

MODULE III

CHAPTER 3

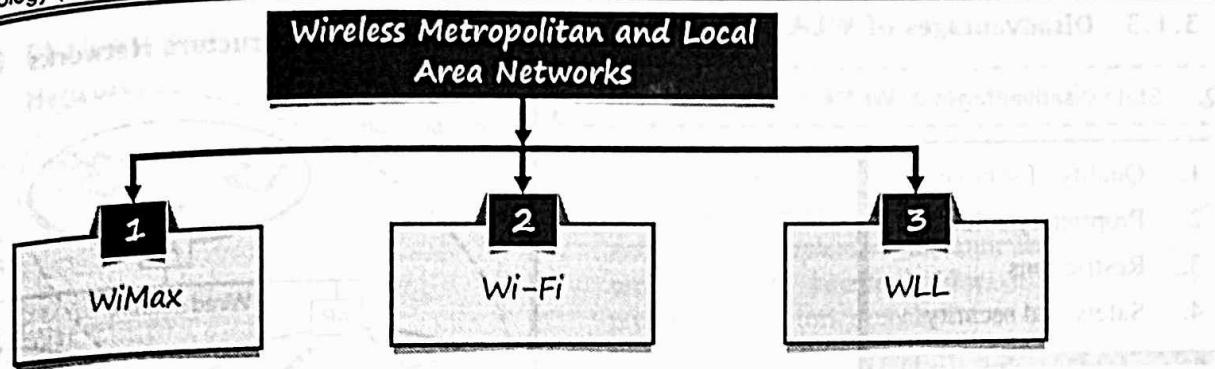
Wireless Metropolitan and Local Area Networks

University Prescribed Syllabus w.e.f Academic Year 2021-2022

IEEE 802.16 (Wi-Max) – Mesh mode, Physical and MAC layer; IEEE 802.11(Wi-Fi) – Architecture, Protocol Stack, Enhancements and Applications.

Self-learning Topics : WLL(Wireless Local Loop).

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3.1 INTRODUCTION

(For reading purpose only...)

- Wireless communication technologies support mobility, portability to the users. Hence it has gained immense popularity in day-to-day life nowadays.
- It has many advantages over wired communication. For example, wired networks rely on cables to connect digital devices together, wireless networks don't rely on wires or cables. This also reduces overall system cost. This also reduces errors due to cable damage.
- Wireless technologies are widely used in both home and business computer networks, for a variety of uses.

3.1.1 Types of Wireless Network Technologies

(For reading purpose only...)

A large number of technologies have been developed to support wireless networking in different scenarios.

- | | |
|---------------|-----------------|
| (i) Bluetooth | (ii) WLAN |
| (iii) WMAN | (iv) MANET |
| (v) Wi-Fi | (vi) Wimax etc. |

- Wireless local area network technologies comprise of fast growing flexible wireless access.
- This is generally useful in smaller organizations like offices, home, or production environments etc.
- WLAN is restricted to the distance parameters. They can be employed in limited ranges like in buildings, campus, single rooms etc.
- They are operated by individuals and not by the large-scale service providers.
- The main aim of WLAN technology is to remove the office cabling and provide the tether less access to internet. It also provides the flexibility to the ad hoc networks. For example, group meetings etc.

3.1.2 Advantages of WLAN

GQ. State advantages of WLAN.

1. Mobility
2. Low implementation cost
3. Installation speed and simplicity
4. Network expansion
5. Reliability
6. Scalability
7. Usage of ISM band

1. Mobility

Internet connectivity can be availed from any location.

2. Low implementation cost

WLANs are easy to set up, relocate, change and manage.

3. Installation speed and simplicity

The speed of installation is fast as no cabling is required. It also simplifies the architecture of the network.

4. Network expansion

Wireless technology reaches where cables cannot be reached.

5. Reliability

WLAN is resistant to different types of cable failures.

6. Scalability

WLAN can be configured in many topologies to meet the needs of the specific applications and installations.

7. Usage of ISM band

WLAN operates in the unlicensed ISM band available for use.

3.1.3 Disadvantages of WLAN

Q. State disadvantages of WLAN.

1. Quality of service
2. Proprietary solutions
3. Restrictions
4. Safety and security

1. Quality of service

- WLAN offer lower quality of service as compared to wired LAN connection.
- This is due to lower bandwidth and limitations in radio transmission.
- WLAN also faces problems like higher data error rates due to interference, higher propagation delays.

2. Proprietary solutions

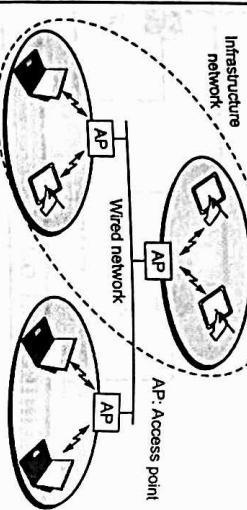
- Due to lack of standardization many companies are offering proprietary solutions.
- But these additional features work only when adapters from the same vendors are used.

3. Restrictions

- WLAN's are limited to low power senders and certain license free frequency bands.
- These frequency bands may not be same worldwide.
- There is lack of government regulations.

4. Safety and security

- Senders and receivers have to be controlled for low radiations.



(sw)Fig. 3.2.1 : Infrastructure networks

3.2.1 Infrastructure Networks

3.2.1(B) Disadvantage of Infrastructure Networks

1. These types of networks may lose flexibility of wireless networks. For example they cannot be used for disaster relief in case of no infrastructure is left.
2. In this the complexity of each node is higher as every node has to implement medium access mechanisms, mechanisms to handle hidden or exposed terminal problems and priority mechanisms.

3.2.2 Ad hoc Networks

3.2.2(A) Advantage of Adhoc Networks

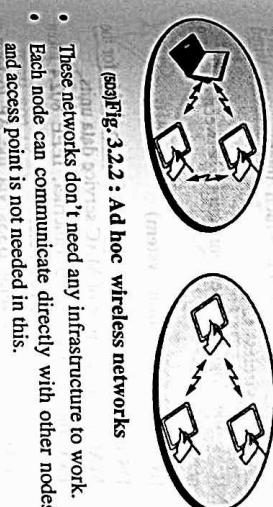
3.2.2(B) Disadvantages of Adhoc Networks

1. Hardware complexity is main issue in ad hoc networks.
2. Two nodes can communicate only when they are in radio range of each other or other nodes are able to forward the data.

3.2.3 Comparison between infrastructure and ad hoc architectures of WLAN

Q. Explain the difference between Ad-hoc Network and infrastructure based wireless networks.

Table 3.2.1 : Comparison between infrastructure and ad hoc architectures of WLAN



(sw)Fig. 3.2.2 : Ad hoc wireless networks

3.2.1(A) Advantages of Infrastructure Networks

3.2.1(A) Advantages of Infrastructure Networks

3.2.2 Ad hoc Networks

3.2.2(A) Advantage of Adhoc Networks

3.2.2(B) Disadvantages of Adhoc Networks

3.2.3 Comparison between infrastructure and ad hoc architectures of WLAN

Q. Explain the difference between Ad-hoc Network and infrastructure based wireless networks.

Table 3.2.1 : Comparison between infrastructure and ad hoc architectures of WLAN

Sl.No.	Parameter	Infrastructure based Network	Ad hoc networks
1.	Working principle	Devices on this type of network all communicate through a single access point, which is generally the wireless router.	Ad-hoc networks don't require a centralized access point. Instead, devices on the wireless network connect directly to each other.
2.	Presence of access point	Required	No
3.	Limitation on range	Wireless routers that acts as access points have higher-power wireless devices so they can cover a larger range	Wireless devices used are generally of low power. Hence range is limited.
4.	Communication with the receiver when it is not in range	If a receiver is out of range of sender, then forwarding of packets is done via access point.	If a receiver is out of range of sender, it will pass the data through other devices on the way.
5.	Hardware	Simple design of individual node	More complex design of individual node as compared to infrastructure based network.
6.	Flexibility	Less flexible	Higher flexibility
7.	Time required for setting up of network	More	Less
8.	Architectural diagram	Refer Fig. 3.2.1	Refer Fig. 3.2.2

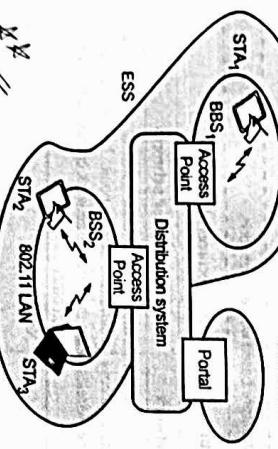
W 3.3 IEEE 802.11 (Wi-Fi)

In the family of IEEE 802.11 many WLAN standards are included. The main goal of this standard is to set up simple and robust WLAN which offers time-bounded and asynchronous services.

