connected components. Then $T_{n,k}=kn^{n-k-1}$.

• High order residue

$$\left[d^{\frac{p-1}{(n,p-1)}} \equiv 1\right]$$

 • Packing and Covering $|\mathsf{maximum}| \mathsf{independent}| \mathsf{set}| + |\mathsf{minimum}| \mathsf{vertex}| \mathsf{cover}| = |V|$

Końig's theorem

|maximum matching| = |minimum vertex cover|

 Dilworth's theorem width = |largest antichain| = |smallest chain decomposition|

• Mirsky's theorem height = |lon

height = |longest chain| = |smallest antichain decomposition| = |minimum anticlique partition|

· Lucas'Theorem

For $n,m\in\mathbb{Z}^*$ and prime P, $\binom{m}{n}\mod P=\Pi\binom{m_i}{n_i}$ where m_i is the i-th digit of m in base P.

• Stirling approximation $n! \approx \sqrt{2\pi n} (\frac{n}{\epsilon})^n e^{\frac{1}{12n}}$

• 1st Stirling Numbers(permutation |P|=n with k cycles) S(n,k)= coefficient of x^k in $\prod_{i=0}^{n-1}(x+i)$ S(n+1,k)=nS(n,k)+S(n,k-1)

 $\bullet \ \ {\it 2nd Stirling Numbers(Partition} \ n \ {\it elements into} \ k \ {\it non-empty set)} \\$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$$

$$S(n+1,k) = kS(n,k) + S(n,k-1)$$

Catalan number

Catalar Humber
$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n-1} \\ \binom{n+m}{n} - \binom{n+m}{n+1} = (m+n)! \frac{n-m+1}{n+1} \quad \text{for} \quad n \geq m$$

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1 \quad \text{and} \quad C_{n+1} = 2(\frac{2n+1}{n+2})C_n$$

$$C_0 = 1 \quad \text{and} \quad C_{n+1} = \sum_{i=0}^n C_i C_{n-i} \quad \text{for} \quad n \geq 0$$

• Extended Catalan number

$$\frac{1}{(k-1)n+1} \binom{kn}{n}$$

• Calculate $c[i-j]+=a[i] \times b[j]$ for a[n],b[m] 1. a=reverse(a); c=mul(a,b); c=reverse(c[:n]); 2. b=reverse(b); c=mul(a,b); c=rshift(c,m-1);

• Eulerian number (permutation $1 \sim n$ with $m \; a[i] > a[i-1]$)

$$A(n,m) = \sum_{i=0}^{m} (-1)^{i} {\binom{n+1}{i}} (m+1-i)^{n}$$

$$A(n,m) = (n-m)A(n-1,m-1) + (m+1)A(n-1,m)$$

Hall's theorem

Let G=(X+Y,E) be a bipartite graph. For $W\subseteq X$, let $N(W)\subseteq Y$ denotes the adjacent vertices set of W. Then, G has a X'-perfect matching (matching contains $X'\subseteq X$) iff $\forall W\subseteq X', |W|\le |N(W)|$.

• Tutte Matrix:

For a graph
$$G=(V,E)$$
, its maximum matching $=\frac{rank(A)}{2}$ where $A_{ij}=((i,j)\in E?(i< j?x_{ij}:-x_{ji}):0)$ and x_{ij} are random numbers.

Erdoš-Gallai theorem

There exists a simple graph with degree sequence $d_1 \geq \cdots \geq d_n$ iff n

$$\sum_{i=1}^n d_i \text{ is even and } \sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k), \forall 1 \leq k \leq n$$

• Euler Characteristic

```
planar graph: V-E+F-C=1 convex polyhedron: V-E+F=2
```

