



RECURSION WITH STRINGS

Q1. Reverse a string using recursion.

I/P: "abcdc"
O/P: "cdcba"

I/P: "babbar"
O/P: "rabbab"

I/P: "abccba"
O/P: "abccba"

Logic: Solve 1 case and identify base cases.

Our function will look something like this:

```
void reverse (string s, int i, int j) {  
    ~~~~~  
    ~~~~~  
}
```

where i & j are the 2 indices to be swapped.

Thus, base case will be when $i > j$.

We solve the given sub-task by swapping elements at i and at j , and then incrementing & decrementing i and j respectively. Call the function again till we hit base case.

Code:

```
#include<iostream>  
using namespace std;  
  
void reverse(string& str, int i, int j) {  
    //base case  
    if(i>j)  
        return ;  
  
    swap(str[i], str[j]);  
    i++;  
    j--;  
  
    //Recursive call  
    reverse(str,i,j);  
}  
  
int main() {  
    string name = "babbar";  
  
    reverse(name, 0, name.length()-1);  
  
    cout << name << endl;  
  
    return 0;  
}
```



Don't forget to pass the string by reference or else we will not be modifying the original string name

Usecase: Let name = "abcde"

① reverse("abcde", 0, 4)

↳ abcde \Rightarrow ebcda
 $i \rightleftharpoons j$

② reverse("ebcda", 1, 3)

↳ ebcda \Rightarrow edcba
 $i \rightleftharpoons j$

③ reverse("edcba", 2, 2)

↳ edcba \Rightarrow edcba
 $i \rightleftharpoons j$

④ reverse("edcba", 3, 1)

↳ Base Case hits.

We can also make
our base case hit
for $i \geq j$

Homework: Improve the code by using up just one
extra variable i instead of 2 (i & j).

Logic: Same as above just use i to get the element
to be swapped. We can use the string's size to
get the element.

First Element : i

Second Element : $n - i - 1$.

Code :

```
1 #include <iostream>
2 using namespace std;
3
4 void reverse(string& s, int i) {
5     if(i >= s.size()/2)
6         return;
7     // swap the i and n-i-1 indexed nodes
8     swap(s[i], s[s.size() - i - 1]);
9     // call for the next index
10    reverse(s, i+1);
11 }
12
13 int main(void)
14 {
15     string s;
16     cin >> s;
17     reverse(s, 0);
18     cout << s << endl;
19     return 0;
20 }
```

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ork\Coding\Recursion\" ; if (\$?) { g++ ReverseString.cp
abcdefg
gfedcba

Q2 Check Palindrome. (If the reversed string is same
as the original string)

I/P : "abccba"

O/P : true

I/P : "abcbde"

O/P : false

Approach 1: Reverse string and check if it's equal to the original string.

Time Complexity : $O(n)$

Space Complexity : $O(n)$

Approach 2: Take 2 pointers i & j pointing to the first and last element. Check if they are equal & then increment and decrement i & j respectively till $i < j$. If at some point $str[i] \neq str[j]$, return false.

Code :

```
bool checkPalindrome(string str, int i, int j) {  
    //base case  
    if(i > j)  
        return true;  
  
    if(str[i] != str[j])  
        return false;  
    else{  
        //Recursive call  
        return checkPalindrome(str, i+1, j-1);  
    }  
}
```

Homework: Write the recursive function using just one extra pointing variable.

Code :

```
1  #include <iostream>  
2  using namespace std;  
3  
4  bool isPalindrome(string& s, int i) {  
5      if(i >= s.size()/2) return true;  
6  
7      if(s[i] != s[s.size()-i-1])  
8          return false;  
9  
10     return isPalindrome(s, i+1);  
11 }  
12  
13 int main(void)  
14 {  
15     string s;  
16     cin >> s;  
17     if(isPalindrome(s, 0)) {  
18         cout << "Palindrome\n";  
19     }  
20     else {  
21         cout << "Not Palindrome\n";  
22     }  
23     return 0;  
24 }
```

(Similar to reverse)
string

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abccba
Palindrome

Note: The time complexity of accessing the size of a `std::string` in C++ is $O(1)$.

Q3. Find $\text{power}(a, b)$ using recursion.

I/P : 2, 4

O/P : 16

Approach: We can represent a^b as follows:

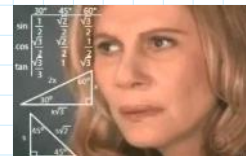
$$a^b = \begin{cases} a \cdot a^{b-1} & \text{if } b \text{ is odd} \\ a^{b/2} \cdot a^{b/2} & \text{if } b \text{ is even} \end{cases}$$

We can use a recursive function with base case being $b=0$ (or $b=1$), because $a^0 = 1$ (or $a^1 = a$)

For a given call, if $b \% 2 == 0$ then return $\text{func}(a, b/2) * \text{func}(a, b/2)$,

else return $a * \text{func}(a, b-1)$ which equals $a * \text{func}(a, b/2) * \text{func}(a, b/2)$.

Confused?



Example: $\text{power}(2, 5)$

$$\begin{aligned} 2^5 &= 2 \times 2^4 \\ 2^4 &= 2^2 \times 2^2 \\ 2^2 &= 2^1 \times 2^1 \\ 2^1 &= 2 \end{aligned}$$

When $\text{power}(2, 5)$ is called, we will perform 2 steps in one call to improve efficiency.

Instead of giving back $2 * \text{power}(2, 4)$

we will calculate $\text{pow}(2, 2)$ before-hand by using $\text{variable} = \text{power}(2, 5/2)$;

Then return $2 * \text{variable} * \text{variable}$.