Syllabus Topic : IEEE 802.11(Wi-Fi) - Architecture**3.3.1 System Architecture of IEEE 802.11**

- Q.** Explain different architectures of WLAN.
Wireless networks can exhibit two different basic system architectures.

- A. Infrastructure based WLAN
B. Ad hoc WLAN



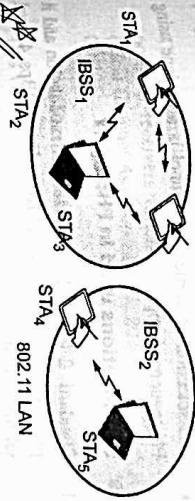
- Refer Fig. 3.3.1. It shows the system architecture of WLAN.

- Q.** It is nothing but a mobile station.

- The STAs in ESS can be mobile or stationary. The STAs in ESS are actually AP. They are part of wired LAN.
- When the BSSs are connected, two stations can communicate directly without AP.
- But if stations are belonging to different BSSs then they have to communicate via AP.

3.3.1(B) Ad hoc WLAN Architecture

802.11 LAN



- Refer Fig. 3.3.2. It shows ad-hoc WLAN architecture.

- In addition to infrastructure-based architectures, IEEE 802.11 also supports ad hoc network between stations.

- It forms IBSS (Independent Basic Service Set).

IBSS (Independent Basic Service Set)

- When all the stations in BSS are mobile and it has no connection with other BSSs, the BSS is known as IBSS.

- It is an ad hoc network. This does not contain AP and it in the network.

- Simplest BSS has only two STAs.

- The association of STA and BSS is entirely dynamic.

- BSS may be isolated, or it may be connected to the backbone Distribution System (DS) through an Access Point (AP).

- This mode of operation is known as ad hoc networking as this type of IEEE 802.11 WLAN is created and maintained without prior administrative arrangement.

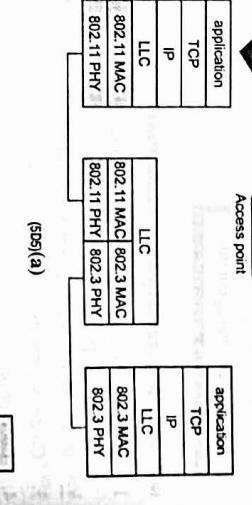
Syllabus Topic : IEEE 802.11(Wi-Fi) -Protocol Stack, Enhancements and Applications

- (i) STA (Station)
(ii) Access Point(AP)
(iii) Portal (PO)
(iv) DS(Distribution System)
(v) BSS (Basic Service Set)
(vi) ESS (Extended Service Set)

- It is a set of two or more BSSs forming a single sub-network.

- ESS configurations consist of multiple BSS cells that can be linked by either wired or wireless backbones known as DS.

- IEEE 802.11 supports ESS configurations in which multiple cells use the same channel.
- If may also use different channel to aggregate throughput.



- When the BSSs are connected, two stations can communicate directly without AP.

- But if stations are belonging to different BSSs then they have to communicate via AP.

- It is typical access point which interconnects wired LAN and wireless LAN.

- It is the logical interconnection between the two networks.

- DS is the backbone network that is responsible for the MAC layer transport of MAC service data units.

- Examples include 802.3 Ethernet, IEEE 802.4 token bus LAN, fiber optic LAN etc.

- A DS connects several BSSs via the AP to form a single network. Thus, it extends the wireless coverage area.

- It handles data transfer between the different APs.

- (v) BSS (Basic Service Set)

- It is the smallest building block of the WLAN system. It consists of some number of stations and access points in the network.

- Simplest BSS has only two STAs.

- The association of STA and BSS is entirely dynamic.

- BSS may be isolated, or it may be connected to the backbone Distribution System (DS) through an Access Point (AP).

- This mode of operation is known as ad hoc networking as this type of IEEE 802.11 WLAN is created and maintained without prior administrative arrangement.

3.4 PROTOCOL ARCHITECTURE OF IEEE 802.11

- (i) ESS (Extended Service Set)

- It is a set of two or more BSSs forming a single sub-network.

- ESS configurations consist of multiple BSS cells that can be linked by either wired or wireless backbones known as DS.

- IEEE 802.11 supports ESS configurations in which multiple cells use the same channel.
- If may also use different channel to aggregate throughput.

3 Management layers

- Management layers are as follows -
- (a) MAC management
- (b) PHY management
- (c) Station management

(a) Tasks performed by MAC Management

1. Supports the association and re-association of a station to an access point
2. Roaming between different access points
3. Controls authentication mechanisms
4. Encryption
5. Synchronization of a station with regard to an access point.
6. Power management to save battery power
7. Maintains the MAC MIB (Management Information Base).

(b) Tasks performed by PHY management

1. Physical MIB (Management Information Base)
2. Channel tuning.
3. Maximum transmit power is 1 W/MHz in US, 100 mW EIRP in Europe and 10 mW/MHz in Japan.
4. Operation at 1 Mbit/s is mandatory while at 2 Mbit/s is optional.
5. 79 Hopping channels for North America and Europe and 23 hopping channels for Japan.
6. Maximum transmit power is 1 W/MHz in US, 100 mW EIRP in Europe and 10 mW/MHz in Japan.
7. FHSS is easier to implement

(c) Tasks performed by station management

1. Interaction with both management layers
2. Responsible for additional higher layer functions like control of bridging, interaction with distribution system in case of access point etc.

M 3.5 PHYSICAL LAYER

Q. Explain the physical layer of IEEE 802.11

- IEEE 802.11 supports three different physical layers.
- One layer based on infra red
- Two layers based on radio transmission
- All PHY variants has CCA provision. It is important in controlling MAC mechanisms and indicates if the medium is current idle.
- PHY layer offers an SAP with 1/2 Mbps transfer rate to the MAC layer.
- Four versions of PHY layer are

 1. FHSS
 2. DSSS
 3. Infra-red
 4. Narrowband microwave LANs

3.5.1 FHSS Frequency Hopping Spread Spectrum

- FHSS is one of the variants of spread spectrum technology. It allows coexistence of multiple networks in the same area.
- It is done by separating different networks by different hopping sequences.

Hopping sequences are generated with the help of PN (pseudo noise) sequence.

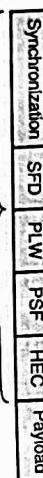
The selection of particular channel is also done using PN sequence hopping pattern.

Specifications used in FHSS PHY

1. Standard 2 level GFSK is used for modulation and it provides transfer rate of 1Mbps. For 2Mbps 4 level GFSK is used.
2. It also uses the 2.4 GHz ISM band.
3. Provides bandwidth of 1 MHz.
4. Operation at 1 Mbit/s is mandatory while at 2 Mbit/s is optional.
5. 79 Hopping channels for North America and Europe and 23 hopping channels for Japan.
6. Maximum transmit power is 1 W/MHz in US, 100 mW EIRP in Europe and 10 mW/MHz in Japan.
7. FHSS is easier to implement

Frame structure used in FHSS PHY layer

- Refer Fig. 3.5.1 it shows the format of IEEE 802.11 PHY frame using FHSS.
- 80 16 12 4 16 Variable bits



(a) Fig. 3.5.1 : Format of IEEE 802.11 PHY frame using FHSS

3.5.2 Direct Sequence Spread Spectrum

- It is another spread spectrum technology used at Physical layer of IEEE 802.11.
- In this method spreading is achieved by code and not by frequency.
- In IEEE 802.11, 11 chip Barker code is used. Its value is $(+1, -1, +1, +1, -1, +1, +1, -1, -1, -1)$.
- The main advantage of this technology is its robustness against interference and its insensitivity to multipath propagation.
- But the implementation is bit complex as compared to FHSS.

