

Models of Calculation Project

The Language Lua

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1 Introduction

In this article, we're going to define axiomatic semantics for coroutines in Lua. We're going to start by defining the basics of Lua, including some basic data types which we'll use in our examples. Then follows a description of coroutines, which is the built-in concurrency tool of Lua. Coroutines in Lua are collaborative, which means a coroutine must explicitly yield its cpu to other coroutines. Only one coroutine ever runs at any time, which means coroutines do not offer parallelism. They still offer a useful abstraction though. As While doesn't offer any kind of concurrency, we believe Lua is sufficiently special.

1.1 Example code

— from <http://lua-users.org/wiki/CoroutinesTutorial>

```
function odd(x)
  print('A: odd', x)
  coroutine.yield(x)
  print('B: odd', x)
end

function even(x)
  print('C: even', x)
  if x==2 then return x end
  print('D: even ', x)
end

co = coroutine.create(
  function (x)
    for i=1,x do
      if i==3 then coroutine.yield(-1) end
      if i % 2 == 0 then even(i) else odd(i) end
    end
  end)

count = 1
while coroutine.status(co) ~= 'dead' do
  print('----', count) ; count = count+1
  errorfree, value = coroutine.resume(co, 5)
  print('E: errorfree, value, status', errorfree, value, coroutine.status(co))
end
```

2 Description of the syntax

The following section describes the grammar of the subset of Lua that we'll be describing. Numerous Lua constructs have been left out for the sake of simplicity, like multiple assignment for example. Whitespace has also been left out of the grammar, as have the optional semicolons.

$$\begin{aligned}
\langle block \rangle &::= \langle stmt \rangle \langle block \rangle \mid '' \\
\langle expr \rangle &::= \langle bool \rangle \mid \langle number \rangle \mid \langle string \rangle \mid \langle table \rangle \mid \langle func \rangle \mid \langle nil \rangle \mid \langle table \rangle \mid \langle co-wrap \rangle \mid \langle lambda \rangle \\
\langle nil \rangle &::= \text{'nil'} \\
\langle bool \rangle &::= \text{'true'} \mid \text{'false'} \\
\langle stmt \rangle &::= \langle if \rangle \mid \langle while \rangle \mid \langle assign \rangle \mid \langle return \rangle \mid \langle expr \rangle \mid \langle co-yield \rangle \mid \langle co-resume \rangle \mid \langle call \rangle \\
\langle assign \rangle &::= \langle var \rangle \text{'='} \langle expr \rangle \\
\langle return \rangle &::= \text{'return'} \langle expr \rangle \\
\langle if \rangle &::= \text{'if'} \langle expr \rangle \text{'then'} \langle block \rangle \langle else-ifs \rangle \langle else \rangle \text{'end'} \\
\langle else-ifs \rangle &::= \text{'elseif'} \langle expr \rangle \text{'then'} \langle block \rangle \langle else-ifs \rangle \mid '' \\
\langle else \rangle &::= \text{'else'} \langle block \rangle \mid '' \\
\langle while \rangle &::= \text{'while'} \langle expr \rangle \text{'do'} \langle block \rangle \text{'end'} \\
\langle co-wrap \rangle &::= \text{'coroutine.wrap'}(\langle func \rangle) \\
\langle co-yield \rangle &::= \langle var \rangle \text{'='} \text{'coroutine.yield'}(\langle expr \rangle) \\
\langle co-resume \rangle &::= \langle var \rangle \text{'='} \langle coroutine \rangle (\langle expr \rangle) \\
\langle call \rangle &::= \langle var \rangle \text{'='} \langle func \rangle (\langle expr \rangle) \\
\langle lambda \rangle &::= \text{'function'} (\langle var \rangle) \langle block \rangle \text{'end'}
\end{aligned}$$

This grammar is still incomplete.

2.1 Expected Problems

3 Description of the semantics

The normal rules will be worked out in the nearby future, as seen in appendix A. We will, however, start some of the harder semantical expressions, namely the thread and the table rules.

3.1 Approach

For this project we have chosen to apply Structural Operational Semantics. This is because with SOS we can easily add extra information to the states.

3.2 Concepts

We will mainly be explaining coroutines, if possible we will expand this with tables, yielding and resuming, pipes and filters. To explain this we will also need to describe functions.

3.2.1 States

Because we are going to use SOS we will look at very small steps concerning the states of the program.

3.2.2 Transitions

There are two types of transitions. We have the form where a statement and a state go to a state: $\langle S, s \rangle \Longrightarrow s'$. Secondly we have the form $\langle S, s \rangle \Longrightarrow \langle S', s' \rangle$. This is from a statement and state to a (new) statements and a (new) state.

