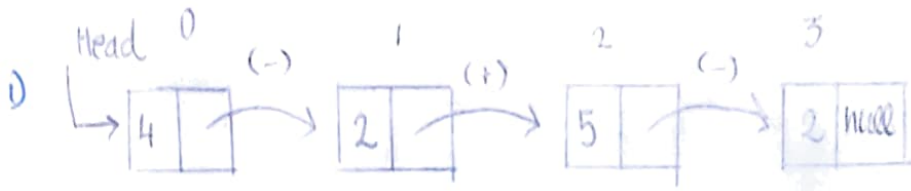


# Linked List

①



algorithm alt-sum(head)

1) If (head == null)

return 0

2) else {

~~while~~

temp = head

count = 0

sum = 0

while (temp != null) {

if (count % 2 == 0) {

sum += temp → data

}

else {

sum -= temp → data

}

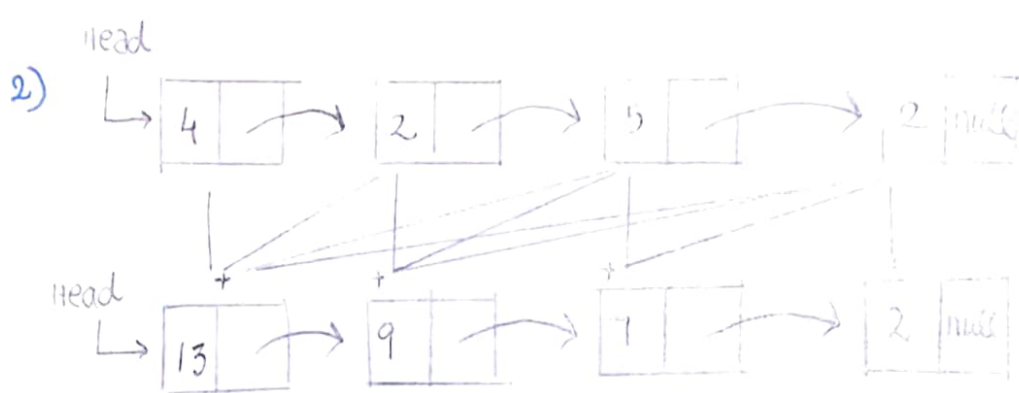
count++;

temp = temp → next

}

3) return sum;

4) end.



Basic logic : for each node sum up the value of current node and the nodes that come after the current node

algorithm add (head)

1) current\_node = head

2) temp = head

~~3) for (; head != null~~

3) for (; current\_node != NULL ; current\_node = current\_node → next) {

this loop is to iterate through the linked list

~~for (temp = head → next~~

to iterate through the next set of nodes that are after the current node.

for (temp = current\_node → next ; temp != null ; temp = temp → next) {

current\_node → data += temp → data

}

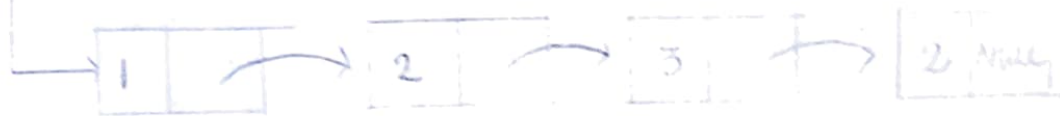
}

4) return head

5) end.

3) head  $x = 2$

2



RESULT:



Algorithm delete\_var (head, x)

IF (head == NULL) // empty list

return 0

while

1) ~~if~~ (head → data == x) // delete in front

head = head → next

3) temp = head

4) while (temp → next → next != NULL) // iterate till last before node

{ if (temp → next → data == x) {

temp → next = temp → next → next

}

else {

temp = temp → next

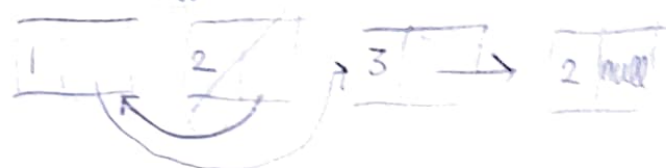
}

}

5) return head

6) end

1st iteration



2nd iteration



Eg1:

4)



ANS:



Eg2:



ANS:



