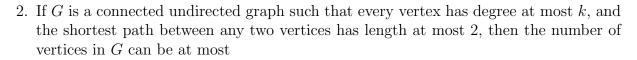
Mathematics for Data Science - 1 Graded Assignment

Week 10

1 MULTIPLE CHOICE QUESTIONS:

- 1. An undirected graph G has 20 vertices and the degree of each vertex is at least 3 and at most 5. Which of the following statements is true regarding the graph G?
 - (a) The minimum number of edges that the graph G can have is 60.
 - (b) The maximum number of edges that the graph G can have is 100.
 - (c) The maximum number of edges that the graph G can have is 60.
 - (d) The minimum number of edges that the graph G can have is 30.

we know that, Sun of degree of all the Volices is twice the number of edges in the graph. * Consider the minimum possible case: Suppose every Vertexp in the graph of has degree 3. Sum of degree of all vertices = 20x3 = 60 > Mirimum number of edges that the graph of can have is $\frac{60}{3} = 30$. * Graider the maximum possible case: Suppose every Voltep in the graph of how degree 5. Sun of degree of all volices = 20×5 = 100 > Maximum number of edges that the graph of can have is 100 = 50. : option (d) is correct.



(a)
$$k^2 - 1$$

(b)
$$k^2 + 1$$

(c)
$$k^2$$

(d)
$$k^2 - k$$

As the degree of every verters is at most 'k', the number of vertices that one adjacent to 'u' can be at most 'k'.

Vertices that one adjacent to 'u' can be at most 'k'.

Also, if we drow a BFS tree starting with verter 'u', then the depth of the tree cannot be more than 2 because the length of the short the depth of the tree cannot be more than 2.

-est Path between any two vertices is at most 2.

V₁ V₂ V₃ V_K

Now each of V, can have at most K-1 adjacent volices because 'u' is already adjacent to each of V?.

BFS tree with at most voltices as be

a, a₂ a_{k-1}

(K-1) + (K-1) + --- + (K-1) (K times.)

: level 2 of BPS tree can have almost K(k-1) Valices. => the graph of can have almost 1 + K + K(k-1)(level 0) (level 1) (level 2)

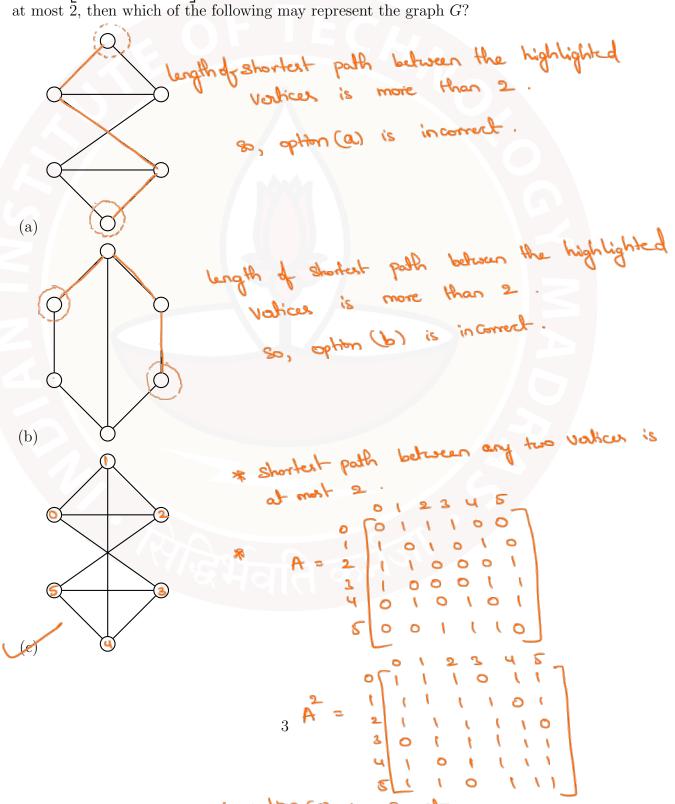
Hence, the graph can have at most k2+1 Volices.

