

# Improved Approximation Algorithms for Bounded-Degree Local Hamiltonians

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**arXiv: 2105.01193**

**Joint work with**

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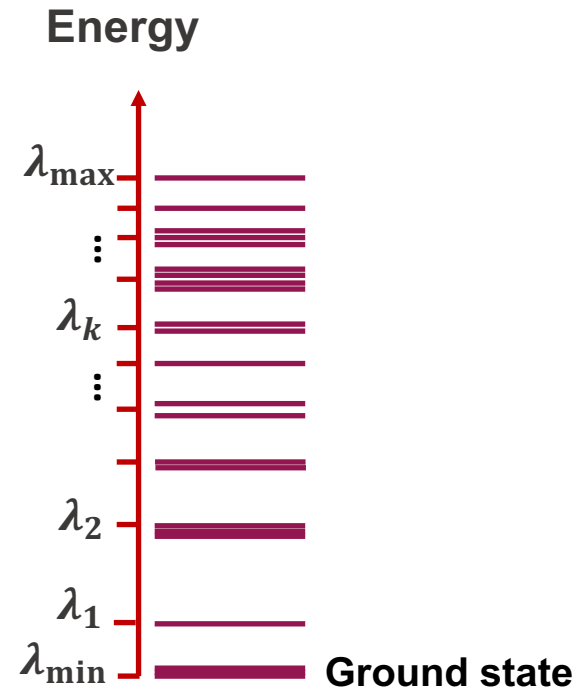
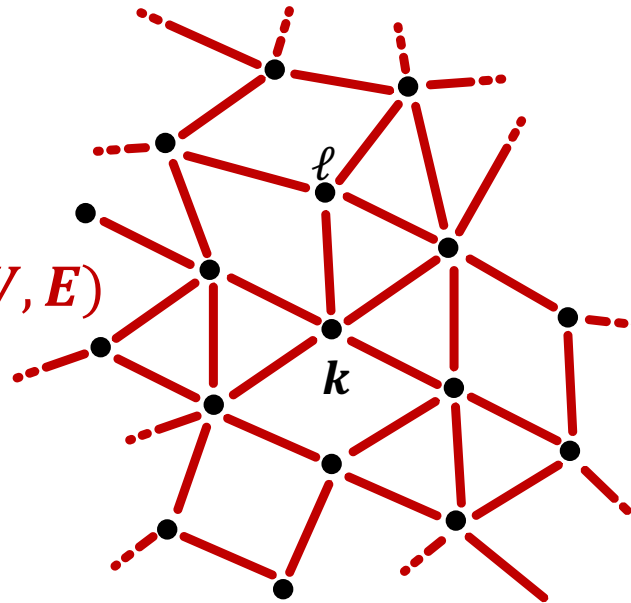
# **Problem Statement and Background**

# Interacting Quantum Systems

Interaction graph  $G = (V, E)$

$|V| = n$  qubits

$|E| = m$  interactions



Local Hamiltonians

$$H = \sum_{\{k,\ell\} \in E} h_{k\ell}$$

Degree- $d$  interaction graph

Ground state of  $H$  captures the low-temperature physics

**Believed to require  $\exp(n)$  resources  
to compute in the worst case**

# **1 Worst-Case Complexity and Rigorous Algorithms**

## **2 Heuristic Quantum Algorithms**

## Worst-Case Complexity

- Although **ground state energy**  $= \lambda_{\min}(H)$ ,  
more convenient to consider estimating

$$\lambda_{\max}(H) = \max_{\psi} \langle \psi | H | \psi \rangle$$

*Equivalent because*

$$\lambda_{\min}(H) = -\lambda_{\max}(-H)$$

- QMA-hard to estimate  $\lambda_{\max}(H)$  with  $\frac{1}{\text{poly}(n)}$  additive error

[Kitaev 1999, Kempe, Kitaev, Regev 2004]

- PCP Theorem: For some constant  $0 < \epsilon < 1$ ,  
remains **NP-hard** to estimate  $\lambda_{\max}$  within additive error  $\epsilon \cdot m$



[Arora, Lund, Motwani, Sudan, Szegedy '98,  
Arora, Safra '98, Dinur '07]

QMA-hard? qPCP conjecture

# Worst-Case Complexity

**Approximation algorithms:** compute estimate  $\hat{\lambda} \leq \lambda_{\max}$  s.t.

$$r = \hat{\lambda} / \lambda_{\max}$$

is as large as possible.