TRACING the algorithm: for eg 2

prev = null

first = head

① st iteration

∴ first = head (A) & first → next = B

while loop is true

first = A

first → next = second → next

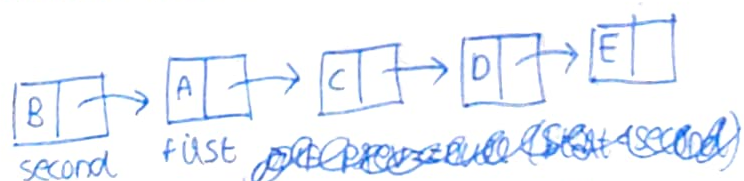
second = B



prev = null

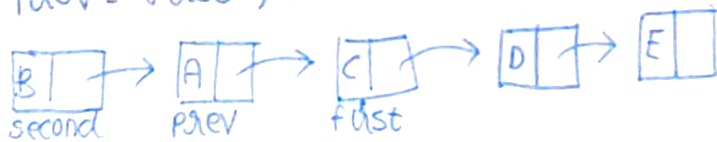
second → next = first

start = ~~first~~  
head → next



done for next iteration purpose

prev = first, first = prev → next



②nd iteration

first = [C | ]

second = [B | ]

prev = [A | ]

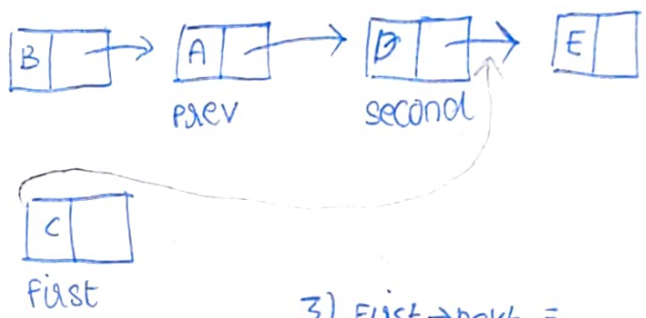
1) second = first → next

⇒ second = [D | ]



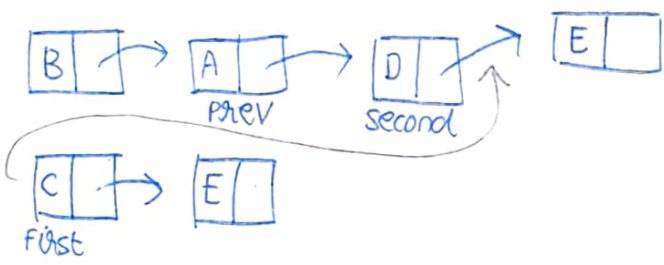
∴ (prev → next ≠ null)

2) prev → next = second



3) first → next =

~~second~~ → next = ~~first~~ second → next



4) second → next = first



5) for next iteration prev = first

first = prev → next



③rd iteration is false because first → next is null  
so return start

Algorithm swap (head)

1) if (head == NULL || head → next == NULL) {  
    return head; }

2) start = head → next

3) prev = NULL



4) first = head

5) while (first != NULL && first → next != NULL) {

    second = first → next;

    if (prev != NULL) {

        prev → next = second

    }

    first → next = second → next

    second → next = first

    prev = first

    first = prev → next

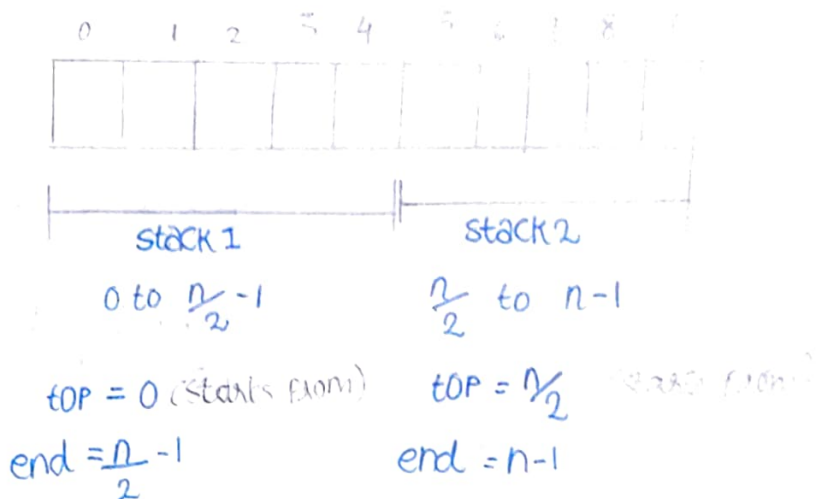
}

6) return start

7) end.



