

# GLEX-Alltoall: gather-scatter-based multi-leader All-to-all Communication on Multi-core Supercomputer

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## Abstract

All-to-all communication is commonly used in parallel applications like FFT. In modern supercomputers, there are multiple cores, NUMAs and network endpoints. These features bring much parallelism. However, there is no method which makes use of the parallelism to improve the all-to-all communication. In this paper, we introduce an optimized NUMA-aware multi-leader all-to-all library which explores the parallelism on network, CPU cores and overlaps the intra- and inter-node communication. The results show that, compared to MPI, our library achieves up to 20x speedup. For application, our method achieves up to 1.75x speedup on peak performance for 16384 cores.

**Keywords:** Collective Communication, Multi-core processor, MPI all-to-all, RDMA, Shared Heap

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## 1. Introduction

Many parallel applications may suffer from global communication. Especially for communication-intensive applications, their time-to-solution and scalability may be affected by global communication. Message Passing Interface (MPI) provides a set of commonly used collective communication. MPI.Alltoall is one of the collective communication where each process will send a different message to all processes. It is broadly used in some parallel applications like Fast Fourier Transform (FFT) [1] and some graph algorithms like MapReduce [2] and Breadth-first search (BFS) [3]. However, each time we double the processes, the all-to-all communication workload is quadrupled. On modern supercomputers, network throughput has a linear relationship with the number of nodes. This brings great challenges to large-scale all-to-all communications.

For multiple-core processes, an effective way is node-aware all-to-all method [4]. It replaces a  $N$  nodes global all-to-all into  $N-1$  times intra-node gather + local transpose + inter-node transpose +  $N-1$  times intra-node gather. This method is very effective for small messages. Because, compared to original method, a node-aware all-to-all reduces the number of inter-node messages from  $(M^N)^2$  to  $N^2$  times ( $M$  is number of processors in each node). The size of the message is increased by  $M^2$  times, which makes effective use of the network bandwidth. In the current supercomputer, a node has multiple CPU cores, NUMA and network endpoints. This architecture brings 4 kinds of parallelism to optimize a node-aware all-to-all method:

- (1) Multiple network endpoints can simultaneously process multiple communication requests.
- (2) Processes in different NUMA can simultaneously access its local memory without contention.
- (3) Multiple processes can simultaneously gather/scatter data and compose communication requests.
- (4) Inter-node communication can be overlapped with intra-node communication.

As we known, no methods combine these parallelism together to improve a node-aware all-to-all collective communication.

In this paper, we proposed a multi-leader node-aware all-to-all method. It uses multiple leaders on different NUMA which open the different network endpoints to gather/scatter data, compose communication requests, and transpose local matrix. It explores the parallelism existing in modern multi-core processor with NUMA memory architecture and multi-port network. For intra-node gather/scatter, we proposed a shared-heap-based remote accessible memory which is similar to intra-node MPI RMA. Inter-node communication is based on Remote Direct Memory Access (RDMA) which provides high throughput and low latency. The results show that, compared to MPI.alltoall, our implementation achieves up to 20x speedup and 4x speedup on average.

## 2. Related Work

From an algorithm perspective: Bruck algorithm [5] is commonly used for small message all-to-all. For mid size messages, isend-irecv algorithm is used. For large messages, linear shift exchange [6], pairwise exchange [7].

When considering the multi-core processors: Cache-oblivious MPI all-to-all (SH-NUMA-CO) based on morton order is proposed to minimize the cache miss rate [8]. For Infiniband and

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multi-core systems, a all-to-all collective (SA-orig) which based  
on shared memory aggregation techniques is proposed in [9].  
For multi-rail QsNet SMP clusters, a shared memory and RDMA  
based all-to-all collectives (elan\_alltoall) is proposed in [10].  
For Intel Many Integrated Core (MIC) architecture, the re-routing  
scheme based all-to-all collective (PAIRWISE-SLR/BRUCK-  
SLR) is proposed in [11]. These works are direct related to  
our work. Table 1 shows the overall design-space for all-to-all  
collective on mult-core processors.

When considering the network topology: A bandwidth-optimal  
all-to-all exchange is proposed for fat-tree network [12]. For  
torus network, a large scale all-to-all is proposed for Blue Gene/L  
Supercomputer [13]. A optimal schedule for all-to-all person-  
alized communication is proposed for multiprocessor systems  
[14]. For Infiniband clusters, their is a topology aware all-to-all  
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