**What is the largest approximation ratio  $r$  achievable with efficient algorithms?**

*Known Algorithms e.g. for*

- **Heisenberg-like interactions:**  $\mathbf{h}_{ij} = \mathbf{I} - \mathbf{X}_i \mathbf{X}_j - \mathbf{Y}_i \mathbf{Y}_j - \mathbf{Z}_i \mathbf{Z}_j$

[Gharibian, Parekh 2019, Anshu, Gosset, Morenz Korol 2020]

- **Positive semidefinite:**  $\mathbf{h}_{ij} \geq 0$

[Gharibian, Kempe 2012]

- **Traceless:**  $\text{Tr}[\mathbf{h}_{ij}] = 0$

[Bravyi, Gosset, König, Temme 2019]

- **Dense or Planar graphs**

[Bansal, Bravyi, Terhal 2009, Gharibian, Kempe 2012, Brandão, Harrow 2014]

# Worst-Case Complexity

**Most of these algorithms compute a quantum state  $|v\rangle$  that**

$$|v\rangle = |v_1\rangle \otimes |v_2\rangle \otimes \cdots \otimes |v_n\rangle$$

**or**

**$|v\rangle =$  tensor product of few-qubit states**

**But ground states may be highly **entangled**,**

**What is the structure of states  
with high approximation ratio?**

# Worst-Case Complexity

**What is the structure of states  
with high approximation ratio?**

**For high degree graphs,  
product states provide good approximations**

Monogamy of Entanglement  
Mean-field Approximation

[Brandão, Harrow 2014]

**For Hamiltonians on **degree- $d$**  graph with  $n$  qubits and  $m$  interactions, there exists  $|v\rangle = |v_1\rangle \otimes |v_2\rangle \otimes \cdots \otimes |v_n\rangle$  s.t.**

$$\lambda_{\max}(H) - \langle v|H|v\rangle \leq O\left(\frac{m}{d^{1/3}}\right)$$

**This work:**

**Extensive** improvement over product states **for bounded-degree graphs** using shallow (low-depth) quantum circuits



# **1 Worst-Case Complexity and Rigorous Algorithms**

## **2 Heuristic Quantum Algorithms**

- Many **heuristic** classical or quantum algorithms for estimating ground state energy
- Ground states could be highly **entangled**  
Potential advantage in using **quantum computers**
- E.g. variationally optimize energy over output states of **shallow (low-depth)** quantum circuits

$$|\psi(\theta)\rangle = U(\theta)|0^n\rangle$$

$\langle\psi(\theta)|H|\psi(\theta)\rangle$   
Measure with quantum computer

$\min_{\theta} \langle\psi(\theta)|H|\psi(\theta)\rangle$   
Optimize with classical computer

- Many **heuristic** classical or quantum algorithms for estimating ground state energy
- Ground states could be highly **entangled**  
Potential advantage in using **quantum computers**
- E.g. variationally optimize energy over output states of **shallow (low-depth)** quantum circuits
  - Can be implemented on **small** quantum computers
  - Some known **limitations** in efficacy

[McClean et al 2018]

[Bravyi, Kliesch, Koenig, Tang 2020]

[Farhi, Gamarnik, Gutmann 2020]

[Bravyi, Gosset, Movassagh 2021]

**Rigorous bounds on the performance of shallow quantum circuits for estimating ground energy?**

# Recap

Many known **rigorous** algorithms output **product states**.