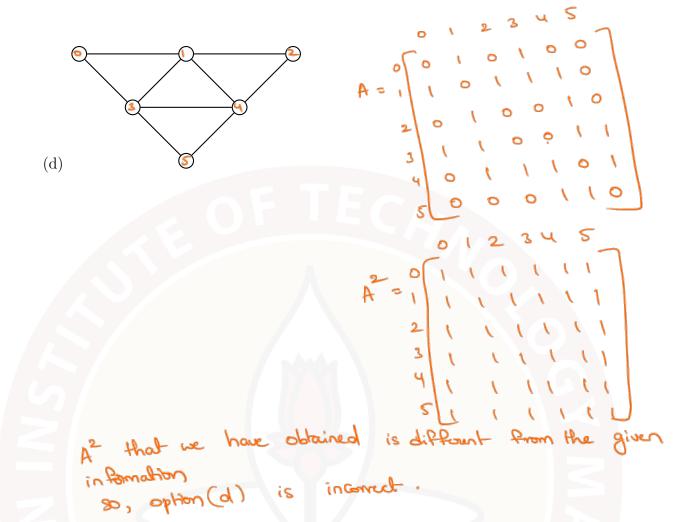
option (b) is Correct.

3. Suppose A is the adjacency matrix of a connected undirected graph G.

1 1						•
	Γ1	1	1	0	1	1
If $A^2 =$	1	1	1	1	0	1
	$\begin{vmatrix} 1 \\ 0 \end{vmatrix}$	1 1 1 1 0 1	1	1	1	0
	0	1	1	1	1	1
	1	0	1	1	1	1
	1	1	0	1	1	1

and the shortest path between any two vertices has length





2 MUTIPLE SELECT QUESTIONS:

4. Suppose in a farewell party of IIT Madras Mathematics department, 60 students were present. As in normal parties, handshaking took place and of course no one shook their own hand. The number of students who have made odd number of handshakes is x. Which of the followings can be a possible value of x?

(b) 13
(c) 21
(d) 28
Sel: We represent this in a graph model, every vertex represents a person and edge between a pair of vertices if the Persons representing those vertices made a handshake.

Now, observe that there are 'x' vertices with even degree with odd degree and '60-x' vertices with even degree.

As we know that number of odd degree vertices in a graph is always even.

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5. Suppose G is a graph with 6 vertices 0, 1, 2, 3, 4, 5 and the adjacency matrix of the

- (a) The graph G is a directed acyclic graph.
- From vertex 4, every other vertex is reachable.
- (c) The longest path in the graph G has length 4, in terms of number of edges.
- (d) The longest path in the graph is $4 \longrightarrow 0 \longrightarrow 2 \longrightarrow 3$

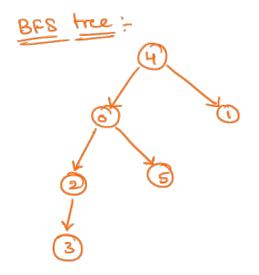
sof: Given, we first draw the graph



As the obtained graph does not have any cycles and also it is a directed graph => q is a directed acyclic graph.

So, option (a) is Correct.

* IP we draw a BFS tree stanking with Votep'y', then we can get the Vortices that are readhable from Votep'y'.



Every Vortup is present in the BFS tree

=> Every Volep is reachable from Verlyp '4'.

so, option (b) is correct.

* we have

from this we find A, A3, A4.

, og

: From AY, we can conclude that there is no path in the graph G that has length 4. So, option (c) is incorrect.

* From A^3 , $y \longrightarrow 3$ is the bongust path in the graph G_1 and the path is $y \longrightarrow 0 \longrightarrow 2 \longrightarrow 3$.

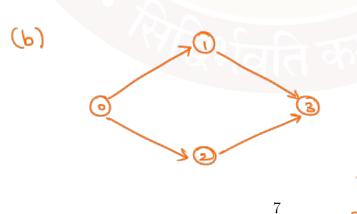
So, option (d) is Garrect.

- 6. Which of the following statements is (are) true?
 - (a) Every directed acyclic graph has a vertex with outdegree 0.
 - (b) In a directed acyclic graph, the longest path between a pair of vertices is always unique.
 - If an undirected graph G is connected, then the graph representing the transitive closure of the graph G is a complete graph.
 - Suppose A is the adjacency matrix of a graph G with n vertices, any non-zero entry (i, j) in the matrix A^k , where k < n, indicates that there is a path of length k from vertex i to vertex j.
 - (a) Suppose every vortage in a DAG has outdegree >0. If the graph has 'n' vortices, then there will be atteast 'n' outward edges (at least one for each value) choose a Vortup 'u', as the outdegree of vertex 'u' is greater NOW,

than o => there is an outward edge (e) from 'u' follow the edge 'e' and proceed to next vertex '3', even the autolignee (10) 70 => there is an outward edge (e) from 19.

Now follow e' and proceed to nept.

