GLEX-Alltoall: gather-scatter-based multi-leader All-to-all Communication on Multi-core Supercomputer

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Abstract

All-to-all communication is commonly used in parallel applications like FFT. In mordern supercomputers, there are multiple cores, NUMAs and network endpoints. These features bring much parallelism. However, there is no method which makes uses the parallelism to improve the all-to-all communication. In this paper, we introduce an optimized NUMA-aware multi-leader all-to-all library which explore the parallelism on network, CPU cores and overlap the intra- and inter-node communication. The results show that, compared to MPI, our library achieves up to 20x speedup. For application, our method achieves up to 1.75x speedup on peak performance for 16384 cores.

Keywords: Collective Communication, Multi-core processor, MPI all-to-all, RDMA, Shared Heap
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1. Introduction

Many parallel applications may suffer from global communication. Especially for communication-intensive applications, their time-to-solution and scalablily may be affected by global communication. Message Passing Interface (MPI) provides a set of commonly used collective communication. MPI_Alltoall is one of the collective communication where each process will send a different message to all processes. It is broadly used in some parallel applications like Fast Fourier Transform (FFT) [1] and some graph algorithms like MapReduce [2] and Breadth-first search (BFS) [3]. However, each time we double the processes, the all-to-all communication workload is quadrupled. On mordern supercomputers, network throughput has a linear relationship with the number of nodes. This brings great challenges to large-scale all-to-all communications.

For multiple-core processes, an effective way is node-aware all-to-all method [4]. It's replace a N nodes global all-to-all into N-1 times intra-node gather + local transpose + inter-node transpose + N-1 times intra-node gather. This method is very effective for small messages. Because, compared to original method, a node-aware all-to-all reduce the number of inter-node messages from $(M^N)^2$ to N^2 times (M is number of processers in each nodes). The size of the message is increased by M^2 times, which makes effective use of the network bandwidth. In the current supercomputer, a node has multiple CPU cores, NUMA and network endpoints. This architecture brings 4 kinds of parallism to optimize a node-aware all-to-all method:

(1) Multiple network endpoints can simultaneously process multiple communication requests.
(2) Processes in different NUMA can simultaneously access it local memory without contention.
(3) Multiple processes can simultaneously gather/scatter data and compose communication requests.
(4) Inter-node communication can be overlapped with intra-node communication.

As we known, no methods combine these parallism togather to improve a node-aware all-to-all collective communication.

In this paper, we proposed a multi-leader node-aware all-to-all method. It using multiple leaders on different NUMA which open the different network endpoints to gather/scatter data, compose communication requests, and transpose local matrix. It explore the parallelism existing in mordern multi-core processor with NUMA memory architecture and multi-port network. For intra-node gather/sactter, we proposed a shared-heap-based remote accessible memory which similar to intra-node MPI RMA. Inter-node communication is based on Remote Direct Memory Access (RDMA) which provide high throughput and low latency. The results show that, compared to MPI_alltoall, our implementation achieves up to 20x speedup and 4x speedup on average.

2. Related Work

From an algorithm perspective: Bruck algorithm [5] is commonly used for small message all-to-all. For mid size messages, isend-irecv algorithm is used. For large messages, linear shift exchange [6], pairwise exchange[7].

When considering the multi-core processors: Cache-oblivious MPI all-to-all (SH-NUMA-CO) based on morton order is proposed to minimize the cache miss rate [8]. For Infiniband and

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Table 1: Overall design-space for all-to-all collective on mulit-core processers

Approach/Metric	SH-NUMA-CO	SA-orig	elan_alltoall	PAIRWISE-SLR and BRUCK-SLR	GLEX-Alltoall (ours)
inter-node implementation	MVAPICH2	MVAPICH	RDMA	MPI-P2P	RDMA
intra-node implementation	Shared Heap	Shared Memory	Shared Memory	MPI-P2P	Shared Heap
Leader-based aggregation support	✔	✔	✔	✔	✔
Multi-leader support	✖	✔	✖	✖	✔
NUMA-aware	✔	✖	✖	✖	✔
Multiple network endpoints	✖	✔	✔	✖	✔
Overlapping Inter/intra-node	✔	✖	✖	✔	✔

multi-core systems, a non-leader all-to-all collective (SA-orig) which based on shared memory aggregation techniques is proposed in [9]. For multi-rail QsNet SMP clusters, a shared memory and RDMA based all-to-all collectives (elan_alltoall) is proposed in [10]. For Intel Many Integrated Core (MIC) architecture, the re-routing scheme based all-to-all collective (PAIRWISE-SLR/BRUCK-SLR) is proposed in [11]. These works are direct related to our work. Table 1 shows the overall design-space comparation between these methods and our method.

Besides, to improve the intra-node communication, several kernel-assistant techniques have been proposed for multi-core processor (LIMIC [12], KNEM [13], XPMEM). Compared to shared memory, these method works well for large message intra-node communication. Because they avoid one memory copy overhead. Shared heap [14] has the same performance with these kernel-assistant techniques. Shared heap do not need extra kernel module.

When considering the network topology: A bandwidth-optimal all-to-all exchange is proposed for fat-tree network [15]. For torus network, a large scale all-to-all is proposed for Blue Gene/L Supercomputer [16]. A optimal schedule for all-to-all personalized communication is proposed for multiprocessor systems [17]. For Infiniband clusters, their is a topology aware all-to-all scheduler which lower the contention by adding extra communication steps [18].

3. Experimental Platforms

The methods in this paper are validated on three different mechine.

(1) HPC-A: Matrix-2000+ CPU, Tianhe series network.
(2) HPC-B: FT-2000+ CPU, Tianhe series network.
(3) HPC-C: Intel E5 CPU, Tianhe series network.
(4) HPC-D: E5-2692 v2 * 2, Tianhe series network.

The Stream bandwidth test results are shown in the Table 2.

Table 2: Memory Bandwidth tested by Stream

Memory Bandwidth tests (MB/s)	Copy (Single-thread)	Copy (Multi-threads)
HPC-A (32 threads)		
HPC-B (32 threads)		
HPC-C (24 threads)		
HPC-D (24 threads)	11280.5	46252.3

4. Motivation

Node-aware all-to-all is a traditional method for optimize small message all-to-all. It include four steps: intra-node gather, local transpose, inter-node all-to-all, intra-node scatter. With the increase in the number of processor cores, multiple CPU cores, NUMA, network endpoints and inter-/intra-node overlapping brings huge parallelism. However, we notice that all traditional method do not take advantage of this parallelism optimize the four steps. For a 32/64 cores CPU, each inter-node message in node-aware method will be 1024/4096 times larger than direct method. If there is one double data between each pair of processes. There are 1024/4096 double data between each pair of nodes. As the single-core memory access bandwidth may not faster than the network bandwidth. Using a single core to gather/scatter/transpose data and initialize communication request may causing communication hotspot happend on a CPU core.

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