V, E, F, C: number of vertices, edges, faces(regions), and components

```
• Burnside Lemma
    |X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|
· Polya theorem
    |Y^x/G| = \frac{1}{|G|} \sum_{g \in G} m^{c(g)}
    m = |Y|: num of colors, c(g): num of cycle
· Cayley's Formula
    Given a degree sequence d_1,\ldots,d_n of a labeled tree, there are
    \frac{(n-2)!}{(d_1-1)!\cdots(d_n-1)!} spanning trees.
• Find a Primitive Root of n:
    n has primitive roots iff n=2,4,p^k,2p^k where p is an odd prime.
    1. Find \phi(n) and all prime factors of \phi(n), says P=\{p_1,...,p_m\}
    2. \forall g \in [2,n), if g^{\frac{\phi(n)}{p_i}} \neq 1, \forall p_i \in P, then g is a primitive root. 3. Since the smallest one isn't too big, the algorithm runs fast.
    4. n has exactly \phi(\phi(n)) primitive roots.
· Taylor series
    f(x) = f(c) + f'(c)(x - c) + \frac{f^{(2)}(c)}{2!}(x - c)^2 + \frac{f^{(3)}(c)}{2!}(x - c)^3 + \cdots
· Lagrange Multiplier
    \min f(x,y), subject to g(x,y)=0
    \frac{\partial f}{\partial x} + \lambda \frac{\partial g}{\partial x} = 0
\frac{\partial f}{\partial y} + \lambda \frac{\partial g}{\partial y} = 0
g(x, y) = 0
• Calculate f(x+n) where f(x) = \sum_{i=0}^{n-1} a_i x^i
    f(x+n) = \sum_{i=0}^{n-1} a_i (x+n)^i = \sum_{i=0}^{n-1} x^i \cdot \frac{1}{i!} \sum_{j=i}^{n-1} \frac{a_j}{j!} \cdot \frac{n^{j-i}}{(j-i)!}
• Bell 數 (有 n 個人, 把他們拆組的方法總數)
    \begin{array}{l} B_0 = 1 \\ B_n = \sum_{k=0}^n s(n,k) \quad (second-stirling) \\ B_{n+1} = \sum_{k=0}^n {n \choose k} B_k \end{array}
· Wilson's theorem
    (p-1)! \equiv -1 (\mod p)
    (p^q!)_p \equiv \begin{cases} 1, & (p=2) \land (q \geq 3), \\ -1, & \text{otherwise.} \end{cases} \pmod{p}^q
· Fermat's little theorem
    a^p \equiv a \pmod{p}
· Euler's theorem
   a^b \equiv \begin{cases} a^{b \bmod \varphi(m)}, \\ a^{b}, \\ \vdots, \end{cases}
                                                  gcd(a, m) = 1,
                                                  \gcd(a,m) \neq 1, b < \varphi(m), \pmod{m}
               a^{(b \mod \varphi(m)) + \varphi(m)}, \quad \gcd(a, m) \neq 1, b \geq \varphi(m).
• 環狀著色(相鄰塗異色)
    (k-1)(-1)^n + (k-1)^n
```

Stringology

8.1 Aho-Corasick AM

```
struct ACM {
 int idx = 0;
 vector<array<int, 26>> tr;
 vector<int> cnt, fail;
 void clear() {
   tr.resize(1, array<int, 26>{});
   cnt.resize(1, 0);
    fail.resize(1, 0);
 }
 ACM() {
   clear();
 int newnode() {
   tr.push_back(array<int, 26>{});
   cnt.push_back(0);
   fail.push_back(0);
   return ++idx;
 void insert(string &s) {
   int u = 0;
    for (char c : s) {
     if (tr[u][c] == 0) tr[u][c] = newnode();
      u = tr[u][c];
   cnt[u]++;
```

```
void build() {
     queue<int> q;
     rep (i, 0, 26) if (tr[0][i]) q.push(tr[0][i]);
     while (!q.empty()) {
      int u = q.front(); q.pop();
       rep (i, 0, 26) {
        if (tr[u][i]) {
           fail[tr[u][i]] = tr[fail[u]][i];
           cnt[tr[u][i]] += cnt[fail[tr[u][i]]];
           q.push(tr[u][i]);
         } else {
          tr[u][i] = tr[fail[u]][i];
        }
      }
    }
   int query(string &s) {
     int u = 0, res = 0;
     for (char c : s) {
      c -= 'a';
      u = tr[u][c];
      res += cnt[u];
     return res;
  }
};
8.2 Double String
```

```
// need zvalue
int ans = 0:
auto dc = [&](auto self, string cur) -> void {
   int m = cur.size();
   if (m <= 1) return;</pre>
   string _s = cur.substr(0, m / 2), _t = cur.substr(m / 2, m);
   self(self, _s);
  self(self, _t);
rep (T, 0, 2) {
     int m1 = _s.size(), m2 = _t.size();
     string s = _t + "$" + _s, t = _s;
reverse(all(t));
     zvalue z1(s), z2(t);
     auto get_z = [&](zvalue &z, int x) -> int {
       if (0 <= x && x < z.z.size()) return z[x];</pre>
       return 0;
     rep (i, 0, m1) if (_s[i] == _t[0]) {
       int len = m1 - i:
       int L = m1 - min(get_z(z2, m1 - i), len - 1),
         R = get_z(z1, m2 + 1 + i);
       if (T == 0) R = min(R, len - 1);
       R = i + R;
       ans += \max(0, R - L + 1);
     ļ
     swap(_s, _t);
     reverse(all(s));
     reverse(all(_t));
};
dc(dc, str);
```