This field is reserved for future use.

3.5.3 Fields in the frame format

1. Synchronization
2. Start Frame Delimiter (SFD)
3. PLCP PDU Length Word (PLW)
4. PLCP Signaling Field (PSF)
5. Header Error Check (HEC)

1. Synchronization
This pattern is used for the synchronization of the potential receivers and signal detection by the CCA (Clear Channel Assessment). It is 80 bit field which is a (010101... Bit pattern).

2. Start Frame Delimiter (SFD)
This is a 16 bit field indicates the start of frame and provide's frame synchronization. The pattern of SFD is 000011001011101.

3. Packet Length Word (PLW)
The 12 bit packet length width shows the length of the payload. The PLW can range between 0 to 4095.

4. PLCP Signalling Field (PSF)
It is 4 bit field that specifies the data rate of the payload following. If all bits are set to zero (0000) it means the lowest data rate (1 Mbit/s). 2 Mbit/s data rate is represented by 0010 bit sequence. Maximum data rate 85 Mbit/s is represented by 1111.

5. Header Error Check (HEC)
16 bit HEC is added to protect the PLCP header. It can recover errors of up to 2 bits, otherwise identify whether PLCP bits are corrupted.

Fields in the frame format

1. Synchronization
2. Start Frame Delimiter (SFD)
3. Signal
4. Service
5. Length
6. Header Error Check (HEC)

1. Synchronization
It is a 128 bit field is used for synchronization, gain setting, energy detection, and for frequency offset compensation. This field consists of scrambled 1 bits.

2. Start frame delimiter (SFD)
This field indicates the starting of a frame and consist the pattern 1111001110100000. This field is used for synchronization at the beginning of the frame.

3. Signal
This field indicates the data rate of the payload. The value 0x0A is for 1 Mbit/s and 0x14 is for 2 Mbit/s, other values are reserved for future use.

4. Service
This field is reserved for future use.

5. Length
This field is reserved for future use.

6. Header Error Check (HEC)
This field is reserved for future use.

Module 3

3.5.4 Specifications

- It also uses 2.4 GHz ISM band
- It offers both 1 Mbit/s and 2 Mbit/s data rate.
- It uses Differential Binary Phase Shift Keying (DBPSK) modulation for 1 Mbit/s transmission and Differential Quadrature Phase Shift Keying (DQPSK) for 2 Mbit/s.
- The symbol rate is 1 MHz and chipping rate is 11 MHz Implementation is complex.
- Provides a better coverage and a more stable signal
- All bits transmitted by the DSSS PHY are scrambled with the polynomial $s(z)=z^7+z^4+1$ for DC blocking and whitening of the spectrum.
- Frame structure for IEEE 802.11 PHY using DSSS
- Refer Fig. 3.5.2.
- The frame consists of two parts
- PLCP part (preamble and header)
- Payload part
- PLCP part is always transmitted at 1Mbps and payload which is MAC data may be transmitted at 2Mbps.
- It is 4 bit field that specifies the data rate of the payload following. If all bits are set to zero (0000) it means the lowest data rate (1 Mbit/s). 2 Mbit/s data rate is represented by 0010 bit sequence. Maximum data rate 85 Mbit/s is represented by 1111.
- It is a 128 bit field is used for synchronization, gain setting, energy detection, and for frequency offset compensation. This field consists of scrambled 1 bits.
- In this method spreading is achieved by code and not by frequency.
- In IEEE 802.11, 11 chip Barker code is used. Its value is $(+1, -1, +1, +1, -1, +1, +1, -1, -1, -1)$.
- The main advantage of this technology is its robustness against interference and its insensitivity to multipath propagation.
- But the implementation is bit complex as compared to FHSS.

- 5. Length**
This 16 bit field is used to indicate the length of a payload in microseconds.
- 6. Header Error Check (HEC)**
HEC is used to protect PLCP header.

3.5.3 Infra-Red

- The PHY layer based on the infrared transmission, makes use of near visible light spectrum at 850-950nm.
- It does not need line of sight between the sender and the receiver. It also works with diffuse light.
- It supports point to multipoint communication.
- Maximum range is about 10m if no sunlight or heat sources interfere with the transmission.
- This is suitable for networks in buildings like classrooms, meeting rooms etc.
- Frequency reuse scheme implementation is very simple.
- A wall is more than enough to isolate one IR based IEEE 802.11 network from another. Infra red cannot penetrate through the walls or other opaque objects.
- Now a days this standard has been obsolete. No devices are available that offer infrared communication.

3.5 Advantages of Infra-red

- Infrared can be more easily secured against eavesdropping as compared to microwave.
- Separate infrared installation can be installed in every room in a building without interference enabling the construction of very large infrared LANs.
- Equipment is relatively simple and inexpensive.

Note

Microwave radio frequencies can be used for voice, data and video transmissions.

They should be coordinated within certain geographic areas to avoid potential interference between the systems.

- In narrowband microwaves, cell configuration approach is employed. In these adjacent cells use non-overlapping frequencies within overall 18 GHz band.
- All transmissions are encrypted so that eavesdropping is avoided. This supports interference free communication.

Narrowband microwave LAN using unlicensed spectrum

- In 1995, Radio LAN has introduced narrowband wireless LAN using unlicensed ISM spectrum.
- It is used for transmission at low power of 0.5 watts or less.
- These Radio LAN products operates on 10 Mbps in the 5.8 MHz band.
- It has the range of 50 m in a semi-open office and 100 m in open office. It makes use of peer-to-peer configuration.

- The Radio LAN product automatically elects one node as the Dynamic Master based on parameters such as location, interference and signal strength.
- The identity of the master can change automatically as conditions change.

- The LAN also includes the dynamic relay function which allows each station to act as a repeater to move data between stations that are out of range of each other.

M.3.6 MEDIUM ACCESS CONTROL LAYER

Q. Explain in Detail IEEE 802.11 MAC sublayer.

- The tasks performed by the MAC layers are

- Control medium access
- Supports roaming
- Authentication
- Power saving mechanisms

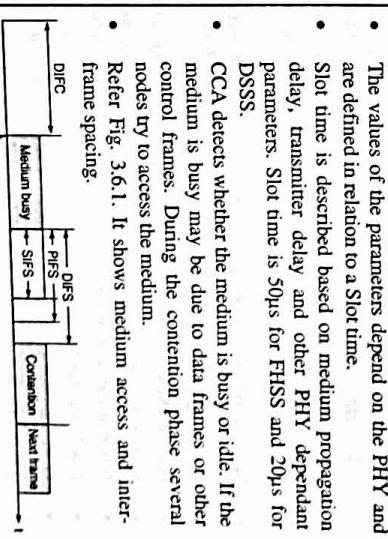
Parameters for controlling the waiting time before medium access

- There are three different parameters that define the priorities of medium access.
- The values of the parameters depend on the PHY and are defined in relation to a Slot time.

- Slot time is described based on medium propagation delay, transmitter delay and other PHY dependent parameters. Slot time is 50μs for FHSS and 20μs for DSSS.
- CCA detects whether the medium is busy or idle. If the medium is busy may be due to data frames or other control frames. During the contention phase several nodes try to access the medium.

- Refer Fig. 3.6.1. It shows medium access and inter-frame spacing.