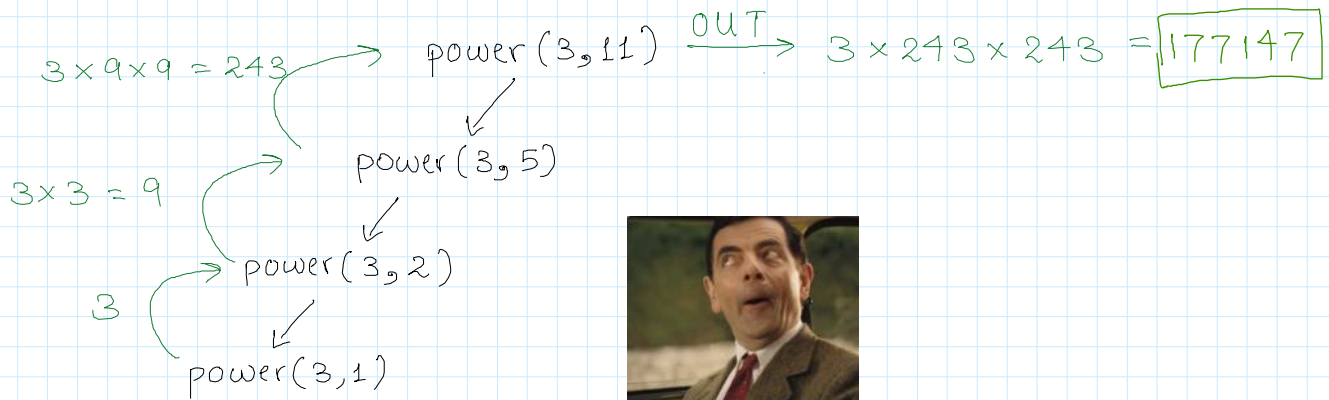
For even power,

return $\text{variable} * \text{variable}$.

Code :

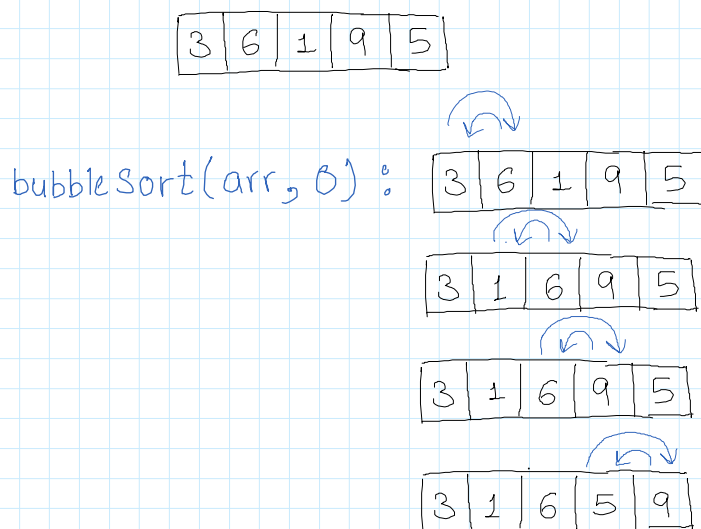
```
int power(int a, int b) {  
    //base case  
    if( b == 0 )  
        return 1;  
  
    if(b == 1)  
        return a;  
  
    //RECURSIVE CALL  
    int ans = power(a, b/2);  
  
    //if b is even  
    if(b%2 == 0) {  
        return ans * ans;  
    }  
    else {  
        //if b is odd  
        return a * ans * ans;  
    }  
}
```

Recursion Tree : $\text{power}(3, 11)$



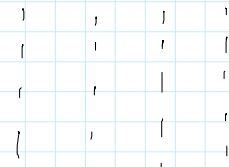
Q4 Bubble Sort using Recursion.

Approach: The recursive function will be called n times, each time it will correctly position the i^{th} largest element correctly. Base case hits when only 1 or no elements are left.



bubbleSort(arr, 1) :

1	3	6	5	9
---	---	---	---	---



bubbleSort(arr, n-1) :

1	3	5	6	9
---	---	---	---	---

Code :

```
void sortArray(int *arr, int n) {
    //base case - already sorted
    if(n == 0 || n == 1) {
        return ;
    }

    //1 case solve karlia - largest element ko end me rakh dega
    for(int i=0; i<n-1; i++) {
        if(arr[i] > arr[i+1]){
            swap(arr[i], arr[i+1]);
        }
    }

    //Recursive Call
    sortArray(arr, n-1);
}
```

Homework:

① selection sort using recursion.

Code :

```
int minIndex(int a[], int i, int j)
{
    if (i == j)
        return i;
    // minimum element's index from a[(i+1)...j]
    int k = minIndex(a, i + 1, j);
    // Comparing with current element and updating answer
    return (a[i] < a[k])? i : k;
}

void recurSelectionSort(int a[], int n, int index = 0)
{
    // Return when starting and size are same
    if (index == n)
        return;
    // minimum element's index in a[index...(n-1)]
    int k = minIndex(a, index, n-1);
    swap(a[k], a[index]);
    // Recursively calling selection sort function
    recurSelectionSort(a, n, index + 1);
}
```

```
23     swap(a[k], a[index]);
24
25     // Recursively calling selection sort function
26     recurSelectionSort(a, n, index + 1);
27 }
28
```

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```
5
3 1 4 5 2
1 2 3 4 5
```

[Refer Here](#) for detailed explanation of Homework Questions.