*How can we **improve** them by applying **quantum circuits**?*

Many **near-term** algorithms use **shallow** quantum circuits

*How can we **rigorously** bound their **performance**?*

# Main Results

# Result: Improving product state approx.

Define variance of a state  $|v\rangle$  by

$$\text{Var}_v(H) = \langle v|H^2|v\rangle - \langle v|H|v\rangle^2$$

Given a degree- $d$  Hamiltonian  $H$  and a product state  $|v\rangle$ , we can **efficiently** compute a **depth- $(d + 1)$**  quantum circuit  $U$  such that the state  $|\psi\rangle = U|v\rangle$  satisfies

$$\langle \psi|H|\psi\rangle \geq \langle v|H|v\rangle + \Omega\left(\frac{\text{Var}_v(H)^2}{d^2 m}\right)$$

- An improvement of  $\Omega(m)$  in estimated energy when

$$\text{Var}_v(H) = \Omega(m) \text{ and } d = O(1).$$

- No improvement when  $|v\rangle$  is an eigenstate of Hamiltonian (e.g. purely classical case)

# Proof Idea of 1<sup>st</sup> Result

Choice of circuit  $U$  for state  $|v\rangle = |v_1\rangle \otimes |v_2\rangle \otimes \cdots \otimes |v_n\rangle$

$$U(\theta) = \bigotimes_{\{i,j\} \in E} e^{i\theta_{ij} P_i P_j} = e^{i \sum_{\{i,j\} \in E} \theta_{ij} P_i P_j}$$

$$\|P_i\| \leq 1, \quad \langle v_i | P_i | v_i \rangle = 0 \quad \forall i \in V$$

- Generalizes **level-1 QAOA**  $P_i = e^{i\beta \sum_{k \in V} X_k} Z_k e^{-i\beta \sum_{k \in V} X_k}$
- **Locally & slightly rotates**  $|v_i\rangle|v_j\rangle$  towards the ground space

Example: Antiferromagnetic Heisenberg Interactions

[Anshu, Gosset, Morenz Korol 2020]

$$H = \sum_{\{i,j\} \in E} w_{ij} h_{ij}$$

$$h_{ij} = \frac{1}{4} (I - X_i X_j - Y_i Y_j - Z_i Z_j) = |\Psi_{ij}\rangle \langle \Psi_{ij}|$$

$$|\Psi_{ij}\rangle = \frac{1}{\sqrt{2}} (|0\rangle_i |1\rangle_j - |1\rangle_i |0\rangle_j)$$

$$e^{-i\theta_{ij} X_i Y_j} |0\rangle_i |1\rangle_j = \cos(\theta_{ij}) |0\rangle_i |1\rangle_j - \sin(\theta_{ij}) |1\rangle_i |0\rangle_j$$



## Bounding improvement in energy

$$U(\boldsymbol{\theta}) = \bigotimes_{\{i,j\} \in E} e^{i\theta_{ij}P_iP_j} = e^{i \sum_{\{i,j\} \in E} \theta_{ij}P_iP_j}$$

$$|\boldsymbol{\psi}\rangle = U(\boldsymbol{\theta})|v\rangle$$

$$\langle \boldsymbol{\psi} | \boldsymbol{h}_{ij} | \boldsymbol{\psi} \rangle = \langle v | U(\boldsymbol{\theta})^\dagger \boldsymbol{h}_{ij} U(\boldsymbol{\theta}) | v \rangle$$

$$= \langle v | \boldsymbol{h}_{ij} | v \rangle - i \theta_{ij} \langle v | [\boldsymbol{P}_i \boldsymbol{P}_j, \boldsymbol{h}_{ij}] | v \rangle + \mathbf{Err} \quad \langle v_i | \boldsymbol{P}_i | v_i \rangle = 0$$


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$$= \langle v|h_{ij}|v\rangle - i\theta_{ij} \langle v|[P_iP_j, h_{ij}]|v\rangle + \mathbf{Err} \quad \langle v_i|P_i|v_i\rangle = 0$$


$$\theta_{k\ell} = \theta_0 \cdot \text{sign}(-i\langle v|[P_kP_\ell, h_{ij}]|v\rangle)$$

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$$|\text{Err}| \leq O(\theta_0^2 d)$$

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$$\geq \langle v|h_{ij}|v\rangle + \theta_0 |\langle v|[P_iP_j, h_{ij}]|v\rangle| - \Omega(\theta_0^2 d)$$

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$$\geq \langle v | h_{ij} | v \rangle + \theta_0 |\langle v | [P_i P_j, h_{ij}] | v \rangle| - \Omega(\theta_0^2 d)$$

There are choices of  $\{P_i\}$  such that for  $\theta_0 \leq O(1/d)$ ,

$$\langle \psi | H | \psi \rangle \geq \langle v | H | v \rangle + \Omega\left(\frac{\text{Var}_v(H)^2}{d^2 m}\right)$$