8.3 Lyndon Factorization

```
| // partition s = w[0] + w[1] + ... + w[k-1],
// w[0] >= w[1] >= ... >= w[k-1]
// each w[i] strictly smaller than all its suffix
// min rotate: last < n of duval_min(s + s)</pre>
// max rotate: last < n of duval_max(s + s)</pre>
// min suffix: last of duval_min(s)
// max suffix: last of duval_max(s + -1)
vector<int> duval(const auto &s) {
   int n = s.size(), i = 0;
   vector<int> pos;
   while (i < n) {
     int j = i + 1, k = i;
     while (j < n \text{ and } s[k] <= s[j]) { // >=}
       if (s[k] < s[j]) k = i; // >
       else k++;
       j++;
     while (i <= k) {
       pos.push_back(i);
```

for (int j = sa[rk[i] - 1];

i + h < n and j + h < n

```
i += j - k;
                                                                               and s[i + h] == s[j + h];) ++h;
    }
                                                                          lcp[rk[i] - 1] = h ? h-- : 0;
  }
                                                                        }
                                                                      }
  pos.push_back(n);
                                                                   };
                                                                          Suffix Array
                                                                    8.6
       Manacher
 8.4
                                                                   | struct SuffixArray {
/* center i: radius z[i * 2 + 1] / 2
                                                                      int n;
  center i, i + 1: radius z[i * 2 + 2] / 2
                                                                      vector<int> suf, rk, S;
                                                                      SuffixArray(vector<int> _S) : S(_S) {
  both aba, abba have radius 2 */
                                                                        n = S.size();
 vector<int> manacher(const string &tmp) { // 0-based
  string s = "%";
                                                                        suf.assign(n, 0);
                                                                        rk.assign(n * 2, -1);
  int l = 0, r = 0;
                                                                        iota(all(suf), 0);
   for (char c : tmp) s += c, s += '%';
                                                                        for (int i = 0; i < n; i++) rk[i] = S[i];
  vector<int> z(ssize(s));
                                                                        for (int k = 2; k < n + n; k *= 2) {
  for (int i = 0; i < ssize(s); i++) {
                                                                          auto cmp = [&](int a, int b) -> bool {
    z[i] = r > i ? min(z[2 * l - i], r - i) : 1;
    while (i - z[i] \ge 0 \& i + z[i] < ssize(s) \& s[i + z[i]]
                                                                            return rk[a] == rk[b] ? (rk[a + k / 2] < rk[b + k / 2])
                                                                                  : (rk[a] < rk[b]);
     == s[i - z[i]])
     ++z[i];
    if(z[i] + i > r) r = z[i] + i, l = i;
                                                                          sort(all(suf), cmp);
                                                                          tmp[suf[0]] = 0;
   return z:
                                                                          for (int i = 1; i < n; i++) {</pre>
| }
                                                                            tmp[suf[i]] = tmp[suf[i - 1]] + cmp(suf[i - 1], suf[i])
 8.5 SA-IS
auto sais(const auto &s) {
                                                                          rk.swap(tmp);
  const int n = (int)s.size(), z = ranges::max(s) + 1;
                                                                        }
  if (n == 1) return vector{0};
                                                                      }
  vector<int> c(z); for (int x : s) ++c[x];
                                                                   |};
  partial_sum(all(c), begin(c));
                                                                    8.7 Z-value
  vector<int> sa(n); auto I = views::iota(0, n);
   vector<bool> t(n); t[n - 1] = true;
                                                                    struct zvalue {
  for (int i = n - 2; i >= 0; i--)
t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
                                                                      vector<int> z
                                                                      int operator[] (const int &x) const {
  auto is_lms = views::filter([&t](int x) {
                                                                        return z[x]:
    return x && t[x] & !t[x - 1];
                                                                      zvalue(string s) {
  auto induce = [8] {
                                                                        int n = s.size();
    for (auto x = c; int y : sa)
                                                                        z.resize(n);
       if (y-- and !t[y]) sa[x[s[y] - 1]++] = y;
                                                                        z[0] = 0;
     for (auto x = c; int y : sa | views::reverse)
                                                                        for (int i = 1, l = 1, r = 0; i < n; i++) {
       if (y-- and t[y]) sa[--x[s[y]]] = y;
                                                                          z[i] = min(z[i - l], max<int>(0, r - i));
                                                                          while (i + z[i] < n \delta s[i + z[i]] == s[z[i]]) z[i]++;
  vector<int> lms, q(n); lms.reserve(n);
                                                                          if (i + z[i] > r) l = i, r = i + z[i];
  for (auto x = c; int i : I | is_lms) {
    q[i] = int(lms.size());
                                                                      }
    lms.push_back(sa[--x[s[i]]] = i);
                                                                   };
  induce(); vector<int> ns(lms.size());
                                                                    9
                                                                          Peppa
  for (int j = -1, nz = 0; int i : sa | is_lms) {
    if (j >= 0) {
                                                                    9.1 LinearSolve
       int len = min({n - i, n - j, lms[q[i] + 1] - i});
                                                                   | // ax + b = 0 \pmod{m}
       ns[q[i]] = nz += lexicographical_compare(
                                                                    std::pair<i64, i64> sol(i64 a, i64 b, i64 m) {
        s.begin() + j, s.begin() + j + len,
                                                                      assert(m > 0);
         s.begin() + i, s.begin() + i + len
                                                                      b *= -1;
      ):
                                                                      i64 x, y;
    }
                                                                      i64 g = exgcd(a, m, x, y);
    j = i;
                                                                      if (g < 0) {
  }
                                                                        g *= -1, x *= -1, y *= -1;
  ranges::fill(sa, 0); auto nsa = sais(ns);
  for (auto x = c; int y : nsa | views::reverse)
                                                                      if (b % g != 0) return {-1, -1};
    y = lms[y], sa[--x[s[y]]] = y;
                                                                      x = x * (b / g) % (m / g);
  return induce(), sa;
                                                                      if(x < 0) {
                                                                        x += m / g;
// sa[i]: sa[i]-th suffix is the
                                                                      }
 // i-th lexicographically smallest suffix.
                                                                      return {x, m / g};
 // lcp[i]: LCP of suffix sa[i] and suffix sa[i + 1].
                                                                   }
 struct Suffix {
  int n;
                                                                    9.2 LinearSolve
   vector<int> sa, rk, lcp;
  Suffix(const auto &s) : n(s.size()),
                                                                   template<typename M = ll>
                                                                    void NTT(vector<M> &P, M w, bool inv = 0) {
    lcp(n - 1), rk(n) {
    vector<int> t(n + 1); // t[n] = 0
                                                                      int n = P.size(), lg = __builtin_ctz(n);
    copy(all(s), t.begin()); // s shouldn't contain 0
                                                                      assert(__builtin_popcount(n) == 1);
                                                                      for (int j = 1, i = 0; j < n - 1; ++j) {
     sa = sais(t); sa.erase(sa.begin());
                                                                        for (int k = n >> 1; k > (i ^= k); k >>= 1); // !!!
     for (int i = 0; i < n; i++) rk[sa[i]] = i;
    for (int i = 0, h = 0; i < n; i++) {</pre>
                                                                        if (j < i) swap(P[i], P[j]);</pre>
       if (!rk[i]) { h = 0; continue; }
```