Diagram illustrating the sequence of events in IEEE 802.11 medium access:



- It makes use of narrowband microwave frequency for radio transmission. It is just wide enough to accommodate signal.
- It can make use of Licensed and Unlicensed (ISM) microwave bands for its operation.

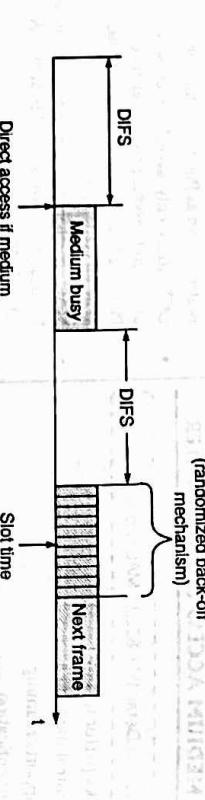
Parameter	Infrared	FHSS	DSSS	Narrowband microwave
No	Data rate (Mbps)	Up to 10	1 to 3	2 to 54
1	Mobility	Stationary/ mobile	Mobile	Stationary/ mobile
2	Range (m)	Up to 25	30-100	30-250
3	Detectability	Negligible	Little	Little
4	Wavelength / Frequency	1:800 to 900 nm	902-928MHz 2.4-2.4835GHz 5.725-5.8GHz	902-928 MHz 2.4-2.4835 GHz 4.9-5.775 GHz 18.825-19.205 GHz
5	Modulation used	ASK	GFSK	QPSK
6	Radiated power	--	61 W	61 W
7	Access method	CSMA	CSMA	Reservation ALOHA, CSMA

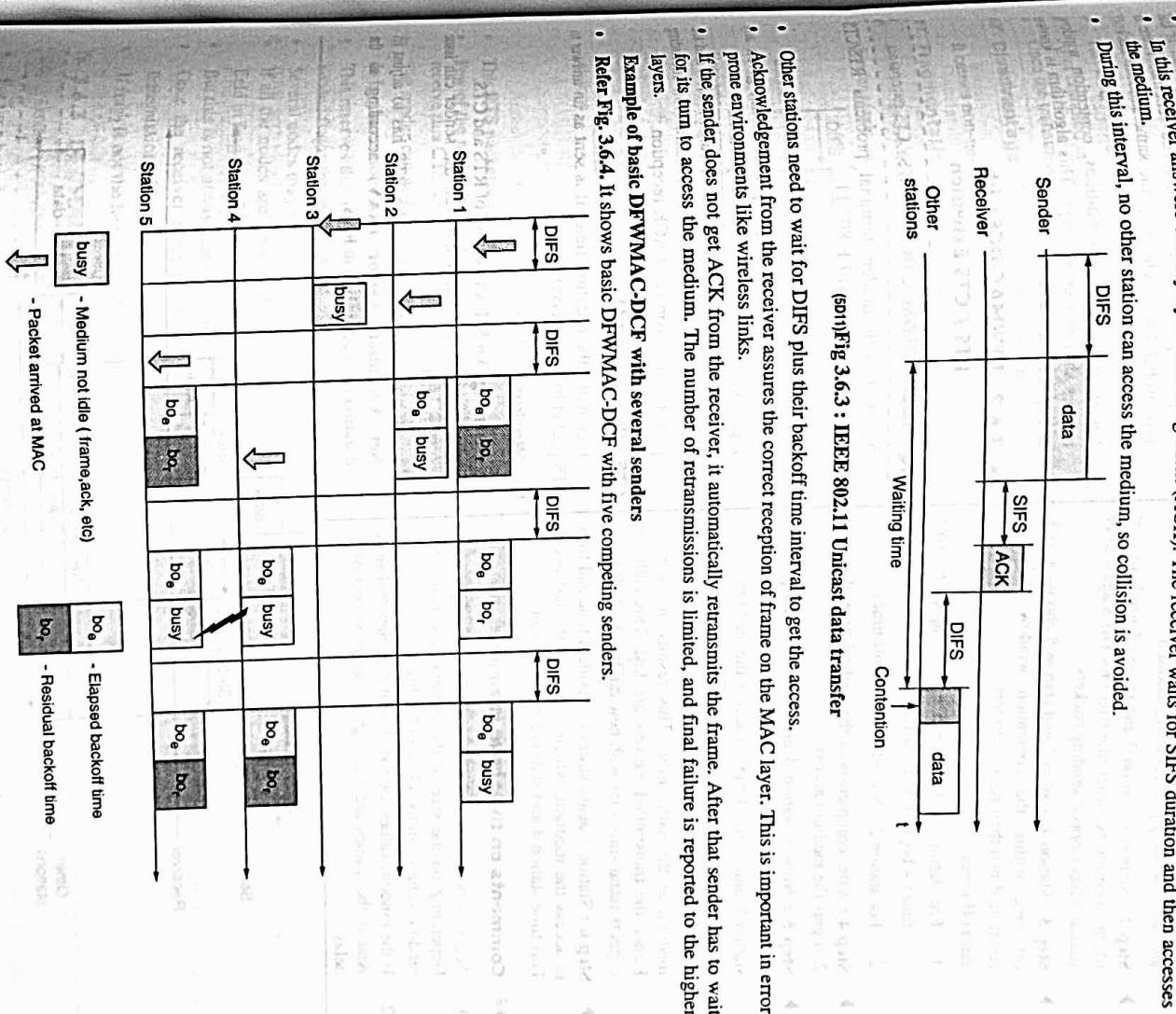
Contention free polling method for time bound service is also known as Point Coordination Function (PCF).

DCF offers only asynchronous service whereas PCF offers both asynchronous and time bounded services.

PCF needs access point to control medium access and to avoid contention.

MAC mechanisms are also known as Distributed Foundation Wireless Medium Access Control (DFWMAC).

- | | |
|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| <ul style="list-style-type: none"> • It has the highest priority. • The shortest waiting time for medium access is defined for short control messages like acknowledgements of data packets or polling responses. • The SIFS for DSSS is 10µs and for FHSS it is 28µs. <p>2. PCF Inter Frame Spacing (PIFS)</p> <ul style="list-style-type: none"> • It has medium priority. • A waiting time between DIFS and SIFS is used for a time bounded service. PIFS is defined as SIFS plus one slot time. | <p>3. DCF Inter Frame Spacing (DIFS)</p> <ul style="list-style-type: none"> • It has lowest priority. • It has the longest waiting time. This is used for asynchronous data service within a contention period. It is defined as SIFS plus two slot times. |
| <p>3.6.1 Basic DFWMAC-DCF using CSMA/CA</p> <p>Q. Explain Basic DFWMAC-DCF using CSMA/CA.</p> <ul style="list-style-type: none"> • The mandatory access mechanism of IEEE 802.11 is based on Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA). • It is random access scheme with carrier sense and collision avoidance through random backoff. • Refer Fig. 3.6.2. It shows basic CSMA/CA mechanism.  <p>Fig. 3.6.2: Basic CSMA/CA mechanism</p> <p>Working of Basic DFWMAC-DCF using CSMA/CA</p> <ul style="list-style-type: none"> • A node can access the medium at once if the medium is idle for the duration of DIFS. It supports short access delay under light load. Additional mechanisms are needed if large number of nodes are needed to access the medium. • If the medium is busy, nodes must wait for DIFS duration before entering the contention phase. • Each node chooses a random backoff time within a contention window. It delays medium access for this random amount of time. • Node keeps on sensing the medium. Once node senses that channel is busy, it has lost this cycle and has to wait for the next chance. If the randomized additional waiting time is over for a node and if the medium is still idle, the node can access medium immediately. • Basic CSMA/CA is not fair. Hence IEEE 802.11 adds a backoff timer to provide fairness. • Each node selects the random waiting time within the range of the contention window. • If one of the stations does not get access to medium in first cycle, it will stop its backoff timer. • It will wait for the channel to be idle again for DIFS and starts the counter again. Once the counter expires, the node accesses the medium. <p>Therefore, deferred stations don't choose a randomized backoff time again, rather they continue to count down. Stations will have to wait only for their backoff timer from the previous cycle.</p> <p>This is the advantage in case of the stations which are waiting for longer time as compared to the stations that are entered recently.</p> | |



Shows basic DFWMAC-DCF with several senders

Time	S1	S2	S3	S4	S5
0	busy				
1		busy			
2			busy		
3				busy	
4					busy
5					

If the sender does not get ACK from the receiver, it automatically retransmits the frame. After that sender has to wait for its turn to access the medium. The number of retransmissions is limited, and final failure is reported to the higher layer.