# Extensions and Tightness

# Result: locally optimal states & tightness

Improved bound:

A product state  $|v\rangle$  is **locally optimal** if for any **single-qubit operator**  $Q$ ,

$$\frac{d}{d\phi} \langle v | e^{-i\phi Q} H e^{i\phi Q} | v \rangle = 0 \quad \text{at} \quad \phi = 0$$

For locally optimal states,

$$\langle \psi | H | \psi \rangle \geq \langle v | H | v \rangle + \Omega \left( \frac{\text{Var}_v(H)^2}{d m} \right)$$

Tightness:

For simple Hamiltonians e.g.  $h_{ij} = Z_i + Z_j$  and

$$|v\rangle = (\cos(\theta) |0\rangle - \sin(\theta) |1\rangle)^{\otimes n}$$

We have

$$\lambda_{\max} - \langle v | H | v \rangle \leq O \left( \frac{\text{Var}_v(H)^2}{d^2 m} \right)$$

# Result: $k$ -local Hamiltonians

## Improvement for $k$ -local Hamiltonians



Given a degree- $d$   $k$ -local Hamiltonian  $H$  and a product state  $|v\rangle$ , we can efficiently compute a shallow quantum circuit  $U$  such that the state  $|\psi\rangle = U|v\rangle$  satisfies

$$\langle\psi|H|\psi\rangle \geq \langle v|H|v\rangle + \Omega\left(\frac{\text{Var}_v(H)^2}{2^{O(k)} d^4 m}\right)$$



# Result: Improving entangled states

Let  $|v\rangle = W|0^n\rangle$  where  $W$  is a quantum circuit of depth  $D$ .

We can efficiently compute a quantum circuit  $U$  such that the state  $|\psi\rangle = U|v\rangle$  satisfies

↓  
Lightcone  $\ell$

$$\langle\psi|H|\psi\rangle \geq \langle v|H|v\rangle + \Omega\left(\frac{\text{Var}_v(H)^2}{2^{O(D)} d^2 m}\right)$$

↓  $\ell^{10}$

- The circuit  $U$  is not constant-depth anymore
- The bound extends to when  $|\psi\rangle$  is the unique ground state of some  $\ell$ -local gapped Hamiltonian

# **Generic Performance and Comparison with Local Classical Algorithms**

# Result: Improvement for random states

Write  $H$  in terms of Pauli operators  $\sigma_1, \sigma_2, \sigma_3$ , and  $\sigma_0 = I$ :

$$H = \sum_{\{i,j\} \in E} \sum_{x,y} f_{xy}^{ij} \sigma_x^i \otimes \sigma_y^j$$

Define

$$\text{quad}(H) = \sum_{\{i,j\} \in E} \sum_{x>0,y>0} \left(f_{xy}^{ij}\right)^2$$

There is an **efficient randomized** algorithm which computes a **depth- $(d+1)$**  quantum circuit  $U$  such that  $|\psi\rangle = U|v\rangle$  satisfies

$$\mathbb{E}_v \langle \psi | H | \psi \rangle \geq \mathbb{E}_v \langle v | H | v \rangle + \Omega \left( \frac{\text{quad}(H)^2}{d m} \right)$$

For triangle-free graphs, we have

$$\mathbb{E}_v \langle \psi | H | \psi \rangle \geq \mathbb{E}_v \langle v | H | v \rangle + \Omega \left( \frac{\text{quad}(H)}{\sqrt{d}} \right)$$

# Result: Local Classical Algorithm

For triangle free graphs, there is an **efficient randomized** algorithm that computes the product state  $|v\rangle$  satisfying