3.2.3 Types

Lua does not have a strong type system. Whenever a value is called the value is passed and the user doesn't need to apply which type this value has. In this model we will restrict ourselves to Coroutines, Tables, Booleans, Strings and Integers as types.

```

Boolean ::= tt|ff
Integer  ::= 0|1|2|...
String   ::= < Char, String > |λ
Char     ::= a|...|z|A|...|Z
Table    ::= Value → Value
Value    ::= Boolean|Integer|String|Char|Table|Coroutine
Coroutine ::= [tobecontinued]

```

3.3 Rules

[ass] $\langle x := a, s \rangle \Longrightarrow \langle [x \mapsto \mathcal{A}[[a]]], s \rangle$

[skip] $\langle \text{skip}, s \rangle \Longrightarrow s$

[comp]
$$\frac{\langle S_1, s \rangle \Longrightarrow \langle S'_1, s' \rangle}{\langle S_1; S_2, s \rangle \Longrightarrow \langle S'_1; S_2, s' \rangle}$$

[iftt] $\langle \text{if } \mathbf{b} \text{ then } S_1 \text{ else } S_2, s \rangle \Longrightarrow \langle S_1, s \rangle$
 where $\mathcal{B}[[\mathbf{b}]]_s = \text{tt}$

[ifff] $\langle \text{if } \mathbf{b} \text{ then } S_1 \text{ else } S_2, s \rangle \Longrightarrow \langle S_2, s \rangle$
 where $\mathcal{B}[[\mathbf{b}]]_s = \text{ff}$

[while] $\langle \text{while } \mathbf{b} \text{ do } S, s \rangle \Longrightarrow \langle \text{if } \mathbf{b} \text{ then } (S; \text{while } \mathbf{b} \text{ do } S) \text{ else skip}, s \rangle$

[resume]
$$\frac{\langle B_I, \text{let}(\text{merge}(\text{env}_I, \text{env}_O) r_I, \mathbb{E}[[a_O]] \text{env}_O) \rangle \Longrightarrow^* \langle r'_I = \text{coroutine.yield}(a_I) B'_I, \text{env}'_I \rangle}{\langle r_O = c(a_O) B_O, \text{env}_O \rangle \Longrightarrow \langle B_O, \text{let}(\text{set}(\text{merge}(\text{env}_O, \text{env}'_I), c, (\text{env}'_I, r'_I, B'_I)), r_O, \mathbb{E}[[a_I]] \text{env}'_I) \rangle}$$

 where $\text{env}_I, r_I, B_I = \text{get}(\text{env}_O, c)$

[wrap] $\langle c = \text{coroutine.wrap}(f), \text{env} \rangle \Longrightarrow \text{set}(\text{env}, c, (\text{env}_f, v_f, B_f))$
 where $\text{env}_f, v_f, B_f = \mathbb{E}[[f]] \text{env}$.

3.4 Environment

$$\begin{aligned}
\text{Env} &= (\langle \text{var} \rangle \rightarrow \text{Val} \cup \text{Ref}, \text{Ref} \rightarrow \text{Val}, \text{Ref}) \\
\text{get}((L, G, \text{ref}_N), v) &= \begin{cases} G(L(v)) & \text{if } L(v) \in \text{Ref} \\ L(v) & \text{else} \end{cases} \\
\text{merge}((L, _, \text{ref}_{N_1}), (_, G, \text{ref}_{N_2})) &= (L, G, \max(\text{ref}_{N_1}, \text{ref}_{N_2})) \\
\text{let}((L, G, \text{ref}_N), v, x) &= \left(\left(v' \mapsto \begin{cases} x & \text{if } v = v' \\ L(v') & \text{else} \end{cases} \right), G, \text{ref}_N \right) \\
\text{set}((L, G, \text{ref}_N), v, c) &= \left(L, \left(r \mapsto \begin{cases} x & \text{if } L(v) = r \\ G(r) & \text{else} \end{cases} \right), \text{ref}_N \right) \\
\text{def}((L, G, \text{ref}_N), v, x) &= \left(\left(v' \mapsto \begin{cases} \text{ref}_N & \text{if } v' = v \\ L(v') & \text{else} \end{cases} \right), \left(r \mapsto \begin{cases} x & \text{if } r = \text{ref}_N \\ G(r) & \text{else} \end{cases} \right), \text{ref}_{N+1} \right) \\
\text{Ref} &= \{\text{ref}_n \mid n \in \mathbb{N}\} \\
\text{new} &= (x \mapsto \text{nil}, x \mapsto \text{nil}, \text{ref}_0) \\
\text{Nil} &= \{\text{nil}\}
\end{aligned}$$

3.5 Expected Problems

We expect some problems with describing how values are passed on. For example, if we want to call to print an Integer, Lua provides us with the `toString` value for that Integer. We need to figure out a good method for describing how Lua does this.

4 Analysis

We will prove that what is printed is actually the thing that needs to be printed.

4.1 Example Code

— from <http://lua-users.org/wiki/CoroutinesTutorial>

```

function odd(x)
  print('A: odd', x)
  coroutine.yield(x)
  print('B: odd', x)
end

function even(x)
  print('C: even', x)
  if x==2 then return x end
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end

co = coroutine.create(
  function (x)
    for i=1,x do
      if i==3 then coroutine.yield(-1) end
      if i % 2 == 0 then even(i) else odd(i) end
    end
  end)

count = 1

```

```
while coroutine.status(co) ~= 'dead' do
  print('---', count) ; count = count+1
  errorfree, value = coroutine.resume(co, 5)
  print('E: errorfree, value, status', errorfree, value, coroutine.status(co))
end
```

4.2 Elaboration

To be added.

Appendix

A Planning

Deadline	Summary
?	Apply feedback
13 June	Finish Analysis
17 June	Submit final version project