Other stations need to wait for DIFS plus their backoff time interval to get the access. Acknowledgement from the receiver assures the correct reception of frame on the MAC layer. This is important in error prone environments like wireless links.

Fig 3.6.3 : IEEE 802.11 Unicast data transfer

The diagram illustrates a probability density function (PDF) for the waiting time in a contention-based system. The x-axis is labeled "Waiting time" and the y-axis is labeled "Contention". The curve shows a peak at zero, representing the highest probability of a short waiting time. As the waiting time increases, the probability decreases rapidly, indicating that long waits are much less likely.

The diagram illustrates the IEEE 802.11 wireless LAN frame structure. It shows a sequence of frames starting with a short inter-frame space (SIFS), followed by a distributed inter-frame space (DIFS) between the SIFS and the data frame. The data frame itself contains a header and payload. The entire sequence is labeled "Other stations".

```

sequenceDiagram
    participant R as Receiver
    participant S as Sender
    R->>S: ACK

```

Sender | data | SIFS |

During this interval, no calls change the discount, so confusion is avoided.

In this receiver, after a short time interval (SIFS), the receiver waits for SIFS duration and then accesses the medium.

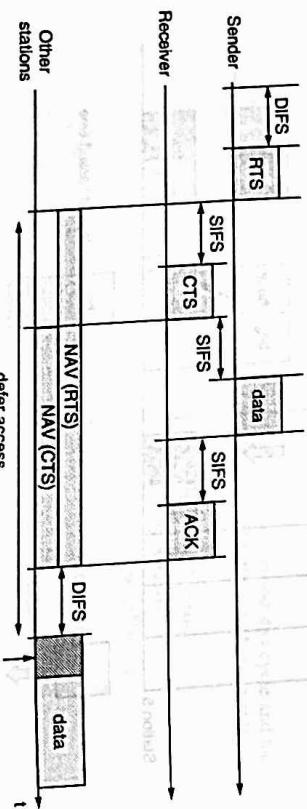
- Let station 3 has the first request from the higher layer to send a packet.
- Step 1 :** Station 3 senses the medium, waits for DIFS and starts sending packets using medium.
- Step 2 :** Station 1, station 2 and station 5 needs to wait till the medium becomes idle (free) for DIFS again after station 3 has stopped sending packets.
- Step 3 :** Station 1, station 2 and station 5 choose a back off time within the contention window and start counting down their backoff timers.

Backoff times

- For station 1 and station 5 = $b_{\text{oe}} \cdot (\text{elapsed backoff time}) + b_{\text{res}} \cdot (\text{residual backoff time})$
 - For station 2 = $b_{\text{oe}} \cdot (\text{elapsed backoff time})$
- Step 4 :** After completion of the backoff time of station 2, it gets the medium access.
- Step 5 :** Now if station 4 also wants to send packets, after DIFS wait time, three stations namely station 1, station 5, and station 4 tries to access the medium.
- As shown in Fig. 3.6.4 the backoff time of station 4 and station 5 is same hence both of them can access the medium at the same time. This results in collision. Hence the transmitted packets are lost. This collision triggers retransmissions with new random backoff time.
- Step 6 :** Station 1 again stores its residual time and tries to access the medium, which is free in the last cycle. This time station 4 and station 5 needs to wait.

Comments on the above example

- Access scheme has problems under heavy or light load. Depending on the size of the contention window, the random values can be too close or too high.
- If the random values are too close it provokes collision. And if the values are too high it causes unnecessary delay.



(contd)Fig. 3.6.5 : IEEE 802.11 hidden node provisions for contention-free access

- NAV specifies the earliest point at which the station can try access the medium.
- On reception of RTS, Receiver answers with CTS after waiting for SIFS interval.
- All the stations receiving distance of sender and receiver are informed that they have to wait for more time before accessing the medium.
- This reserves the medium for single sender. Hence it is also known as virtual reservation scheme.
- After completion of SIFS duration, sender can send the data. Receivers send the ACK after SIFS waiting time.
- Once the transmission is completed, the NAV in each node marks the medium as free and standard cycle starts again.

3.6.2 DFWMAC-DCF with RTS / CTS Extension

- The contention window doubles up to maximum each time the collision occurs.
- The larger the contention window, the greater is the resolution power of the randomized scheme. There is less probability of choosing the same random backoff time value if the contention window is large.
- Under the light load conditions, contention window assures shorter access delays. This algorithm is known as exponential backoff.

Q. Explain DFWMAC-DCF with RTS / CTS extension.

To deal with hidden terminal problems RTS/CTS protocols are defined in IEEE 802.11.

RTS (request to send)

- It is a control packet. This packet is not given any higher priority compared to other packets.
- It includes the receiver of the data transmission and the duration required for the data transmission including packet transmission and ACK reception.

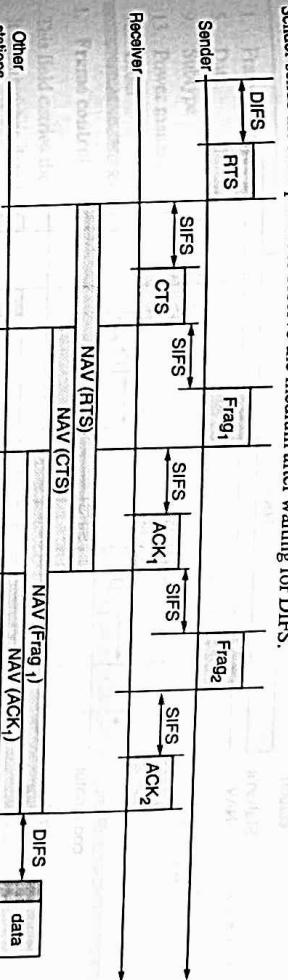
CTS (clear to send)

- It contains the duration field. It is sent as an answer to RTS packet by the receiver.

Working

Refer Fig. 3.6.6. It shows mechanism for fragmentation of user data.

- Sender sends the RTS packet to reserve the medium after waiting for DIFS.
- After waiting for SIFS, CTS is sent by receiver.
- Other stations receive the RTS and CTS.
- NAV (RTS) and NAV (CTS) are updated.
- After waiting for SIFS, Frag1 is transmitted.
- ACK1 is received by sender.
- After waiting for SIFS, Frag2 is transmitted.
- ACK2 is received by sender.
- After waiting for SIFS, data is transmitted.



(contd)Fig. 3.6.6 : IEEE 802.11 fragmentation of user data

- This RTS packet includes the duration for the transmission of the first fragment and the corresponding ACK.
- Few of the nodes which receive this RTS can adjust their NAV accordingly. Receiver answers with CTS. Few receivers receive this CTS and they adjust their NAV accordingly.
- The sender can now send first data frame frag1 after waiting for SIFS interval.
- In fragmentation mode, another duration value is included in frag1.
- This reserves the medium for the said duration of the transmission following. This following transmission comprises of second fragment and its ACK.
- Several nodes may receive this reservation and set their NAV accordingly.
- If all the nodes are static and transmission conditions have not changed, then the set of nodes receiving the duration field in frag1 should be same as the set of nodes that has received the initial RTS.
- But this is not always the case due to mobility in the network.
- Once the receiver receives frag1 it answers with ACK after SIFS interval including the reservation for the next transmission of frag2.
- If frag2 is not the last frame of transmission, process continues.