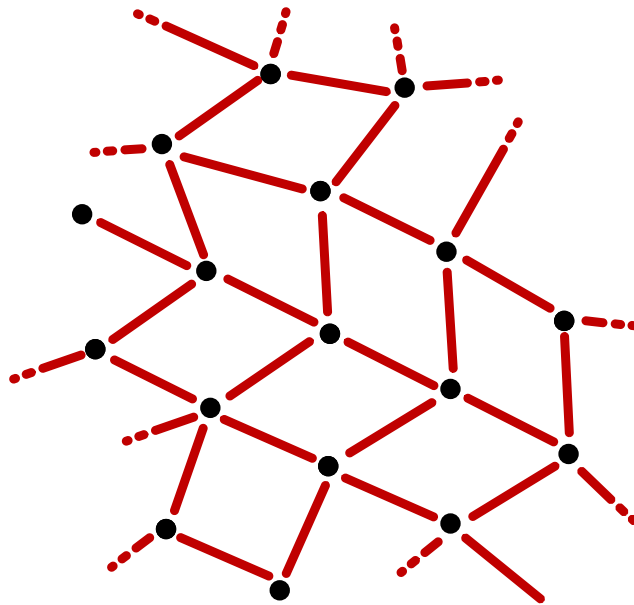
$$\mathbb{E}_v \langle v | H | v \rangle \geq \frac{1}{4} \text{Tr}(H) + \Omega\left(\frac{\text{quad}(H)}{\sqrt{d}}\right)$$

Similar to [Hastings '19, Harrow, Montanaro '17, Barak et al '15]

# Result: Local Classical Algorithm

## Local Classical Algorithm

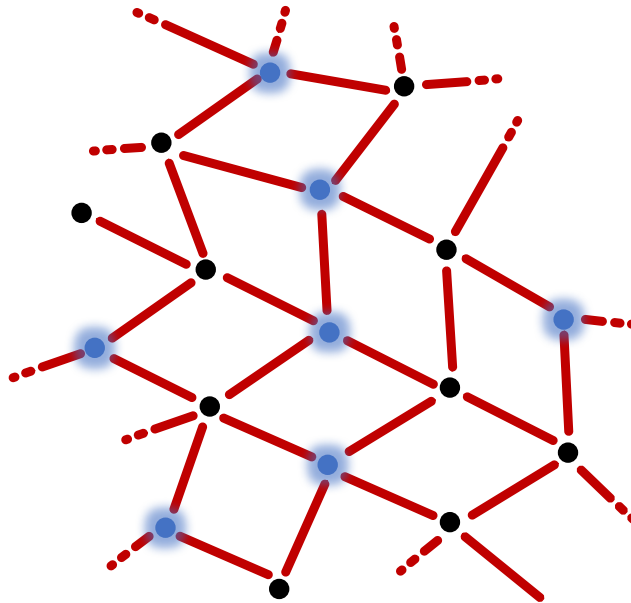
- **Assign i.i.d states to all vertices uniformly at random**



# Result: Local Classical Algorithm

## Local Classical Algorithm

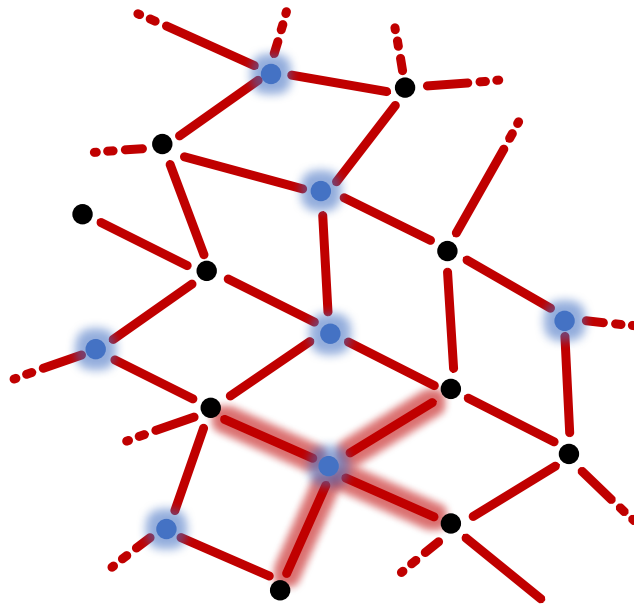
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- Randomly divide vertices into two sets  $\{\bullet\}$ ,  $\{\bullet\}$