* After n steps, we will definely end up with a cycle which contradicts the fact that the graph is a directed acyclic graph. Hence, Every Verley in a DAG Cannot have outdegree 70 => Those will be a Verter which has outdegree O. So, option (a) is correct



this is a directed acydic graph. bengest path from @ to 3 Can be $\langle 0 \longrightarrow 1 \longrightarrow 3 \rangle$: Engest path may not be unique.

- (C) the graph representing the transitive closure of any given graph G is Gontructed by adding edges to the graph G such that if those is a path from varley 'à to varley 'b' in G then we add edge (a,b) to the graph G.
 - 80, if a graph 'H' is Gamacted, then we know that there will be a path from any vertup "i' to any other vertex" i' in the graph.

=> Hence every edge (?ij) is added to the graph 'H'
making 'H' a complete graph.

... the graph representing the transitive closure of a connected graph is complete.

80, option (c) is correct.

(d) we defined A^{K} in this way (given in the bectures).

option (d) is correct.

USE THE FOLLOWING INFORMATION FOR QUESTIONS [7-8]:

Shreya needs to perform 10 tasks namely $\{A, B, C, D, \dots J\}$. Some tasks needs to be performed after performing a particular task. In the below table, column 1 shows the tasks and column 2 shows the sets of tasks that can be performed only after performing the particular task.

$egin{array}{c ccccccccccccccccccccccccccccccccccc$	Ļ
$E = \{G\}$	
$G = \{J\}$	
$\begin{array}{c c} H & \{G,J\} \\ \hline I & \{J\} \\ \hline \end{array}$	
$J \mid \{\}$	

7. Which of the following sequences may represent the possible order in which Shreya can perform the tasks?

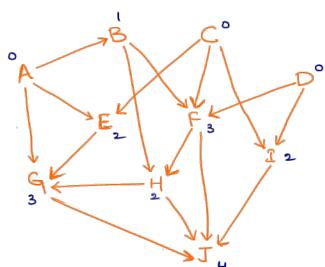
(a) A, C, B, D, E, I, F, H, G, J

(A, D, C, B, E, I, F, H, G, J)

(c) C, A, D, E, B, I, F, G, H, J

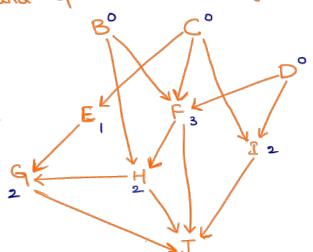
Sel: Draw a directed graph 'K' that represents the given data such that each vortex represents a task and a directed edge from vortex 'i' to vortex 'j' if task 'j' can be Performed only after task ".". * observe that the obtained graph is a DAG (Directed then in the Same Bider she can perform the tasks. acydic graph). Now, if we find . our aim is to find topological sequence of the obtained &DAG.

* Find in degree of each Verlep in the graph.



Now, we have three volices (A,C,D) that has indegree O. i we can choose any one of them i.e, she can perform any

Suppose we choose volue 'A', remove it from the DAG and update the indegree of each of the remaining vertex.

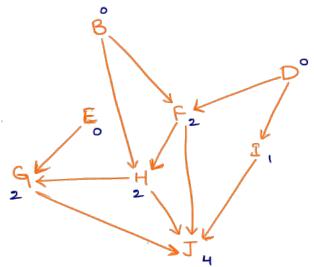


* indegree of Vortex B has become 0 and indegree (E) = 1, indegree $(G_1) = 2$.

* Repeat the process till the last Vortex is removed and added to the topological sequence.

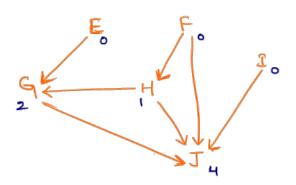
Now, we have again three Volices (B,C,D) having indegree

So we can choose any one of them. Suppose we choose vistep'c'. : Remove Volter 'C' from the graph and update the indegree of each of the Vertex, also add 'C' to the Topological Sequence.

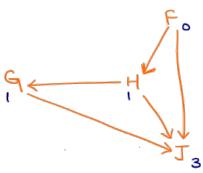




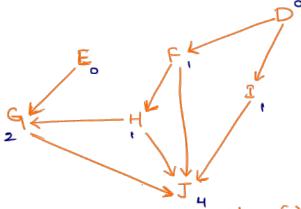
* Topological sequence:



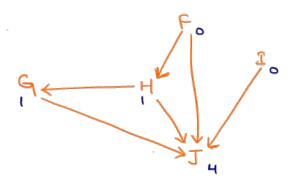
* Topdogical Sequence >

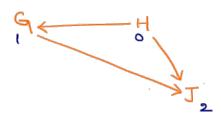


* topological Sequence:



* indegree (F) = 1, indegree (H) = 1





* indegree (H)=0, indegree of (J)=2

G J

* indegree (G) = 0, indegree (J)=1

* Topological Sequence: A,C,B,D,E,I,F,H, * indegree (I) = 0

* Topological Sequence:
A, C, B, D, E, I, F, H, G,

* Finally we remove 'J' and add it to the topological Sequence.