3.6.3 DFWMAC-PCF with Polling

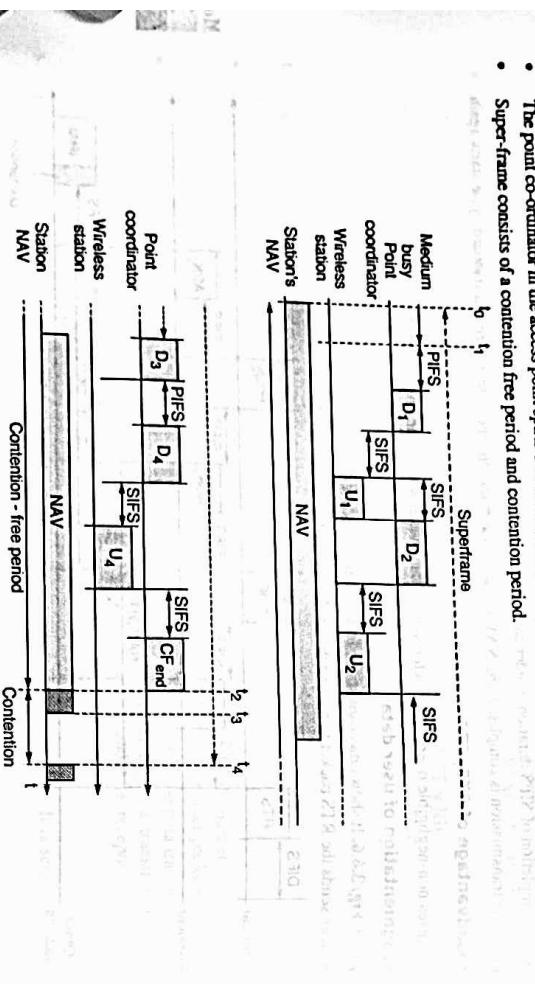
- Q. Explain DFWMAC-PCF with Polling.
- For time bounded service, the standard PCF (Point Coordination Function) is specified on the top of the standard DCF mechanisms.

Wireless Technology (MU-Sem 6-IT)

(Wireless Metropolitan and Local Area Networks) ... Page no (3-16)

- PCF makes use of an access point that controls medium access and polls the single nodes.
- But ad hoc networks cannot use this function and hence it does not provide QoS rather supports best effort services in IEEE 802.11 WLANs.

- Refer Fig. 3.6.7. It shows contention free access using polling mechanisms. Many stations with their NAVs are shown in Fig. 3.6.7.
- The point co-ordinator in the access point splits the access time into super frame periods.
- Super-frame consists of a contention free period and contention period.



Wireless Technology (MU-Sem 6-IT)

Note
 DS : Distribution System.
 AP : Access Point.
 DA : Destination Address.
 SA : Source Address
 BSSID : Basic Service Set Identifier.
 RA : Receiver Address.
 TA : Transmitter Address.

Address 1 identifies the physical receiver. Every station access point or wireless node filters on address 1.

Address 2 represents the transmitter of a frame.
 Address 3 and Address 4 are mainly necessary for the logical assignment of frames.

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 Address 3 and Address 4 are mainly necessary for the logical assignment of frames.

- ▶ 13. Power management : Value 1 indicates the station goes in power save mode, 0 represents the station remains active.
- ▶ 14. More data : This field indicates a receiver that sender has more data to send than the current frame.
- ▶ 15. Wired Equivalent Privacy (WEP) : Indicates that the standard security mechanism of IEEE 802.11 is used.

- ▶ 16. Order : Value 1 indicates the received frames must be processed in strict order.

- ▶ Fig. 3.6.9 shows three different types of control packets: Acknowledgement packet, RTS, and CTS packet.

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3. **CTS (Clear to Send) packet**

- CTS frame copies the transmitter address from the RTS packet into its receiver address field.

- After reading the duration field, it subtracts the time to send the CTS and SIFS and writes the results into its own duration field.

- This beacon frame is used within BSS for conveying timing information.

- This time stamp is used by node to adjust its local clock.

- For synchronization, node is not required to hear to every beacon frame, but internal clocks should be adjusted periodically.

- If the medium is busy, the beacon frame is deferred. Therefore, the transmission of the beacon frame is not always periodic.

- MAC management controls all the functions related to system integration. This means the integration of wireless station into BSS, formation of an ESS, synchronization of stations etc.

- Following are the function groups

1. Synchronization

2. Power management

3. Roaming

4. Management information base (MIB)

5. Synchronization

- It includes the functions related to finding of wireless LAN, synchronization of internal clocks, generation of beacon signals etc.

- Each node of an 802.11 network maintains a internal clock. To synchronize the clocks of all nodes, IEEE 802.11 specifies a Timing Synchronization Function (TSF).

- These synchronized clocks are needed for :

1. Power management

2. Coordination of the PCF and

3. Synchronization in FHSS hopping sequence

Synchronization process for infrastructure-based networks

- In infrastructure-based networks, an access point controls the synchronization process. It is done by transmitting beacon frame periodically.

- Other wireless nodes adjust their local clocks according to beacon time stamp.

- Refer Fig. 3.7.1. It shows the beacon transmission in a busy 802.11 infrastructure network.

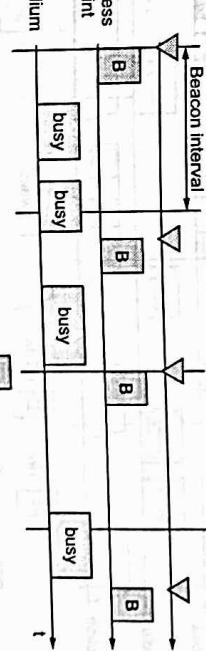
- AP cannot always assure to send its beacon B signal periodically if the medium is busy.

- However, AP always tries to schedule transmissions according to expected beacon interval. It is known as target beacon transmission time.

- The timestamp of the beacon indicates the actual transmission time and not the scheduled time.

- Synchronization process for ad hoc networks.

- In this each node maintains its own synchronization timer and starts transmission of a beacon frame after the beacon interval.



(See) Fig. 3.7.1 : Beacon transmission in a busy 802.11 infrastructure network

- Using PCF the local timer of the node can predict the start of a super frame. It means the start of the contention free and the contention period.

- Beacon frame**

- It contains the time stamp or other management information used for power management and roaming.

- This beacon frame is used within BSS for conveying timing information.

- This time stamp is used by node to adjust its local clock.

- For synchronization, node is not required to hear to every beacon frame, but internal clocks should be adjusted periodically.

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- For synchronization, node is not required to hear to every beacon frame, but internal clocks should be adjusted periodically.

- If the medium is busy, the beacon frame is deferred. Therefore, the transmission of the beacon frame is not always periodic.

- Refer Fig. 3.7.2. It shows an example where multiple stations try to send their beacon frames.

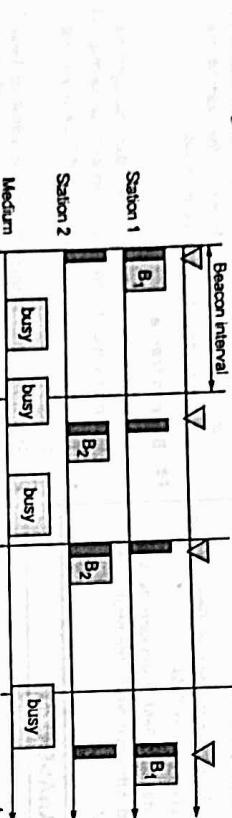


Fig. 3.7.2 : Beacon transmission in a busy 802.11 ad-hoc network

- But the standard backoff algorithm is used so that only one beacon frame is used at a time.

- All the other nodes adjust their internal clocks according to received beacon and suppress their beacons for this cycle. In case of collision, beacon is lost. In this case the beacon interval is shifted slightly.