# Result: Local Classical Algorithm

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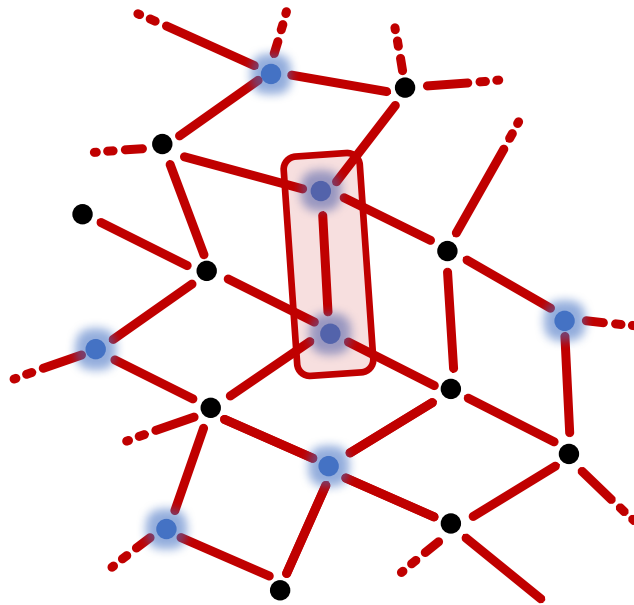
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# Result: Local Classical Algorithm

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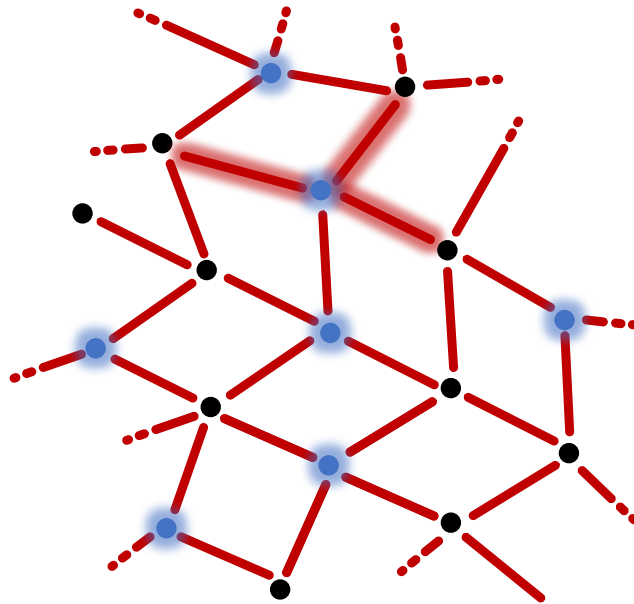




# Result: Local Classical Algorithm

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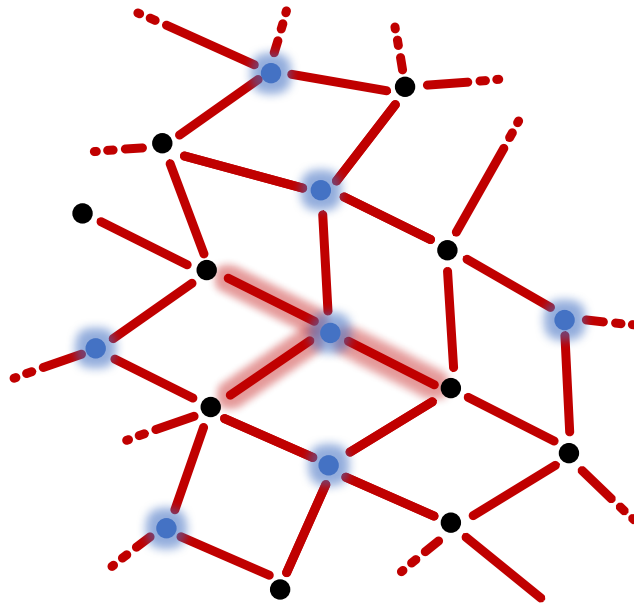
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# Result: Local Classical Algorithm

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**Classical local** algorithms may achieve the **same scaling** with product states.

But their output can be **further improved** by our shallow circuit

We also saw

For **locally optimal states**,

$$\langle \psi | H | \psi \rangle \geq \langle v | H | v \rangle + \Omega \left( \frac{\text{Var}_v(H)^2}{d_m} \right)$$

Better energy improvement can be achieved with **structured** initial states.

