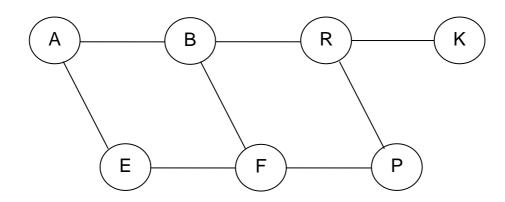
Technical Assignment 2

Part II. Data Structure and Algorithms

In the following graph, nodes are represented as alphabets.



Depth First Search (DFS)

A stack is used. PUSH and POP operations of a stack are used.

PATH 1:

Step 1:

Visited: Result: -

K	R	В	Α	Е	F	Р
0	0	0	0	0	0	0

-
-
-
-
-

Stack

Step 2:

Visited: Result: K

K	R	В	Α	Е	F	Р
1	0	0	0	0	0	0

-
-
-
-
K

Stack

Step 3:	_	
Visited: Re	esult : K R	-
K R B A E F P		- D
1 1 0 0 0 0 0		R K
		Stack
		Stack
Step 4: P and B are the adjacent nodes of I	R. For this PATH, we consider B.	
Visited: Res	sult : K R B	-
K R B A E F P		 B
1 1 1 0 0 0 0		R
		K
		Stack
Step 5: F and A are the adjacent nodes of E	B. For this PATH, we consider A.	
Visited: Res	sult : K R B A	-
violitod .	Jan. 10 10 D A	A
K R B A E F P		В
1 1 1 1 0 0 0	_	R
	L	K
		Stack
Step 6: B and E are the adjacent nodes of	A. For this PATH, we consider E as B has already been	
visited.		Е
Visited: Res	sult : K R B A E	Α
1.0		В

Stack

R

Result of PATH 1 - KRBAE

1

PATH 2:

Visited: Result: -

K	R	В	Α	Е	F	Р
0	0	0	0	0	0	0

-	
-	
-	
-	
-	_

Stack

Step 2:

Visited: Result: K

K	R	В	Α	Е	F	Р
1	0	0	0	0	0	0

ı
-
-
-
K

Stack

Step 3:

Visited: Result: K R

K	R	В	Α	Е	F	Р
1	1	0	0	0	0	0

ı
-
•
R
K

Stack

Step 4: P and B are the adjacent nodes of R. For this PATH, we consider P as we have already considered B in PATH 1.

Visited: Result: KRP

K	R	В	Α	Е	F	Р
1	1	0	0	0	0	1

-	
-	
Р	
R	
K	

Stack

Step 5: R and F are the adjacent nodes of P. For this PATH, we consider F as R has already been visited.

Visited: Result: KRPF

K	R	В	Α	Е	F	Ρ
1	1	0	0	0	1	1

F
Р
R
K

Stack

Step 6: B and E are the adjacent nodes of F. For this PATH, we consider E as it is the final destination.

Visited:

	K	R	В	Α	Е	F	Р
Ī	1	1	0	0	1	1	1

	E	
Posult · K P D F F	F	
Result: KRPFE	Р	
	R	
	K	

Stack

Result of PATH 2 - KRPFE

PATH 3:

Sieb i	Step	1
--------	------	---

Visited: Result: -

K	R	В	Α	Е	F	Р
0	0	0	0	0	0	0

-	
-	
-	
-	
-	

Stack

Step 2:

Visited: Result: K

K	R	В	Α	Е	F	Р
1	0	0	0	0	0	0

-
-
1
-
K

Stack

Step 3:

Visited: Result: K R

K	R	В	Α	Е	F	Р
1	1	0	0	0	0	0

-
ı
-
R
K

Stack

Step 4: P and B are the adjacent nodes of R. For this PATH, we consider B.

Visited: Result: KRB

K	R	В	Α	Е	F	Р
1	1	1	0	0	0	0

-
ı
В
R
K

Stack

Step 5: F and A are the adjacent nodes of B. For this PATH, we consider F as we have already used A in PATH 1.

Visited: Result: KRBF

K	R	В	Α	Е	F	Р
1	1	1	0	0	1	0

-	
F	
В	
R	
K	

Stack

Step 6: B and E are the adjacent nodes of A. For this PATH, we consider E as B has already been_

visited.

Visited:

Result:	: KR	BFE
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K	R	В	Α	Е	F	Р
1	1	1	0	1	1	0

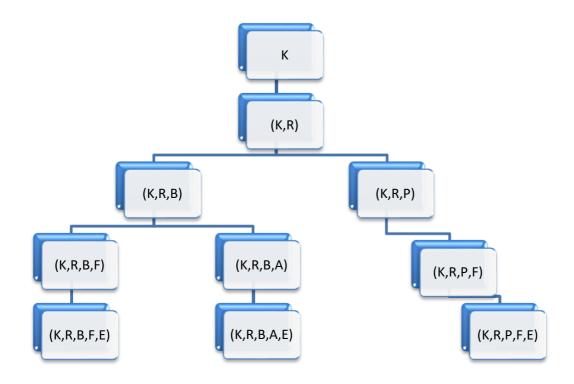
Stack

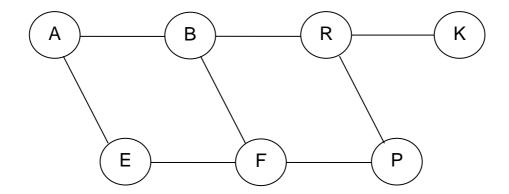
E F

B R K

Result of PATH 3 - KRBFE

DFS Tree





Breadth First Search (BFS)

A Queue is used. ENQUEUE and DEQUEUE are the operations used.

Consider the RED colored element in the QUEUE as DEQUEUE element and the one in GREEN to be the current node.

Step 1:		
Result : K	Queue :	K
Step 2:		
Result: KR	Queue :	K R
Step 3:		
Result: KRB	Queue :	K
		R B
Step 4:		
Result: KRBP	Queue:	K

Step 5:

Result: KRBPA Queue: K

R B P A

R B

Step 6: Result: KRBPAF	Queue :	K R B P A
Step 7:		
Result: KRBPAF	Queue :	K R B P A F
Step 8:		
Result: KRBPAF	Queue :	K R B P A
Step 9:		
Result: KRBPAFE	Queue :	K R B P A
Step 10 :		
		K
		R
		В

Α

Result: KRBPAFE Queue: F

Final BFS - KRBPAFE

BFS Tree