2. Power management

- In IEEE 802.11 power management, the transceiver is switched off for the duration whenever it is not needed.
- This is simple in case of transmitter as the triggering device it itself whereas it is difficult in case of receivers.
- The receiver can't know in advance when to switch off.
- Switching off the transceiver should be transparent to existing protocols and should be flexible enough to support different applications.
- Longer off periods can save the battery life but on the other hand it reduces the average throughput and vice versa.
- There are two states involved in power saving
 - Sleep and Awake
 - Buffeting of data in senders
- If the station is asleep, the sender intending to communicate with the power saving station has to buffer the data.
- The sleeping station has to wake up periodically and needs to stay awake for certain time. During this interval all the senders can announce the destinations of their buffered packets.
- Once the station detects that it is the destination for the buffered packet, it needs to stay awake till transmission takes place. TSF (Timing Synchronization Function) is needed for waking up at the right moment. All the stations need to be awake at the same time.

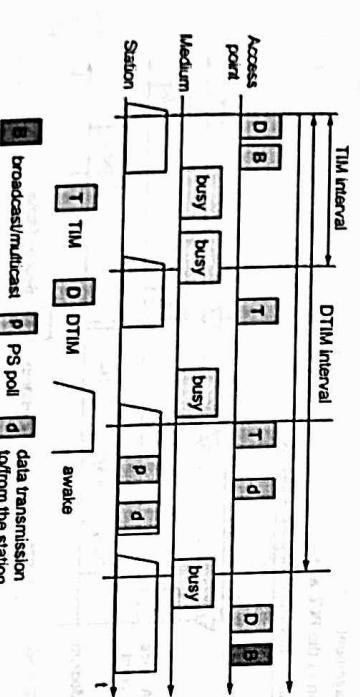


Fig. 3.7.3 : Power management in IEEE 802.11 Infrastructure networks

- management in infrastructure-based architectures. It is simpler. The AP buffers all the frames intended for the power saving stations.
- Along with beacon frame AP also transmits DTIM (time indication map). DTIM contains the list of stations for which unicast data frames are buffered in the access point.
- TSF ensures that sleeping stations will wake up periodically and will listen to the beacon frame and TM (Timing Synchronization Function).
- For multicast/broadcast transmission, nodes have to be awake always.
- A frame is also sent from the node to the access point. For this also nodes need to be awake.
- Refer Fig. 3.7.3. It shows power management in IEEE 802.11 infrastructure networks.

- Fig. 3.7.3 shows power management in IEEE 802.11 infrastructure networks. Fig. 3.7.3 shows power management in IEEE 802.11 infrastructure networks. It shows one access point and one station. The state of the medium is also shown in this.
- AP sends the beacon frame at each beacon interval. The beacon interval is same as that of TIM interval.
- The AP also maintains Delivery Traffic Indication Map (DTIM). This is used for sending multicast/broadcast frames. DTIM is multiple of TIM.
- As shown in the Fig. 3.7.3, all stations are awake to receive the broadcast frame.
- Once the station receives this frame, it again goes into sleeping mode. Once again, all stations wake up before the next TIM transmission.
- If the medium is busy, TIM is delayed. Hence the station stays awake. The AP has nothing to send and therefore station again goes to sleep mode.

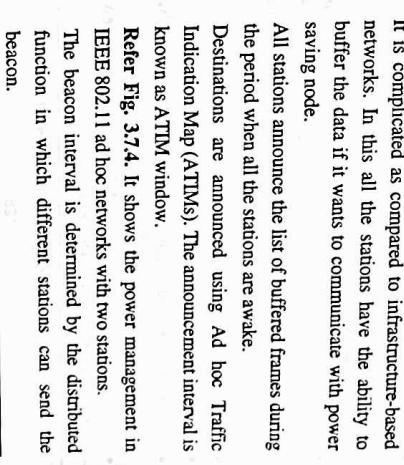


Fig. 3.7.4 : Power management in IEEE 802.11 ad hoc networks with two stations

- At the next TIM interval, AP indicates that the station is the destination for the buffered frame and hence station answers to this with PS (Power Saving) poll. It then stays awake for reception.
- AP then transmits the data for the station. Station acknowledges the receipt and may also send some data.
- Now the AP has large broadcast data to send at the next DTIM interval. This is again delayed as the medium is busy.
- Based on interval thresholds a station may stay awake if the sleeping period is too short. This mechanism clearly shows the trade off between short delays in station access and saving battery power.
- The shorter the TIM interval, shorter is the delay but lower is the power saving effect.
- Power management in ad hoc networks**
- It is complicated as compared to infrastructure-based networks. In this all the stations have the ability to buffer the data if it wants to communicate with power saving node.
- All stations announce the list of buffered frames during the period when all the stations are awake.
- Destinations are announced using Ad hoc Traffic Indication Map (ATIMs). The announcement interval is known as ATIM window.
- Refer Fig. 3.7.4. It shows the power management in IEEE 802.11 ad hoc networks with two stations.
- The beacon interval is determined by the distributed function in which different stations can send the beacon.

- Due to this all the stations wake up at the same time and they remain awake for ATM interval.

- If no frame is buffered to them, these stations goes to sleep. This happens in first two steps.

- In the third step, station 1 has data buffered for station 2. It is shown by the ATM transmitted by ATM from station 1.

- Station 2 acknowledges this ATM and stays awake for transmission.

- Once the ATM window is complete, station 1 can transmit and station 2 acknowledges its receipt. In this case stations stay awake for next beacon.

Disadvantages

- More ATM transmissions results in collision.

- The access delays of large networks are difficult to predict.

- QoS guarantees can't be given under heavy load conditions.

Roaming

- Wireless networks in the building require more than single AP to cover all the rooms. AP has the transmission range of 10-20m. each floor of building also need one AP.

- Moving between the AP is known as roaming.

- The steps of roaming procedure are as follows

- If the current link quality is detected to be poor by station, it starts scanning for another AP.

- It is actually search of another BSS or setting up of new BSS in case of ad hoc networks.

- The only difference in the signal field. In this the rate is encoded in multiples of 100kbps.

- Table 3.8.1 represents the relation between the field value and data rate.

Table 3.8.1 : Field value and data rates

Field value	Data rate
0xA	1Mbps
0x4	2 Mbps
0x37	5.5 Mbps
0x6E	11 Mbps

- Ds then updates its database of current location of stations. It also informs old AP that the station is no longer exist in its BSS.

IEEE 802.11b

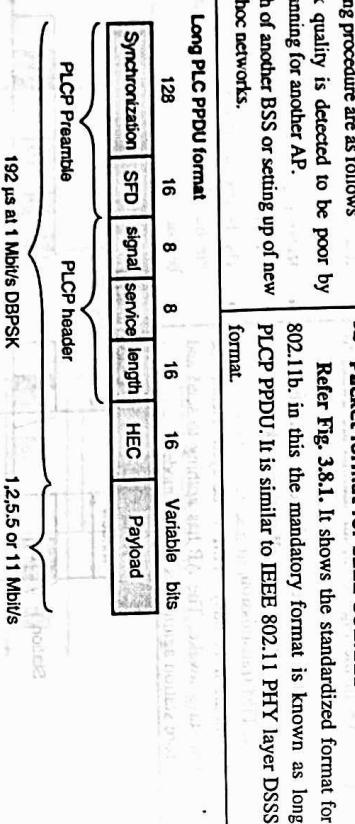
- This standard only defines new PHY layer. All the MAC schemes, MAC managements schemes etc. are same as that of IEEE 802.11.

Features

- 1. Data rate : 11, 5.5, 2 or 1 Mbps
- 2. Lower data rates use Barker codes.
- 3. Modulation used: DBPSK or DQPSK or CCK

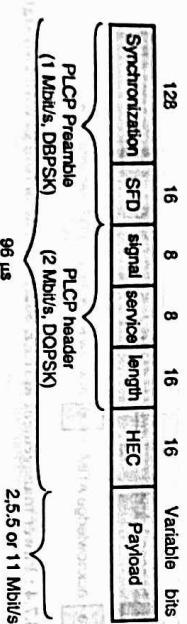
Packet format for IEEE 802.11b

- Refer Fig. 3.8.1. It shows the standardized format for 802.11b, in this the mandatory format is known as long PLCP PPDU. It is similar to IEEE 802.11 PHY layer DSSS format.