# Open Questions

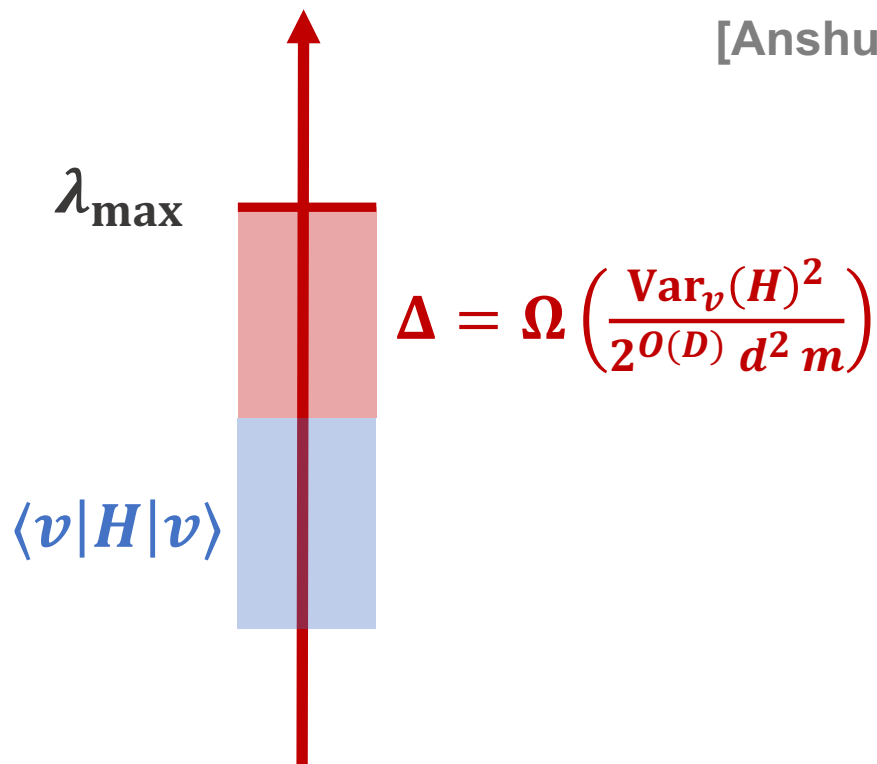
# Open questions

- Limitations on energy of states generated with **low-depth circuits**

$$\langle v|H|v\rangle \leq \lambda_{\max} - \Omega\left(\frac{\text{Var}_v(H)^2}{2^{O(D)} d^2 m}\right)$$

Examples of Hamiltonians with almost **NLTS** property?

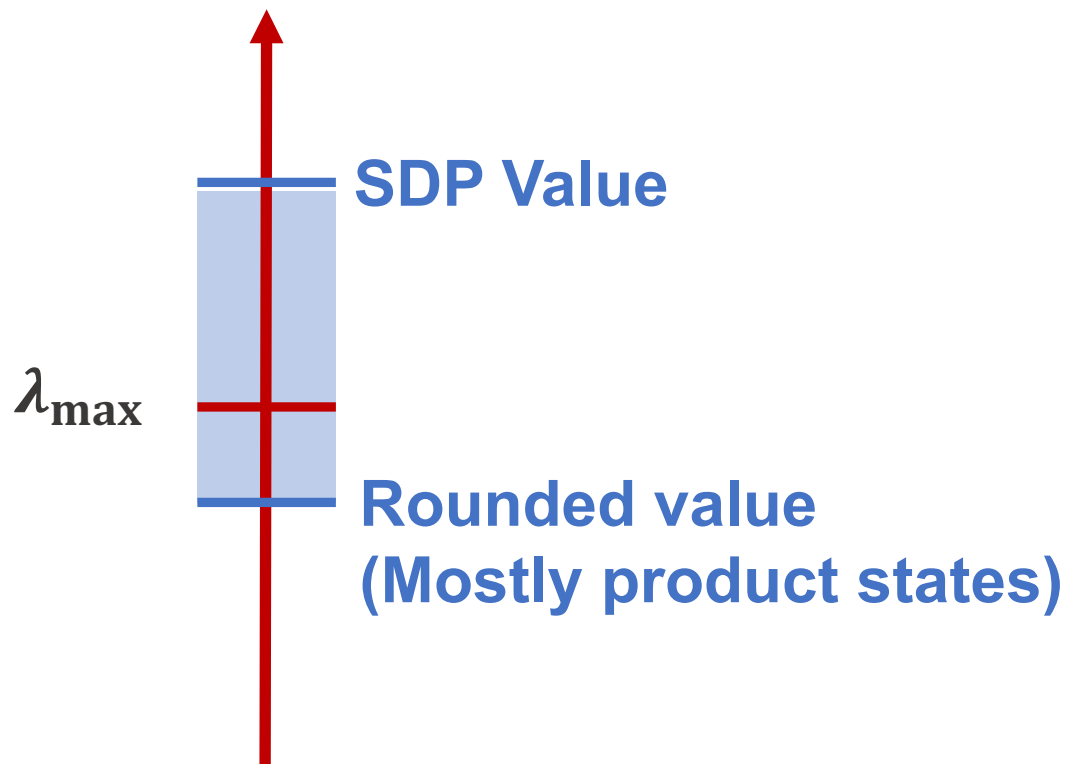
[Anshu, Nirkhe 2021]