1) Implementing 2 stacks in a single array



algorithm 2stacks()

1) initialize an array of size  $n$

2)  $\text{top1} = -1$

3)  $\text{top2} = \frac{n}{2}$

4) to push into first stack do the following operations

(i) check if  $\text{top1} \neq \frac{n}{2} - 1$   
 $\text{top1}++;$   
 Push (element)  
 $\text{top1}++;$

(ii) else print stack 1 overflow

5) to pop from first stack

(i) if  $(\text{top1} \neq -1)$   
 $\text{POP}()$   
 $\text{top1}--;$

(ii) else print stack 1 underflow

7) to pop from stack 2

(i) if  $(\text{top2} \neq \frac{n-1}{2})$   
 $\text{POP}()$   
 $\text{top2}--;$

(ii) else print stack 2 underflow

8) end

6) to push into second stack

(i) if  $(\text{top2} \neq n)$   $\text{top2}++;$  Push (element),  $\text{top2}++;$

(ii) else print stack 2 overflow

### 5) CONVERT INFIX TO POSTFIX IN STACKS

INPUT:  $A+B*C+D$       OUTPUT:  $ABC*+D+$

LOGIC:

Precedence of operators

- |        |               |
|--------|---------------|
| 1) ( ) | associativity |
| 2) ^   | R to L        |
| 3) / * | L to R        |
| 4) + - | L to R        |

$A+B*C+D$

1)  $A+BC*+D$

2)  $ABC*++D$

3)  $ABC*+D+$

RULES FOR CONVERSION USING STACK

- write operand directly as output
- If operator, perform push operation based on the following conditions

(i) If the precedence of the incoming operator is higher than the top operator in the stack then push it into stack

(ii) else POP and keep checking again

(iii) If there is same precedence like + and - check if L to R or R to L. If L to R pop and check else push



INPUT  $A+B * C+D$

Input	Stack	Postfix
A	-	A
+	<div style="border: 1px solid black; padding: 2px; display: inline-block;">+</div>	A
B	<div style="border: 1px solid black; padding: 2px; display: inline-block;">+</div>	AB
*	<div style="border: 1px solid black; padding: 2px; display: inline-block;">* + <small>* &gt; +</small></div>	AB
C	<div style="border: 1px solid black; padding: 2px; display: inline-block;">* +</div>	ABC
+	<div style="border: 1px solid black; padding: 2px; display: inline-block;">* + POP</div>	ABC* +
D	<div style="border: 1px solid black; padding: 2px; display: inline-block;">+ + + <small>L to R so POP</small></div>	ABC* +
D	<div style="border: 1px solid black; padding: 2px; display: inline-block;">+</div>	ABC* + D
empty	pop	ABC* + D +

### Algorithm intopost (expression)

- 1) Iterate through every character and perform the following operations
  - 2) If the character is an operand display as output
  - 3) else  
Check the Priority of the current operator and the operator in the top of the stack
  - 4) If the incoming priority is greater, then push operator into stack and goto step 6
  - 5) else pop ~~from~~ from stack and display as output and repeat step 3
  - 6) Repeat from step 2 for every character
  - 7) If the expression is fully completed and the stack is not empty, pop out all the operators from the stack
- end.

### 3) well formedness of brackets

$\{ () \}$  ✓       $\{ ( \} )$  ✗

### Algorithm brackets (expression)

- 1) Iterate through every character of the expression
- 2) If the character is an opening bracket, then push into stack

3) Else if the character is a closing bracket  
check if the top of the stack is an  
equivalent opening bracket

⑥ ⑥

4) If yes then POP the top of stack and  
continue from step 2

5) If not, then return that the brackets are  
not balanced and exit

6) If the expression is fully completed and  
stack is emptied return that the brackets are  
balanced.

7) end.

4) Range count problem

5
4
3
2
1

Var = 2

Output  $\rightarrow 4$

because 4 numbers in  
the stack are greater  
than or equal to 2

Algorithm range-count (stack, var)

1) ~~if stack~~

1) count = 0

2) while (isEmpty (stack) == False)

~~pop stack~~ pop (stack)

if (stack[top]  $\geq$  var) {

~~count~~ count++;

}

top--;

3) return count;

6) end