[Fig. 3.8.1 : IEEE 802.11b PHY packet formats]

Short PLCP PPDU format(Optional)



[Fig. 3.8.2 : IEEE 802.11b non-overlapping channel selection]

- Features

- 1. Data rate : 54 Mbps

- 2. Makes use of OFDM

- 3. Maximum transmit power : 200mW EIRP for lower frequency band (indoor use) 1W EIRP for higher frequency band (for indoor and outdoor use)
- 4. Number of subcarriers used : 52 (48 data+4 pilot)
- 5. Spacing between the subcarriers: 312.5kHz

- Modulation used : BPSK, QPSK, 16-QAM or 64-QAM

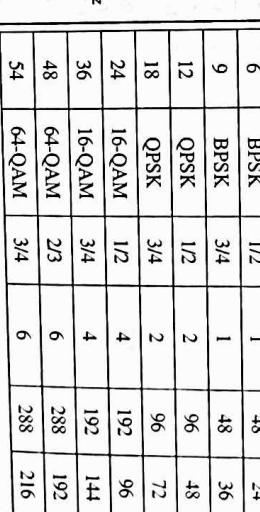
- Different coding schemes are used for different modulation schemes used. Refer Table 3.9.1.

Table 3.9.1 : Rate dependent parameters for IEEE 802.11a

Data rate (Mbit/s)	Modulation	Coding rate	Coded bits per subcarriers	Coded bits per symbol	Data bits per symbol	OFDM symbol
6	BPSK	1/2	1	48	24	
9	BPSK	3/4	1	48	36	
12	QPSK	1/2	2	96	48	
18	QPSK	3/4	2	96	72	
24	16-QAM	1/2	4	192	96	
36	16-QAM	3/4	4	192	144	
48	64-QAM	2/3	6	288	192	
54	64-QAM	3/4	6	288	216	

OFDM in IEEE 802.11a

[Fig. 3.9.1 : Usage of OFDM in IEEE 802.11a]



[Fig. 3.8.1 : IEEE 802.11b PHY packet formats]

- IEEE 802.11a use fixed symbol rate of 250000 symbols per second independent of data rate.
- It shows 52 subcarriers are equally spaced around the center frequency (26 on right and 26 on left). The spacing between the subcarriers is 312.5 kHz.
- Subcarriers with numbers -21, -7, 7 and 21 are used as pilot signals.
- IEEE 802.11 a physical layer PDU

Refer Fig. 3.9.2.

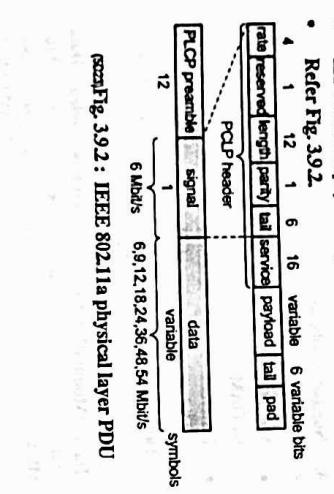


Fig. 3.9.2: IEEE 802.11a physical layer PDU

3.9.1 Comparison of Different IEEE 802.11x Standards

GQ. Compare various IEEE 802.11x standards. (a/b/g/n etc.)

Table 3.9.2 : Comparison of different IEEE 802.11x standards

IEEE 802.11 standard	Description	Frequency spectrum	Band width (MHz)	Data rate	Modulation used	Indoor range (m)	Outdoor range (m)	Year
IEEE 802.11	Standard for WLAN operation	2.4GHz ISM band	20	2Mbps	DSSS,FHSS	20	100	1997
IEEE802.11a	Standard for WLAN infrastructure (UNII) band	5GHz unlicensed national	20	Maximum 54 Mbps	OFDM	35	120	1999
IEEE802.11b	Standard for WLAN operation	2.4GHz ISM band	20	Max 11 Mbps	DSSS	38	140	1999
IEEE802.11g	High rate extension	2.4GHz ISM band	20	Maximum 54 Mbps	DSSS, OFDM	38	140	2003
IEEE 802.11n	PhysicalMAC: enhancement to increase throughput	2.4 GHz, 5GHz	24,40	600 Mbps	OFDM	70	250	2009

Syllabus Topic : IEEE 802.16(WiMAX) – Mesh Mode

HM 3.10 WMAN / IEEE 802.16 / WiMAX

- Wireless Metropolitan Area Network (MAN) is the name trademarked by the IEEE 802.16 Working Group. It is invented in the year 2001.
- This group focuses on Broadband Wireless Access Standards for its wireless metropolitan area network standard. The standard is popularly known as WiMAX.

3.10.1 WiMAX

Q. Explain in detail Wi-Max Technology.

1. Operating frequency : 10 – 66 GHz

3. Modulation used : QPSK, 16 QAM and 64 QAM

2. Bit rate : 32 - 134 Mbps at 28 MHz

4. Bandwidth : maximum 28 MHz

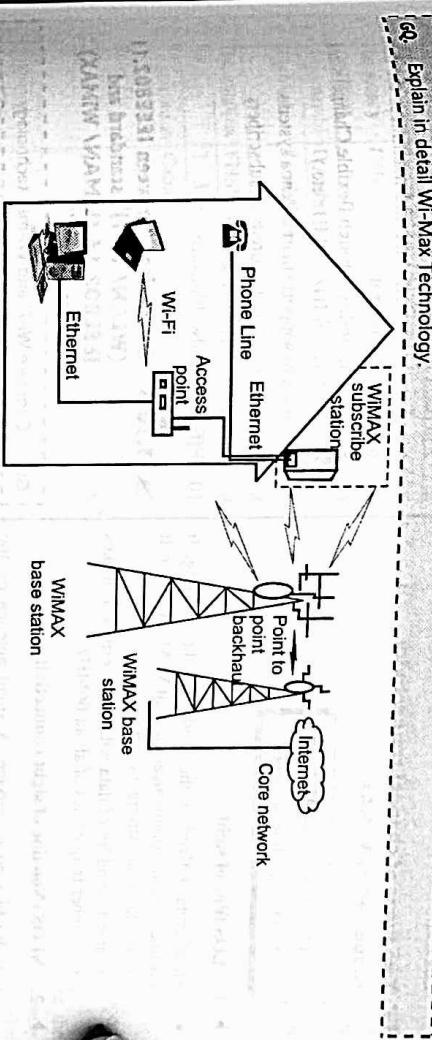


Fig. 3.10.1 : Architecture of WiMAX

- WiMAX is abbreviated for worldwide interoperability for microwave access.
- It is designed to provide broadband Internet access from fixed or mobile devices via antennas. It is the last mile solution for wireless broadband access. It is an alternative to cable and DSL (Digital Subscriber's Line). Hence it is said that WiMAX is IP centric service provided over wide area.
- Subscriber stations communicate with base-stations that are connected to a core network.
- It is a good alternative to fixed line networks. It is simple to build and relatively inexpensive.
- Refer Fig. 3.10.1. It shows general architecture of WiMAX.
- WiMAX supports very high bandwidth solutions where large spectrum deployments (i.e. >10 MHz) are desired. It also reduces the cost as no cables are involved.
- The Architecture of the WiMAX comprises of :
 - 1. Core network
 - 2. BS (Base station)
 - 3. SS (Subscriber station)
 - 4. TE (Terminal equipment)
- 1. Core network
 - It is the standard internet network. It provides the platform for the broadband connectivity.
- 2. BS (Base station)
 - It is the WiMAX cell site. They are the towers with antennas equipped over it. It provides connectivity with the public network. The communication between the subscriber station and the base station is two-way.
 - Uplink connection is from SS to BS and downlink connection is from BS to SS.
 - A single BS can cover upto 3000 Sq. miles of area.