# Open questions

- **Rounding to entangled states in SDP relaxations?**

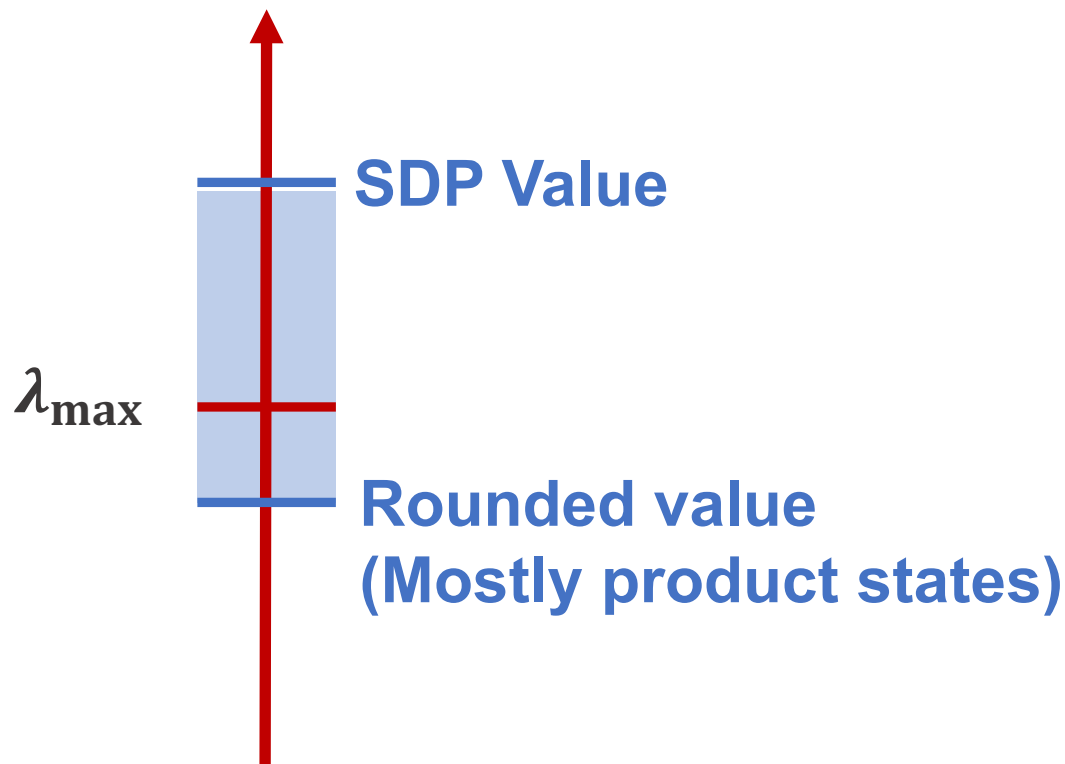
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# Open questions

- **Rounding to entangled states in SDP relaxations?**  
[Parekh, Thompson 2020, Anshu, Gosset, Morenz Korol 2020]
- **More theoretical study of near-term algorithms for estimating ground-state energy**



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