vector<M> ws = $\{inv ? M\{1\} * fpow(w, MOD - 2, MOD) : w\};$

rep (i, 1, lg) ws.pb(ws[i - 1] * ws[i - 1] % MOD);

```
reverse(all(ws));
   rep (i, 0, lg) {
     for (int k = 0; k < n; k += 2 << i) {
       M base = M\{1\};
        rep (j, k, k + (1 << i)) {
          auto t = base * P[j + (1 << i)] % MOD;</pre>
          auto u = P[j];
          P[j] = (u + t) \% MOD;
          P[j + (1 << i)] = (u - t + MOD) % MOD;
          base = base * ws[i] % MOD;
     }
   }
   if (inv) rep (i, 0, n) P[i] = P[i] * fpow(n, MOD - 2, MOD) %
 const int N = 1 << 20;</pre>
 const auto w = fpow(3, (MOD - 1) / N, MOD);
 9.3 FractionSearch
// Binary search on Stern-Brocot Tree
 // Parameters: n, pred
 // n: Q_n is the set of all rational numbers whose denominator
      does not exceed n
 // pred: pair<i64, i64> -> bool, pred({0, 1}) must be true
 // Return value: {{a, b}, {x, y}}
 // a/b is bigger value in Q_n that satisfy pred()
 // x/y is smaller value in Q_n that not satisfy pred()
 // Complexity: 0(log^2 n)
 using Pt = pair<i64, i64>;
 Pt operator+(Pt a, Pt b) { return {a.ff + b.ff, a.ss + b.ss}; }
 Pt operator*(i64 a, Pt b) { return {a * b.ff, a * b.ss}; }
 pair<pair<i64, i64>, pair<i64, i64>> FractionSearch(i64 n,
      const auto &pred) {
   pair<i64, i64> low{0, 1}, hei{1, 0};
   while (low.ss + hei.ss <= n) {</pre>
     bool cur = pred(low + hei);
     auto &fr{cur ? low : hei}, &to{cur ? hei : low};
     u64 L = 1, R = 2;
     while ((fr + R * to).ss \le n \text{ and } pred(fr + R * to) == cur)
     {
L *= 2;
'= 2;
       R *= 2;
     while (L + 1 < R) {
       u64 M = (L + R) / 2;
        ((fr + M * to).ss \le n \text{ and } pred(fr + M * to) == cur ? L :
       R) = M;
     fr = fr + L * to;
   }
   return {low, hei};
| }
 9.4 Triangular
    • Cosine Law (餘弦定理)
      c^{2} = a^{2} + b^{2} - 2ab\cos C
b^{2} = a^{2} + c^{2} - 2ac\cos B
a^{2} = b^{2} + c^{2} - 2bc\cos A
    • Weierstrass Substitution (t-代換)
      設 t = \tan \frac{\theta}{2},則有:
      \sin\theta = \frac{2\tilde{t}}{1+t^2},\quad \cos\theta = \frac{1-t^2}{1+t^2},\quad d\theta = \frac{2}{1+t^2}\,dt
    • Brahmagupta's Formula (海龍公式, 四邊形版本)
      若四邊形為圓內接,邊長 a,b,c,d,半周長 s=\frac{a+b+c+d}{2},則:
      A = \sqrt{(s-a)(s-b)(s-c)(s-d)}
一般四邊形 (Bretschneider's formula):
      A = \sqrt{(s-a)(s-b)(s-c)(s-d) - abcd\cos^2\left(\frac{A+C}{2}\right)}
```