· Topological Sequence: A, C, B, D, E, I, F, H, G, J.

NOTE: The obtained topological Sequence is not unique because instead of stooting with A, we can also start with Verkep'c's volter

- : Shreya Can perform tosks in the order AICIBIDIEIIFIHIGIT
 Hence, option (a) is correct.
- Similarly, A,D,C,B,E,I,F,H,G,J and D,C,A,B,E,I,F,H,G,J are also a possible topological sequence which means shreya can perform tanks in that order too.

 tence, options (b),(d) are correct.
 - * Option (c) is incorrect because took 'G' Carnot be portamed before took 'H'. .: that is not a persible order.

3 NUMERICAL ANSWER TYPE:

8. If each task takes 5 minutes to complete and she performs all the independent tasks simultaneously, then the time(in minutes) taken by Shreya to complete all the tasks is Sel:- we compute the bengest path to each vortex in the DAG (that use got in the previous problem). * we first initialize languat-path-to (i) = 0 for each vortupi'i' in indegree (i) for each vorlep'in the DAG. langest-path-to (?). indegree (?) Now, we find a vertup'u' in the DAG which has indegree 0. * we remove Vortup'u' from the graph and update indegree(i) and longest-path-to(i) for every vertex'i' that is adjount to Vortep 'u'.

* we update langust-path-to(i) as tengent-path-to(i) = Max of langust-path-to(i), 1+ langust-path-to(u) for

* Repeat the process by Anding a new Vertep '19' that has indegree 0, removing it from the grouph and update indegree(i) and removing it from the bright-path-to(i) till all the Vertices are removed from the graph.

Finally, after removing all the Vortices and updating bengut-puth-to(1) for all Vertices, we get

langust-path-to(A) = langust-path-to(C) = langust-path-to(O) = 0.

langust-path-to(B) = langust-path-to(E) = langust-path-to(I) = 1.

langust-path-to(F) = 2.

langust-path-to(H) = 3.

langust-path-to(G) = 4.

langust-path-to(G) = 5.

She can first perform tarks 'A', C', D' at same time and next tarks 'B', E', I' at some time followed by tark 'F' then by tark 'H'
then by 'G' and finally tark 'J'

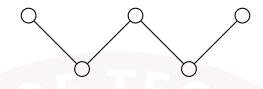
then by 'G' and finally tark 'J'

=> She takes 5 minutes (for tanks A,C,D) + 5 minutes (for tanks B,E,Z)
+ 5 minutes (tank F) + 5 minutes (tank H) + 5 minutes (tank G)
+ 5 minutes (tank T)

= 5+5+5+5+5 = 30 minutes.

Answer - 30

9. Suppose R is a relation defined on a set S and it is represented by a graph G that is shown below.

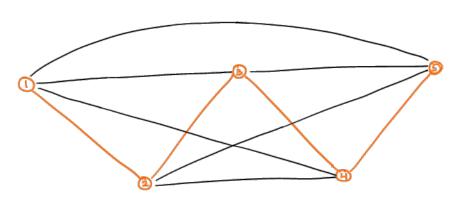


Find the number of edges that need to be added to the graph G such that the new graph obtained after adding the edges represents a transitive relation.

+ we have : R = { (1,2), (2,3), (3,4), (4,5)} The transitive clasure of R is obtained by repeatedly adding (a_1c) to R for each $(a_1b) \in R + (b_1c) \in R$. (1,2) & (2,3) one in R => (1,3) is added to R (2,3) & (3,4) in R \Rightarrow (2,4) is added to R \otimes (3,4) & (4,5) in R \Rightarrow (3,5) is added to R \otimes Similarly, Now again from 1, 2, 23, (1,3) & (3,4) one in R => (1,4) is added to R (1,3) & (3,5) one in R => (1,5) is added to R (2,4) & (4,5) are in R > (2,5) is added to R. .. the set which represents the transitive desure of R is S (1,2), (2,3), (1,3), (3,4), (2,4), (1,4), (4,5), (1,5),

(2,5), (3,5)

So the graph representing this transitive desure of R is



. . 6 edges need to be added to the given graph G.

NOTE: If a graph is Connected, then the transitive clasure of that graph is a complete graph.